# Threading strategy and program overview

This program demonstrates a multi-threaded packet sniffing application that captures packets from a specified network interface using libpcap and libpthreads. The program utilizes a thread pool design to distribute packet analysis work among a fixed number of worker threads.

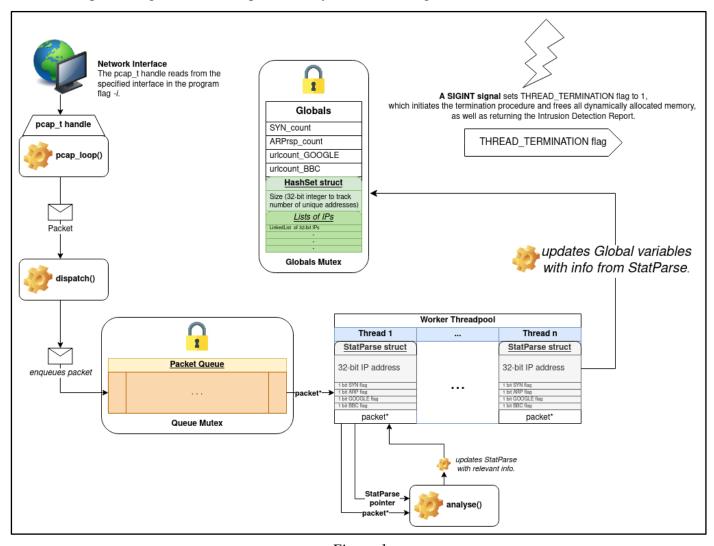


Figure 1

Figure 1 provides a diagrammatical overview of the programs execution, starting from the network interface in the upper left, and terminating upon detection of a SIGINT signal as represented in the upper right.

#### Program Overview:

#### 1. Packet Handling:

The pcap\_loop() function captures packets continuously, invoking the dispatch() function for each packet, which enqueues the packet's header in binary format to a FIFO work queue. This is performed under lock of the queue

*mutex*. Once enqueued, the pthread\_cond\_broadcast() function is invoked which wakes all sleeping threads.

## 2. Thread work and packet analysis:

- o Worker threads are responsible for packet analysis. They either wait for packets in the queue, or sleep until there is a packet in the queue, and then analyse the packet data using the analyse () function. After analysis, the thread updates global variables after obtaining a lock on the *globals mutex*.
- o A StatParse struct in each thread tracks detected attacks and relevant information for each packet and is reset at the beginning of each analysis.
- o A termination signal (SIGINT) triggers the termination process, signaling worker threads to complete their tasks and clean up resources.

#### 3. Data Structures:

- o Queue in queue.c implements a basic FIFO queue used to store packet information for worker threads.
- o HashMap in hashmap.c implements a simple hash set to store unique IP addresses involved in SYN attacks.

# Justification for the Thread-pool model

The chosen threading model involves a thread pool with a fixed number of worker threads (NUMTHREADS). This approach was selected for several reasons:

## 1. Efficiency:

Due to the very large number of packet throughput on any given interface, thread creation for each packet would cause excessive overhead, and therefore a fixed pool of threads can more efficiently handle packet analysis.

#### 2. Resource Utilization:

Resource exhaustion is limited by the number of concurrent threads to a manageable level, based on a system's capabilities. The threads also share resources from the process stack, such as the packet queue, the global variables, and the hashset of unique IP addresses.

## 3. Concurrency:

Multi-threading allows simultaneous processing of packets, maximizing the utilization of multi-core systems. Thread switching is also largely faster than context switching, and as such a multithreaded model is preferred to multi-process.

### 4. Scalability:

The number of threads can also be adjusted based on system resources and performance requirements. Since the DCS virtual machine only offers 1 core, 4 threads were decided as an estimate on the number of cores for an actual system running the program.

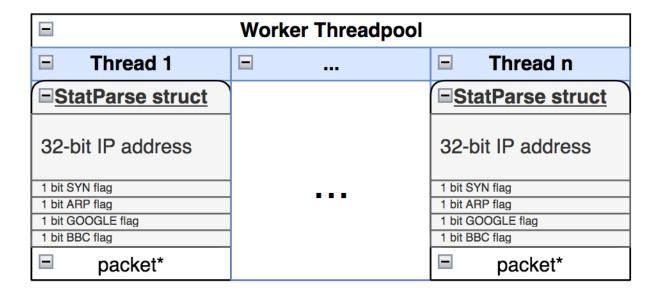


Figure 2: The thread-pool and thread stack

The stack of each thread is represented diagrammatically in Figure 2. Each thread has a StatParse struct, defined in analysis.h. This struct allocates memory for an IP as a 32 bit unsigned integer, and 4 x 1-bit flags of type uint8\_t (or char). It also contains a pointer to the most recently dequeued packet data. The StatParse struct utilizes 1-bit flag fields to not only reduce the amount of memory that the program interacts with, but also allows for the program to perform bit setting efficiently using bitwise operations. The need for each thread to perform it's packet analysis with computational efficiency is noticed when observing the number of packets per second possibly present in modern day ethernet routers. The IEEE 802 Standards Committee recently addressed for example, 100 and up to 400 Gbps ethernet speeds<sup>1</sup>, which can mean millions of packets per second. As such, reducing the overhead of operations performed in the threads' analysis program are vital.

The following code snippet exhibits the utilisation of bitwise operations to improve computational speeds:

```
uint8_t flags = tcp_header->th_flags;
if ((flags & TH_SYN) && !(flags & TH_ACK)) {
    parse->SYN |= 1;
}
```

<sup>1</sup> [3]

As a SYN attack is comprised via the flooding of packets with the SYN TCP flag ON, and the ACK flag OFF, these must be checked by the program. To do this efficiently, the BITWISE AND operation is performed on the flag field to ensure that this is true.

To set the SYN bit in the StatParse struct to 1, instead of performing SYN++, SYN+=1 or SYN=1, SYN | =1 is performed instead. Figure 3 and Figure  $4^2$  compare the x86 assembly instructions of SYN++ and SYN | =1.



Figure 3: x86-64 assembly instructions for SYN++ compiled in gcc 13.2 highlighted in red

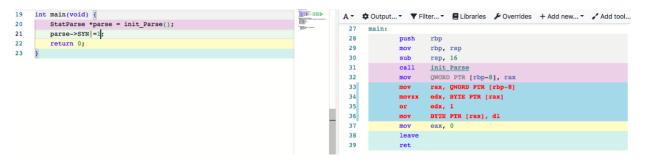


Figure 4: x86-64 assembly instructions for SYN | =1 compiled in gcc 13.2 highlighted in red

As we can see, SYN++ requires 12 assembly instructions, whereas SYN = 1 requires only 4. Furthermore, SYN=1 compiles to the exact same instructions as SYN = 1, performing the same assembly operation: or [REG], 1 in order to set the bit. Using SYN = 1 ensures that if a user was using a different compiler, they would still be performing the same 4 instructions whereas SYN=1 may give different results.

A hash set was used to store addresses as it would be largely space inefficient to store use a dynamic array in a very high throughput network, such as a datacentre. It is therefore optimal to implement a hashing function which both achieves low computational cost and even

 $<sup>^{2}[1]</sup>$ 

distribution of hash results. I used a 32-bit hash function created by Thomas Mueller of H2 Database which was shown to have an avalanche effect better than MurmurHash3, which uses computationally cheap *xorshift-multiply-xorshift* operations, as well as using the same multiplicative constant twice<sup>3</sup>, for the best performance.

# **Testing**

To test the program, numerous tests were performed. On top of successfully passing all tests provided within the coursework specification<sup>4</sup>, SYN packet stress tests were performed on quantities of 30000, 100000 and 1000000 packets. All tests successfully detected all SYN packets and unique IP addresses with no memory leaks on program termination as displayed through valgrind --leak-check=full.

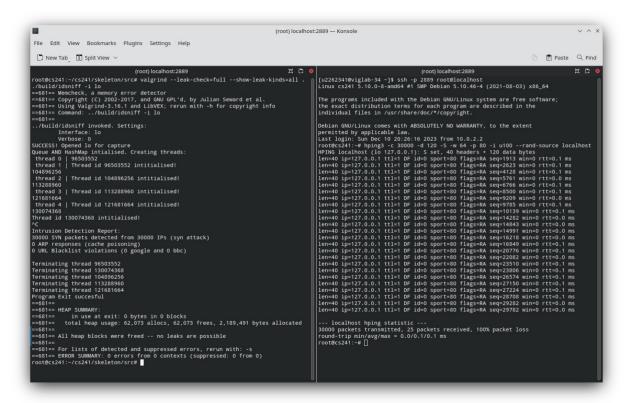


Figure 5: 30000 packet stress test with no memory leaks possible.

 $<sup>^{3}[4]</sup>$ 

<sup>&</sup>lt;sup>4</sup> [2]

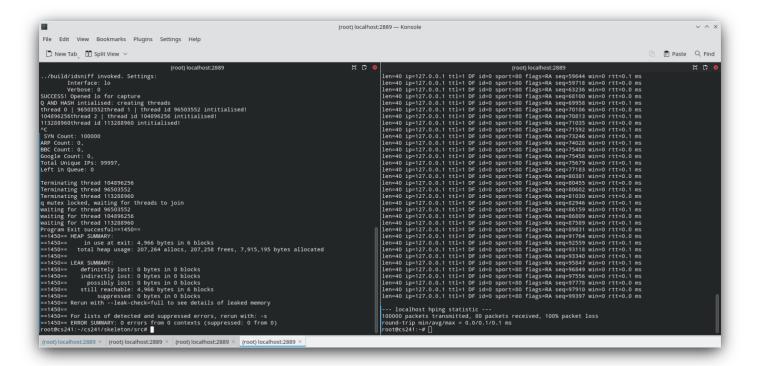


Figure 6: 1000000 packet stress test with no memory leaks

# Bibliography

- [1] "Compiler Explorer," 2023. [Online]. Available: https://godbolt.org/.
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