

Assignment 2

Due Date: September 22, 2025

1 - Lagrange Interpolation and Error Analysis

Let $f \in C^{n+1}[a, b]$, and let P_n be its Lagrange interpolating polynomial at the distinct nodes $x_0, x_1, \dots, x_n \in [a, b]$. The interpolation error is:

$$R_n(x) = f(x) - P_n(x) = \frac{f^{(n+1)}(\xi)}{(n+1)!} \prod_{i=0}^n (x - x_i), \quad \xi \in (a, b)$$

Suppose $f(x) = e^x$ on the interval $[0, 1]$, for equally spaced nodes $x_i = \frac{i}{n}$.

Question 1.1

Derive an explicit bound for the maximum error $\max_{x \in [0,1]} |R_n(x)|$.

Proof. To bound $|R_n(x)|$, we start with the error formula:

$$|R_n(x)| = \left| \frac{f^{(n+1)}(\xi)}{(n+1)!} \prod_{i=0}^n (x - x_i) \right|, \quad \xi \in (0, 1)$$

Note $\frac{1}{(n-1)!}$ is a constant. We can bound $R_n(x)$ by bounding $f^{(n+1)}(\xi)$ and $|\prod_{i=0}^n (x - x_i)|$ separately.

Now notice that $f^{(n+1)}(x) = f(x) = e^x$ for all n . Since e^x is strictly increasing on $[0, 1]$, we have that

$$f^{(n+1)}(\xi) = \max_{x \in [0,1]} |f^{(n+1)}(x)| = f(1) = e \quad (1)$$

Next, we need to bound $|\prod_{i=0}^n (x - x_i)|$. Let $\omega_n(x) = \prod_{i=0}^n (x - x_i)$. Note that $\omega_n(x)$ is a polynomial of degree $n+1$. Since $\omega_n(x)$ is continuous on the compact interval $[0, 1]$, it attains its maximum at some point in $[0, 1]$. We may write that:

$$|\omega_n(x)| \leq W_n \quad (2)$$

Where $W_n = \max_{x \in [0,1]} |\omega_n(x)| = |\omega_n(x^*)|$ for some $x^* \in [0, 1]$.

Now, we can use (1) and (2) to bound $|R_n(x)|$ from above:

$$|R_n(x)| \leq \frac{e}{(n+1)!} W_n$$

Since $R_n(x)$ is continuous on the compact interval $[0, 1]$, $\sup_{x \in [0,1]} R_n(x) = \max_{x \in [0,1]} R_n(x)$.

Since the supremum is the least upper bound, we have that:

$$\max_{x \in [0,1]} |R_n(x)| \leq \frac{e}{(n+1)!} W_n$$

Which gives us an explicit bound for the remainder. □

Question 1.2

Show the asymptotic decay of this error as $n \rightarrow \infty$.

Proof. From the previous part, we have that:

$$\max_{x \in [0,1]} |R_n(x)| \leq \frac{e}{(n+1)!} W_n$$

Let us first make a more crude estimate of W_n . Note that for any $x \in [0, 1]$ and for each i , we have:

$$|x - x_i| \leq 1$$

Since there are $n + 1$ terms in the product, we can write:

$$|\omega_n(x)| = \left| \prod_{i=0}^n (x - x_i) \right| \leq 1^{n+1} = 1$$

Therefore, we get the crude bound:

$$W_n = \max_{x \in [0,1]} |\omega_n(x)| \leq 1$$

Plugging this into our error bound, we get:

$$\max_{x \in [0,1]} |R_n(x)| \leq \frac{e}{(n+1)!}$$

Taking a limit as $n \rightarrow \infty$, we have:

$$\lim_{n \rightarrow \infty} \max_{x \in [0,1]} |R_n(x)| \leq \lim_{n \rightarrow \infty} \frac{e}{(n+1)!} = 0$$

Therefore, we have shown that the error decays to 0 as $n \rightarrow \infty$. □

2 - Interpolation Programming Exercise

Consider the function:

$$f(x) = \frac{1}{1 + 20x}, \quad x \in [-1, 1]$$

Question 2.1

Construct the interpolation polynomial $P_n(x)$ at equidistant nodes for $n = 5$.

Solution. The following methods were used to compute this polynomial:

Listing 1: 2.1 Python

```

1 import numpy as np
2
3 def f(x):
4     return 1 / (1 + 20 * x**2)
5
6 def lagrange_coefficients(nodes):
7     x = nodes
8     num_nodes = len(nodes)
9     # Add zeroth divided differences
10    dd_table = np.array([[f(xi) for xi in nodes]])
11
12    # Calculate divided difference table
13    for i in range(1, num_nodes):
14        ith_dd = np.zeros(num_nodes)
15
16        for j in range(num_nodes - i):
17            # Calculate ith divided differences
18            ith_dd[j] = (dd_table[i - 1, j + 1] - dd_table[i - 1, j]) / (
19                x[j + i] - x[j]
20            )
21
22        # Append the ith divided difference row to the table
23        dd_table = np.vstack([dd_table, ith_dd])
24
25    # Extract coefficients (first column of the table)
26    a = np.array([dd_table[i, 0] for i in range(dd_table.shape[0])])
27    return a
28
29 def generate_lagrange(nodes, degree):
30     # Get coefficients
31     a = lagrange_coefficients(nodes)
32     # Start with function and constant term
33     equation = f"P_{{{degree}}}(x) = {a[0]}"
34     w = []
35     # Build (x - xi) terms

```

```

36     for xi in nodes:
37         if abs(xi) <= 1e-14:
38             w.append(f'({xi})')
39         elif xi < 0:
40             w.append(f"(x + {abs(xi)})")
41         else:
42             w.append(f"(x - {xi})")
43     b = ""
44
45     # Build polynomial string
46     for i in range(1, len(a)):
47         for j in range(i):
48             # Multiply (x - xi) terms
49             b += w[j]
50             # Multiply (x - xi) product with current coefficient and add term
51             if a[i] >= 0:
52                 equation += f" + {a[i]}{b}"
53             else:
54                 equation += f" - {abs(a[i])}{b}"
55         b = ""
56
57     return equation
58
59 # Sample Execution
60 if __name__ == "__main__":
61     # Generate 6 nodes for 5th degree polynomial
62     nodes = np.linspace(-1, 1, 6)
63     equation = generate_equation(nodes, 5)
64     with open("./plots_2/q2_2/lagrange_equations.txt", "w") as file:
65         file.write(equation + "\n")

```

The interpolation polynomial $P_5(x)$ at equidistant nodes for $n = 5$ is given by:

$$\begin{aligned}
 P_5(x) = & 0.047619047619047616 \\
 & + 0.18583042973286878(x + 1.0) \\
 & + 1.1227255129694158(x + 1.0)(x + 0.6) \\
 & - 2.0647825525874324(x + 1.0)(x + 0.6)(x + 0.19999999999999996) \\
 & + 1.2904890953671473(x + 1.0)(x + 0.6)(x + 0.19999999999999996) \\
 & (x - 0.200000000000000018) \\
 & - 3.552713678800501e - 15(x + 1.0)(x + 0.6)(x + 0.19999999999999996) \\
 & (x - 0.200000000000000018)(x - 0.60000000000000001)
 \end{aligned}$$

Question 2.2

Plot $f(x)$ and $P_n(x)$, and report the behavior of the interpolation error as n increases from 5 to 10 and 20.

Solution. In addition to the functions defined in the previous part, the following functions were used to help plot $f(x)$ and $P_n(x)$:

Listing 2: 2.2 Python

```

1  import numpy as np
2  import matplotlib.pyplot as plt
3  # I created a Figure class for figure management.
4  # See figure.py at end of this pdf for implementation.
5  from figure import Figure
6
7  def calculate_lagrange(nodes, a, x):
8      # Start with constant term
9      y = a[0]
10     # Build (x - xi) terms
11     w = [(x - xi) for xi in nodes]
12     # Temporary variable to hold (x - xi) product
13     b = 1
14
15     for i in range(1, len(a)):
16         for j in range(i):
17             # Multiply (x - xi) terms
18             b *= w[j]
19         # Multiply (x - xi) product with current coefficient
20         y += a[i] * b
21         b = 1
22
23     return y
24
25 def calculate_lagrange_output(nodes, x_coords=[]):
26     # Get coefficients
27     a = lagrange_coefficients(nodes)
28     # If no x_coords provided, use nodes as x_coords
29     if len(x_coords) == 0:
30         x_coords = nodes
31     # Calculate y coordinates for each x coordinate
32     y_coordinates = np.array([calculate_lagrange(nodes, a, x) for x in x_coords])
33
34     return y_coordinates
35
36 # Sample Execution
37 if __name__ == "__main__":

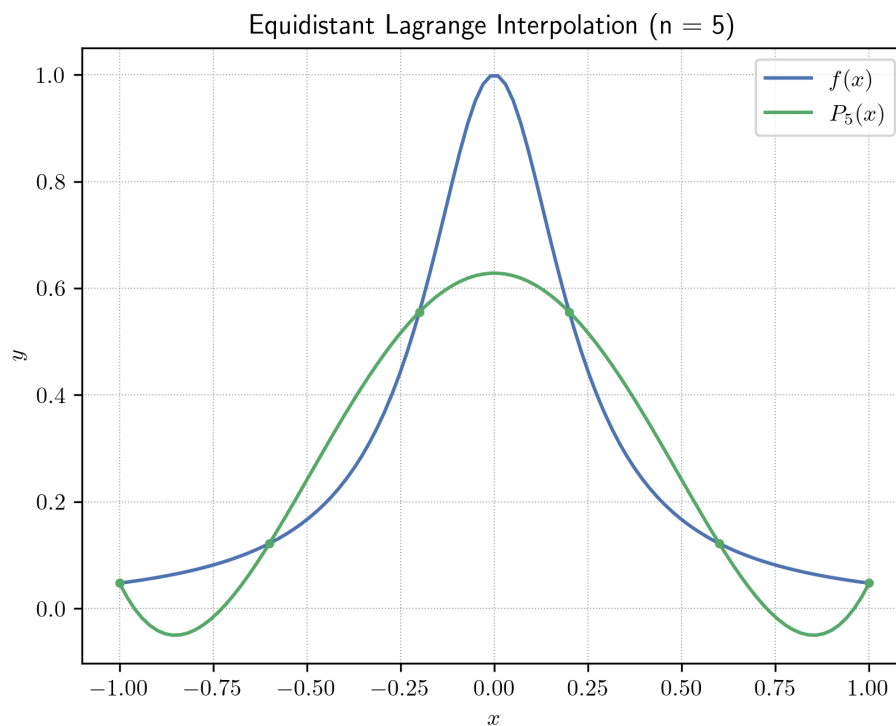
```

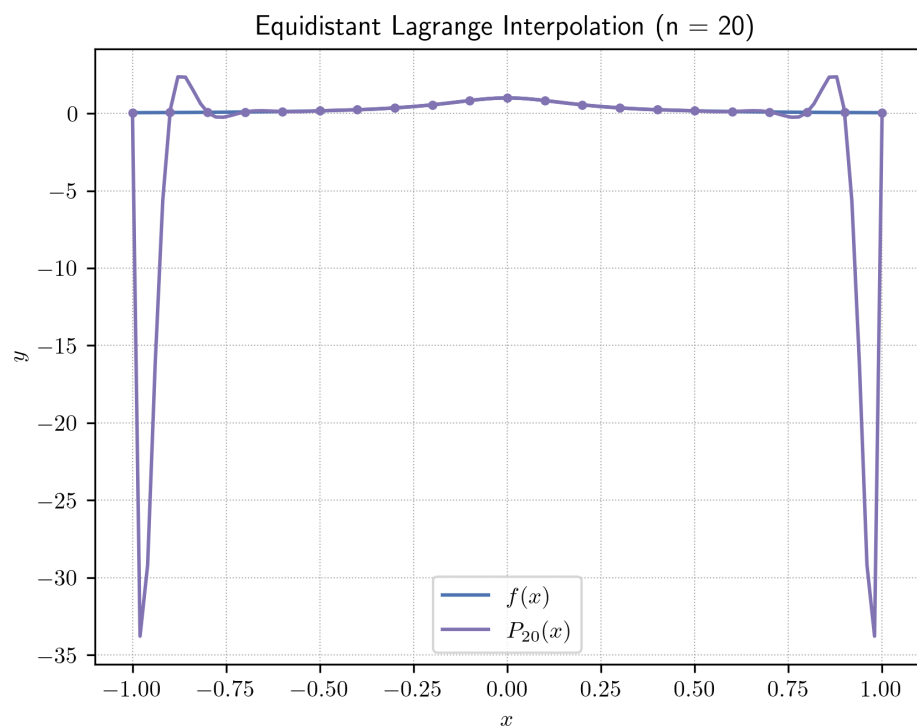
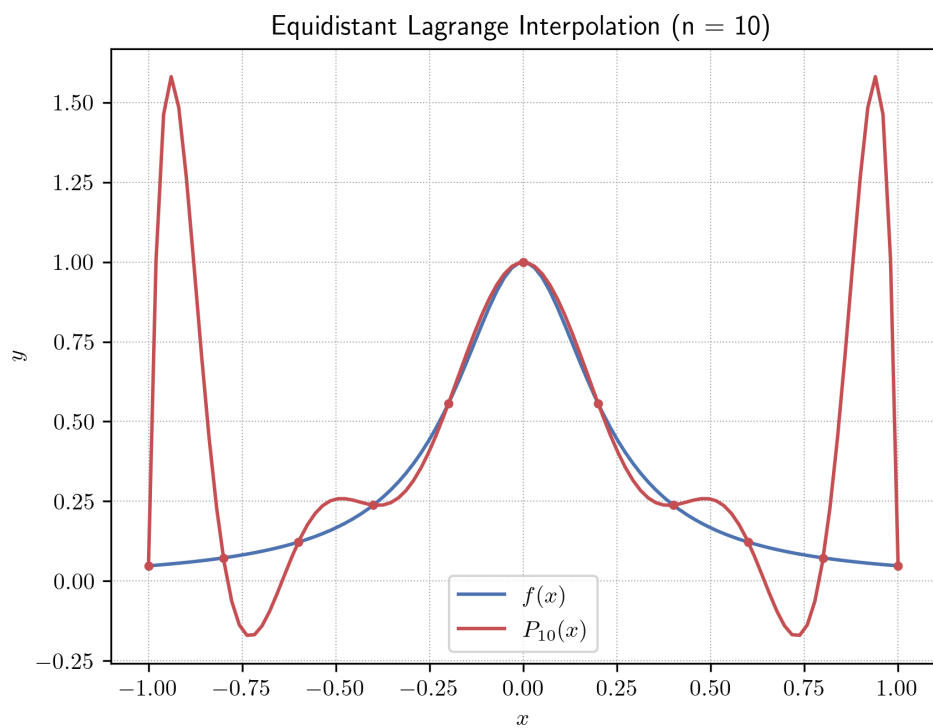
```

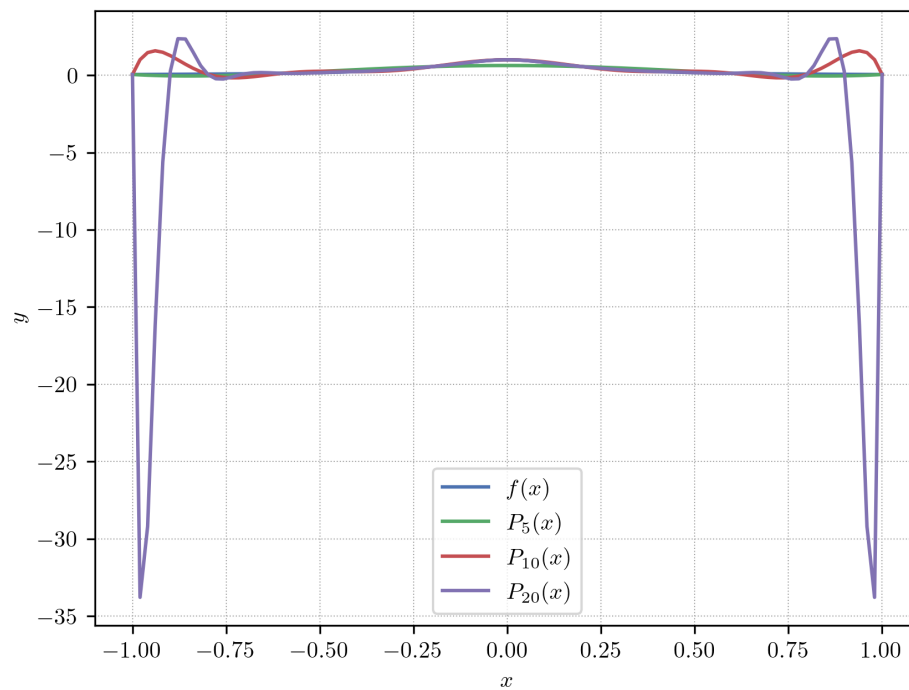
38 fig_f.title = "Equidistant Lagrange Interpolation"
39 x = np.linspace(-1, 1, 100)
40 f_x = f(x)
41 nodes_x = np.linspace(-1, 1, 6)
42 nodes_y = calculate_lagrange_output(nodes_x)
43 p_5 = calculate_lagrange_output(nodes_x, x)
44
45 # Figure attributes:
46 # x_sets, y_sets, labels, colors(indexes from default rotation list),
47 # markers, linestyle, title, xlabel, ylabel
48 # single or multiple axes can be added during initialization
49 fig_f = Figure(x, f_x, r"$f(x)$")
50 fig_p5 = Figure(x, p_5, r"$P_5(x)$", 1)
51 fig_nodes = Figure(nodes_x, nodes_y, "", 1, ".", "")
52
53 # Returns a new Figure that merges fig_f, fig_p5, and fig_nodes.
54 # The title of the new figure is title + " (n = 5)"
55 fig = fig_f.copy().merge([fig_p5, fig_nodes], title + " (n = 5)")
56
57 # Returns the matplotlib figure and saves png with 300 dpi
58 fig.get_figure("./plots_2/q2_2/p5.png")
59 plt.show()

```

The following plots show $f(x)$ and $P_n(x)$ for $n = 5, 10, 20$:







As n increases from 5 to 10 and 20, we see that the interpolation error decreases towards the center of the interval $[-1, 1]$. However, the error increases significantly near the edges of the interval, demonstrating Runge's phenomenon.

Question 2.3

Repeat the interpolation using Chebyshev nodes, and compare the results with the equidistant-node case.

Solution. In addition to the functions defined in the previous parts, the following function was used to generate the Chebyshev nodes:

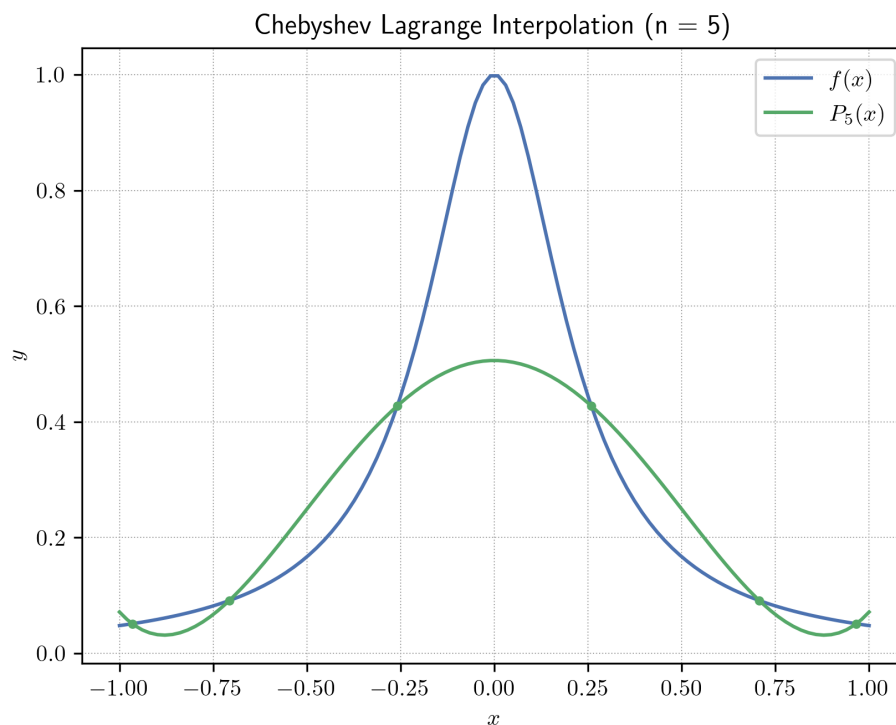
Listing 3: 2.3 Python

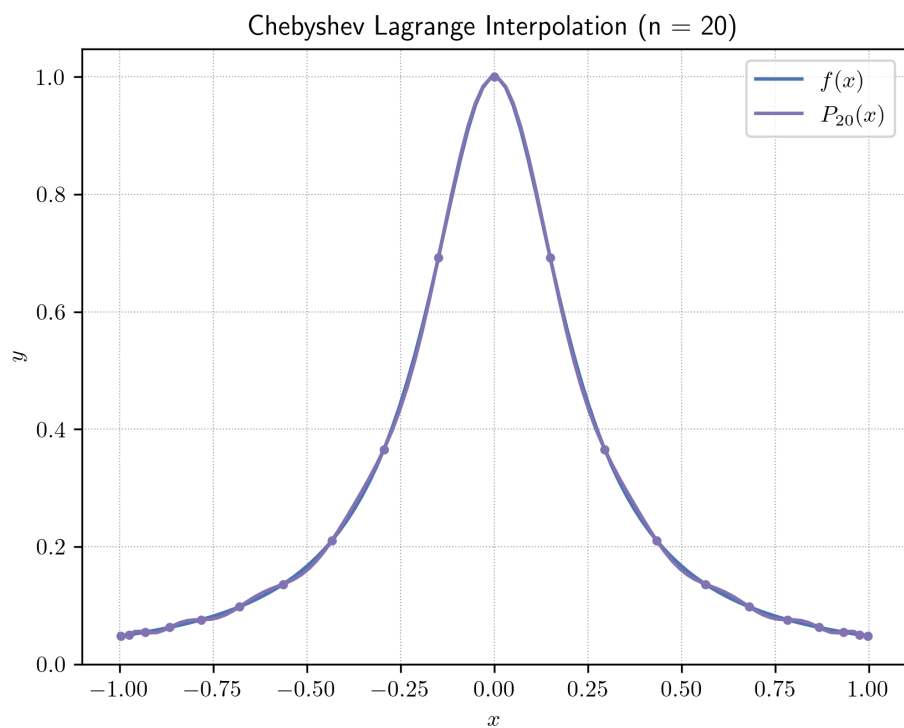
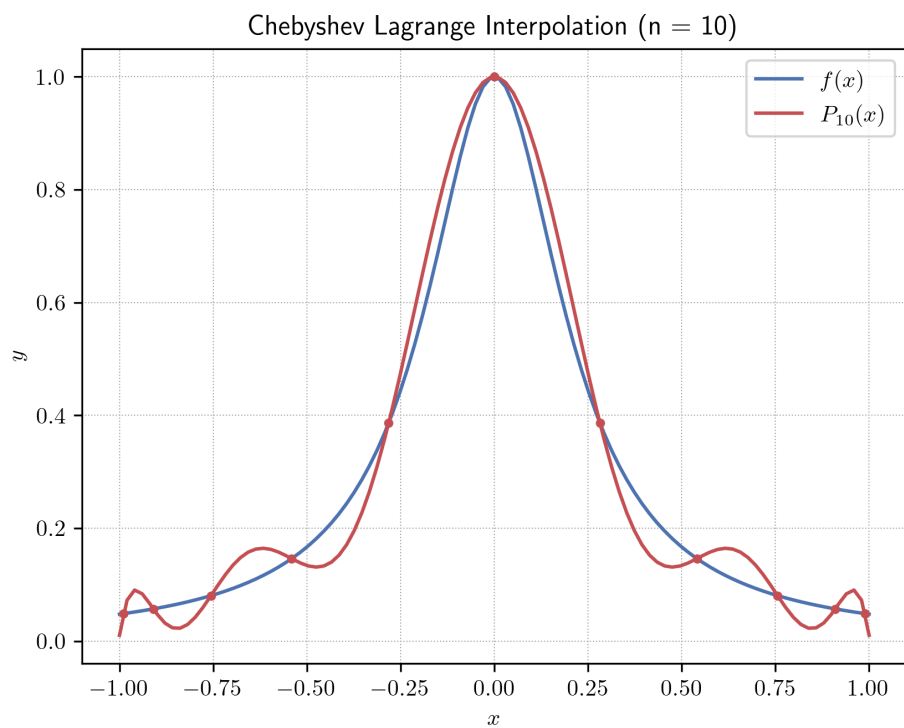
```

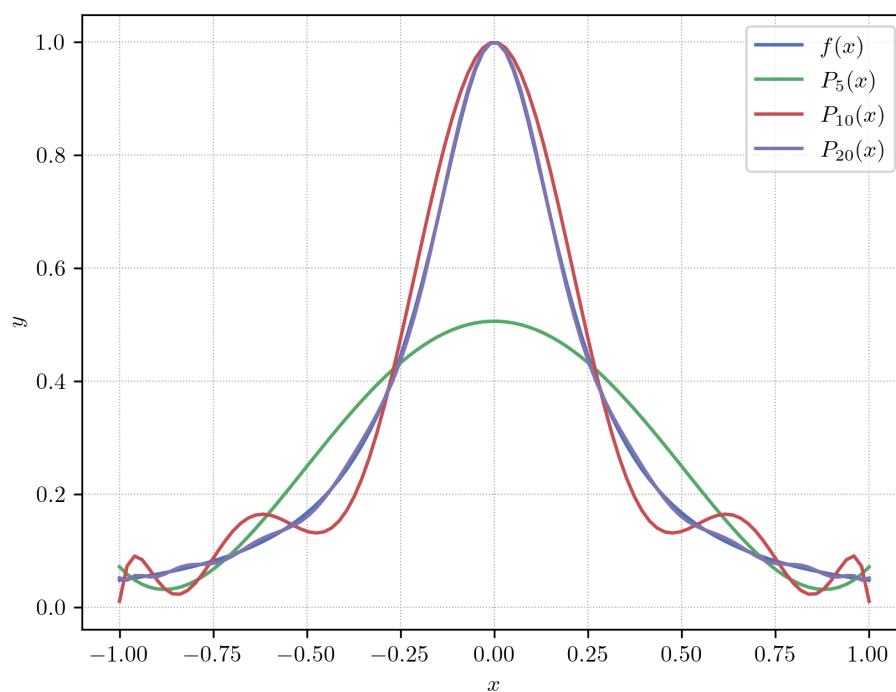
1 import numpy as np
2
3 def chebyshev_nodes(n):
4     nodes = np.array(
5         [(math.cos((2 * k - 1) * math.pi / (2 * n))) for k in range(1, n + 1)]
6     )
7
8     return nodes

```

The following plots show $f(x)$ and $P_n(x)$ for Chebyshev nodes with $n = 5, 10, 20$:







Comparing these plots to those with equidistant nodes, we see that using Chebyshev nodes significantly reduces the interpolation error across the entire interval $[-1, 1]$. The oscillations near the edges of the interval are much less pronounced, demonstrating that Chebyshev nodes help mitigate Runge's phenomenon and provide a more accurate approximation of $f(x)$.

3 - Chebyshev Polynomials and Their Roots

The Chebyshev polynomials of the first kind is defined by:

$$T_n(x) = \cos(n \arccos x), \quad x \in [-1, 1]$$

Question 3.1

Prove that the numbers

$$x_k = \cos\left(\frac{2k-1}{2n}\pi\right), \quad k = 1, 2, \dots, n$$

are the roots of $T_n(x)$.

Proof. We want to find the values of x such that $T_n(x) = 0$. By the definition of Chebyshev polynomials, we have:

$$T_n(x) = \cos(n \arccos x)$$

Setting $T_n(x) = 0$, we get:

$$T_n(x) = \cos(n \arccos x) = 0$$

The cosine function is zero at odd multiples of $\frac{\pi}{2}$. So $T_n(x)$ can only equal zero when $n \arccos x$ is an odd multiple of $\frac{\pi}{2}$. Therefore, we can write:

$$n \arccos x = \frac{(2k-1)\pi}{2}, \quad k \in \mathbb{Z}$$

Note that for $n = 0$, $T_0(x) = \cos(0 \cdot \arccos(x)) = 1$ which has no roots. So consider $n \geq 1$. Dividing both sides by n , we have:

$$\arccos x = \frac{(2k-1)\pi}{2n}, \quad k \in \mathbb{Z}$$

Taking the cosine of both sides, we get:

$$x = \cos\left(\frac{(2k-1)\pi}{2n}\right), \quad k \in \mathbb{Z}$$

Therefore, the set of x defined above are the roots of $T_n(x)$ for $n \geq 1$. Note that if we restrict k to the integers $1, 2, \dots, n$, we exactly get that:

$$x_k = \cos\left(\frac{2k-1}{2n}\pi\right), \quad k = 1, 2, \dots, n$$

are the roots of $T_n(x)$. □

Question 3.2

Show that these roots are distinct and lie in the interval $(-1, 1)$.

Proof. To show that these roots are distinct, we need to show that for $k_1 \neq k_2$, we have $x_{k_1} \neq x_{k_2}$. Assume $k_1, k_2 \in \{1, 2, \dots, n\}$ such that $k_1 \neq k_2$ and $k_1 < k_2$. Then, assuming for fixed $n \geq 1$:

$$x_{k_1} = \cos\left(\frac{2k_1 - 1}{2n}\pi\right), \quad x_{k_2} = \cos\left(\frac{2k_2 - 1}{2n}\pi\right)$$

Note that $\frac{2k_1 - 1}{2n}\pi$ and $\frac{2k_2 - 1}{2n}\pi$ are distinct angles in the interval $(0, \pi)$. Since the cosine function is strictly decreasing on the interval $[0, \pi]$, and since $\frac{2k_1 - 1}{2n}\pi < \frac{2k_2 - 1}{2n}\pi$ are distinct, we have:

$$x_{k_1} = \cos\left(\frac{2k_1 - 1}{2n}\pi\right) > \cos\left(\frac{2k_2 - 1}{2n}\pi\right) = x_{k_2}$$

By the law of trichotomy, $x_{k_1} \neq x_{k_2}$. Therefore, the roots x_k are distinct for $k = 1, 2, \dots, n$.

Next, we need to show that these roots lie in the interval $(-1, 1)$. Note that from above, we have shown that x_k is a strictly decreasing sequence. Therefore, we only need to show that the largest root x_1 is less than 1 and the smallest root x_n is greater than -1 .

First, let us show that $x_1 < 1$:

$$x_1 = \cos\left(\frac{2(1) - 1}{2n}\pi\right) = \cos\left(\frac{\pi}{2n}\right)$$

Note that $\frac{\pi}{2n} \in (0, \frac{\pi}{2}]$. In this interval, the range of the cosine function is $[0, 1)$. Therefore, we have:

$$x_1 = \cos\left(\frac{\pi}{2n}\right) < 1 \tag{1}$$

Next, let us show that $x_n > -1$:

$$x_n = \cos\left(\frac{2n - 1}{2n}\pi\right) = \cos\left(\pi - \frac{\pi}{2n}\right)$$

Note that $\pi - \frac{\pi}{2n} \in [\frac{\pi}{2}, \pi)$. In this interval, the range of the cosine function is $(-1, 0]$. Therefore, we have:

$$x_n = \cos\left(\pi - \frac{\pi}{2n}\right) > -1 \tag{2}$$

Combining (1) and (2), we have shown that:

$$-1 < x_n < x_k < x_1 < 1$$

Therefore, x_k are distinct roots that lie in the interval $(-1, 1)$ for $k = 1, 2, \dots, n$. □

Question 3.3

Write your own code to generate $T_n(x)$ using the recursion formula.

Solution. The following code was used to generate the Chebyshev polynomials::

Listing 4: 3.3 Python

```

1  import sympy as smp
2
3  def chebyshev_poly(n):
4      # Base Cases
5      if n == 0:
6          return 1
7      elif n == 1:
8          return smp.symbols("x")
9      # Recursive Case
10     else:
11         return 2 * smp.symbols('x') * chebyshev_poly(n - 1) - chebyshev_poly(n
- 2)
12
13 def generate_chebyshev(n):
14     x = smp.symbols('x')
15     poly = smp.expand(chebyshev_poly(n))
16     function = f"T_{{{n}}}({x})"
17     equation = function + " = " + smp.latex(poly)
18     return equation
19
20 # Sample Execution
21 if __name__ == "__main__":
22     with open("./plots_2/q3_3/chebyshev_polynomials.txt", "w") as file:
23         for n in range(6):
24             equation = generate_chebyshev(n)
25             file.write(equation + "\n")

```

The Chebyshev polynomials $T_n(x)$ for $n = 0, 1, \dots, 5$ are given by:

$$T_0(x) = 1$$

$$T_1(x) = x$$

$$T_2(x) = 2x^2 - 1$$

$$T_3(x) = 4x^3 - 3x$$

$$T_4(x) = 8x^4 - 8x^2 + 1$$

$$T_5(x) = 16x^5 - 20x^3 + 5x$$

4 - Lagrange Interpolation for Nonsmooth Function

Let $f(x) = |x|$ on the interval $[-1, 1]$.

Question 4.1

Construct the interpolation polynomial $P_n(x)$ for equidistant nodes when n is even.

Question 4.2

Show that $P_n(x)$ is an even polynomial.

Question 4.3

Investigate analytically (for small n) how well $P_n(x)$ approximates $f(x)$.

Question 4.4

Discuss why convergence may be slower for nonsmooth functions.