# ARONDIGHT'S STANDARD CODE LIBRARY\*

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 $<sup>{\</sup>rm *https://www.github.com/footoredo/Arondight}$ 

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# 1 计算几何

```
1 int sign(DB x) {
     return (x > eps) - (x < -eps);
 4 DB msart(DB x) {
    return sign(x) > 0 ? sqrt(x) : 0;
   struct Point {
    DB x. v:
    Point rotate(DB ang) const { // 逆时针旋转 ang 弧度
       return Point(cos(ang) * x - sin(ang) * y,
11
           cos(ang) * y + sin(ang) * x);
13
14
     Point turn90() const { // 逆时针旋转 90 度
15
       return Point(-y, x);
16
    Point unit() const {
       return *this / len();
19
20 };
21 DB dot(const Point& a, const Point& b) {
     return a.x * b.x + a.y * b.y;
23 }
   DB det(const Point& a, const Point& b) {
     return a.x * b.y - a.y * b.x;
   #define cross(p1,p2,p3) ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
   #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
   |bool isLL(const Line& 11, const Line& 12, Point& p) { // 直线与直线交点
    DB s1 = det(12.b - 12.a, 11.a - 12.a),
31
        s2 = -det(12.b - 12.a, 11.b - 12.a);
    if (!sign(s1 + s2)) return false;
    p = (11.a * s2 + 11.b * s1) / (s1 + s2);
    return true:
35 }
   bool onSeg(const Line& 1, const Point& p) { // 点在线段上
     return sign(det(p - 1.a, 1.b - 1.a)) == 0 \&\& sign(dot(p - 1.a, p - 1.b)) <= 0;
38 }
39 Point projection(const Line & 1, const Point& p) {
     return 1.a + (1.b - 1.a) * (dot(p - 1.a, 1.b - 1.a) / (1.b - 1.a).len2());
41 }
42 DB disToLine(const Line& 1, const Point& p) { // 点到 * 直线 * 距离
     return fabs(det(p - 1.a, 1.b - 1.a) / (1.b - 1.a).len());
44 }
45 DB disToSeg(const Line& 1, const Point& p) { // 点到线段距离
    return sign(dot(p - 1.a, 1.b - 1.a)) * sign(dot(p - 1.b, 1.a - 1.b)) == 1?
     \rightarrow disToLine(1, p) : std::min((p - 1.a).len(), (p - 1.b).len());
47 }
48 // 圆与直线交点
49 bool isCL(Circle a, Line 1, Point& p1, Point& p2) {
50 DB x = dot(1.a - a.o. 1.b - 1.a).
       y = (1.b - 1.a).len2(),
```

```
d = x * x - y * ((1.a - a.o).len2() - a.r * a.r);
 52
     if (sign(d) < 0) return false;
      Point p = 1.a - ((1.b - 1.a) * (x / y)), delta = (1.b - 1.a) * (msqrt(d) / y);
     p1 = p + delta; p2 = p - delta;
     return true;
 57 }
 58 //圆与圆的交面积
 59 DB areaCC(const Circle& c1, const Circle& c2) {
     DB d = (c1.o - c2.o).len():
     if (sign(d - (c1.r + c2.r)) >= 0) return 0;
     if (sign(d - std::abs(c1.r - c2.r)) <= 0) {</pre>
       DB r = std::min(c1.r, c2.r);
64
       return r * r * PI:
 65
     DB x = (d * d + c1.r * c1.r - c2.r * c2.r) / (2 * d),
       t1 = acos(x / c1.r), t2 = acos((d - x) / c2.r);
     return c1.r * c1.r * t1 + c2.r * c2.r * t2 - d * c1.r * sin(t1):
 69 }
 70 // 圆与圆交点
 71 bool isCC(Circle a, Circle b, P& p1, P& p2) {
     DB s1 = (a.o - b.o).len():
     if (sign(s1 - a.r - b.r) > 0 \mid | sign(s1 - std::abs(a.r - b.r)) < 0) return false;
 74
     DB s2 = (a.r * a.r - b.r * b.r) / s1;
     DB aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
     P \circ = (b.o - a.o) * (aa / (aa + bb)) + a.o;
     P delta = (b.o - a.o).unit().turn90() * msgrt(a.r * a.r - aa * aa);
     p1 = o + delta, p2 = o - delta;
 79
     return true;
 80 | }
    // 求点到圆的切点,按关于点的顺时针方向返回两个点
    bool tanCP(const Circle &c, const Point &p0, Point &p1, Point &p2) {
     double x = (p0 - c.o).len2(), d = x - c.r * c.r;
     if (d < eps) return false; // 点在圆上认为没有切点
     Point p = (p0 - c.o) * (c.r * c.r / x);
     Point delta = ((p0 - c.o) * (-c.r * sqrt(d) / x)).turn90();
     p1 = c.o + p + delta;
     p2 = c.o + p - delta;
 89
     return true:
 90
    // 求圆到圆的外共切线, 按关于 c1.o 的顺时针方向返回两条线
    vector<Line> extanCC(const Circle &c1, const Circle &c2) {
     vector<Line> ret:
     if (sign(c1.r - c2.r) == 0) {
       Point dir = c2.o - c1.o;
 96
       dir = (dir * (c1.r / dir.len())).turn90();
       ret.push_back(Line(c1.o + dir, c2.o + dir));
 98
        ret.push back(Line(c1.o - dir, c2.o - dir));
 99
      } else {
100
        Point p = (c1.0 * -c2.r + c2.o * c1.r) / (c1.r - c2.r);
101
        Point p1, p2, q1, q2;
102
        if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) {
103
         if (c1.r < c2.r) swap(p1, p2), swap(q1, q2);
104
          ret.push back(Line(p1, q1));
105
         ret.push back(Line(p2, q2));
106
```

```
107
     return ret;
109 }
110 // 求圆到圆的内共切线,按关于 c1.o 的顺时针方向返回两条线
111 std::vector<Line> intanCC(const Circle &c1, const Circle &c2) {
      std::vector<Line> ret:
     Point p = (c1.0 * c2.r + c2.o * c1.r) / (c1.r + c2.r);
113
114
      Point p1, p2, q1, q2;
      if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) { // 两圆相切认为没有切线
115
       ret.push back(Line(p1, q1));
116
117
       ret.push_back(Line(p2, q2));
118
    }
119
     return ret:
120 }
    bool contain(vector<Point> polygon, Point p) { // 判断点 p 是否被多边形包含,包括落在边界上
      int ret = 0, n = polygon.size();
     for(int i = 0; i < n; ++ i) {
123
124
       Point u = polygon[i], v = polygon[(i + 1) % n];
       if (onSeg(Line(u, v), p)) return true; // Here I guess.
126
       if (sign(u.v - v.v) \le 0) swap(u, v);
127
       if (sign(p.y - u.y) > 0 \mid | sign(p.y - v.y) \le 0) continue;
       ret += sign(det(p, v, u)) > 0;
128
129
130
     return ret & 1;
131 | }
132 // 用半平面 (a1,a2) 的逆时针方向去切凸多边形
    std::vector<Point> convexCut(const std::vector<Point>&ps, Point q1, Point q2) {
134
      std::vector<Point> qs; int n = ps.size();
135
      for (int i = 0; i < n; ++i) {
       Point p1 = ps[i], p2 = ps[(i + 1) % n];
136
137
       int d1 = crossOp(q1,q2,p1), d2 = crossOp(q1,q2,p2);
138
       if (d1 \ge 0) qs.push back(p1);
139
       if (d1 * d2 < 0) qs.push back(isSS(p1, p2, q1, q2));
140
    }
141
     return qs;
142 }
143 // 求凸包
    std::vector<Point> convexHull(std::vector<Point> ps) {
     int n = ps.size(); if (n <= 1) return ps;</pre>
      std::sort(ps.begin(), ps.end());
146
147
      std::vector<Point> qs;
     for (int i = 0; i < n; qs.push back(ps[i ++]))</pre>
148
149
       while (qs.size() > 1 && sign(det(qs[qs.size() - 2], qs.back(), ps[i])) \le 0)
150
         qs.pop back();
151
     for (int i = n - 2, t = gs.size(); i \ge 0; gs.push back(ps[i --]))
       while ((int)qs.size() > t && sign(det(qs[qs.size() - 2], qs.back(), ps[i])) \le 0)
          qs.pop back();
     return qs;
```

```
1.1 凸包
1 // 凸包中的点按逆时针方向
2 struct Convex {
   int n:
   std::vector<Point> a, upper, lower;
   void make_shell(const std::vector<Point>& p,
```

```
std::vector<Point>& shell) { // p needs to be sorted.
7
       clear(shell); int n = p.size();
8
       for (int i = 0, j = 0; i < n; i++, j++) {
9
        for (; j \ge 2 \&\& sign(det(shell[j-1] - shell[j-2]),
10
                p[i] - shell[j-2])) \le 0; --j) shell.pop back();
11
         shell.push back(p[i]);
12
      }
13
14
     void make convex() {
15
       std::sort(a.begin(), a.end());
16
       make shell(a, lower);
17
       std::reverse(a.begin(), a.end());
18
       make shell(a, upper);
19
       a = lower; a.pop back();
20
       a.insert(a.end(), upper.begin(), upper.end());
21
       if ((int)a.size() >= 2) a.pop back();
22
       n = a.size():
23
24
     void init(const std::vector<Point>& a) {
25
       clear(a); a = a; n = a.size();
26
       make convex():
27
28
     void read(int _n) { // Won't make convex.
       clear(a); n = n; a.resize(n);
30
       for (int i = 0: i < n: i++)
31
        a[i].read():
32
    }
33
     std::pair<DB, int> get tangent(
        const std::vector<Point>& convex, const Point& vec) {
35
       int 1 = 0, r = (int)convex.size() - 2:
36
       assert(r >= 0);
37
       for (; 1 + 1 < r; ) {
38
        int mid = (1 + r) / 2;
39
        if (sign(det(convex[mid + 1] - convex[mid], vec)) > 0)
40
          r = mid:
41
         else 1 = mid;
42
      }
43
      return std::max(std::make pair(det(vec, convex[r]), r),
44
           std::make_pair(det(vec, convex[0]), 0));
45
46
     int binary_search(Point u, Point v, int 1, int r) {
47
       int s1 = sign(det(v - u, a[1 \% n] - u)):
48
       for (; 1 + 1 < r;) {
49
        int mid = (1 + r) / 2;
50
         int smid = sign(det(v - u, a[mid % n] - u));
51
        if (smid == s1) l = mid:
52
         else r = mid:
53
      }
54
       return 1 % n;
55
     // 求凸包上和向量 vec 叉积最大的点, 返回编号, 共线的多个切点返回任意一个
     int get_tangent(Point vec) {
58
       std::pair<DB, int> ret = get_tangent(upper, vec);
59
       ret.second = (ret.second + (int)lower.size() - 1) % n;
60
       ret = std::max(ret, get tangent(lower, vec));
```

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```
return ret.second:
62
    // 求凸包和直线 u, v 的交点, 如果不相交返回 false, 如果有则是和 (i, next(i)) 的交点, 交在
     → 点上不确定返回前后两条边其中之一
    bool get intersection(Point u, Point v, int &i0, int &i1) {
65
      int p0 = get tangent(u - v), p1 = get tangent(v - u);
      if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p1] - u)) <= 0) {
        if (p0 > p1) std::swap(p0, p1);
        i0 = binarv search(u, v, p0, p1);
        i1 = binary search(u, v, p1, p0 + n);
        return true;
71
      }
      else return false:
72
74 };
```

# 1.2 三角形的心

```
1 Point inCenter(const Point &A, const Point &B, const Point &C) { // 内心
     double a = (B - C).len(), b = (C - A).len(), c = (A - B).len(),
      s = fabs(det(B - A, C - A)).
      r = s / p;
    return (A * a + B * b + C * c) / (a + b + c);
 7 | Point circumCenter(const Point &a. const Point &b. const Point &c) { // 外心
    Point bb = b - a, cc = c - a:
    double db = bb.len2(), dc = cc.len2(), d = 2 * det(bb, cc);
    return a - Point(bb.y * dc - cc.y * db, cc.x * db - bb.x * dc) / d;
11 | }
12 Point othroCenter(const Point &a, const Point &b, const Point &c) { // 垂心
    Point ba = b - a, ca = c - a, bc = b - c;
     double Y = ba.y * ca.y * bc.y,
         A = ca.x * ba.y - ba.x * ca.y,
          x0 = (Y + ca.x * ba.y * b.x - ba.x * ca.y * c.x) / A,
         y0 = -ba.x * (x0 - c.x) / ba.y + ca.y;
    return Point(x0, y0);
```

#### 1.3 半平面交

```
16 | } else {
17
       return sign(det(a, b)) > 0;
18
19 }
20 bool operator < (const Line &10, const Line &11) {
     if (sameDir(10, 11)) {
      return 11.include(10.a);
23
      return (10.b - 10.a) < (11.b - 11.a):
25
26
   bool check(const Line &u, const Line &v, const Line &w) { return w.include(intersect(u,
28 | vector<Point> intersection(vector<Line> &1) {
     sort(l.begin(), l.end());
30
     deque<Line> q;
     for (int i = 0; i < (int)1.size(); ++i) {</pre>
      if (i && sameDir(l[i], l[i - 1])) {
33
         continue:
34
35
      while (q.size() > 1 && !check(q[q.size() - 2], q[q.size() - 1], 1[i])) q.pop back();
       while (q.size() > 1 && !check(q[1], q[0], l[i])) q.pop front();
       q.push_back(l[i]);
37
38
39
     while (q.size() > 2 \&\& !check(q[q.size() - 2], q[q.size() - 1], q[0])) q.pop_back();
     while (a.size() > 2 && !check(a[1], a[0], a[a.size() - 1])) a.pop front():
41
     vector<Point> ret:
     for (int i = 0; i < (int)q.size(); ++i) ret.push back(intersect(q[i], q[(i + 1) %
43
     return ret:
44
```

#### 1.4 圆交面积及重心

```
struct Event {
     Point p;
     double ang:
     int delta;
     Event (Point p = Point(0, 0), double ang = 0, double delta = 0) : p(p), ang(ang),

    delta(delta) {}
6 }:
   bool operator < (const Event &a, const Event &b) {
     return a.ang < b.ang;</pre>
   void addEvent(const Circle &a, const Circle &b, vector<Event> &evt, int &cnt) {
     double d2 = (a.o - b.o).len2(),
12
          dRatio = ((a.r - b.r) * (a.r + b.r) / d2 + 1) / 2,
13
          pRatio = sqrt(-(d2 - sqr(a.r - b.r)) * (d2 - sqr(a.r + b.r)) / (d2 * d2 * 4));
     Point d = b.o - a.o, p = d.rotate(PI / 2),
15
         q0 = a.o + d * dRatio + p * pRatio,
16
         q1 = a.o + d * dRatio - p * pRatio;
     double ang0 = (q0 - a.o).ang(),
17
18
          ang1 = (q1 - a.o).ang();
     evt.push back(Event(q1, ang1, 1));
     evt.push_back(Event(q0, ang0, -1));
```

```
21 | cnt += ang1 > ang0:
23 bool issame(const Circle &a, const Circle &b) { return sign((a.o - b.o).len()) == 0 &&
      \hookrightarrow sign(a.r - b.r) == 0: }
24 bool overlap(const Circle &a, const Circle &b) { return sign(a.r - b.r - (a.o -
      \rightarrow b.o).len()) >= 0: }
25 bool intersect(const Circle &a, const Circle &b) { return sign((a.o - b.o).len() - a.r -
26 Circle c[N]:
27 double area[N]; // area[k] -> area of intersections >= k.
28 Point centroid[N];
29 bool keep[N];
30 void add(int cnt. DB a. Point c) {
     area[cnt] += a:
     centroid[cnt] = centroid[cnt] + c * a;
33 }
   void solve(int C) {
     for (int i = 1: i <= C: ++ i) {
           area[i] = 0:
           centroid[i] = Point(0, 0);
37
38
     for (int i = 0; i < C; ++i) {
       int cnt = 1:
       vector<Event> evt;
       for (int j = 0; j < i; ++j) if (issame(c[i], c[j])) ++cnt;
       for (int i = 0: i < C: ++i) {
         if (j != i && !issame(c[i], c[j]) && overlap(c[j], c[i])) {
           ++cnt;
         }
       for (int j = 0; j < C; ++j) {
         if (j != i && !overlap(c[i], c[i]) && !overlap(c[i], c[j]) && intersect(c[i],
           addEvent(c[i], c[j], evt, cnt);
51
         }
52
       if (evt.size() == Ou) {
         add(cnt, PI * c[i].r * c[i].r, c[i].o):
55
       } else {
         sort(evt.begin(), evt.end());
         evt.push back(evt.front());
         for (int j = 0; j + 1 < (int)evt.size(); ++j) {
           cnt += evt[j].delta;
           add(cnt, det(evt[j].p, evt[j + 1].p) / 2, (evt[j].p + evt[j + 1].p) / 3);
61
           double ang = evt[j + 1].ang - evt[j].ang;
           if (ang < 0) {
             ang += PI * 2:
                    if (sign(ang) == 0) continue;
                    add(cnt, ang * c[i].r * c[i].r / 2, c[i].o +
                       Point(sin(ang1) - sin(ang0), -cos(ang1) + cos(ang0)) * (2 / (3 *
      \hookrightarrow ang) * c[i].r));
           add(cnt, -sin(ang) * c[i].r * c[i].r / 2, (c[i].o + evt[i].p + evt[i + 1].p) /
      \rightarrow 3);
```

#### 1.5 三维向量绕轴旋转

```
1 // 三维绕轴旋转, 大拇指指向 axis 向量方向, 四指弯曲方向转 w 弧度
  Point rotate(const Point& s, const Point& axis, DB w) {
    DB x = axis.x, y = axis.y, z = axis.z;
    DB s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
       cosw = cos(w), sinw = sin(w);
    DB a[4][4];
    memset(a, 0, sizeof a):
    a[3][3] = 1;
    a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
    a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
    a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
    a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
    a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
    a[1][2] = v * z * (1 - cosw) / s1 + x * sinw / ss1;
    a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
    a[2][1] = v * z * (1 - cosw) / s1 - x * sinw / ss1;
    a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;
    DB ans [4] = \{0, 0, 0, 0\}, c[4] = \{s.x, s.y, s.z, 1\};
    for (int i = 0: i < 4: ++ i)
      for (int j = 0; j < 4; ++ j)
        ans[i] += a[j][i] * c[j];
21
    return Point(ans[0], ans[1], ans[2]);
23 }
```

# 1.6 三维凸包

```
inline P cross(const P& a, const P& b) {
    return P(
        a.y * b.z - a.z * b.y,
        a.z * b.x - a.x * b.z
        a.x * b.y - a.y * b.x
          ):
   inline DB mix(const P& a, const P& b, const P& c) {
    return dot(cross(a, b), c):
11 }
12
   inline DB volume(const P& a, const P& b, const P& c, const P& d) {
    return mix(b - a, c - a, d - a):
15 }
16
17 struct Face {
   int a. b. c:
   inline Face() {}
    inline Face(int _a, int _b, int _c):
```

```
21
       a(a), b(b), c(c) {}
     inline DB area() const {
       return 0.5 * cross(p[b] - p[a], p[c] - p[a]).len();
23
24
25
     inline P normal() const {
       return cross(p[b] - p[a], p[c] - p[a]).unit();
27
     __inline DB dis(const P& p0) const {
28
      return dot(normal(), p0 - p[a]);
30
31 };
   std::vector<Face> face. tmp: // Should be O(n).
   int mark[N][N], Time, n;
    inline void add(int v) {
37
    ++ Time:
38
     clear(tmp);
     for (int i = 0; i < (int)face.size(); ++ i) {</pre>
       int a = face[i].a, b = face[i].b, c = face[i].c;
       if (sign(volume(p[v], p[a], p[b], p[c])) > 0) {
         mark[a][b] = mark[b][a] = mark[a][c] =
43
           mark[c][a] = mark[b][c] = mark[c][b] = Time;
44
       }
       else {
46
         tmp.push back(face[i]);
47
       }
48
     clear(face); face = tmp;
     for (int i = 0: i < (int)tmp.size(): ++ i) {</pre>
       int a = face[i].a, b = face[i].b, c = face[i].c;
       if (mark[a][b] == Time) face.emplace back(v, b, a);
52
53
       if (mark[b][c] == Time) face.emplace back(v, c, b);
       if (mark[c][a] == Time) face.emplace back(v, a, c);
       assert(face.size() < 500u);</pre>
56
57 }
58
   void reorder() {
     for (int i = 2; i < n; ++ i) {
61
       P \text{ tmp} = cross(p[i] - p[0], p[i] - p[1]);
       if (sign(tmp.len())) {
63
         std::swap(p[i], p[2]);
         for (int j = 3; j < n; ++ j)
65
           if (sign(volume(p[0], p[1], p[2], p[j]))) {
             std::swap(p[j], p[3]);
             return:
      }
70
71 }
72
   void build convex() {
     reorder():
     clear(face);
```

```
face.emplace_back(0, 1, 2);
face.emplace_back(0, 2, 1);
for (int i = 3; i < n; ++ i)
    add(i);
}
</pre>
```

# 2 数论

### $2.1 O(m^2 \log n)$ 求线性递推数列第 n 项

```
Given a_0, a_1, \ldots, a_{m-1}

a_n = c_0 \times a_{n-m} + \cdots + c_{m-1} \times a_{n-1}

Solve for a_n = v_0 \times a_0 + v_1 \times a_1 + \cdots + v_{m-1} \times a_{m-1}
```

```
void linear recurrence(long long n, int m, int a[], int c[], int p) {
     long long v[M] = \{1 \% p\}, u[M << 1], msk = !!n;
     for(long long i(n); i > 1; i >>= 1) {
      msk <<= 1;
     for(long long x(0); msk; msk >>= 1, x <<= 1) {
       fill n(u, m << 1, 0);
       int b(!!(n & msk));
       x = b;
10
       if(x < m) {
11
         u[x] = 1 \% p;
12
       }else {
13
         for(int i(0); i < m; i++) {
14
           for(int j(0), t(i + b); j < m; j++, t++) {</pre>
15
             u[t] = (u[t] + v[i] * v[j]) % p;
16
           }
17
18
         for(int i((m << 1) - 1); i >= m; i--) {
19
           for(int j(0), t(i - m); j < m; j++, t++) {
20
             u[t] = (u[t] + c[j] * u[i]) % p;
21
22
        }
23
      }
24
       copy(u, u + m, v);
     //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] * a[m - 1].
27
     for(int i(m): i < 2 * m: i++) {
28
       a[i] = 0;
29
       for(int j(0); j < m; j++) {
30
         a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31
32
33
     for(int j(0); j < m; j++) {
34
       b[j] = 0;
       for(int i(0): i < m: i++) {
36
         b[i] = (b[i] + v[i] * a[i + j]) % p;
37
      }
38
    for(int j(0); j < m; j++) {
      a[i] = b[i];
41
```

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42 | }

```
2.2 求逆元
```

```
void ex_gcd(long long a, long long b, long long &x, long long &y) {
   if (b == 0) {
      x = 1;
      y = 0;
      return;
}
long long xx, yy;
ex_gcd(b, a % b, xx, yy);
y = xx - a / b * yy;
x = yy;
}
long long inv(long long x, long long MODN) {
   long long inv_x, y;
ex_gcd(x, MODN, inv_x, y);
   return (inv_x % MODN + MODN) % MODN;
}
```

# 2.3 中国剩余定理

```
1 // 返回 (ans, M), 其中 ans 是模 M 意义下的解
 2 std::pair<long long, long long> CRT(const std::vector<long long>& m, const

    std::vector<long long>& a) {
    long long M = 1, ans = 0;
     int n = m.size();
    for (int i = 0; i < n; i++) M *= m[i];
    for (int i = 0: i < n: i++) {
       ans = (ans + (M / m[i]) * a[i] % M * inv(M / m[i], m[i])) % M: // 可能需要大整数相乘
     →取模
     return std::make pair(ans, M);
   // 模数不互质的情况
   bool solve(int n, std::pair<long long, long long> input[],
13
                    std::pair<long long, long long> &output) {
     output = std::make pair(1, 1);
     for (int i = 0; i < n; ++i) {
       long long number, useless;
       // euclid(a, b, x, y)
17
18
       euclid(output.second, input[i].second, number, useless);
19
       long long divisor = std::__gcd(output.second, input[i].second);
       if ((input[i].first - output.first) % divisor) return false;
20
21
       number *= (input[i].first - output.first) / divisor;
22
       fix(number, input[i].second); // fix 成正的
23
       output.first += output.second * number;
24
       output.second *= input[i].second / divisor;
       fix(output.first, output.second);
26
27
    return true;
28 }
```

# 2.4 魔法 CRT

```
1 // MOD is the given module
   // Do not depend on LL * LL % LL
   inline int CRT(int *a) {
     static int x[N]:
     for (int i = 0; i < N; i ++) {
      x[i] = a[i];
      for (int j = 0; j < i; j ++) {
        int t = (x[i] - x[i] + mod[i]) \% mod[i]:
        if (t < 0) t += mod[i];</pre>
10
        x[i] = 1LL * t * Inv[i][i] % mod[i];
11
12
     int sum = 1, ret = x[0] % MOD;
13
     for (int i = 1; i < N; i ++) {
      sum = 1LL * sum * mod[i - 1] % MOD;
16
      ret += 1LL * x[i] * sum % MOD:
17
      if (ret >= MOD) ret -= MOD;
18
19
    return ret;
20 }
21
   for (int i = 0; i < N; i ++)
    for (int j = i + 1; j < N; j ++) {
23
      Inv[i][j] = fpw(mod[i], mod[j] - 2, mod[j]);
24
```

#### 2.5 素性测试

```
int strong_pseudo_primetest(long long n,int base) {
       long long n2=n-1.res:
       int s=0:
       while (n2\%2==0) n2>>=1,s++;
       res=powmod(base,n2,n);
       if((res==1)||(res==n-1)) return 1:
7
       s--:
       while(s>=0) {
           res=mulmod(res,res,n);
10
           if(res==n-1) return 1;
11
12
13
       return 0; // n is not a strong pseudo prime
14 }
15 int isprime(long long n) {
     static LL testNum[]={2,3,5,7,11,13,17,19,23,29,31,37};
     static LL lim[]={4,0,1373653LL,25326001LL,25000000000LL,2152302898747LL,
      \rightarrow 3474749660383LL,341550071728321LL,0,0,0,0);
     if(n<2||n==3215031751LL) return 0:
19
     for(int i=0;i<12;++i){
20
       if(n<lim[i]) return 1;</pre>
21
       if(strong pseudo primetest(n,testNum[i])==0) return 0;
22
    }
23
     return 1;
24
```

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# 2.6 质因数分解

```
1 int ansn; LL ans[1000];
 2 LL func(LL x,LL n) { return(mod mul(x,x,n)+1)%n; }
 3 LL Pollard(LL n){
    LL i,x,v,p;
     if(Rabin Miller(n)) return n;
     if(!(n&1)) return 2;
     for(i=1;i<20;i++){
       x=i; y=func(x,n); p=gcd(y-x,n);
       while (p==1) {x=func(x,n); y=func(func(y,n),n); p=gcd((y-x+n)\%n,n)\%n;}
       if(p==0||p==n) continue;
       return p;
12
13 | }
   void factor(LL n){
    LL x;
    x=Pollard(n);
    if(x==n){ ans[ansn++]=x; return; }
    factor(x), factor(n/x):
```

#### 2.7 线下整点

#### 2.8 原根相关

- 1. 模 m 有原根的充要条件:  $m = 2, 4, p^a, 2p^a$ , 其中 p 是奇素数;
- 2. 求任意数 p 原根的方法: 对  $\phi(p)$  因式分解, 即  $\phi(p) = p_1^{r_1} p_2^{r_2} \cdots p_k^{r_k}$ , 若恒成立:

$$g^{\frac{p-1}{g}} \neq 1 \pmod{p}$$

那么 g 就是 p 的原根。

3. 若模 m 有原根,那么它一共有  $\Phi(\Phi(m))$  个原根。

### 3 代数

#### 3.1 快速傅里叶变换

```
int prepare(int n) {
   int len = 1;
   for (; len <= 2 * n; len <<= 1);
   for (int i = 0; i < len; i++) {
      e[0][i] = Complex(cos(2 * pi * i / len), sin(2 * pi * i / len));
      e[1][i] = Complex(cos(2 * pi * i / len), -sin(2 * pi * i / len));
}
return len;
}
void DFT(Complex *a, int n, int f) {</pre>
```

```
for (int i = 0, j = 0; i < n; i++) {
11
12
      if (i > j) std::swap(a[i], a[j]);
13
       for (int t = n >> 1; (j = t) < t; t >>= 1);
14
15
     for (int i = 2; i <= n; i <<= 1)
16
      for (int j = 0; j < n; j += i)
17
        for (int k = 0; k < (i >> 1); k++) {
18
           Complex A = a[j + k];
19
           Complex B = e[f][n / i * k] * a[j + k + (i >> 1)];
20
           a[j + k] = A + B;
21
           a[i + k + (i >> 1)] = A - B;
22
        }
     if (f == 1) {
23
      for (int i = 0; i < n; i++)
25
         a[i].a /= n;
26
27 }
```

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#### 3.2 快速数论变换

```
1 // meminit(A, l, r) 是将数组 A 的 [l, r) 清 O。
  // memcopy(target, source, l, r) 是将 source 的 [l, r) 复制到 target 的 [l, r)
   #define meminit(A, 1, r) memset(A + (1), 0, sizeof(*A) * ((r) - (1)))
   #define memcopy(B, A, 1, r) memcpy(B, A + (1), sizeof(*A) * ((r) - (1)))
   void DFT(int *a, int n, int f) { // 封闭形式, 常数小 (10^7) 跑 2.23 秒)
    for (register int i = 0, j = 0; i < n; i++) {
      if (i > j) std::swap(a[i], a[j]);
8
      for (register int t = n \gg 1; (j ^= t) < t; t >>= 1);
9
10
     for (register int i = 2; i \le n; i \le 1) {
11
      static int exp[MAXN];
12
      \exp[0] = 1; \exp[1] = fpm(PRT, (MOD - 1) / i);
13
      if (f == 1) \exp[1] = fpm(\exp[1], MOD - 2);
14
      for (register int k = 2; k < (i >> 1); k++) {
15
        \exp[k] = 111 * \exp[k - 1] * \exp[1] % MOD;
16
17
      for (register int j = 0; j < n; j += i) {
18
        for (register int k = 0; k < (i >> 1); k++) {
19
          register int &pA = a[j + k], &pB = a[j + k + (i >> 1)];
20
          register int A = pA, B = 111 * pB * exp[k] % MOD;
21
          pA = (A + B) \% MOD;
22
          pB = (A - B + MOD) \% MOD;
23
24
      }
25
26
    if (f == 1) {
27
      register int rev = fpm(n, MOD - 2, MOD);
28
      for (register int i = 0; i < n; i++) {
29
        a[i] = 111 * a[i] * rev % MOD:
30
31
33 // 在不写高精度的情况下合并 FFT 所得结果对 MOD 取模后的答案
34 // 值得注意的是,这个东西不能最后再合并,而是应该每做一次多项式乘法就 CRT 一次
35 int CRT(int *a) {
```

```
static int x[3];
     for (int i = 0; i < 3; i++) {
       x[i] = a[i];
       for (int j = 0; j < i; j++) {
         int t = (x[i] - x[j] + FFT[i] \rightarrow MOD) \% FFT[i] \rightarrow MOD;
41
         if (t < 0) t += FFT[i] -> MOD;
         x[i] = 1LL * t * inv[j][i] % FFT[i] -> MOD;
43
    }
44
     int sum = 1, ret = x[0] % MOD;
     for (int i = 1; i < 3; i ++) {
       sum = 1LL * sum * FFT[i - 1] \rightarrow MOD % MOD;
       ret += 1LL * x[i] * sum % MOD;
       if(ret >= MOD) ret -= MOD;
50
    }
51
    return ret;
52 | }
53 for (int i = 0; i < 3; i++) // inv 数组的预处理过程, inverse(x, p) 表示求 x 在 p 下逆元
     for (int j = 0; j < 3; j++)
       inv[i][j] = inverse(FFT[i] -> MOD, FFT[j] -> MOD);
```

#### 3.3 自适应辛普森积分

```
namespace adaptive simpson {
     template<typename function>
     inline double area(function f, const double &left, const double &right) {
       double mid = (left + right) / 2;
       return (right - left) * (f(left) + 4 * f(mid) + f(right)) / 6;
     template<typename function>
     inline double simpson(function f, const double &left, const double &right, const

    double &eps, const double &area sum) {
       double mid = (left + right) / 2;
11
       double area_left = area(f, left, mid);
12
       double area_right = area(f, mid, right);
       double area total = area left + area right;
       if (fabs(area total - area sum) <= 15 * eps) {
15
         return area_total + (area_total - area_sum) / 15;
16
       return simpson(f, left, right, eps / 2, area left) + simpson(f, mid, right, eps / 2,
17
      → area right):
19
     template<typename function>
     inline double simpson(function f, const double &left, const double &right, const
21

    double &eps) {
22
       return simpson(f, left, right, eps, area(f, left, right));
23
24 }
```

# 3.4 单纯形

```
const double eps = 1e-8;
// max{c * x | Ax <= b, x >= 0} 的解, 无解返回空的 vector, 否则就是解.
vector<double> simplex(vector<vector<double> > &A, vector<double> b, vector<double> c) {
```

```
vector < vector < double > D(n + 2, vector < double > (m + 1));
     vector<int> ix(n + m);
     for(int i = 0; i < n + m; i++) {
       ix[i] = i;
9
     for(int i = 0; i < n; i++) {
11
       for(int j = 0; j < m - 1; j++) {
12
         D[i][i] = -A[i][i]:
13
14
       D[i][m-1]=1;
15
       D[i][m] = b[i];
       if (D[r][m] > D[i][m]) {
16
17
         r = i:
18
       }
19
20
21
     for(int j = 0; j < m - 1; j++) {
22
       D[n][i] = c[i];
23
     D[n + 1][m - 1] = -1;
24
25
     for(double d: :) {
26
       if (r < n) 
         swap(ix[s], ix[r + m]);
27
28
         D[r][s] = 1. / D[r][s];
29
         for(int i = 0: i <= m: i++) {
30
           if (j != s) {
31
             D[r][i] *= -D[r][s];
32
33
34
         for(int i = 0; i \le n + 1; i++) {
35
           if (i != r) {
36
             for(int j = 0; j \le m; j++) {
37
               if (i != s) {
38
                 D[i][j] += D[r][j] * D[i][s];
39
40
41
             D[i][s] *= D[r][s];
42
           }
43
         }
       r = -1, s = -1:
       for(int j = 0; j < m; j++) {
47
         if (s < 0 || ix[s] > ix[j]) {
48
           if (D[n + 1][j] > eps || D[n + 1][j] > -eps && D[n][j] > eps) {
             s = j;
50
           }
51
         }
52
       }
53
       if (s < 0) {
54
         break:
55
56
       for(int i = 0; i < n; i++) {
57
         if (D[i][s] < -eps) {</pre>
58
           if (r < 0 \mid | (d = D[r][m] / D[r][s] - D[i][m] / D[i][s]) < -eps
```

int n = A.size(), m = A[0].size() + 1, r = n, s = m - 1;

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```
|| d < eps && ix[r + m] > ix[i + m]) {
             r = i;
62
63
         }
64
       }
       if (r < 0) {
         return vector<double> ():
68
69
    if (D[n + 1][m] < -eps) {
       return vector<double> ():
72
73
     vector<double> x(m - 1);
     for(int i = m; i < n + m; i++) {</pre>
       if (ix[i] < m - 1) {
         x[ix[i]] = D[i - m][m];
78
79
    }
    return x;
```

# 4 字符串

# 4.1 后缀数组

```
const int MAXN = MAXL * 2 + 1:
   int a[MAXN], x[MAXN], y[MAXN], c[MAXN], sa[MAXN], rank[MAXN], height[MAXN];
   void calc sa(int n) {
    int m = alphabet, k = 1;
    memset(c, 0, sizeof(*c) * (m + 1));
    for (int i = 1: i \le n: ++i) c[x[i] = a[i]]++:
     for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
     for (int i = n; i; --i) sa[c[x[i]]--] = i;
     for (; k <= n; k <<= 1) {
       for (int i = n - k + 1; i \le n; ++i) y[i - n + k] = i;
       for (int i = 1; i \le n; ++i)
13
        if (sa[i] > k) y[++tot] = sa[i] - k;
       memset(c, 0, sizeof(*c) * (m + 1));
       for (int i = 1; i <= n; ++i) c[x[i]]++;
       for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
17
       for (int i = n; i; --i) sa[c[x[y[i]]]--] = y[i];
       for (int i = 1; i \le n; ++i) y[i] = x[i];
       tot = 1; x[sa[1]] = 1;
       for (int i = 2; i <= n; ++i) {
         if (\max(sa[i], sa[i-1]) + k > n || y[sa[i]] != y[sa[i-1]] || y[sa[i] + k] !=
      \rightarrow v[sa[i - 1] + k]) ++tot:
         x[sa[i]] = tot;
23
24
       if (tot == n) break; else m = tot;
27 void calc_height(int n) {
```

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#### 4.2 后缀自动机

```
1 static const int MAXL = MAXN * 2; // MAXN is original length
2 static const int alphabet = 26; // sometimes need changing
3 int 1, last, cnt, trans[MAXL][alphabet], par[MAXL], sum[MAXL], seq[MAXL], mxl[MAXL],
     → size[MAXL]; // mxl is maxlength, size is the size of right
   char str[MAXL]:
   inline void init() {
    l = strlen(str + 1); cnt = last = 1;
    for (int i = 0; i \le 1 * 2; ++i) memset(trans[i], 0, sizeof(trans[i]));
    memset(par, 0, sizeof(*par) * (1 * 2 + 1));
    memset(mxl, 0, sizeof(*mxl) * (1 * 2 + 1)):
     memset(size, 0, sizeof(*size) * (1 * 2 + 1)):
11 }
   inline void extend(int pos, int c) {
    int p = last, np = last = ++cnt;
     mxl[np] = mxl[p] + 1; size[np] = 1;
     for (; p && !trans[p][c]; p = par[p]) trans[p][c] = np;
     if (!p) par[np] = 1;
17
     else {
18
       int q = trans[p][c];
       if (mxl[p] + 1 == mxl[q]) par[np] = q;
20
       else {
21
         int nq = ++cnt;
22
         mxl[nq] = mxl[p] + 1;
23
         memcpy(trans[nq], trans[q], sizeof(trans[nq]));
24
         par[nq] = par[q];
         par[np] = par[q] = nq;
26
         for (; trans[p][c] == q; p = par[p]) trans[p][c] = nq;
27
28
29 }
   inline void buildsam() {
     for (int i = 1; i <= 1; ++i) extend(i, str[i] - 'a');</pre>
    memset(sum, 0, sizeof(*sum) * (1 * 2 + 1));
     for (int i = 1; i <= cnt; ++i) sum[mxl[i]]++;
     for (int i = 1; i <= 1; ++i) sum[i] += sum[i - 1];
    for (int i = cnt; i; --i) seq[sum[mxl[i]]--] = i;
    for (int i = cnt; i; --i) size[par[seq[i]]] += size[seq[i]];
37
```

#### 4.3 EX 后缀自动机

```
inline void add_node(int x, int &last) {
  int lastnode = last;
```

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```
if (c[lastnode][x]) {
       int nownode = c[lastnode][x];
       if (l[nownode] == l[lastnode] + 1) last = nownode;
         int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
         for (int i = 0; i < alphabet; ++i) c[auxnode][i] = c[nownode][i];</pre>
         par[auxnode] = par[nownode]; par[nownode] = auxnode;
         for (; lastnode && c[lastnode][x] == nownode; lastnode = par[lastnode]) {
           c[lastnode][x] = auxnode:
13
         last = auxnode;
14
       }
    } else {
       int newnode = ++cnt: l[newnode] = l[lastnode] + 1:
       for (; lastnode && !c[lastnode][x]; lastnode = par[lastnode]) c[lastnode][x] =
       if (!lastnode) par[newnode] = 1;
       else {
         int nownode = c[lastnode][x];
         if (l[lastnode] + 1 == l[nownode]) par[newnode] = nownode;
21
22
23
           int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
24
           for (int i = 0; i < alphabet; ++i) c[auxnode][i] = c[nownode][i];</pre>
           par[auxnode] = par[nownode]; par[nownode] = par[newnode] = auxnode;
25
           for (: lastnode && c[lastnode][x] == nownode: lastnode = par[lastnode]) {
             c[lastnode][x] = auxnode:
28
           }
29
        }
30
       }
31
       last = newnode:
32
33 }
```

#### 4.4 后缀树

- 1. 边上的字符区间是左闭右开区间;
- 如果要建立关于多个串的后缀树,请用不同的分隔符,并且对于每个叶子结点,去掉和它父亲的连边上出现的第一个分隔符之后的所有字符;

# 4.5 回文自动机

```
int nT, nStr, last, c[MAXT][26], fail[MAXT], r[MAXN], l[MAXN], s[MAXN];
int allocate(int len) {
    l[nT] = len;
    r[nT] = 0;
    fail[nT] = 0;
    memset(c[nT], 0, sizeof(c[nT]));
    return nT++;
    }
    void init() {
    nT = nStr = 0;
    int newE = allocate(0);
    int newO = allocate(-1);
    last = newE;
    fail[newE] = newO;
    fail[newO] = newE;
    s[0] = -1;
```

```
17 | }
18 void add(int x) {
     s[++nStr] = x;
     int now = last;
     while (s[nStr - l[now] - 1] != s[nStr]) now = fail[now];
     if (!c[now][x]) {
      int newnode = allocate(l[now] + 2), &newfail = fail[newnode];
24
       newfail = fail[now];
       while (s[nStr - l[newfail] - 1] != s[nStr]) newfail = fail[newfail]:
26
       newfail = c[newfail][x]:
       c[now][x] = newnode:
27
28
29
    last = c[now][x]:
    r[last]++:
31 }
32
   void count() {
    for (int i = nT - 1; i \ge 0; i--) {
34
      r[fail[i]] += r[i];
35
   }
36 }
```

# 5 数据结构

#### 5.1 KD-Tree

```
long long norm(const long long &x) {
       // For manhattan distance
       return std::abs(x):
       // For euclid distance
5
       return x * x;
6
   struct Point {
       int x, y, id;
10
       const int& operator [] (int index) const {
11
12
           if (index == 0) {
13
               return x:
14
          } else {
15
               return y;
16
           }
17
       }
18
19
       friend long long dist(const Point &a, const Point &b) {
20
           long long result = 0;
21
           for (int i = 0; i < 2; ++i) {
22
               result += norm(a[i] - b[i]);
23
24
           return result;
25
26 } point[N];
28 struct Rectangle {
      int min[2], max[2];
```

```
31
       Rectangle() {
32
           min[0] = min[1] = INT MAX; // sometimes int is not enough
           max[0] = max[1] = INT MIN;
33
34
       }
35
36
       void add(const Point &p) {
           for (int i = 0; i < 2; ++i) {
37
38
               min[i] = std::min(min[i], p[i]);
39
               max[i] = std::max(max[i], p[i]);
40
           }
41
       }
42
43
       long long dist(const Point &p) {
44
           long long result = 0;
45
           for (int i = 0; i < 2; ++i) {
               // For minimum distance
46
               result += norm(std::min(std::max(p[i], min[i]), max[i]) - p[i]);
47
               // For maximum distance
               result += std::max(norm(max[i] - p[i]), norm(min[i] - p[i]));
50
51
           return result;
52
53 };
54
   struct Node {
       Point seperator;
57
       Rectangle rectangle;
58
       int child[2]:
59
60
       void reset(const Point &p) {
61
           seperator = p;
62
           rectangle = Rectangle();
63
           rectangle.add(p);
           child[0] = child[1] = 0;
   } tree[N << 1];</pre>
67
68
   int size, pivot;
69
   bool compare(const Point &a, const Point &b) {
       if (a[pivot] != b[pivot]) {
71
72
           return a[pivot] < b[pivot];</pre>
73
       }
74
       return a.id < b.id;
75 | }
76
   // 左閉右開: build(1, n + 1)
   int build(int 1, int r, int type = 1) {
79
       pivot = type;
       if (1 >= r) {
80
81
           return 0;
83
       int x = ++size;
       int mid = 1 + r \gg 1;
```

```
std::nth element(point + 1, point + mid, point + r, compare);
 85
 86
        tree[x].reset(point[mid]);
 87
        for (int i = 1; i < r; ++i) {
 88
            tree[x].rectangle.add(point[i]);
 89
 90
        tree[x].child[0] = build(1, mid, type ^ 1);
 91
        tree[x].child[1] = build(mid + 1, r, type ^ 1);
 92
        return x:
 93
 94
    int insert(int x, const Point &p, int type = 1) {
        pivot = type;
        if (x == 0) {
 97
 98
            tree[++size].reset(p);
 99
            return size;
100
       }
        tree[x].rectangle.add(p);
101
102
        if (compare(p, tree[x].seperator)) {
            tree[x].child[0] = insert(tree[x].child[0], p, type ^ 1);
103
104
105
            tree[x].child[1] = insert(tree[x].child[1], p, type ^ 1);
106
       }
107
        return x;
108
109
110 // For minimum distance
111 // For maximum: 下面递归 query 时 0, 1 换顺序;< and >;min and max
    void query(int x, const Point &p, std::pair<long long, int> &answer, int type = 1) {
        pivot = type:
114
        if (x == 0 || tree[x].rectangle.dist(p) > answer.first) {
115
            return;
       }
116
117
        answer = std::min(answer.
118
                 std::make pair(dist(tree[x].seperator, p), tree[x].seperator.id));
119
        if (compare(p, tree[x].seperator)) {
120
            query(tree[x].child[0], p, answer, type ^ 1);
121
            query(tree[x].child[1], p, answer, type ^ 1);
122
            query(tree[x].child[1], p, answer, type ^ 1);
123
124
            query(tree[x].child[0], p, answer, type ^ 1);
       }
125
126
    }
127
    std::priority_queue<std::pair<long long, int> > answer;
129
    void query(int x, const Point &p, int k, int type = 1) {
131
        pivot = type;
        if (x == 0 || (int)answer.size() == k && tree[x].rectangle.dist(p) >
       133
            return:
       }
134
135
        answer.push(std::make_pair(dist(tree[x].seperator, p), tree[x].seperator.id));
136
        if ((int)answer.size() > k) {
137
            answer.pop();
138
       }
```

```
5.2 Treap
   struct Node{
     int mn, key, size, tag;
     bool rev;
     Node* ch[2];
     Node(int mn, int key, int size): mn(mn), key(key), size(size), rev(0), tag(0){}
     void downtag():
     Node* update(){
       mn = min(ch[0] \rightarrow mn, min(key, ch[1] \rightarrow mn));
       size = ch[0] \rightarrow size + 1 + ch[1] \rightarrow size;
       return this:
11
12 };
   typedef pair<Node*, Node*> Pair;
   Node *null, *root;
   void Node::downtag(){
     if(rev){
       for(int i = 0; i < 2; i++)
         if(ch[i] != null){
18
19
            ch[i] -> rev ^= 1:
            swap(ch[i] -> ch[0], ch[i] -> ch[1]);
21
22
       rev = 0:
23
    }
24
     if(tag){
25
       for(int i = 0; i < 2; i++)
26
         if(ch[i] != null){
27
            ch[i] -> key += tag;
            ch[i] -> mn += tag;
29
            ch[i] -> tag += tag;
30
         }
31
       tag = 0;
32
33 }
     static int s = 3023192386:
     return (s += (s << 3) + 1) & (^{\circ}0u >> 1);
37 }
38 bool random(int x, int y){
     return r() \% (x + y) < x;
40 }
41 Node* merge(Node *p, Node *q){
     if(p == null) return q;
     if(q == null) return p;
    p -> downtag();
    q -> downtag();
```

```
if(random(p -> size, q -> size)){
        p \rightarrow ch[1] = merge(p \rightarrow ch[1], q);
48
        return p -> update();
49
      }else{
50
        q \rightarrow ch[0] = merge(p, q \rightarrow ch[0]);
51
        return q -> update();
52
53
   Pair split(Node *x. int n){
     if(x == null) return make pair(null, null);
      x -> downtag();
      if(n \le x \rightarrow ch[0] \rightarrow size){
        Pair ret = split(x \rightarrow ch[0], n);
        x \rightarrow ch[0] = ret.second;
60
        return make pair(ret.first, x -> update());
61
     Pair ret = split(x \rightarrow ch[1], n - x \rightarrow ch[0] \rightarrow size - 1);
62
     x \rightarrow ch[1] = ret.first;
     return make_pair(x -> update(), ret.second);
65
   pair<Node*, Pair> get segment(int 1, int r){
     Pair ret = split(root, l - 1);
     return make_pair(ret.first, split(ret.second, r - 1 + 1));
70
   int main(){
     null = new Node(INF, INF, 0):
     null \rightarrow ch[0] = null \rightarrow ch[1] = null;
73
    root = null;
74 | }
```

#### 5.3 Link/cut Tree

```
1 inline void reverse(int x) {
    tr[x].rev ^= 1; swap(tr[x].c[0], tr[x].c[1]);
3 }
   inline void rotate(int x. int k) {
    int y = tr[x].fa, z = tr[y].fa;
       tr[x].fa = z; tr[z].c[tr[z].c[1] == y] = x;
       tr[tr[x].c[k ^ 1]].fa = y; tr[y].c[k] = tr[x].c[k ^ 1];
9
       tr[x].c[k ^ 1] = y; tr[y].fa = x;
10 }
11
12 inline void splay(int x, int w) {
     int z = x: pushdown(x):
14
     while (tr[x].fa != w) {
15
       int y = tr[x].fa; z = tr[y].fa;
16
      if (z == w) {
17
         pushdown(z = y); pushdown(x);
18
         rotate(x, tr[y].c[1] == x);
19
         update(y); update(x);
20
      } else {
21
         pushdown(z); pushdown(y); pushdown(x);
22
         int t1 = tr[y].c[1] == x, t2 = tr[z].c[1] == y;
23
         if (t1 == t2) rotate(y, t2), rotate(x, t1);
```

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```
else rotate(x, t1), rotate(x, t2);
         update(z); update(y); update(x);
26
       }
27
    }
28
    update(x);
    if (x != z) par[x] = par[z], par[z] = 0;
31
   inline void access(int x) {
     for (int y = 0; x; y = x, x = par[x]) {
33
       splay(x, 0);
       if (tr[x].c[1]) par[tr[x].c[1]] = x, tr[tr[x].c[1]].fa = 0;
       tr[x].c[1] = y; par[y] = 0; tr[y].fa = x; update(x);
38 }
39
   inline void makeroot(int x) {
     access(x); splay(x, 0); reverse(x);
42 }
   inline void link(int x, int y) {
     makeroot(x); par[x] = y;
47
   inline void cut(int x, int y) {
     access(x): splav(v, 0):
    if (par[y] != x) swap(x, y), access(x), splay(y, 0);
51
    par[y] = 0;
52 }
   inline void split(int x, int y) { // x will be the root of the tree
    makeroot(y); access(x); splay(x, 0);
```

#### 5.4 树状数组查询第 k 小元素

```
1 inline void reverse(int x) {
     tr[x].rev ^= 1; swap(tr[x].c[0], tr[x].c[1]);
 3 }
   inline void rotate(int x, int k) {
    int y = tr[x].fa, z = tr[y].fa;
       tr[x].fa = z; tr[z].c[tr[z].c[1] == y] = x;
       tr[tr[x].c[k ^ 1]].fa = y; tr[y].c[k] = tr[x].c[k ^ 1];
       tr[x].c[k ^ 1] = y; tr[y].fa = x;
10 }
11
   inline void splay(int x, int w) {
    int z = x; pushdown(x);
     while (tr[x].fa != w) {
       int y = tr[x].fa; z = tr[y].fa;
16
       if (z == w) {
         pushdown(z = y); pushdown(x);
         rotate(x, tr[y].c[1] == x);
         update(v): update(x):
       } else {
```

```
21
         pushdown(z); pushdown(y); pushdown(x);
22
         int t1 = tr[y].c[1] == x, t2 = tr[z].c[1] == y;
23
         if (t1 == t2) rotate(v, t2), rotate(x, t1);
24
         else rotate(x, t1), rotate(x, t2);
25
         update(z); update(y); update(x);
26
27
    }
28
    update(x);
    if (x != z) par[x] = par[z], par[z] = 0;
30 }
31
   inline void access(int x) {
    for (int y = 0; x; y = x, x = par[x]) {
       splay(x, 0);
      if (tr[x].c[1]) par[tr[x].c[1]] = x, tr[tr[x].c[1]].fa = 0;
       tr[x].c[1] = y; par[y] = 0; tr[y].fa = x; update(x);
37
38
39
   inline void makeroot(int x) {
    access(x); splay(x, 0); reverse(x);
42 }
43
   inline void link(int x, int y) {
    makeroot(x); par[x] = y;
46 }
47
   inline void cut(int x, int y) {
     access(x); splay(y, 0);
    if (par[y] != x) swap(x, y), access(x), splay(y, 0);
    par[v] = 0;
52 }
53
   inline void split(int x, int y) { // x will be the root of the tree
    makeroot(y); access(x); splay(x, 0);
56 }
```

# 6 图论

#### 6.1 基础

```
struct Graph { // Remember to call .init()!
       int e, nxt[M], v[M], adj[N], n;
       bool base:
       __inline void init(bool _base, int _n = 0) {
           assert(n < N);
           n = _n; base = _base;
           e = 0; memset(adj + base, -1, sizeof(*adj) * n);
       __inline int new_node() {
10
           adi[n + base] = -1;
11
           assert(n + base + 1 < N);
12
           return n++ + base:
13
14
       __inline void ins(int u0, int v0) { // directional
```

```
assert(u0 < n + base && v0 < n + base);
v[e] = v0; nxt[e] = adj[u0]; adj[u0] = e++;
assert(e < M);

}
__inline void bi_ins(int u0, int v0) { // bi-directional ins(u0, v0); ins(v0, u0);
};
};</pre>
```

```
6.2 KM
1 struct KM {
    // Truly O(n^3)
    // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时边权取负, 但不能连的还是 -INF, 使用时先对 1
     → -> n 调用 hungary() , 再 get ans() 求值
    int w[N][N]:
    int lx[N]. lv[N]. match[N]. wav[N]. slack[N]:
    bool used[N];
    void init() {
      for (int i = 1; i \le n; i++) {
        match[i] = 0:
        lx[i] = 0:
        lv[i] = 0;
11
        wav[i] = 0;
13
14
    }
     void hungary(int x) {
      match[0] = x;
      int j0 = 0;
17
18
       for (int j = 0; j \le n; j++) {
        slack[j] = INF;
20
        used[j] = false;
21
       }
22
23
       do {
24
        used[j0] = true;
        int i0 = match[j0], delta = INF, j1 = 0;
        for (int j = 1; j <= n; j++) {
          if (used[i] == false) {
27
28
            int cur = -w[i0][j] - lx[i0] - lv[j];
            if (cur < slack[j]) {</pre>
30
              slack[j] = cur;
              way[j] = j0;
32
33
             if (slack[j] < delta) {</pre>
34
              delta = slack[i]:
35
              j1 = j;
36
37
          }
38
        for (int j = 0; j \le n; j++) {
           if (used[i]) {
            lx[match[j]] += delta;
            lv[i] -= delta:
44
           else slack[j] -= delta;
```

```
45
46
         j0 = j1;
47
      } while (match[j0] != 0);
48
49
       do {
         int j1 = way[j0];
50
51
         match[j0] = match[j1];
52
         j0 = j1;
      } while (i0):
54
55
     int get ans() {
      int sum = 0:
57
       for(int i = 1; i <= n; i++) {
        if (w[match[i]][i] == -INF); // 无解
60
        if (match[i] > 0) sum += w[match[i]][i];
61
62
      return sum;
63
64 } km;
```

#### 6.3 点双连通分量

bcc.forest is a set of connected tree whose vertices are chequered with cut-vertex and BCC.

```
const bool BCC_VERTEX = 0, BCC EDGE = 1:
   struct BCC { // N = NO + MO. Remember to call init(&raw graph).
     Graph *g, forest; // g is raw graph ptr.
    int dfn[N], DFN, low[N];
    int stack[N], top;
     int expand_to[N];  // Where edge i is expanded to in expaned graph.
     // Vertex i expaned to i.
     int compress to[N]: // Where vertex i is compressed to.
     bool vertex_type[N], cut[N], compress_cut[N], branch[M];
     //std::vector<int> BCC component[N]; // Cut vertex belongs to none.
     inline void init(Graph *raw graph) {
12
      g = raw graph;
13
14
     void DFS(int u, int pe) {
       dfn[u] = low[u] = ++DFN; cut[u] = false;
       if (!~g->adj[u]) {
16
17
         cut[u] = 1;
18
         compress to[u] = forest.new node();
19
         compress cut[compress to[u]] = 1;
20
21
       for (int e = g->adj[u]; ~e; e = g->nxt[e]) {
22
         int v = g \rightarrow v[e];
         if ((e ^ pe) > 1 && dfn[v] > 0 && dfn[v] < dfn[u]) {
23
24
           stack[top++] = e:
25
           low[u] = std::min(low[u], dfn[v]);
26
27
         else if (!dfn[v]) {
28
           stack[top++] = e; branch[e] = 1;
29
           DFS(v. e):
30
           low[u] = std::min(low[v], low[u]);
```

```
if (low[v] >= dfn[u]) {
32
             if (!cut[u]) {
33
               cut[u] = 1;
               compress to[u] = forest.new node();
35
               compress cut[compress to[u]] = 1;
36
             int cc = forest.new node();
             forest.bi ins(compress to[u], cc);
39
             compress cut[cc] = 0:
             //BCC component[cc].clear();
41
42
               int cur e = stack[--top];
               compress to [expand to [cur el] = cc:
               compress to[expand to[cur e^1]] = cc;
               if (branch[cur e]) {
                 int v = g \rightarrow v[cur e];
                 if (cut[v])
                   forest.bi ins(cc, compress to[v]);
                    //BCC component[cc].push back(v);
                    compress to[v] = cc;
51
52
53
             } while (stack[top] != e);
55
56
         }
57
       }
58
     void solve() {
       forest.init(g->base):
       int n = g->n;
       for (int i = 0; i < g > e; i + +) {
63
         expand to[i] = g->new node();
64
       memset(branch, 0, sizeof(*branch) * g->e);
       memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0;
       for (int i = 0; i < n; i++)
         if (!dfn[i + g->base]) {
           top = 0;
           DFS(i + g->base, -1);
72
   } bcc:
   bcc.init(&raw graph);
76 bcc.solve():
  // Do something with bcc.forest ...
```

# 6.4 边双连通分量

```
struct BCC {
Graph *g, forest;
int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N], top, dfs_clock;
// tot[] is the size of each BCC, belong[] is the BCC that each node belongs to
pair<int, int > ori[M]; // bridge in raw_graph(raw node)
bool is_bridge[M];
```

```
inline void init(Graph *raw graph) {
8
       g = raw graph;
9
       memset(is bridge, false, sizeof(*is bridge) * g -> e);
10
       memset(vis + g \rightarrow base, 0, sizeof(*vis) * g \rightarrow n);
11
12
     void tarjan(int u, int from) {
13
       dfn[u] = low[u] = ++dfs clock; vis[u] = 1; stack[++top] = u;
14
       for (int p = g -> adj[u]; ~p; p = g -> nxt[p]) {
15
         if ((p ^ 1) == from) continue:
16
         int v = g \rightarrow v[p];
17
         if (vis[v]) {
18
            if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
19
         } else {
20
            tarjan(v, p);
21
            low[u] = min(low[u], low[v]);
22
            if (low[v] > dfn[u]) is bridge[p / 2] = true;
23
24
       7
25
       if (dfn[u] != low[u]) return;
26
       tot[forest.new node()] = 0;
27
28
         belong[stack[top]] = forest.n;
29
         vis[stack[top]] = 2;
30
         tot[forest.n]++;
31
         --top:
32
       } while (stack[top + 1] != u):
33
     }
34
     void solve() {
       forest.init(g -> base);
       int n = g \rightarrow n:
37
       for (int i = 0; i < n; ++i)
38
         if (!vis[i + g -> base]) {
39
            top = dfs clock = 0;
40
            tarjan(i + g \rightarrow base, -1);
41
42
       for (int i = 0; i < g -> e / 2; ++i)
43
         if (is bridge[i]) {
44
            int e = forest.e:
45
            forest.bi_ins(belong[g \rightarrow v[i * 2]], belong[g \rightarrow v[i * 2 + 1]], g \rightarrow w[i * 2]);
46
            ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i * 2]);
47
            ori[e + 1] = make pair(g -> v[i * 2], g -> v[i * 2 + 1]);
48
49
50 } bcc;
```

#### 6.5 最小树形图

```
const int MAXN,INF;// INF >= sum( W_ij )
int from[MAXN + 10] [MAXN * 2 + 10],n,m,edge[MAXN + 10] [MAXN * 2 + 10];
int sel[MAXN * 2 + 10],fa[MAXN * 2 + 10],vis[MAXN * 2 + 10];
int getfa(int x){if(x == fa[x]) return x; return fa[x] = getfa(fa[x]);}
void liuzhu(){ // 1-base: root is 1, answer = (sel[i], i) for i in [2..n]
fa[1] = 1;
for(int i = 2; i <= n; ++i){
    sel[i] = 1; fa[i] = i;</pre>
```

```
for(int j = 1; j <= n; ++j) if(fa[j] != i)</pre>
         if(from[i][i] = i, edge[sel[i]][i] > edge[i][i]) sel[i] = j;
    }
11
12
     int limit = n;
13
     while(1){
14
       int prelimit = limit; memset(vis, 0, sizeof(vis)); vis[1] = 1;
       for(int i = 2; i <= prelimit; ++i) if(fa[i] == i && !vis[i]){</pre>
         int j = i; while(!vis[j]) vis[j] = i, j = getfa(sel[j]);
         if(i == 1 || vis[i] != i) continue: vector<int> C: int k = i:
18
         do C.push back(k), k = getfa(sel[k]); while(k != j);
19
         ++limit:
         for(int i = 1; i <= n; ++i){
           edge[i][limit] = INF. from[i][limit] = limit:
21
22
23
         fa[limit] = vis[limit] = limit;
24
         for(int i = 0; i < int(C.size()); ++i){</pre>
25
           int x = C[i], fa[x] = limit;
           for(int j = 1; j \le n; ++ j)
27
             if(edge[j][x] != INF && edge[j][limit] > edge[j][x] - edge[sel[x]][x]){
28
               edge[j][limit] = edge[j][x] - edge[sel[x]][x];
29
               from[i][limit] = x:
30
             }
31
         for(int j=1;j<=n;++j) if(getfa(j)==limit) edge[j][limit] = INF;</pre>
33
         sel[limit] = 1:
34
         for(int i = 1: i \le n: ++i)
35
           if(edge[sel[limit]][limit] > edge[j][limit]) sel[limit] = j;
36
37
       if(prelimit == limit) break:
38
39
     for(int i = limit; i > 1; --i) sel[from[sel[i]][i]] = sel[i];
40 }
```

#### 66 帯花树

```
vector<int> link[maxn];
 2 int n,match[maxn],Queue[maxn],head,tail;
   int pred[maxn].base[maxn].start.finish.newbase;
   bool InQueue[maxn], InBlossom[maxn];
   void push(int u){ Queue[tail++]=u;InQueue[u]=true; }
   int pop(){ return Queue[head++]; }
   int FindCommonAncestor(int u.int v){
    bool InPath[maxn];
    for(int i=0;i<n;i++) InPath[i]=0;</pre>
    while(true){ u=base[u];InPath[u]=true;if(u==start) break;u=pred[match[u]]; }
    while(true){ v=base[v]:if(InPath[v]) break:v=pred[match[v]]: }
12
    return v:
13 }
   void ResetTrace(int u){
     while(base[u]!=newbase){
16
17
       v=match[u];
       InBlossom[base[u]]=InBlossom[base[v]]=true;
       u=pred[v]:
       if(base[u]!=newbase) pred[u]=v:
21
```

```
22 | }
   void BlossomContract(int u,int v){
     newbase=FindCommonAncestor(u,v);
25
     for (int i=0;i<n;i++)
     InBlossom[i]=0;
27
     ResetTrace(u);ResetTrace(v);
     if(base[u]!=newbase) pred[u]=v;
     if(base[v]!=newbase) pred[v]=u;
30
     for(int i=0:i<n:++i)</pre>
31
     if(InBlossom[base[i]]){
32
       base[i]=newbase;
33
       if(!InQueue[i]) push(i);
34
35 }
36
   bool FindAugmentingPath(int u){
37
     bool found=false;
     for(int i=0;i<n;++i) pred[i]=-1,base[i]=i;</pre>
     for (int i=0;i<n;i++) InQueue[i]=0;</pre>
     start=u;finish=-1; head=tail=0; push(start);
     while(head<tail){
42
       int u=pop();
43
       for(int i=link[u].size()-1;i>=0;i--){
44
         int v=link[u][i]:
45
         if(base[u]!=base[v]&&match[u]!=v)
46
           if(v==start||(match[v]>=0&&pred[match[v]]>=0))
47
             BlossomContract(u.v):
48
           else if(pred[v]==-1){
49
             pred[v]=u;
50
             if(match[v]>=0) push(match[v]);
51
             else{ finish=v: return true: }
52
53
54
     return found:
56
   void AugmentPath(){
    int u=finish.v.w:
59
    while(u>=0){ v=pred[u]; w=match[v]; match[v]=u; match[u]=v; u=w; }
60 }
   void FindMaxMatching(){
    for(int i=0;i<n;++i) match[i]=-1;</pre>
    for(int i=0:i<n:++i) if(match[i]==-1) if(FindAugmentingPath(i)) AugmentPath();</pre>
63
64 }
```

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#### 6.7 Dominator Tree

```
vector<int> prec[N], succ[N];
vector<int> ord;
int stamp, vis[N];
int num[N];
int fa[N];
void dfs(int u) {
   vis[u] = stamp;
   num[u] = ord.size();
   ord.push_back(u);
```

```
for (int i = 0; i < (int)succ[u].size(); ++i) {</pre>
       int v = succ[u][i]:
       if (vis[v] != stamp) {
         fa[v] = u;
13
         dfs(v);
15
       }
16
    }
   int fs[N]. mins[N]. dom[N]. sem[N]:
   int find(int u) {
     if (u != fs[u]) {
21
       int v = fs[u];
       fs[u] = find(fs[u]):
       if (mins[v] != -1 && num[sem[mins[v]]] < num[sem[mins[u]]]) {
24
         mins[u] = mins[v];
25
26
    }
27
    return fs[u];
   void merge(int u, int v) { fs[u] = v; }
   vector<int> buf[N]:
31 int buf2[N]:
   void mark(int source) {
     ord.clear();
     ++stamp;
     dfs(source):
     for (int i = 0; i < (int)ord.size(); ++i) {</pre>
       int u = ord[i];
37
       fs[u] = u, mins[u] = -1, buf2[u] = -1;
39
     for (int i = (int)ord.size() - 1; i > 0; --i) {
       int u = ord[i], p = fa[u];
41
       sem[u] = p:
42
       for (int j = 0; j < (int)prec[u].size(); ++j) {</pre>
43
         int v = prec[u][i];
         if (use[v] != stamp) continue;
         if (num[v] > num[u]) {
47
           find(v): v = sem[mins[v]]:
48
         if (num[v] < num[sem[u]]) {</pre>
           sem[u] = v;
         }
51
52
       }
53
       buf[sem[u]].push_back(u);
54
       mins[u] = u;
       merge(u, p);
       while (buf[p].size()) {
         int v = buf[p].back();
57
58
         buf[p].pop back();
         find(v):
         if (sem[v] == sem[mins[v]]) {
           dom[v] = sem[v];
         } else {
           buf2[v] = mins[v];
64
```

# 6.8 无向图最小割

```
1 int cost[maxn] [maxn], seq[maxn], len[maxn], n, m, pop, ans;
2 bool used[maxn];
   void Init(){
    int i,j,a,b,c;
     for(i=0;i<n;i++) for(j=0;j<n;j++) cost[i][j]=0;
     for(i=0;i<m;i++){</pre>
       scanf("%d %d %d",&a,&b,&c); cost[a][b]+=c; cost[b][a]+=c;
     pop=n; for(i=0;i<n;i++) seq[i]=i;
10 }
11 | void Work(){
     ans=inf; int i,j,k,l,mm,sum,pk;
     while(pop > 1){
       for(i=1;i<pop;i++) used[seq[i]]=0; used[seq[0]]=1;</pre>
14
15
       for(i=1;i<pop;i++) len[seq[i]]=cost[seq[0]][seq[i]];</pre>
16
       pk=0; mm=-inf; k=-1;
17
       for(i=1;i<pop;i++) if(len[seq[i]] > mm){ mm=len[seq[i]]; k=i; }
18
       for(i=1;i<pop;i++){</pre>
19
         used[seq[l=k]]=1;
20
         if(i==pop-2) pk=k;
21
         if(i==pop-1) break;
22
         mm=-inf:
23
         for(j=1;j<pop;j++) if(!used[seq[j]])</pre>
            if((len[seq[j]]+=cost[seq[1]][seq[j]]) > mm)
24
25
              mm=len[seq[j]], k=j;
       }
26
27
28
       for(i=0;i<pop;i++) if(i != k) sum+=cost[seq[k]][seq[i]];</pre>
29
       ans=min(ans.sum):
30
       for(i=0;i<pop;i++)</pre>
31
         cost[seq[k]][seq[i]]=cost[seq[i]][seq[k]]+=cost[seq[pk]][seq[i]];
32
       seq[pk]=seq[--pop];
33
34
     printf("%d\n",ans);
```

#### 6.9 重口味费用流

```
int S, T, totFlow, totCost;
int dis[N], slack[N], visit[N];
int modlable () {
```

```
int delta = INF:
       for (int i = 1; i <= T; i++) {
           if (!visit[i] && slack[i] < delta) delta = slack[i];</pre>
           slack[i] = INF:
11
       if (delta == INF) return 1;
       for (int i = 1; i <= T; i++)
           if (visit[i]) dis[i] += delta;
15 }
16
   int dfs (int x, int flow) {
      if (x == T) {
           totFlow += flow:
           totCost += flow * (dis[S] - dis[T]);
21
           return flow;
      }
22
23
       visit[x] = 1;
       int left = flow:
       for (int i = e.last[x]; ~i; i = e.succ[i])
           if (e.cap[i] > 0 && !visit[e.other[i]]) {
27
               int y = e.other[i];
28
               if (dis[y] + e.cost[i] == dis[x]) {
                   int delta = dfs (y, min (left, e.cap[i]));
                   e.cap[i] -= delta;
                   e.cap[i ^ 1] += delta:
                   left -= delta;
                   if (!left) { visit[x] = 0; return flow; }
33
                   slack[v] = min (slack[v], dis[v] + e.cost[i] - dis[x]):
37
38
       return flow - left:
39 }
   pair <int, int> minCost () {
       totFlow = 0; totCost = 0;
       fill (dis + 1, dis + T + 1, 0):
       do {
45
               fill (visit + 1, visit + T + 1, 0);
           } while (dfs (S, INF)):
      } while (!modlable ());
49
       return make_pair (totFlow, totCost);
50 }
```

```
6.10 2-SAT

int S, T, totFlow, totCost;

int dis[N], slack[N], visit[N];

int modlable () {
    int delta = INF;
    for (int i = 1; i <= T; i++) {
        if (!visit[i] && slack[i] < delta) delta = slack[i];
}</pre>
```

```
slack[i] = INF;
10
      }
11
       if (delta == INF) return 1;
12
       for (int i = 1; i \le T; i++)
13
           if (visit[i]) dis[i] += delta;
14
       return 0:
15 }
16
   int dfs (int x. int flow) {
       if (x == T) {
19
           totFlow += flow;
           totCost += flow * (dis[S] - dis[T]);
21
           return flow:
22
      }
23
       visit[x] = 1;
24
       int left = flow;
       for (int i = e.last[x]; ~i; i = e.succ[i])
26
           if (e.cap[i] > 0 && !visit[e.other[i]]) {
27
               int y = e.other[i];
28
               if (dis[y] + e.cost[i] == dis[x]) {
29
                   int delta = dfs (y, min (left, e.cap[i]));
30
                   e.cap[i] -= delta;
31
                   e.cap[i ^ 1] += delta;
32
                   left -= delta;
33
                   if (!left) { visit[x] = 0: return flow: }
34
35
                   slack[y] = min (slack[y], dis[y] + e.cost[i] - dis[x]);
36
37
38
       return flow - left:
39
40
   pair <int, int> minCost () {
       totFlow = 0: totCost = 0:
       fill (dis + 1, dis + T + 1, 0);
       do {
45
           do {
46
               fill (visit + 1, visit + T + 1, 0);
47
           } while (dfs (S, INF));
48
       } while (!modlable ());
       return make_pair (totFlow, totCost);
```

# 7 其他

# 7.1 Dancing Links

```
struct Node {
   Node *1, *r, *u, *d, *col;
   int size, line_no;
   Node() {
       size = 0; line_no = -1;
       l = r = u = d = col = NULL;
   }
} *root;
```

```
void cover(Node *c) {
     c->1->r = c->r; c->r->1 = c->1;
     for (Node *u = c->d; u != c; u = u->d)
       for (Node *v = u - r; v != u; v = v - r) {
14
         v->d->u = v->u:
15
         v->u->d = v->d;
16
         -- v->col->size;
17
18 }
19
   void uncover(Node *c) {
     for (Node *u = c -> u : u != c : u = u -> u) {
21
       for (Node *v = u->1; v != u; v = v->1) {
23
         ++ v->col->size;
24
         v\rightarrow u\rightarrow d = v;
25
         v\rightarrow d\rightarrow u = v:
26
27
    }
28
     c->l->r = c; c->r->l = c;
29 }
30
   std::vector<int> answer;
   bool search(int k) {
     if (root->r == root) return true:
     Node *r = NULL:
     for (Node *u = root->r; u != root; u = u->r)
       if (r == NULL || u->size < r->size)
     if (r == NULL || r->size == 0) return false;
39
     else {
40
       cover(r);
41
       bool succ = false;
42
       for (Node *u = r -> d: u != r && !succ: u = u -> d) {
         answer.push_back(u->line_no);
         for (Node *v = u \rightarrow r; v != u; v = v \rightarrow r) // Cover row
45
            cover(v->col):
46
         succ |= search(k + 1):
47
         for (Node *v = u->1; v != u; v = v->1)
48
            uncover(v->col);
49
         if (!succ) answer.pop_back();
50
51
       uncover(r);
52
       return succ;
53
54 | }
55
   bool entry[CR][CC];
   Node *who[CR][CC];
   int cr, cc;
   void construct() {
     root = new Node();
     Node *last = root:
     for (int i = 0; i < cc; ++ i) {
```

```
64
        Node *u = new Node():
 65
        last->r = u; u->l = last;
 66
        Node *v = u; u->line_no = i;
 67
        last = u;
 68
        for (int j = 0; j < cr; ++ j)
 69
          if (entry[j][i]) {
 70
            ++ u->size;
 71
            Node *cur = new Node();
 72
            who[i][i] = cur:
 73
            cur->line no = j;
 74
            cur->col = u;
 75
            cur->u = v; v->d = cur;
 76
            v = cur:
 77
 78
        v->d = u; u->u = v;
 79
      last->r = root; root->l = last;
 81
      for (int j = 0; j < cr; ++ j) {
       Node *last = NULL;
 83
        for (int i = cc - 1; i \ge 0; -- i)
          if (entry[j][i]) {
 84
 85
            last = who[j][i];
86
            break;
87
 88
        for (int i = 0; i < cc; ++ i)
 89
          if (entrv[i][i]) {
 90
            last->r = who[j][i];
            who[j][i]->1 = last;
91
 92
            last = who[j][i];
 93
 94
 95
 96
97
    void destruct() {
      for (Node *u = root->r; u != root; ) {
       for (Node *v = u->d; v != u; ) {
100
          Node *nxt = v->d;
101
          delete(v):
102
          v = nxt;
103
104
        Node *nxt = u->r;
105
        delete(u): u = nxt:
106
107
      delete root;
108 }
```

#### 7.2 蔡勒公式

0 for Sunday. Day and month is 1-based.

```
int zeller(int y,int m,int d) {
   if (m<=2) y--,m+=12; int c=y/100; y%=100;
   int w=((c>>2)-(c<<1)+y+(y>>2)+(13*(m+1)/5)+d-1)%7;
   if (w<0) w+=7; return(w);
}</pre>
```

# 8 技巧

# 8.1 真正的释放 STL 容器内存空间

```
template <typename T>

__inline void clear(T& container) {
    container.clear(); // 或者删除了一堆元素
    T(container).swap(container);
}
```

#### 8.2 无敌的大整数相乘取模

Time complexity O(1).

```
1 // 需要保证 x 和 y 非负
2 long long mult(long long x, long long y, long long MODN) {
3 long long t = (x * y - (long long)((long double)x / MODN * y + 1e-3) * MODN) % MODN;
4 return t < 0 ? t + MODN : t;
5 }
```

# 8.3 无敌的读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化
 2 // 用 isdigit() 会小幅变慢
 3 // 返回 false 表示读到文件尾
 4 namespace Reader {
       const int L = (1 << 15) + 5;
       char buffer[L], *S, *T;
       __inline bool getchar(char &ch) {
          if (S == T) {
              T = (S = buffer) + fread(buffer, 1, L, stdin);
              if (S == T) {
           ch = EOF;
          return false;
15
       ch = *S++;
       return true;
17
       __inline bool getint(int &x) {
       char ch; bool neg = 0;
       for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^= ch == '-';
21
       if (ch == EOF) return false;
22
       x = ch - '0';
      for (; getchar(ch), ch >= '0' && ch <= '9'; )</pre>
        x = x * 10 + ch - '0';
      if (neg) x = -x;
       return true;
27
```

# 8.4 梅森旋转算法

High quality pseudorandom number generator, twice as efficient as rand() with -02. C++11 required.

```
1 #include <random>
2
```

```
3 int main() {
4    std::mt19937 g(seed); // std::mt19937_64
5    std::cout << g() << std::endl;
6 }</pre>
```

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# 9 提示

### 9.1 控制 cout 输出实数精度

```
std::cout << std::fixed << std::setprecision(5);</pre>
```

#### 9.2 vimrc

```
set nu
set sw=4
set ts=4
set ts=4
syntax on
set cindent
```

# 9.3 让 make 支持 c + 11

In .bashrc or whatever:

```
export CXXFLAGS='-std=c++11 -Wall'
```

#### 9.4 tuple 相关

#### 9.5 线性规划转对偶

```
maximize \mathbf{c}^T \mathbf{x}
subject to \mathbf{A} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq 0 \iff \text{minimize } \mathbf{y}^T \mathbf{b}
subject to \mathbf{y}^T \mathbf{A} \geq \mathbf{c}^T, \mathbf{y} \geq 0
```

#### 9.6 32-bit/64-bit 随机素数

32-bit	64-bit
73550053	1249292846855685773
148898719	1701750434419805569
189560747	3605499878424114901
459874703	5648316673387803781
1202316001	6125342570814357977
1431183547	6215155308775851301
1438011109	6294606778040623451
1538762023	6347330550446020547
1557944263	7429632924303725207
1981315913	8524720079480389849

#### 9.7 NTT 素数及其原根

Prime	Primitive root
1053818881	7
1051721729	6
1045430273	3
1012924417	5
1007681537	3

```
9.8 Java Hints
 1 import iava.io.*:
   import java.util.*;
   import java.math.*;
   public class Main {
    static int get(char c) {
       if (c <= '9')
         return c - '0';
       else if (c \le 'Z')
         return c - 'A' + 10;
11
12
         return c - 'a' + 36;
13
     static char get(int x) {
       if (x <= 9)
         return (char)(x + '0');
16
       else if (x \le 35)
         return (char)(x - 10 + 'A'):
20
         return (char)(x - 36 + 'a');
21
     static BigInteger get(String s, BigInteger x) {
22
23
       BigInteger ans = BigInteger.valueOf(0), now = BigInteger.valueOf(1);
       for (int i = s.length() - 1; i >= 0; i--) {
         ans = ans.add(now.multiply(BigInteger.valueOf(get(s.charAt(i)))));
25
26
         now = now.multiplv(x):
27
      }
28
       return ans;
29
     public static void main(String [] args) {
       Scanner cin = new Scanner(new BufferedInputStream(System.in));
32
       for (; ; ) {
33
         BigInteger x = cin.nextBigInteger();
         if (x.compareTo(BigInteger.valueOf(0)) == 0)
           break:
```

```
36
        String s = cin.next(), t = cin.next(), r = "";
37
        BigInteger ans = get(s, x).mod(get(t, x));
38
        if (ans.compareTo(BigInteger.valueOf(0)) == 0)
39
40
        for (; ans.compareTo(BigInteger.valueOf(0)) > 0;) {
41
          r = get(ans.mod(x).intValue()) + r;
42
           ans = ans.divide(x);
43
44
        System.out.println(r):
45
      }
46
47
49 // Arrays
50 int a[];
51 .fill(a[, int fromIndex, int toIndex], val); | .sort(a[, int fromIndex, int toIndex])
52 // String
53 String s:
   .charAt(int i); | compareTo(String) | compareToIgnoreCase () | contains(String) |
55 length () | substring(int 1, int len)
56 // BigInteger
57 | .abs() | .add() | bitLength () | subtract () | divide () | remainder () |
     58 pow(int) | multiply () | compareTo () |
59 gcd() | intValue () | longValue () | isProbablePrime(int c) (1 - 1/2^c) |
60 nextProbablePrime () | shiftLeft(int) | valueOf ()
61 // BigDecimal
62 ROUND CEILING | ROUND DOWN FLOOR | ROUND HALF DOWN | ROUND HALF EVEN | ROUND HALF UP |
   .divide(BigDecimal b, int scale , int round mode) | doubleValue () | movePointLeft(int)
     64 setScale(int scale , int round mode) | stripTrailingZeros ()
65 // StringBuilder
66 StringBuilder sb = new StringBuilder ();
67 | sb.append(elem) | out.println(sb)
```