

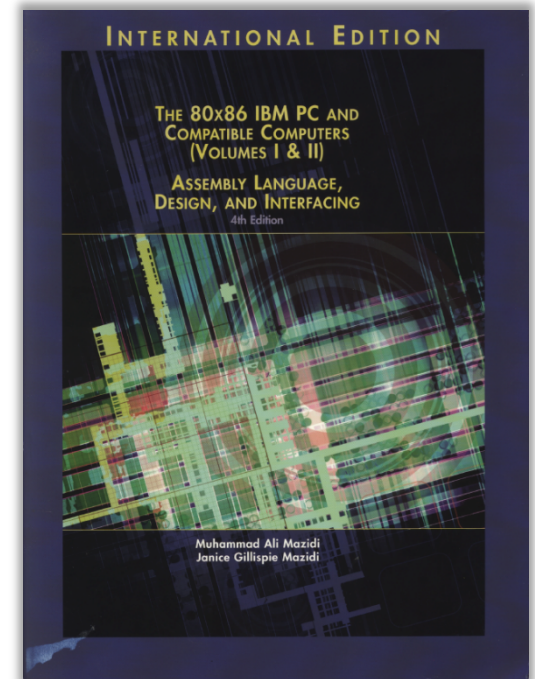
Lecture 04: Assembly Language Programming (1)

Reference Book:

The 80x86 IBM PC and Compatible Computers

Chapter 2

Assembly Language Programming



Programming Languages

- Machine language

- Binary code for CPU but not human beings

How to convert your program in low/high-level languages into machine language?

- *Low-level language*: deals with the internal structure of a CPU
 - Hard to program, poor portability but very efficient
 - BASIC, Pascal, C, Fortran, Perl, TCL, Python, ...
 - *High-level languages*: do not have to be concerned with the internal details of a CPU
 - Easy to program, good portability but less efficient

Assembly Language Programs

- A series of *statements*
 - *Assembly language instructions*
 - | Perform the real work of the program
 - *Directives (pseudo-instructions)*
 - | Give instructions **for the *assembler*** program about how to translate the program into machine code.
- Consists of multiple segments
 - But CPU can access only one data segment, one code segment, one stack segment and one extra segment (***Why?***)

Form of an statement

- `[label:] mnemonic [operands] [;comment]`
 - *label* is a reference to this statement
 - Rules for names: each label must be unique; letters, 0-9, (?), (.), (@), (_, and (\$) ; first character cannot be a digit; less than 31 characters
 - ":" is needed if it is an instruction otherwise omitted
 - ";" leads a comment, the assembler omits anything on this line following a semicolon

Shell of a Real Program

- Full segment definition (old fashion)
 - See an example later
- Simplified segment definition

```
;THE FORM OF AN ASSEMBLY LANGUAGE PROGRAM
;NOTE: USING SIMPLIFIED SEGMENT DEFINITION

.MODEL SMALL
.STACK 64
.DATA
DATA1 DB 52H
DATA2 DB 29H
SUM DB ?
.CODE
MAIN PROC FAR ;this is the program entry point
MOV AX,@DATA ;load the data segment address
MOV DS,AX ;assign value to DS
MOV AL,DATA1 ;get the first operand
MOV BL,DATA2 ;get the second operand
ADD AL,BL ;add the operands
MOV SUM,AL ;store the result in location SUM
MOV AH,4CH ;set up to return to DOS
INT 21H ;
MAIN ENDP
END MAIN ;this is the program exit point
```

Figure 2-1. Simple Assembly Language Program

Model Definition

■ The **MODEL** directive

- Selects the size of the memory model
- SMALL: code \leq 64KB, data \leq 64KB
- MEDIUM: data \leq 64KB, code $>$ 64KB
- COMPACT: code \leq 64KB, data $>$ 64KB
- LARGE: data $>$ 64KB but single set of data $<$ 64KB, code $>$ 64KB
- HUGE: data $>$ 64KB, code $>$ 64KB
- TINY: code + data $<$ 64KB

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Figure 2-1. Simple Assembly Language Program

Simplified Segment Definition

- Simplified segment definition
 - **.CODE, .DATA, .STACK**
 - Only three segments can be defined
 - Automatically correspond to the CPU's CS, DS, SS
 - DOS determines the CS and SS segment registers automatically. DS (and ES) has to be manually specified.

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Figure 2-1. Simple Assembly Language Program

Segments All at a Glance

- Stack segment
- Data segment
 - Data definition
- Code segment
 - Write your statements
 - Procedures definition
 - label **PROC** [**FAR**|**NEAR**]
 - label **ENDP**
 - Entrance proc should be FAR

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Figure 2-1. Simple Assembly Language Program

Full Segment Definition

- Full segment definition
label **SEGMENT**
label **ENDS**
 - You name those labels
 - as many as needed
 - DOS assigns CS, SS
 - Program assigns DS (manually load data segments) and ES

```
DaSeg1 segment
    str1 db 'Hello World! $'
DaSeg1 ends

StSeg segment
    dw 128 dup(0)
StSeg ends

CoSeg segment
    start proc far
        assume cs:CoSeg, ss:StSeg

        mov ax, DaSeg1    ; set segment registers:
        mov ds, ax
        mov es, ax

        call subr          ; call subroutine

        mov ah, 1          ; wait for any key....
        int 21h

        mov ah, 4ch         ; exit to operating system.
        int 21h
    start endp

    subr proc

        mov dx, offset str1
        mov ah, 9
        int 21h            ; output string at ds:dx

        ret
    subr endp

CoSeg ends

    end start    ; set entry point and stop the assembler.
```

Program Execution

- Program starts from the entrance
 - Ends whenever calls 21H interruption with AH = 4CH
- Procedure caller and callee
 - **CALL** *procedure*
 - **RET**

```
DaSeg1 segment
    str1 db 'Hello World! $'
DaSeg1 ends
```

```
StSeg segment
    dw 128 dup(0)
```

```
CoSeg segment
    start proc far
        assume cs:CoSeg, ss:StSeg

        mov ax, DaSeg1      ; set segment registers:
        mov ds, ax
        mov es, ax

        call subr            ; call subroutine

        mov ah, 1           ; wait for any key....
        int 21h

        mov ah, 4ch         ; exit to operating system.
        int 21h

    start endp

    subr proc

        mov dx, offset str1
        mov ah, 9
        int 21h             ; output string at ds:dx

        ret
    subr endp

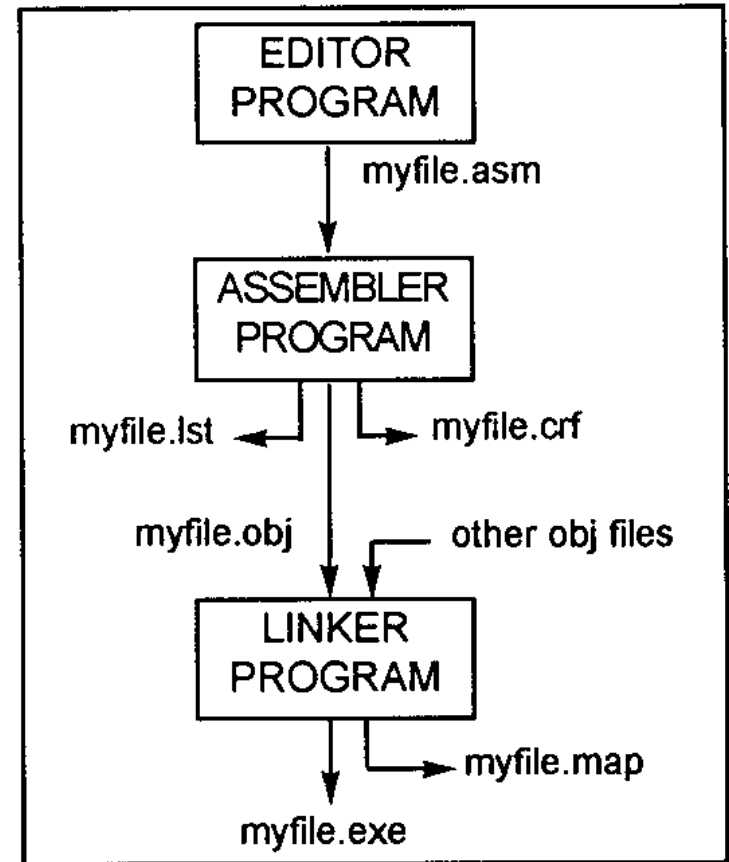
CoSeg ends

    end start                ; set entry point and stop the assembler.
```

Build up Your Program

C>MASM A:MYFILE.ASM <enter>

C>LINK A:MYFILE.OBJ <enter>



Control Transfer Instructions

■ Range

■ **SHORT**, *intra*segment

- | IP changed: one-byte range (-128~127)

■ **Near**, *intra*segment

- | IP changed: two-bytes range (-32768~32767)
- | If control is transferred within the same code segment

■ **FAR**, *inter*segment

- | CS and IP all changed
- | If control is transferred outside the current code segment

■ Jumps

■ CALL statement

Conditional Jumps

- Jump according to the value of the flag register
- Short jumps

Mnemonic	Condition Tested	“Jump IF ...”
JA/JNBE	$(CF = 0) \text{ and } (ZF = 0)$	above/not below nor zero
JAE/JNB	$CF = 0$	above or equal/not below
JB/JNAE	$CF = 1$	below/not above nor equal
JBE/JNA	$(CF \text{ or } ZF) = 1$	below or equal/not above
JC	$CF = 1$	carry
JE/JZ	$ZF = 1$	equal/zero
JG/JNLE	$((SF \text{ xor } OF) \text{ or } ZF) = 0$	greater/not less nor equal
JGE/JNL	$(SF \text{ xor } OF) = 0$	greater or equal/not less
JL/JNGE	$(SF \text{ xor } OF) = 1$	less/not greater nor equal
JLE/JNG	$((SF \text{ xor } OF) \text{ or } ZF) = 1$	less or equal/not greater
JNC	$CF = 0$	not carry
JNE/JNZ	$ZF = 0$	not equal/not zero
JNO	$OF = 0$	not overflow
JNP/JPO	$PF = 0$	not parity/parity odd
JNS	$SF = 0$	not sign
JO	$OF = 1$	overflow
JP/JPE	$PF = 1$	parity/parity equal
JS	$SF = 1$	sign

Unconditional Jumps

- JMP [**SHORT**|**NEAR**|**FAR** PTR] *label*
- Near by default

Subroutines & CALL Statement

- Range
 - **NEAR**: procedure is defined within the same code segment with the caller
 - **FAR**: procedure is defined outside the current code segment of the caller
- **PROC & ENDP** are used to define a subroutine
- **CALL** is used to call a subroutine
 - **RET** is put at the end of a subroutine
 - *Difference between a far and a near call?*

Calling a NEAR proc

- ✓ The CALL instruction and the subroutine it calls are in the same segment.
- ✓ Save the current value of the IP on the stack.
- ✓ load the subroutine's offset into IP (nextinst + offset)

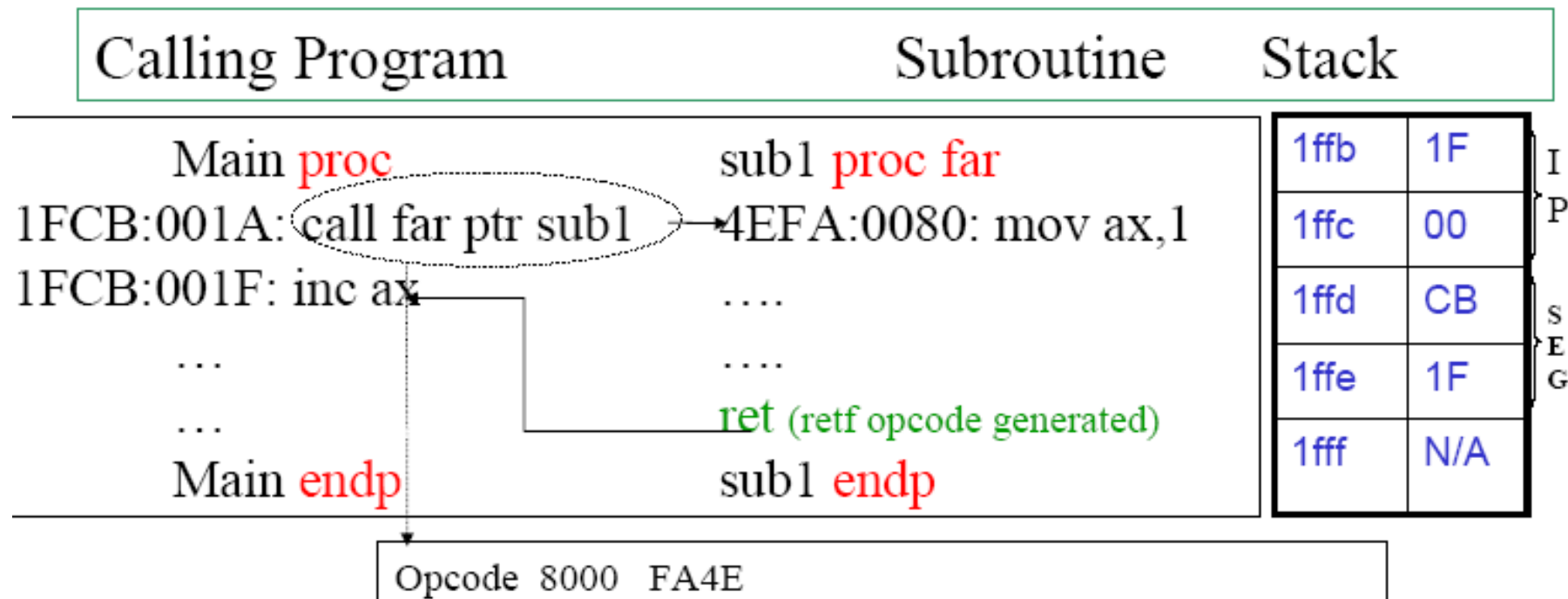
Calling Program	Subroutine	Stack
-----------------	------------	-------

Main proc	sub1 proc
001A: call sub1	0080: mov ax,1
001D: inc ax	...
.	ret
Main endp	sub1 endp

1ffd	1D
1ffe	00
1fff	(not used)

Calling a FAR proc

- ✓ The CALL instruction and the subroutine it calls are in the “Different” segments.
- ✓ Save the current value of the CS and IP on the stack.
- ✓ Then load the subroutine’s CS and offset into IP.



Data Types & Definition

- CPU can process either 8-bit or 16 bit ops
 - What if your data is bigger?
- Directives
 - **ORG:** indicates the beginning of the offset address
 - E.g., `ORG 10H`
 - Define variables:
 - **DB:** allocate byte-size chunks
 - E.g., `x DB 12` | `y DB 23H,48H` | `z DB 'Good Morning!'` |
`str DB "I'm good!"`
 - **DW, DD, DQ**
 - **EQU:** define a constant
 - E.g., `NUM EQU 234`
 - **DUP:** duplicate a given number of characters
 - E.g., `x DB 6 DUP(23H)` | `y DW 3 DUP(0FF10H)`

More about Variables

- For variables, they may have names
 - E.g., `luckyNum DB 27H`, `time DW 0FFFFH`
- Variable names have three attributes:
 - **Segment value** } Logical address
 - **Offset address** }
 - **Type:** how a variable can be accessed (e.g., DB is byte-wise, DW is word-wise)
- Get the segment value of a variable
 - Use **SEG** directive (E.g., `MOV AX, SEG luckyNum`)
- Get the offset address of a variable
 - Use **OFFSET** directive, or **LEA** instruction
 - E.g., `MOV AX, OFFSET time`, or `LEA AX, time`

More about Labels

■ Label definition:

■ Implicitly:

■ E.g., `AGAIN: ADD AX, 03423H`

■ Use **LABEL** directive:

■ E.g., `AGAIN LABEL FAR`

`ADD AX, 03423H`

■ Labels have three attributes:

- **Segment value:**
 - **Offset address:**
 - **Type:** range for jumps, NEAR, FAR
- } Logical address

More about the PTR Directive

- Temporarily change the type (range) attribute of a variable (label)
 - To guarantee that both operands in an instruction match
 - To guarantee that the jump can reach a label
- E.g.,

```
DATA1  DB  10H,20H,30H  ;  
DATA2  DW  4023H,0A845H  
  
.....  
MOV    BX, WORD PTR DATA1 ; 2010H -> BX  
MOV    AL, BYTE PTR DATA2 ; 23H -> AL  
MOV    WORD PTR [BX], 10H ; [BX], [BX+1] ← 0010H  
  
E.g.,    JMP FAR PTR aLabel
```

.COM Executable

- One segment in total
 - Put data and code all together
 - Less than 64KB

```
TITLE  PROG2-4 COM PROGRAM TO ADD TWO WORDS
PAGE   60,132
CODSG  SEGMENT
      ORG 100H
      ASSUME CS:CODSG,DS:CODSG,ES:CODSG
;-----THIS IS THE CODE AREA
PROGCODE PROC NEAR
      MOV  AX,DATA1    ;move the first word into AX
      MOV  SUM,AX      ;move the sum
      MOV  AH,4CH      ;return to DOS
      INT  21H
PROGCODE ENDP
;-----THIS IS THE DATA AREA
DATA1  DW  2390
DATA2  DW  3456
SUM     DW  ?
;-----
CODSG  ENDS
      END  PROGCODE
```

```
TITLE  PROG2-5 COM PROGRAM TO ADD TWO WORDS
PAGE   60,132
CODSG  SEGMENT
      ASSUME CS:CODSG,DS:CODSG,ES:CODSG
      ORG 100H
START: JMP  PROGCODE   ;go around the data area
;-----THIS IS THE DATA AREA
DATA1  DW  2390
DATA2  DW  3456
SUM     DW  ?
;-----THIS IS THE CODE AREA
PROGCODE: MOV  AX,DATA1    ;move the first word into AX
          ADD  AX,DATA1    ;add the second word
          MOV  SUM,AX      ;move the sum
          MOV  AH,4CH
          INT  21H
;-----
CODSB  ENDS
      END  START
```