



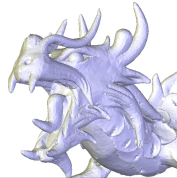
Lyon 1

## Mesh and Computational Geometry

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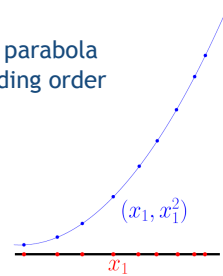
M2 ID3D  
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## Incremental Delaunay Algorithm

- Complexity analysis
- Worst case :
  - Points distributed on a parabola and inserted in descending order of abscissa.

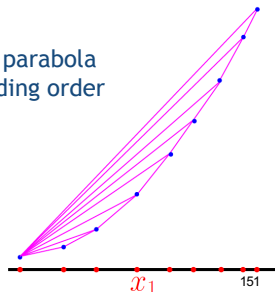


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## Incremental Delaunay Algorithm

- Complexity analysis
- Worst case :
  - Points distributed on a parabola and inserted in descending order of abscissa.



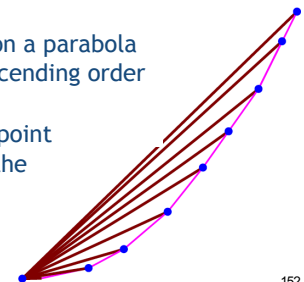
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## Incremental Delaunay Algorithm

- Complexity analysis
- Worst case :
  - Points distributed on a parabola and inserted in descending order of abscissa.
  - Each new inserted point conflicts with ALL the triangles

$$\Omega(n^2)$$



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## Incremental Delaunay Algorithm

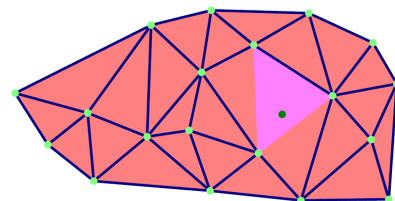
- Complexity analysis
- In average :
  - Complexity dependent on the strategy used to locate the triangle containing the point to be inserted

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## Location strategies

- Exhaustive search among all the triangles



$$O(n)$$

O. Devillers

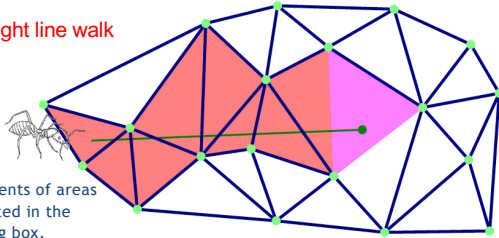
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### Location strategies

- Search from a first randomly selected triangle

Straight line walk



3n elements of areas distributed in the bounding box, therefore  $\sqrt{3n}$  elements encountered on average on a straight line

$O(\sqrt{n})$  on average

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### Location strategies

- Search from a first randomly selected triangle
  - Straight walk
  - Requires a predicate of segments intersection

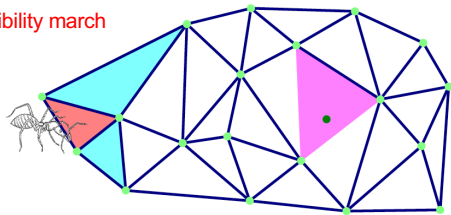
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### Location strategies

- Search from a first randomly selected triangle
  - Some minor deviations from the straight line walk

Visibility march



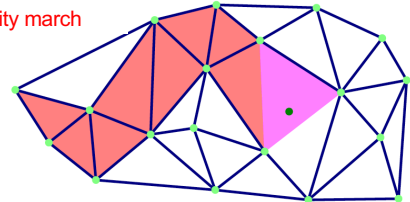
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### Location strategies

- Search from a first randomly selected triangle
  - Some minor deviations from the straight line walk

Visibility march



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### Location strategies

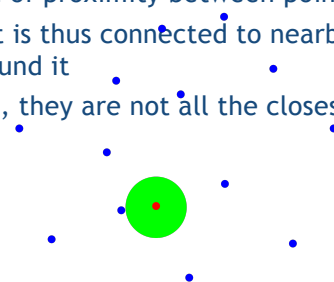
- Search from a first randomly selected triangle
  - Visibility walking
  - Only requires an orientation predicate to find the next triangle to walk in

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### Delaunay and proximity in space

- Delaunay triangulation allows to model the notion of proximity between points
- Each point is thus connected to nearby points around it
- Be careful, they are not all the closest!

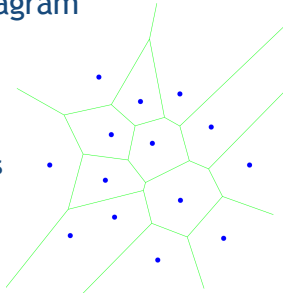


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### Voronoi diagram

- Voronoi's cell of a site  $P_i$  is the set of points closer to this site than to other sites



$$V_i = \{P \in \mathbb{R}^k \text{ t. que } PP_i < PP_j \text{ pour tout } j \neq i\}$$

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### Voronoi diagram

- Given a set  $E$  of points in  $\mathbb{R}^k$ , the partitioning of  $\mathbb{R}^k$  into cells composed of points having the same nearest neighbour in  $E$  is called a Voronoi diagram of  $E$

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### Voronoi diagram

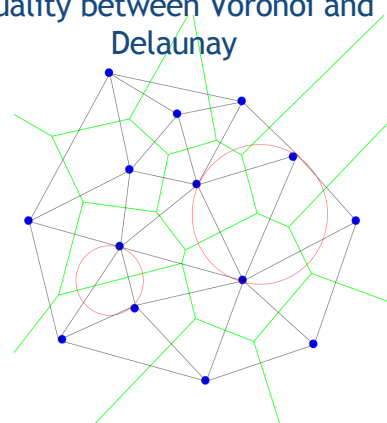
- Possible construction:
  - $V_i$  : intersection of half-spaces  $h_{ij}^i$  where  $h_{ij}$  is the mediator of segment  $P_iP_j$  and  $h_{ij}^i$  is the half-space delimited by  $h_{ij}$  containing  $P_i$

In practice we will proceed differently!

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### Duality between Voronoi and Delaunay



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### Duality between Voronoi and Delaunay

- Each Voronoi vertex is located at the center of the circumscribed circle of a Delaunay triangle
- Two Voronoi vertices are connected if they are associated with adjacent triangles

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### Coordinates of the centre of the circle circumscribed to a triangle ABC

- Useful for displaying the Voronoi diagram
- 1st possibility:
  - Write the equation of the mediator for each edge  
ex: For the edge AB, set of points  $M$  such that  $MA^2 = MB^2$
  - Solving a system of 2 equations with 2 unknowns (it is enough to take 2 mediators)
  - Numerically unstable

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## Coordinates of the centre of the circle circumscribed to a triangle ABC

### • 2nd possibility :

- Let's consider the angles

$$\hat{A} = \widehat{CAB} \quad \hat{B} = \widehat{ABC} \quad \hat{C} = \widehat{BCA}$$

- Then the barycentric coordinates of the center H of the circumscribed circle with respect to A, B and C are elegantly expressed :

$$H(\tan \hat{B} + \tan \hat{C}, \tan \hat{C} + \tan \hat{A}, \tan \hat{A} + \tan \hat{B})$$

- Reminder :

$$\tan(\widehat{ABC}) = \frac{\sin(\widehat{ABC})}{\cos(\widehat{ABC})} = \frac{\text{sign}((\vec{BC} \times \vec{BA}) \cdot \vec{k}) \|\vec{BC} \times \vec{BA}\|}{\vec{BC} \cdot \vec{BA}} \quad 167$$

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## Coordinates of the centre of the circle circumscribed to a triangle ABC

$$H = \text{Barycenter}((A, \tan \hat{B} + \tan \hat{C}), (B, \tan \hat{C} + \tan \hat{A}), (C, \tan \hat{A} + \tan \hat{B}))$$

- Reminder :

$$\tan(\widehat{ABC}) = \frac{\sin(\widehat{ABC})}{\cos(\widehat{ABC})} = \frac{\text{sign}((\vec{BC} \times \vec{BA}) \cdot \vec{k}) \|\vec{BC} \times \vec{BA}\|}{\vec{BC} \cdot \vec{BA}}$$

$$\text{Barycenter}((A, \alpha a), (B, \alpha b), (C, \alpha c))$$

$$= \text{Barycenter}((A, a), (B, b), (C, c))$$

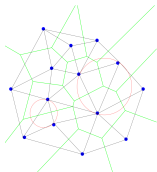
- Ensure to have no more denominators in the expression of your barycentric coordinates (normalization performed afterwards)

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## Duality between Voronoi and Delaunay

- Each Delaunay vertex is dual to one Voronoi cell
- Each Delaunay edge is dual to a Voronoi edge
- Each Voronoi vertex is dual to a Delaunay triangle
- What about Delaunay, Voronoi and their duality in 3D?

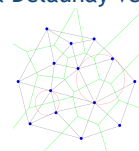


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## Duality between Voronoi and Delaunay

- Which data structure for Voronoi?
  - Walking around a Voronoi face is performed by walking through the faces/edges incident at a Delaunay vertex.
  - To move from one Voronoi cell to an adjacent cell is like moving from a Delaunay vertex to an adjacent vertex.

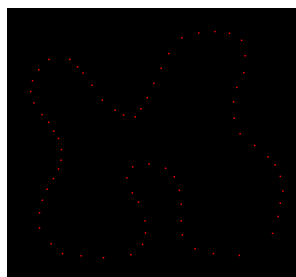


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## Use of Voronoi Center for line reconstruction from a 2D point set (Crust)

- Let consider a set of points sampled on a line
- The points are not ordered
- How to approximate the input line with a polygonal line?

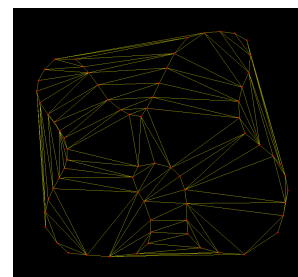


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## Use of Voronoi Center for line reconstruction from a point set (Crust)

- If the sampling is dense enough, Delaunay encloses a good candidate
- How to remove the edges crossing the shape?

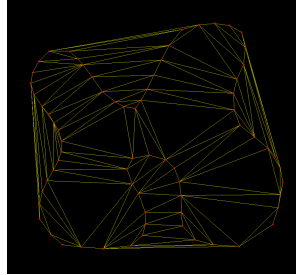


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### Use of Voronoi Center for line reconstruction from a point set (Crust)

- Add points as far as possible from the input line
  - Points located on the internal and external skeleton of the shape
  - Centers of maximal empty circles

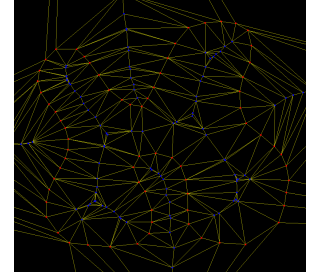


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### Use of Voronoi Center for line reconstruction from a point set (Crust)

- The Voronoi centers are close to the shape skeleton
  - Let's add them to break edges that cross the shape while preserving the boundaries

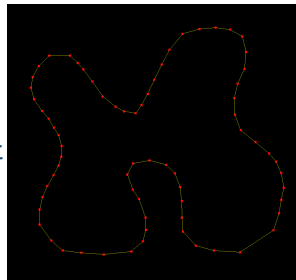


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### Use of Voronoi Center for line reconstruction from a point set (Crust)

- Keep the edges that join initial input points
- Correctness of the algorithm if the input point set is locally denser than a given proportion  $\epsilon$  of the distance to the skeleton ( $\epsilon$ -sampling)

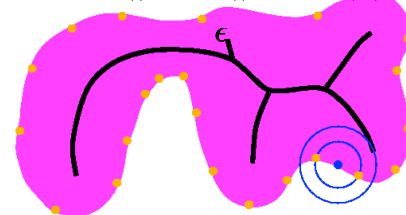


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### Local Feature Size (lfs) and $\epsilon$ -sampling

- Given  $\epsilon$ , an  $\epsilon$ -sampling of a shape is a set of samples  $P_i$  such that for each  $x$  there is a  $i$  such that  $\|x - P_i\| \leq \epsilon lfs(x)$



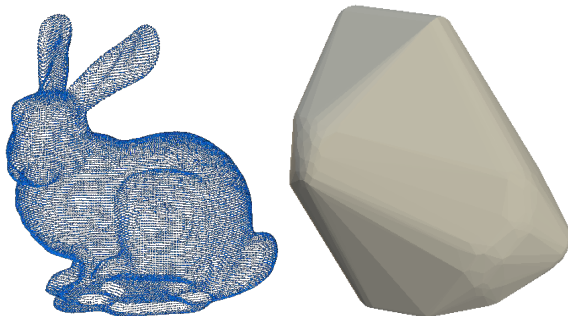
Sampling density locally proportional to  $1/lfs$   
Measure of thickness and curvature

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### Crust in 3D?

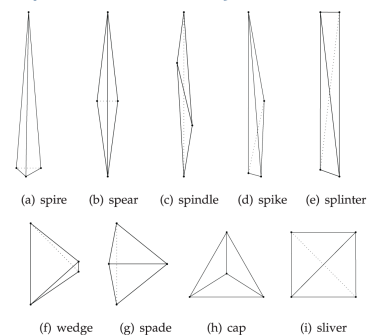
In 3D the Delaunay triangulation provides a tetrahedrisation of the shape



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### Crust in 3D?

- Zoology of Delaunay tetrahedra based on point samples from an object's surface

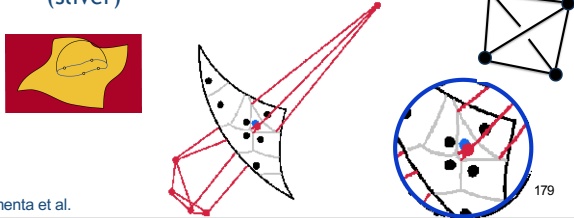


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## Crust in 3D

- Some Voronoi centers can be far from the skeleton and close to the surface...
  - No control over the position of the centers of flat tetrahedrons: 4 neighboring points almost cocyclic may have their center near the surface (sliver)



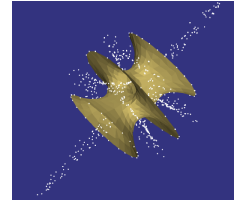
Amenta et al.

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## 3D Crust

- Need to filter Voronoi centers



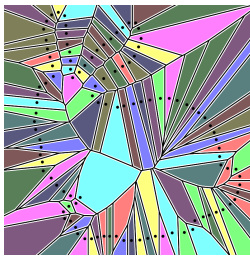
- Consider only the poles!
  - ie. Voronoi vertices certified to be far from the surface by one of the point samples

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## Poles

- In presence of a dense and non noisy sampling
  - long and thin Voronoi cells,
  - direction similar to the normal to the surface.

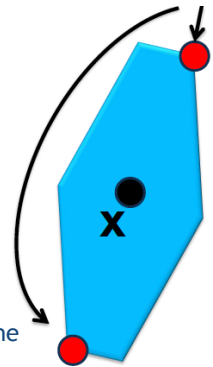


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## Poles

- Let  $V_x$  be the Voronoi's cell of a point  $x$
- Positive pole  $p^+$  : Voronoi Vertex of  $V_x$  further away from  $x$
- Vector pole  $x p^+$  : approximation of the normal direction at  $x$
- The negative pole  $p^-$  : farthest vertex of  $V_x$  in the opposite direction to the vector  $x p^+$

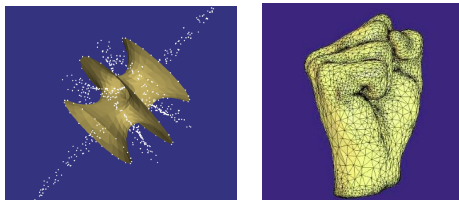


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## Crust in 3D

- Adding poles in the 3D triangulation



Amenta et al 98

- Reconstructed surface composed of Delaunay faces relying on 3 input point samples

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