

# 1 Introduction

A *graph*  $G$  consists of a set of *vertices*  $V(G)$  and a set of *edges*  $E(G)$

Each edge is associated with one or two vertices, its *ends*

An edge is a *loop* if the ends are the same, a *link* if the ends are different

Graphs are assumed to be *finite* - both  $V(G)$  and  $E(G)$  are finite sets

A *path* is a sequence  $(a_0, A_1, a_1, A_2, a_2, \dots, A_n, a_n)$  of vertices  $a_i$  and edges  $A_j$  such that:

1. If  $1 \leq i \leq n$  the ends of  $A_i$  are  $a_{i-1}$  and  $a_i$
2. If  $1 \leq i \leq n$  then  $a_{i-1} = a_i$  iff  $A_i$  is a loop

If all terms in path are distinct, then the path is *simple*

If all terms are distinct except that  $a_0 = a_n$ , then the path is *circular*

$x, y \in V(G)$  are *connected* if there is a path from  $x$  to  $y$  in  $G$

Connectedness is an equivalence relation on  $V(G)$  so that if  $V(G) \neq \emptyset$  then it can be partitioned into disjoint non-null subsets  $V_1, \dots, V_k$  such that two vertices of  $G$  are connected iff they are in the same  $V_i$

The  $G[V_i]$  are the *connected components* of  $G$ , they are edge- and vertex-disjoint and cover  $G$

The number of components of  $G$  is denoted  $p_0(G)$

A graph is *connected* iff  $p_0(G) = 0$  or  $1$ , the first case arising if  $G$  is the empty graph

A connected graph with no circular path is a tree

$$\alpha_0(G) = |G| \text{ and } \alpha_1(G) = e(G)$$

Let  $Q_n$  be a finite set of  $n > 0$  elements

$f : V(G) \rightarrow Q_n$  is an *n-colouring* of  $G$  if each edge  $xy$  of  $G$  has  $f(x) \neq f(y)$

The number of *n-colourings* of  $G$  wrt  $Q_n$  is denoted by  $P(G, n)$

If  $V(G) = \emptyset$  then we say  $P(G, n) = 1$ , also  $P(G, n) = 0$  if  $G$  contains a loop

When  $G$  is loopless,  $P(G, n)$  is a polynomial in  $n$  of degree  $|G|$

For planar graphs,  $P(G, n)$  is called the *chromatic polynomial* of  $G$

$$P(G, n) = \sum_S (-1)^{e(S)} n^{p_0(S)} \quad (\text{summing over spanning subgraphs } S \text{ of } G)$$

An *orientation* of  $G$  distinguishes one end of each edge  $A$  as positive,  $p(A)$ , and one as negative,  $q(A)$

If  $A$  is a loop, then  $p(A) = q(A)$  otherwise  $p(A) \neq q(A)$

If  $a \in V(G)$  and  $A \in E(G)$  then  $\eta(A, a) = 0$  if  $A$  is a loop or  $a$  is not an end of  $A$ .  
Otherwise,  $\eta(A, a) = 1$  or  $-1$  depending as whether  $a$  is the positive or negative end of  $A$

A mapping  $f$  of  $V(G)$  or  $E(G)$  into  $Q_n$  is a *0-chain* or *1-chain* respectively *on*  $G$  *over*  $Q_n$

If  $V(G) = \emptyset$  then there is just one 0-chain on  $G$  over  $Q_n$

If  $E(G) = \emptyset$  then there is just one 1-chain on  $G$  over  $Q_n$

If  $h$  is a 0-chain on  $G$  over  $Q_n$  its *coboundary*,  $\delta h$  is the 1-chain on  $G$  over  $Q_n$  satisfying

$$(\delta h)(A) = \sum_a \eta(A, a)h(a) \quad (2)$$

for each  $A \in E(G)$ , equivalently

$$(\delta h)(A) = h(p(A)) - h(q(A)) \quad (2a)$$

If  $g$  is a 1-chain, its *boundary*  $\delta g$  is the 0-chain satisfying

$$(\delta g)(a) = \sum_A \eta(A, a)g(A) \quad (3)$$

for each  $a \in V(G)$

We call  $g$  a *1-cycle* on  $G$  over  $Q_n$  if  $\delta g \equiv \mathbf{0}$

## 2 Colour-coboundaries and colour-cycles

A *colour-coboundary* or *colour-cycle* on  $G$  over  $Q_n$  is a 1-chain  $g$  on  $G$  over  $Q_n$  which is a coboundary or a 1-cycle respectively and which satisfies  $g(A) \neq 0$  for each  $A \in E(G)$

The number of colour-coboundaries of  $G$  over  $Q_n$  is denoted  $\theta(G, n)$

The number of colour cycles on  $G$  over  $Q_n$  is denoted  $\phi(G, n)$

$\theta(G, n)$  and  $\phi(G, n)$  are independent of orientation

If  $e(G) = 0$  then we say  $\theta(G, n) = \phi(G, n) = 1$

By (2a), the colour-coboundaries on  $G$  over  $Q_n$  are the coboundaries of the  $n$ -colourings of  $G$

Also,  $\delta h_1 = \delta h_2$  for 0-chains  $h_1, h_2$  iff  $h_1(a) - h_2(a)$  is constant in each component of  $G$ , for all  $A \in E(G)$ :

$$h_1(p(A)) - h_1(q(A)) = h_2(p(A)) - h_2(q(A)) \iff h_1(p(A)) - h_2(p(A)) = h_1(q(A)) - h_2(q(A))$$