

Practice Sheet 2

CSE 112 Computer Organization

The instructions supported by the ISA are mentioned in the table below. **The ISA has 16 General purpose registers: r0 to r15**

Name	Semantics	Syntax
Add	Performs $\text{reg1} = \text{reg2} + \text{reg3}$	<code>add reg1 reg2 reg3</code>
subi	Performs $\text{reg1} = \text{reg2} - \text{Imm}$	<code>sub reg1 reg2 \$Imm</code>
mvi	Performs $\text{reg1} = \text{Imm}$	<code>mov reg1 \$Imm</code>
Mov	Performs $\text{reg1} = \text{reg2}$	<code>mov reg1 reg2</code>
Branch not equal	Branch to addr if $\text{reg1} \neq \text{reg2}$	<code>bneq reg1 reg2 addr</code>
Branch and link	Jumps to label after saving the return address to r1.	<code>brl label</code>
Push	Pushes the data stored in reg1 onto the stack	<code>push reg1</code>
Pop	Pops the data stored on the top of the stacks into reg1	<code>pop reg1</code>
mul	Performs $\text{reg1} = \text{reg2} * \text{reg3}$	<code>mul reg1 reg2 reg3</code>

Apart from the above instructions, the assembler and the operating system support the following subroutines:

Name	Semantics	Syntax
Input	Reads immediate data from user into reg	in reg
Output	Prints str on the console	out "str"

Caller-callee conventions:

The following are the caller callee convention:

- There are 15 registers r0 to r15.
- r15 - program counter.
- r0 - stack pointer.
- r1 - link register and return address
- r2 - return value.
- r3 and r4 holds the first and second argument to the callee
- The stack is automatically managed by push and pop.
- All the registers from r1-r7 are caller saved. On the other hand, registers r8-r14 are callee saved.
- Whenever the branch and link instruction is used, the return address is stored in r1 and the program counter jumps to the given label.

Q1: Convert the following high level code into assembly language. Follow the caller-callee conventions mentioned above.

You can only use callee saved registers for storing variables in bar functions and caller saved registers for foo function for storing variables.

```

int baz(int a,int b)
{
    return a+b;
}

int bar()
{
    int a = 10;
    int b = 100;
    int c = 1000;

```

```

        int d = baz(a,b);
        return a+b+c+d;
    }

int foo()    // Use only caller saved registers
{
    int a = 10;
    int b = 100;
    int c = bar();
    int d = baz(a,b);
    return a+b+c+d;
}

int main()
{
    return foo();
}

```

Q2. Write a function myfunc which computes factorial of a number n passes as an argument in assembly language. Use ISA as provided in Q1.

a. Use iterative call

Assumptions for part b:

1. The number whose factorial is to be calculated (n) is present in r5.

b. Use recursive call

Assumptions for part b:

1. Return Address for myfunc is present in link register r1.
2. The number whose factorial is to be calculated (n) is present in register stack.

A factorial of a number n is:

$n \times (n-1) \times (n-2) \times \dots \times 1$

Solution Q1:

baz:

```
add r2 r3 r4    // Add the arguments and return it
```

```
mov r15 r1      // Return
```

bar:

```
push r8         // Push callee saved register
```

```
push r9         // Push callee saved register
```

```
push r10        // Push callee saved register
```

```
push r11        // Push callee saved register
```

```
mvi r8 #10      // r8 is a
```

```
mvi r9 #100     // r9 is b
```

```
mvi r10 #1000   // r10 is c
```

```
mov r3 r8       // Prepare first argument of baz
```

```
mov r4 r9       // Prepare second argument of baz
```

```
push r1         // Push caller saved register
```

```
brl baz        // call baz function
```

```
pop r1         // Pop caller saved register
```

```
mov r11 r2      // r11 is d
```

```
mvi r2, #0     // Initialize sum with 0
```

```
add r2 r2 r8    // Add a
```

```
add r2 r2 r9    // Add b
```

```
add r2 r2 r10   // Add c
```

```
add r2 r2 r11   // Add d
```

```
pop r11        // Pop callee saved register
```

```
pop r10        // Pop callee saved register
```

```
pop r9         // Pop callee saved register
```

```

    pop r8            // Pop callee saved register
    mov r15 r1        // Jump back to caller function

foo:  push r1          // Push caller saved register
      brl bar          // Push caller saved register
      pop r1           // Push caller saved register
      mov r6 r2        // r6 is c

      mvi r3 $10       // r3 is a, send the first argument to baz
      mvi r4 $100      // r4 is b, send the second argument to baz
      push r1           // Push caller saved register
      push r3           // Push caller saved register
      push r4           // Push caller saved register
      push r6           // Push caller saved register
      brl baz          // Call baz function
      pop r6           // Pop caller saved register
      pop r4           // Pop caller saved register
      pop r3           // Pop caller saved register
      pop r1           // Pop caller saved register

      add r2 r2 r6      // Prepare return value
      add r2 r2 r3      // Prepare return value
      add r2 r2 r4      // Prepare return value
      mov r15 r1        // Jump back to caller function

main: push r1          // Push caller saved register
      brl foo          // Call foo function
      pop r1           // Pop caller saved register
      mov r15 r1        // Jump back to caller function

```

Solution Q2 (a)

```
myfunc:    mvi r6 $1
           mvi r7 $0
           Bneq r5 r7 MulLoop
           mvi r2 $0
           Pop r1
           Mov r15 r1
MulLoop:   Mul r6 r6 r5           // n is saved in r5
           subi r5 r5 $1
           Bneq r5 r7 MulLoop
           Pop r1
           Mov r15 r1
```

Solution Q2 (b):

```
myfunc:    pop r5
           Push r1
           Mvi r7 $0
           Bneq r5 r7 Done
myfunc2:   pop r5
           Push r1
           mvi r7 $1
           Bneq r5 r7 else
Done:      mvi r2 $1
           Pop r1
           Mov r15 r1
Else:      mov r8 r5
           subi r5 $1
           push r8
           push r5
```

brl myfunc2

pop r8

mul r2 r2 r8

Pop r1

Mov r15 r1