

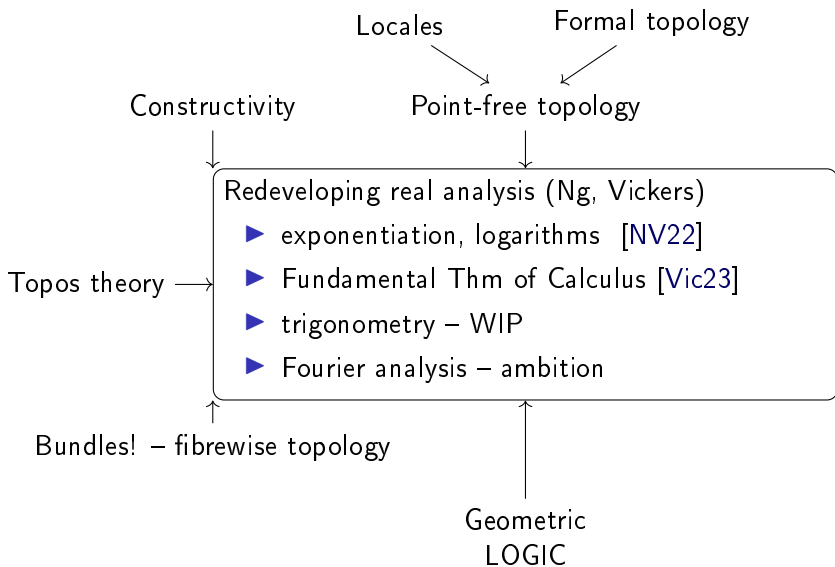
Real analysis via logic

Steve Vickers

School of Computer Science
University of Birmingham

British Logic Colloquium, Birmingham, 7 September 2024

- ▶ Topology and continuity as *logical* phenomena (geometric logic).
- ▶ Now applying logical approach to real analysis.
- ▶ Gives fresh perspective on the analysis.



Topology, continuity as *logical* phenomena

Usual approach: maths is discrete

Topological space (point-set) = set + extra structure

Continuous map = function respecting that structure

Geometric logic: “all” maths is continuous

Discrete maths of sets very restricted – eg \mathbb{Q} is a set, real line \mathbb{R} isn't. Can still access \mathbb{R} , by taking logical, *point-free* approach –

Topological space = logical theory (point = model of theory)

Continuity of map = definition respects logical constraints

Example: real line \mathbb{R} as logical theory

Signature

For each rational $q \in \mathbb{Q}$: two propositional symbols $[\cdot < q]$, $[q < \cdot]$.

Axioms – eg

$$\begin{aligned} [\cdot < q'] \vdash [\cdot < q] \quad (\text{if } q' < q) \\ [\cdot < q] \wedge [q < \cdot] \vdash \perp \\ \top \vdash [q < \cdot] \vee [\cdot < r] \quad (\text{if } q < r) \\ : \end{aligned}$$

Model x is real number as *Dedekind cut*:

Specify truth values $[x < q]$ and $[q < x]$ for every q ,
ie which rationals are bigger than x , which are smaller.

Example: real line \mathbb{R} as logical theory

Signature

For each rational $q \in \mathbb{Q}$: two propositional symbols $[\cdot < q]$, $[q < \cdot]$.

Axioms – eg

$$\begin{aligned} & [\cdot < q'] \vdash [\cdot < q] \quad (\text{if } q' < q) \\ & [\cdot < q] \wedge [q < \cdot] \vdash \perp \\ & \top \vdash [q < \cdot] \vee [\cdot < r] \quad (\text{if } q < r) \\ & \quad : \\ & [\cdot < q] \vdash \bigvee_{q' < q} [\cdot < q'] \end{aligned}$$

Model x is real number as *Dedekind cut*:

Specify truth values $[x < q]$ and $[q < x]$ for every q ,
ie which rationals are bigger than x , which are smaller.

Defining maps

Think of maps $f: \mathbb{R} \rightarrow \mathbb{R}$ in a style of programming languages.

Let $x: \mathbb{R}$

:

:

$f(x): \mathbb{R} := \dots$

:

Declare formal parameter x .

Do some auxiliary calculations.

Define result $f(x)$ as model:

- ▶ specify truth values $[f(x) < q]$, $[q < f(x)]$,
- ▶ prove that axioms hold.

eg absolute value $|\cdot|: \mathbb{R} \rightarrow \mathbb{R}$

Let $x: \mathbb{R}$

$[|x| < q] := [x < q] \wedge [-q < x]$

$[q < |x|] := [q < x] \vee [x < -q]$

... and prove axioms

Inside the box, in scope of x , is a *different mathematics!*

1. Lots of non-standard truth values

$[x < q]$, $[q < x]$ (for each rational q)

Each expresses *where* (ie for which models x) something is true.

2. Continuity = different logic

Continuity: inverse image of open is open

$f^{-1}([x < q]) = [f(x) < q]$ is made from truth values of form $[x < r]$ and $[s < x]$ using \wedge and \vee .

Similarly for $[q < f(x)]$.

We want continuity, therefore restrict mathematics inside box to limit how we construct $f(x)$.

Geometricity

Pure logic

Restricted formulae: $\forall, \wedge, =, \exists$.

Axioms as sequents:

formula \vdash_{context} formula

Context = finite stock of free variables, with implicit \forall .

Geometricity

Pure logic

Restricted formulae: $\vee, \wedge, =, \exists$.

Axioms as sequents:

formula \vdash_{context} formula

Context = finite stock of free variables, with implicit \forall .

Infinite \vee can often be avoided by using \exists with an infinite set. eg

$$[\cdot < q] \vdash_{q:\mathbb{Q}} (\exists q':\mathbb{Q})(q' < q \wedge [\cdot < q']) \text{ for } [\cdot < q] \vdash \bigvee_{q' < q} [\cdot < q'].$$

Corresponding maths

Restricted maths of sets:

Disjoint unions, quotients, finite products, equalizers, free algebras.

Function spaces Y^X , powersets $\mathcal{P}X$, the real line \mathbb{R} are *not sets!* They must be dealt with as spaces.

Technicalities

- ▶ The “maths inside the box” is the geometric fragment of the internal mathematics of the classifying topos $\mathcal{S}[X]$.
“Map” = geometric morphism.
- ▶ “Classifying topos” is slippery constructively – depends on choice of a base topos \mathcal{S} . To avoid that dependency, work without infinite disjunctions. [Vic17]

[Vic99] shows the technique in action in domain theory.

[Vic07] explains how standard topos results (eg [MLM92]) arrive at this point of view.

[Vic22] gives a more up-to-date discussion.

Logical manipulations I: Decomposing theories

Theory of Dedekind reals =

theory of *lower reals* to account for $[q < \cdot]$

+ theory of *upper reals* to account for $[\cdot < q]$

+ two axioms to relate them

Good strategy for point-free analysis (eg exp, log, integration)

1. Deal with lower and upper cases separately,
2. then combine them.

Can provide fresh insights – eg Ostrowski's Theorem in number theory (Ng [Ng22, NV]).

Logical manipulations II: Building up theories

Let $x:\mathbb{R}$

:

:

Space $\mathbb{T}(x) := \dots$

:

Declare formal parameter x .

Do some auxiliary calculations.

Define space $\mathbb{T}(x)$ as theory:

► define signature (as set)

► define axiom set (with appropriate structure).

Geometricity suggests $\mathbb{T}(x)$ depends *continuously* on x .

Logically – $(\mathbb{T}(x))_{x:\mathbb{R}}$ defines extension of theory of reals.

Models = pairs (x, y) , $x:\mathbb{R}$, y model of $\mathbb{T}(x)$.

Dependent type theory: write $\sum_{x:\mathbb{R}} \mathbb{T}(x)$.

Forgetful map $\sum_{x:\mathbb{R}} \mathbb{T}(x) \rightarrow \mathbb{R}$, $(x, y) \mapsto x$, makes a *bundle* over \mathbb{R} .

None of this works satisfactorily in point-set topology

Can't describe bundle as continuously indexed family of spaces.

Logical approach works better!

Logical manipulations III: Modal logic \rightarrow hyperspaces

Hyperspace¹ = space of subspaces

eg Use \Box modality to construct new theory, models are to be subspaces W .

$\Box\phi$ assigned value true for those W in which every point has ϕ true in old theory.

Clearly $\Box(\phi \wedge \psi) \equiv \Box\phi \wedge \Box\psi$.

For \mathbb{R} , signature has elements of the form $\Box\perp$, $\Box[q < \cdot]$, $\Box[\cdot < r]$ and $\Box([q < \cdot] \vee [\cdot < r])$.

Under suitable axioms, model = compact subspace of \mathbb{R} .

Application eg: Heine-Borel Theorem, closed intervals $[x, y]$ are compact

Using hyperspace, can demonstrate that $[x, y]$ depends *continuously* on x and y .

Other hyperspaces available; works for all spaces

¹Point-free hyperspaces aka *powerlocales*. For geometricity see [Vic04].

Conclusion

Logic makes topology work
... better than *topology* does!

See [NV22] “Point-free construction of real exponentiation” for introduction to putting this into practice.

Bibliography I

- [MLM92] S. Mac Lane and I. Moerdijk, *Sheaves in geometry and logic*, Springer-Verlag, 1992.
- [Ng22] Ming Ng, *Adelic geometry via topos theory*, Ph.D. thesis, School of Computer Science, University of Birmingham, 2022.
- [NV] Ming Ng and Steven Vickers, *A point-free look at Ostrowski's Theorem and absolute values*, Submitted for publication. Archived at arXiv:2308.14758.
- [NV22] ———, *Point-free construction of real exponentiation*, Logical Methods in Computer Science **18** (2022), no. 3, 15:1–15:32, DOI 10.46298/lmcs-18(3:15)2022.
- [Vic99] Steven Vickers, *Topical categories of domains*, Mathematical Structures in Computer Science **9** (1999), 569–616.

Bibliography II

- [Vic04] ———, *The double powerlocale and exponentiation: A case study in geometric reasoning*, Theory and Applications of Categories **12** (2004), 372–422, Online at <http://www.tac.mta.ca/tac/index.html#vol12>.
- [Vic07] ———, *Locales and toposes as spaces*, Handbook of Spatial Logics (Marco Aiello, Ian E. Pratt-Hartmann, and Johan F.A.K. van Benthem, eds.), Springer, 2007, pp. 429–496.
- [Vic17] ———, *Arithmetic universes and classifying toposes*, Cahiers de topologie et géométrie différentielle catégorique **58** (2017), no. 4, 213–248.
- [Vic22] ———, *Generalized point-free spaces, pointwise*, <https://arxiv.org/abs/2206.01113>, 2022.

Bibliography III

- [Vic23] ———, *The fundamental theorem of calculus point-free, with applications to exponentials and logarithms*, <https://arxiv.org/abs/2312.05228>, 2023.