Locators point-free

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Locators (Booij [Boo20]); supervised by Martín Escardó

A Dedekind real x is defined by -

- ightharpoonup which rationals q are less than it (q < x)
- ightharpoonup which rationals r are greater than it (x < r)

x has the *locatedness* property:

$$q < r \vdash_{q,r:\mathbb{Q}} (q < x) \lor (x < r)$$

A *locator* for x is structure that, for all pairs q < r, chooses a valid one out of the properties q < x and x < r.

Thus with a locator we have

$$q < r \vdash_{q,r:\mathbb{Q}} (q < x) + (x < r)$$

 $\ensuremath{\mathbb{R}}$ is space of (Dedekind) reals.

 $\mathbb{R}^{\mathfrak{L}}$ is space of reals equipped with locators.

 $\operatorname{pr}_1:\mathbb{R}^\mathfrak{L}\to\mathbb{R}$ is projection.

Analysis with locators (Booij)

Given a locator for x -

can compute decimal expansion without using Axiom of Choice. Locator supplies the choices needed.

Analysis with locators -

tracks the usual real analysis.

Constructions on \mathbb{R} (including arithmetic, limits) lifted to $\mathbb{R}^{\mathfrak{L}}$.

At the end, decimal expansion can be output.

Univalent type theory (UTT) v. point-free topology

Booij's work was in setting of UTT.

- ► Allows for *dis*continuous functions.
- Type constructors include arbitrary function types.
- Not all reals have locators, so analysis with locators does not automatically restrict to real analysis.

We transfer it to point-free topology.

- All types are topologized, all maps continuous.
- ► Not all function types exist. ②
- ▶ $\operatorname{pr}_1: \mathbb{R}^{\mathfrak{L}} \to \mathbb{R}$ is a *point-free surjection*, and a coequalizer. Hence real analysis can be recovered from analysis with locators. \mathfrak{S}

Point-free topology

- The points of a space are described as the models of a geometric theory.
- ➤ Continuity of maps is ensured by restricting to geometric constructions colimits, finite limits, free algebras.
- A bundle over a space X is a space Y_x constructed (geometrically) out of points x of X.
- ► A set X is a discrete space.
- A function is a map between sets.

Point-free topology in a context

Let x be a point of X {

In the scope of that declaration x : X, we are working

- ▶ in a geometric mathematics
- \blacktriangleright that includes a *generic* point x of X.

It is equivalent to working internally in the topos of sheaves over X.

Anything we do here with x can be transported to any specific point a of X by substitution.

(That's the universal property of classifying toposes.)

Let's define a space Y_x .

It constructs a bundle over X.

The substitution $Y_x[a/x]$ gives the fibre over a.

}

Bundle space Y is dependent sum $\Sigma_{x:X} Y_x$

X is a geometric theory.

 Y_{x} is a geometric theory, defined using ingredients of model of X.

Combine them to get theory Y.

Its models are pairs (x, y), x : X and $y : Y_x$

– just as they should be for a dependent sum $\sum_{x:X} Y_x$.

Projection $p: Y \to X$ forgets y.

Fibre over a is space of points y of $Y_x[a/x]$.

Locators using a pullback

Given a real x and rationals q < r –

two truthvalues $\phi_I = (q < x)$ and $\phi_r = (x < r)$, disjunction is \top . Write Λ for space of such pairs (ϕ_I, ϕ_r) of truthvalues, $Q_<$ for set of pairs of rationals q < r. Get $\mathbb{R} \times Q_< \to \Lambda$, so $\mathbb{R} \to \Lambda^{Q_<}$. (Sets are locally compact, so exponentiable.)

Structure of A

Three classical solutions of $\phi_{\it I} \lor \phi_{\it r} = \top$:

For constructive solutions, Λ is the ideal completion of this 3-element poset.

Locators using a pullback

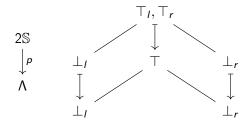
At x and q < r, locator provides element of $\phi_I + \phi_r$

Two more truth values

- $ightharpoonup \phi_I' = \phi_I \text{ chosen}$
- $\phi_r' = \phi_r$ chosen

Write 2S for space of quadruples $(\phi_I', \phi_I, \phi_r', \phi_r)$

- 4 classical solutions.



Locators using a pullback

A real equipped with a locator is -

- ► a real x
- ightharpoonup a map $f: Q_{<}
 ightarrow 2\mathbb{S}$
- If $f(q < r) = (\phi'_I, \phi_I, \phi'_r, \phi_r)$ then $\phi_I = (q < x)$ and $\phi_r = (x < r)$

The following diagram is a pullback.

$$\mathbb{R}^{\mathfrak{L}} \xrightarrow{\operatorname{pr}_{1}} (2\mathbb{S})^{Q_{<}} \downarrow^{p^{Q_{<}}} \mathbb{R} \xrightarrow{x \mapsto (q < r) \mapsto (q < x, x < r)} \Lambda^{Q_{<}}$$

Open surjections

A map $f: Y \to X$ is surjective iff –

its inverse image function f^{-1} is one-one on opens.

This does *not* mean that every point of X has a pre-image.

It is a *conservativity* principle

- certain facts about X can be validly derived from facts about Y.

Surjections are better behaved if they are open

- ie direct images of opens are open (Joyal & Tierney [JT84])

Open surjections are preserved by pullback.

An open surjection is the coequalizer of its kernel pair.

Open surjections

For a map $f: Y \to X$:

f open $\Leftrightarrow f$ fibrewise overt f open surjection $\Leftrightarrow f$ fibrewise positive overt

A property holds "fibrewise" for f iff it holds for the fibre Y_x over the generic point.

This only makes sense if the property is geometric – so if it holds for the generic fibre, it holds for all fibres.

Overt spaces

A space X is *overt* if it has a positivity predicate Pos on its opens.

- ▶ Think: Pos(U) means U non-empty.
- It suffices to define Pos on a base of opens.
- Pos must respect covers: if Pos(U), and U is covered by a family $\{V_i\}$, then we must have $Pos(V_i)$ for some i.
- ▶ Each basic U must be covered by the family $\{\top \mid Pos(U)\}$.

The lower hyperspace (powerlocale) P_L

Points of $P_L X = \text{overt}$, "weakly closed" subspaces of X.

X is itself overt iff it appears in $P_L X$, as the top point \top . "Top" means right adjoint to !.

$$1 \xrightarrow{\frac{\top}{\top}} \mathsf{P}_L X$$

- \triangleright P_L is a geometric construction (Vickers [Vic04])
- Hence overtness is a geometric property
- ▶ and we can talk about "fibrewise overtness".
- ► Also, P_L preserves coreflexive equalizers.

Fibrewise overtness of $f: Y \rightarrow X$

Given x: X, we get two coreflexive equalizer diagrams:

$$Y_{X} \xrightarrow{i_{X}} Y \xrightarrow{y \mapsto \langle x, y \rangle} X \times Y ,$$

$$P_{L} Y_{X} \xrightarrow{P_{L} i_{X}} P_{L} Y \xrightarrow{P_{L}(y \mapsto \langle x, y \rangle)} P_{L}(X \times Y) .$$

Now apply $\Sigma_{x:X}$ and simplify to get another equalizer.

$$(\mathsf{P}_L/X)\ Y \longrightarrow X \times \mathsf{P}_L\ Y \xrightarrow[\pi_2\,\mathsf{P}_L\langle f,\mathsf{Id}_Y\rangle]{\tau} \mathsf{P}_L(X\times Y)$$

au is strength, $(P_L/X) Y = \sum_{x:X} Y_x$.



$f: Y \to X$ is fibrewise overt iff it is open

In context of x:X:

Overtness of Y_x is map $1 \to P_L Y_x$ with certain properties – and can get it from map $1 \to P_L Y$.

Applying $\Sigma_{x:X}$:

Fibrewise overtness of f is a map $X \to P_L Y$ with certain properties.

f is open iff (Vickers [Vic95])

 $-P_L f$ has a right adjoint with certain properties This structure is equivalent to that for fibrewise overtness.

f open iff fibrewise overt

Given that, further calculation shows surjection iff fibrewise positive.

$p^{Q_{<}}$ is fibrewise positive overt $(p \colon 2\mathbb{S} \to \Lambda)$

More generally: if X is a set with decidable equality, then

$$p^X \colon (2\mathbb{S})^X \to \Lambda^X$$
 is fibrewise positive overt

Point of Λ^X is decomposition $X = X_I \cup X_r$.

Fibre = set $X_I + X_r$.

For overtness, must say which opens are positive.

Suffices to do this for basic opens

- and check that covers are respected.

If $U \leq \bigvee_i V_i$, and U is positive, then one of the V_i s must be positive.

Subasic opens x_l, y_r for $x \in X_l, y \in X_r$.

For $x \in X_l \cap X_r$ have

$$x_l \wedge x_r \vdash \bot$$

 $\top \vdash x_l \vee x_r$

Basic opens (finite meets of subbasics) -

$$B = \{ (S_I, S_r) \in \mathcal{F}X_I \times \mathcal{F}X_r \mid (\forall x \in S_I)(\forall y \in S_r)x \neq y \}$$
$$(S_I, S_r) = \bigwedge_{x \in S_I} x_I \wedge \bigwedge_{y \in S_r} y_r$$
$$B \text{ is a set!}$$

Positivity predicate: every basic in B is positive.

Respecting covers

$$(S_I, S_r) \vdash \bigvee (\{(S_I \cup \{x\}, S_r), (S_I, S_r \cup \{x\})\} \cap B)$$

Suppose $(S_l, S_r) \in B, x \in X$. Want at least one of $(S_l \cup \{x\}, S_r)$, $(S_l, S_r \cup \{x\})$ in B. Immediate if $x \in S_l \cup S_r$. Otherwise (and condition is decidable), use $x \in X_l \cup X_r$.



$\operatorname{pr}_1\colon \mathbb{R}^\mathfrak{L} o \mathbb{R}$ is an open surjection

- We now know p^X is fibrewise overt. (We've defined the positivity predicate.)
- ▶ It's fibrewise positive because (\emptyset, \emptyset) , which is \top , is positive.
- ► Hence it's an open surjection.
- Now put $X = Q_{<}$
- ightharpoonup and pull back to get pr_1 .

Open surjections are coequalizers of their kernel pairs

$K = \text{kernel pair of } pr_1$

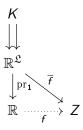
Point = real x equipped with two locators.

 \overline{f} depends only on x

- gives same result for any locators Then \overline{f} factors uniquely via pr_1 .

In defining f, can pretend every real has a locator

- even though not all of them do.



Examples of surjections

All these are triqotient maps (Plewe [Ple97]).

They are coequalizers of their kernel pairs, and preserved by pullback.

$$\operatorname{pr}_1 \colon \mathbb{R}^{\mathfrak{L}} \to \mathbb{R}$$

open surjection

$$2^{\omega} \rightarrow [-1,1]$$

$$(s_i)_{i=1}^{\infty} \mapsto \sum_{i=1}^{\infty} s_i/2^i \text{ (where } s_i = \pm 1)$$

proper surjection (Vickers [Vic17])

$$\lim : \operatorname{Cauchy}_f X \to \overline{X}$$

Limits of Cauchy sequences in a generalized metric space X.

- triquotient (Vickers [Vic98])

Taking $X = \mathbb{Q}$ gives surjection to \mathbb{R} .

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