

## Tutorial -3

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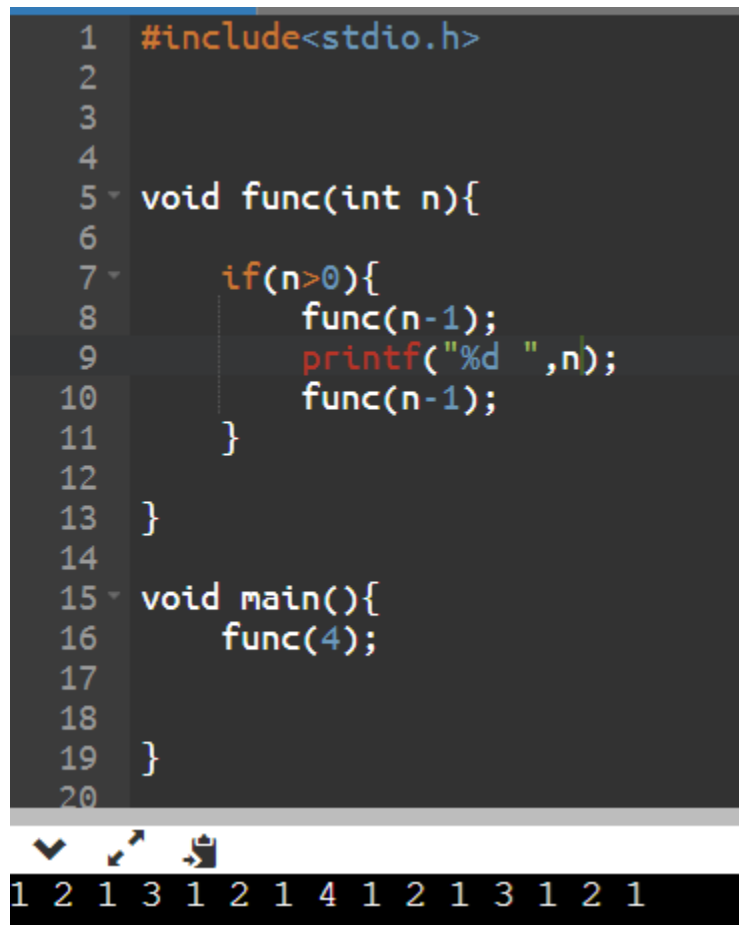
In this tutorial, we will go through the concepts of recursion, time and space complexities, and how to solve recurrence relations

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### Recursion:

#### Question 1)

```
1  #include<stdio.h>
2
3
4
5  void func(int n){
6
7      if(n>0){
8          func(n-1);
9          printf("%d ",n);
10         func(n-1);
11     }
12
13 }
14
15 void main(){
16     func(4);
17
18
19 }
20
```



1 2 1 3 1 2 1 4 1 2 1 3 1 2 1

### Explanation:

Draw recursion tree and traverse in top to bottom, left to right fashion

## Question 2)

```
1  #include<stdio.h>
2  int c=0;
3
4  void func(int n){
5
6      if(n==0){
7          return;
8      }
9      c++;
10     func(n/10);
11
12 }
13
14 void main(){
15     |
16     func(123456789);
17     printf("Value of c is %d",c);
18
19 }
20
21
```

### Output:

Value of c is 9

### Explanation:

Draw recursion tree and traverse in top to bottom fashion. Final value of c=9, after that we reach the base case n=0 and so return.

## Efficiency of Algorithms:

### Question 1)

Find the Time Complexity and Space Complexity of the following snippet:

```
int a = 0, b = 0;
for (i = 0; i < N; i++) {
    a = a + rand(); //rand() returns a random number
}
for (j = 0; j < M; j++) {
    b = b + rand();
}
```

### Answer:

Time Complexity =  $O(m+n)$

Space Complexity =  $O(1)$  (constant number of variables)

### Question 2)

Find the Time Complexity of the following snippet:

```
int a = 0;
for (i = 0; i < N; i++) {
    for (j = N; j > i; j--) {
        a = a + i + j;
    }
}
```

### Answer:

code runs total no of times  
=  $N + (N - 1) + (N - 2) + \dots 1 + 0$   
=  $N * (N + 1) / 2$   
=  $1/2 * N^2 + 1/2 * N$   
 $O(N^2)$  times.

### Question 3)

Find the Time Complexity of the following snippet:

```
int a = 0, i = N;
while (i > 0) {
    a += i;
    i /= 2;
}
```

**Answer:**

$O(\log n)$

### Question 4)

Find the Time Complexity of the following function:

```
int fun(int n) {
    for (int i = 1; i <= n; i++) {
        for (int j = 1; j < n; j += i) {
            // Some  $O(1)$  task
        }
    }
}
```

**Answer:**

$O(n * \log n)$

**Explanation:**

## Harmonic Series:

$$\sum_{k=1}^n \frac{1}{k} = 1 + \frac{1}{2} + \dots + \frac{1}{n} \approx \ln n$$

For i = 1, the inner loop is executed n times.  
For i = 2, the inner loop is executed approximately n/2 times.  
For i = 3, the inner loop is executed approximately n/3 times.  
For i = 4, the inner loop is executed approximately n/4 times.  
.....  
.....  
For i = n, the inner loop is executed approximately n/n times.

The total time complexity of the above algorithm is  $(n + n/2 + n/3 + \dots + n/n)$   
 $= n * (1/1 + 1/2 + 1/3 + \dots + 1/n)$   
[Recall the complexity of harmonic series]

Hence, the time complexity of fun is  $O(n \log n)$

### Question 5)

Find the Time Complexity of the following function:

```
void demo()
{
    int i, j;
    for (i=1; i<=n; i++)
        for (j=1; j<=log(i); j++)
            System.out.println("Hello World");
}
```

### Answer:

$\Theta(\log 1) + \Theta(\log 2) + \Theta(\log 3) + \dots + \Theta(\log n)$   
 $= \Theta(\log n!)$   
 $= \Theta(n \log n)$

## Solving Recurrence Relations:

### Question 1)

Solve the following recurrence:

$$T(n) = 2T(n/2) + n$$

$$T(1) = 0$$

**Answer:**

$$\begin{aligned} T(n) &= 2T(n/2) + n \\ &= 2[2T(n/2^2) + n/2] + n = 2^2 T(n/2^2) + 2n \\ &= 2^2 [2T(n/2^3) + n/2^2] + 2n = 2^3 T(n/2^3) + 3n \\ &\dots\dots\dots \\ &= 2^k T(n/2^k) + kn \\ \text{If } n/2^k &= 1 \Rightarrow n = 2^k \Rightarrow k = \log n \text{ (with base 2)} \\ \text{Complexity} &= O(n \log n) \end{aligned}$$

### Question 2)

Find time complexity of the recursive fibonacci method

```
public static int rFib( int n) {  
    if (n == 0 || n == 1)  
        return 1;  
    else  
        return rFib (n-1) + rFib (n-2);  
}
```

**Answer:**

$O(2^n)$

### Few Basic Questions:

Find the time complexity of the given expressions?

**Q1)**  $3n \log n + 2n^{1.8}$                       **Ans:**  $O(n^{1.8})$

**Q2)**  $8 \log n + 4 \log \log n$                       **Ans:**  $O(\log n)$

**Q3)**  $3n + 2(\log n)^2 + 4 \log n$                       **Ans:**  $O(n)$

**Q4)** An algorithm takes 0.5 ms for input size 100. How large a problem can be solved in 1 min if the running time is Linear (assume low-order terms are negligible)

**Answer:**

$x/100 = 60,000/0.5$   
solving for x gives an input size of 12,000,000

**Q5)** An algorithm takes 0.5 ms for input size 100. How large a problem can be solved in 1 min if the running time is  $O(n \log n)$  (assume low-order terms are negligible)

**Answer:**

$x \log x / 100 \log 100 = 60,000/0.5$   
solving for x gives an input size of 3,656,807

### Homework Questions:

**Q1)** An algorithm takes 0.5 ms for input size 100. How large a problem can be solved in 1 min if the running time is Quadratic (assume low-order terms are negligible)

**Q2)** An algorithm takes 0.5 ms for input size 100. How large a problem can be solved in 1 min if the running time is Cubic (assume low-order terms are negligible)