

Propositional and Predicate Logics

Solutions (2016)

1 Models and Entailment in Propositional Logic

1. (a) Truth table for $A \wedge \neg B \models A \vee B$:

A	B	$A \wedge \neg B$	$A \vee B$
0	0	0	0
0	1	0	1
1	0	1	1
1	1	0	1

The entailment is **true**.

- (b) Truth table for $A \vee B \models A \wedge \neg B$:

A	B	$A \vee B$	$A \wedge \neg B$
0	0	0	0
0	1	1	0
1	0	1	1
1	1	1	0

The entailment is **false**.

- (c) Truth table for $A \Leftrightarrow B \models A \Rightarrow B$:

A	B	$A \Leftrightarrow B$	$A \Rightarrow B$
0	0	1	1
0	1	0	1
1	0	0	0
1	1	1	1

The entailment is **true**.

- (d) Truth table for $(A \Leftrightarrow B) \Leftrightarrow C \models A \vee \neg B \vee \neg C$:

A	B	C	$(A \Leftrightarrow B) \Leftrightarrow C$	$A \vee \neg B \vee \neg C$
0	0	0	0	1
0	0	1	1	1
0	1	0	1	1
0	1	1	0	0
1	0	0	1	1
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

The entailment is **true**.

- (e) Truth table for $(\neg A \wedge B) \wedge (A \Rightarrow B)$:

A	B	$(\neg A \wedge B) \wedge (A \Rightarrow B)$
0	0	0
0	1	1
1	0	0
1	1	0

The expression is **satisfiable**.

- (f) Truth table for $(\neg A \wedge B) \wedge (A \Leftrightarrow B)$:

A	B	$(\neg A \wedge B) \wedge (A \Leftrightarrow B)$
0	0	0
0	1	0
1	0	0
1	1	0

The expression is **not satisfiable**.

2. In the following, let $Q = 2^{100}$.

- (a) The expression $A_{31} \wedge \neg A_{76}$ is satisfied by 1 models out of 4 possible for the variables A_{31}, A_{76} . For all 100 variables A_1, A_2, \dots, A_{100} , the answer is thus $\boxed{\frac{1}{4}Q}$.
- (b) The expression $A_{44} \wedge A_{49} \wedge A_{78}$ is satisfied by 1 models out of 8 possible for the variables A_{44}, A_{49}, A_{78} . For all 100 variables A_1, A_2, \dots, A_{100} , the answer is thus $\boxed{\frac{1}{8}Q}$.
- (c) The expression $A_{44} \vee A_{49} \vee A_{78}$ is satisfied by 7 models out of 8 possible for the variables A_{44}, A_{49}, A_{78} . For all 100 variables A_1, A_2, \dots, A_{100} , the answer is thus $\boxed{\frac{7}{8}Q}$.
- (d) The expression $A_{70} \Rightarrow \neg A_{92}$ is satisfied by 3 models out of 4 possible for the variables A_{70}, A_{92} . For all 100 variables A_1, A_2, \dots, A_{100} , the answer is thus $\boxed{\frac{3}{4}Q}$.
- (e) The expression $(A_7 \Leftrightarrow A_{72}) \wedge (A_{83} \Leftrightarrow A_{84})$ is satisfied by 4 models out of 16 possible for the variables $A_7, A_{72}, A_{83}, A_{84}$. For all 100 variables A_1, A_2, \dots, A_{100} , the answer is thus $\boxed{\frac{4}{16}Q}$.
- (f) The expression $\neg A_9 \wedge \neg A_{19} \wedge A_{37} \wedge A_{50} \wedge A_{68} \wedge A_{73} \wedge A_{79} \wedge A_{81}$ is satisfied by 1 models out of 256 possible for the variables $A_{19}, A_{37}, A_{50}, A_{68}, A_{73}, A_{79}, A_{81}, A_9$. For all 100 variables A_1, A_2, \dots, A_{100} , the answer is thus $\boxed{\frac{1}{256}Q}$.

Another way to look at this is to realize that the 8 variables $A_9, A_{19}, A_{37}, A_{50}, A_{68}, A_{73}, A_{79}, A_{81}$ all have their values “fixed” by the expression, so that the number of possible models is reduced from 2^{100} to $\boxed{2^{100-8} = 2^{92} = \frac{1}{256}Q}$.

3. Table 1 shows the 16 possible models. There are $16 = 2^4$ possibilities because we ignore the Wumpus and only consider whether there are pits in the four adjacent rooms $[3, 1], [3, 2], [3, 3]$ and $[4, 4]$.

The 6th column of the table shows the models that are consistent with the knowledge base (KB), where the state of the KB is given in the assignment text. The 7th, 8th and 9th columns show the truth values of respectively α_1, α_2 and α_3 .

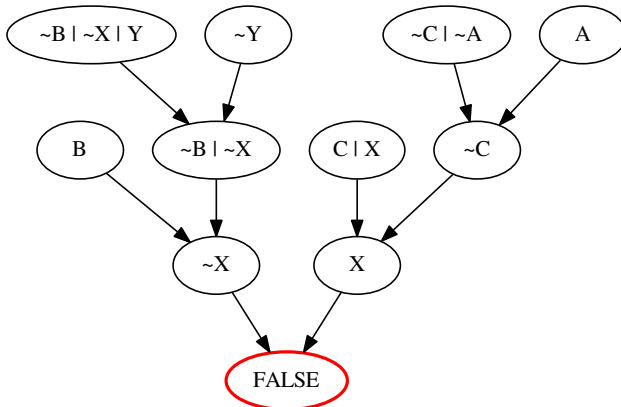
The KB is only true in three models: 10, 11 and 12. α_1 and α_3 are both true in all three of these models, thus both sentences are entailed by the KB. α_2 is true in model 11 and 12, but not in model 10. The KB does therefore *not* entail α_2 .

Table 1: 16 models for a restricted view of the Wumpus World, where KB is the current state of the knowledge base after visiting [4, 1], [4, 2] and [4, 3]. α_1 = “There is a pit in [3, 1]”. α_2 = “There is a pit in [3, 3]”. α_3 = “There is a pit in [3, 3] or [4, 4]”.

Index	Pits				KB	α_1	α_2	α_3
	P_{31}	P_{32}	P_{33}	P_{44}				
1	0	0	0	0	0	0	0	0
2	0	0	0	1	0	0	0	1
3	0	0	1	0	0	0	1	1
4	0	0	1	1	0	0	1	1
5	0	1	0	0	0	0	0	0
6	0	1	0	1	0	0	0	1
7	0	1	1	0	0	0	1	1
8	0	1	1	1	0	0	1	1
9	1	0	0	0	0	1	0	0
10	1	0	0	1	1	1	0	1
11	1	0	1	0	1	1	1	1
12	1	0	1	1	1	1	1	1
13	1	1	0	0	0	1	0	0
14	1	1	0	1	0	1	0	1
15	1	1	1	0	0	1	1	1
16	1	1	1	1	0	1	1	1

2 Resolution in Propositional Logic

- $\neg A \vee (B \wedge C) \equiv (B \vee \neg A) \wedge (C \vee \neg A)$
 - $\neg(A \Rightarrow B) \wedge \neg(C \Rightarrow D) \equiv \neg B \wedge A \wedge \neg D \wedge C$
 - $\neg(A \Rightarrow B) \vee \neg(C \Rightarrow D) \equiv (\neg D \vee \neg B) \wedge (C \vee \neg B) \wedge (\neg D \vee A) \wedge (C \vee A)$
 - $(A \Rightarrow B) \Leftrightarrow C \equiv (B \vee \neg A \vee \neg C) \wedge (\neg B \vee C) \wedge (A \vee C)$
- The following figure shows one possible example of a resolution.



3. (a) The truth table is as follows:

<i>Drinks</i>	<i>Food</i>	<i>Party</i>	$((Food \vee Drinks) \Rightarrow Party) \Rightarrow (\neg Party \Rightarrow \neg Food)$
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

The expression evaluates to true in all cases, so the sentence is **valid**.

- (b) Converting the left-hand side (LHS) to CNF:

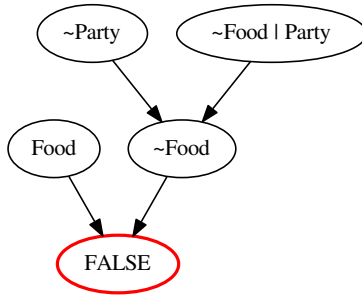
- i. $(Food \vee Drinks) \Rightarrow Party$
- ii. $Party \vee \neg(Food \vee Drinks)$
- iii. $Party \vee (\neg Food \wedge \neg Drinks)$
- iv. $(Party \vee \neg Food) \wedge (Party \vee \neg Drinks)$

Converting the right-hand side (RHS) to CNF:

- i. $\neg Party \Rightarrow \neg Food$
- ii. $\neg Food \vee \neg(\neg Party)$
- iii. $Party \vee \neg Food$

We see that in CNF, the RHS consists simply of one of the clauses of the LHS. The full sentence, $LHS \Rightarrow RHS$, must therefore necessarily be true.

- (c) Proof-by-contradiction by negating the expression, converting to CNF and performing resolution. CNF of negated expression: $Food \wedge \neg Party \wedge (\neg Food \vee Party) \wedge (\neg Drinks \vee Party)$



3 Representations in First-Order Logic

1. (a) $PlayedCharacter(ChristianBale,Batman) \wedge PlayedCharacter(GeorgeClooney,Batman) \wedge PlayedCharacter(ValKilmer,Batman)$
- (b) $\forall c: \neg PlayedCharacter(Bale,c) \vee \neg PlayedCharacter(Ledger,c)$
- (c) $\forall m: CharacterInMovie(Batman,m) \wedge Directed(Nolan,m) \rightarrow PlayedInMovie(ChristianBale,m)$
- (d) $\exists m: CharacterInMovie(TheJoker,m) \wedge CharacterInMovie(Batman,m)$

(e) $\exists m: \text{PlayedInMovie}(\text{KevinCostner}, m) \wedge \text{Directed}(\text{KevinCostner}, m)$

(f) $\forall m: \text{PlayedInMovie}(\text{GeorgeClooney}, m) \rightarrow \neg (\text{PlayedInMovie}(\text{Tarantino}, m) \vee \text{Directed}(\text{Tarantino}, m))$

This is equivalent to:

$\forall m: \neg (\text{PlayedInMovie}(\text{GeorgeClooney}, m) \wedge \text{PlayedInMovie}(\text{Tarantino}, m)) \wedge \neg (\text{PlayedInMovie}(\text{GeorgeClooney}, m) \wedge \text{Directed}(\text{Tarantino}, m))$

And also

$\forall m: (\text{PlayedInMovie}(\text{Tarantino}, m) \vee \text{Directed}(\text{Tarantino}, m)) \rightarrow \neg \text{PlayedInMovie}(\text{GeorgeClooney}, m)$

(g) $\exists m: \text{PlayedInMovie}(\text{UmaThurman}, m) \wedge \text{Directed}(\text{Tarantino}, m)$

2. (a) $\forall x, y: \text{Divisible}(x, y) \leftrightarrow \exists z: (z \leq x) \wedge (x = z \times y)$

(b) $\forall x: \text{Even}(x) \leftrightarrow \text{Divisible}(x, 2)$

(c) $\forall x: \text{Odd}(x) \leftrightarrow \neg \text{Divisible}(x, 2)$

(d) $\forall x: \text{Odd}(x) \leftrightarrow \exists y: \text{Even}(y) \wedge (x = y + 1)$

(e) $\forall x: \text{Prime}(x) \leftrightarrow \forall y: \neg(x = y) \rightarrow \neg \text{Divisible}(x, y)$

3. List of predicates:

- $\text{PersonDNA}(p, d)$: predicate. Person p has DNA d .
- $\text{Derived}(d_1, d_2)$: predicate. DNA d_1 is derived from DNA d_2
- $\text{ParentPerson}(p, q)$ predicate. Person q is a parent (father/mother) of person p .
- $=$: predicate. Compare two persons, true if they are the same person.

$\forall p, d: \text{PersonDNA}(p, d) \rightarrow [\forall q: \neg(q = p) \rightarrow \neg \text{PersonDNA}(q, d)]$

$\wedge [\forall q, d_2: \text{ParentPerson}(p, q) \wedge \text{PersonDNA}(q, d_2) \rightarrow \text{Derived}(d, d_2)]$

4 Resolution in First-Order Logic

1. (a) Answer: $\theta = \{x/\text{Plato}\}$

(b) Answer: $\theta = \{y/\text{TheRepublic}\}$

(c) Answer: $\theta = \{x/\text{Peter}, y/\text{Metaphysics}\}$

(d) Answer: Impossible. You can not unify x to two different atoms ("Fear And Trembling" and "Kierkegaard")

(e) Answer: $\theta = \{\text{Author}(y)/\text{Kant}, y/\text{CritiqueOfPureReason}\}$

2. (a) $\text{Philosopher}(c_x) \wedge \text{StudentOf}(c_y, c_x)$

Where c_y and c_x are Skolem constants substituting variables x and y .

(b) $\forall y, x: \text{Philosopher}(x) \wedge \text{StudentOf}(y, x) \rightarrow [\text{Book}(S_z(x, y)) \wedge \text{Write}(x, S_z(x, y)) \wedge \text{Read}(y, S_z(x, y))]$

Where $S_z(x, y)$ is a Skolem function substituting variable z .

3. (a) We start with the CNF form of the SuperActor applying inference rule and Skolemization

- $\forall x: \text{SuperActor}(x) \leftrightarrow [\exists m: \text{PlayedInMovie}(x,m) \wedge \text{Directed}(x,m)]$

Break the double-implication into 2 conjoined implications

$$\begin{aligned} \forall x : (&\text{SuperActor}(x) \rightarrow [\exists m : \text{PlayedInMovie}(x, m) \wedge \text{Directed}(x, m)]) \\ &\wedge ([\exists m : \text{PlayedInMovie}(x, m) \wedge \text{Directed}(x, m)] \rightarrow \text{SuperActor}(x)) \end{aligned} \quad (1)$$

$$\begin{aligned} \Leftrightarrow \forall x : (&\neg \text{SuperActor}(x) \vee [\exists m : \text{PlayedInMovie}(x, m) \wedge \text{Directed}(x, m)]) \\ &\wedge (\neg [\exists m : \text{PlayedInMovie}(x, m) \wedge \text{Directed}(x, m)] \vee \text{SuperActor}(x)) \end{aligned} \quad (2)$$

$$\begin{aligned} \Leftrightarrow \forall x : (&\neg \text{SuperActor}(x) \vee [\exists m : \text{PlayedInMovie}(x, m) \wedge \text{Directed}(x, m)]) \\ &\wedge ([\forall m : \neg \text{PlayedInMovie}(x, m) \vee \neg \text{Directed}(x, m)] \vee \text{SuperActor}(x)) \end{aligned} \quad (3)$$

$$\begin{aligned} \Rightarrow \forall x : (&\neg \text{SuperActor}(x) \vee [\text{PlayedInMovie}(x, F[x]) \wedge \text{Directed}(x, F[x])]) \\ &\wedge ([\forall m : \neg \text{PlayedInMovie}(x, m) \vee \neg \text{Directed}(x, m)] \vee \text{SuperActor}(x)) \end{aligned} \quad (4)$$

$$\begin{aligned} \Rightarrow (&\neg \text{SuperActor}(x) \vee [\text{PlayedInMovie}(x, F[x]) \wedge \text{Directed}(x, F[x])]) \\ &\wedge ([\neg \text{PlayedInMovie}(x, m) \vee \neg \text{Directed}(x, m)] \vee \text{SuperActor}(x)) \end{aligned} \quad (5)$$

$$\begin{aligned} \Rightarrow (&\neg \text{SuperActor}(x) \vee \text{PlayedInMovie}(x, F[x])) \wedge (\neg \text{SuperActor}(x) \vee \text{Directed}(x, F[x])) \\ &\wedge (\neg \text{PlayedInMovie}(x, m) \vee \neg \text{Directed}(x, m) \vee \text{SuperActor}(x)) \end{aligned} \quad (6)$$

Now let us look at the second and third formula

- $\forall m: \text{Directed}(\text{Tarantino}, m) \leftrightarrow \text{PlayedInMovie}(\text{UmaThurman}, m)$

Break the double-implication into 2 conjoined implications

$$\begin{aligned} \forall m : (&\text{Directed}(\text{Tarantino}, m) \rightarrow \text{PlayedInMovie}(\text{UmaThurman}, m)) \\ &\wedge (\text{PlayedInMovie}(\text{UmaThurman}, m) \rightarrow \text{Directed}(\text{Tarantino}, m)) \end{aligned} \quad (7)$$

$$\begin{aligned} \Leftrightarrow \forall m : (&\neg \text{Directed}(\text{Tarantino}, m) \vee \text{PlayedInMovie}(\text{UmaThurman}, m)) \\ &\wedge (\neg \text{PlayedInMovie}(\text{UmaThurman}, m) \vee \text{Directed}(\text{Tarantino}, m)) \end{aligned} \quad (8)$$

$$\begin{aligned} \Rightarrow (&\neg \text{Directed}(\text{Tarantino}, m) \vee \text{PlayedInMovie}(\text{UmaThurman}, m)) \\ &\wedge (\neg \text{PlayedInMovie}(\text{UmaThurman}, m) \vee \text{Directed}(\text{Tarantino}, m)) \end{aligned} \quad (9)$$

- $\exists m: \text{PlayedInMovie}(\text{UmaThurman}, m) \wedge \text{PlayedInMovie}(\text{Tarantino}, m)$

$$\begin{aligned} \exists m : &\text{PlayedInMovie}(\text{UmaThurman}, m) \wedge \text{PlayedInMovie}(\text{Tarantino}, m) \\ \Rightarrow &\text{PlayedInMovie}(\text{UmaThurman}, c) \wedge \text{PlayedInMovie}(\text{Tarantino}, c) \end{aligned} \quad (10)$$

Now we add the hypothesis $\neg \text{SuperActor}(\text{Tarantino})$ and apply resolution rule until we achieve a contradiction.

$$\neg \text{SuperActor}(\text{Tarantino}) \quad (11)$$

$$\text{PlayedInMovie}(\text{UmaThurman}, c) \quad (12)$$

$$\text{PlayedInMovie}(\text{Tarantino}, c) \quad (13)$$

$$\neg \text{PlayedInMovie}(\text{UmaThurman}, m) \vee \text{Directed}(\text{Tarantino}, m) \quad (14)$$

$$\neg \text{Directed}(\text{Tarantino}, m) \vee \text{PlayedInMovie}(\text{UmaThurman}, m) \quad (15)$$

Unification of 12, 14 and 15

$$\neg \text{PlayedInMovie}(\text{UmaThurman}, c) \vee \text{Directed}(\text{Tarantino}, c) \quad (16)$$

$$\neg \text{Directed}(\text{Tarantino}, c) \vee \text{PlayedInMovie}(\text{UmaThurman}, c) \quad (17)$$

Conjunction of 12 and 16

$$\Rightarrow \text{PlayedInMovie}(\text{UmaThurman}, c) \wedge \text{Directed}(\text{Tarantino}, c) \quad (18)$$

$$\Rightarrow \text{Directed}(\text{Tarantino}, c) \quad (19)$$

Conjunction of 19 and 13

$$\Rightarrow \text{PlayedInMovie}(\text{Tarantino}, c) \wedge \text{Directed}(\text{Tarantino}, c) \quad (20)$$

Given from 6

$$\neg \text{PlayedInMovie}(x, m) \vee \neg \text{Directed}(x, m) \vee \text{SuperActor}(x) \quad (21)$$

Conjunction of 20 and 21

$$\Rightarrow \text{PlayedInMovie}(\text{Tarantino}, c) \wedge \text{Directed}(\text{Tarantino}, c) \wedge \text{SuperActor}(x) \quad (22)$$

$$\Rightarrow \text{SuperActor}(x) \quad (23)$$

$$\Rightarrow \mathbf{F} \quad (24)$$

- (b) Translate the information given in FOL into English (or Norwegian) and describe in high level the reasoning you could apply in English to have the same result (in other words, describe a proof of the result in natural language).

A SuperActor is someone that is a director and an actor in the same film. Uma Thurman is performing in a movie if and only if Tarantino is the director. There is a movie that has Uma Thurman and Tarantino as actors in it. We know that everytime Uma Thurman is playing a character in a movie Tarantino is directing. So in this specific movie Tarantino is director and actor. So we can conclude that he is an SuperActor.