



https://skepu.github.io

Portable Programming of Heterogeneous Parallel Systems with SkePU

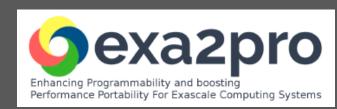
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Updated slides at: https://skepu.github.io/tutorials/escience21





Acknowledgments

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Details can be found in **August Ernstsson's Licentiate thesis**https://doi.org/10.3384/lic.diva-170194 and in the following open-access article:

August Ernstsson, Johan Ahlqvist, Stavroula Zouzoula, Christoph Kessler: "SkePU 3: Portable High-Level Programming of Heterogeneous Systems and HPC Clusters." *International Journal of Parallel Programming*, Springer, May 2021 (online), print version to appear. https://link.springer.com/article/10.1007/s10766-021-00704-3

Many **others** have contributed to SkePU over the years, too, see the SkePU web page for acknowledgments.

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Agenda

Introduction to Skeleton Programming and SkePU Overview

SkePU Fundamentals

SkePU API

Examples

SkePU Advanced Topics







https://skepu.github.io

Introduction to Skeleton Programming and SkePU Overview

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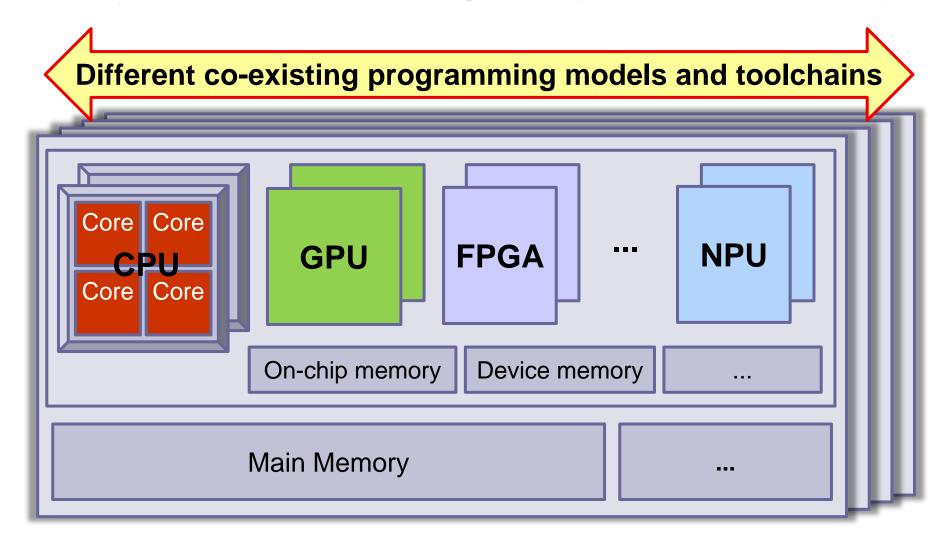
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Parallelism and Heterogeneity are here to stay ...

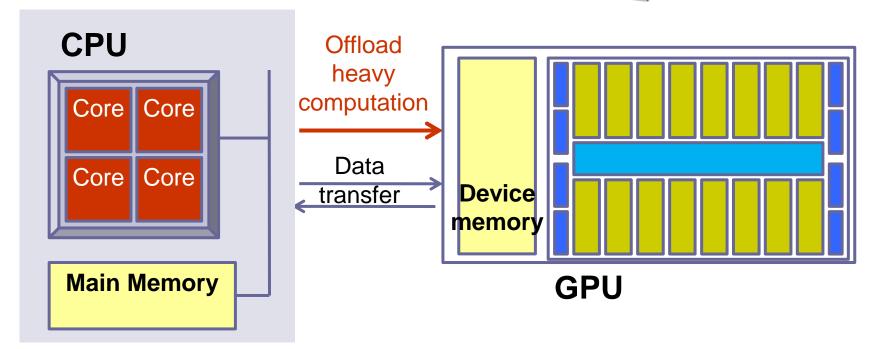


→ Software Productivity ??

Example:



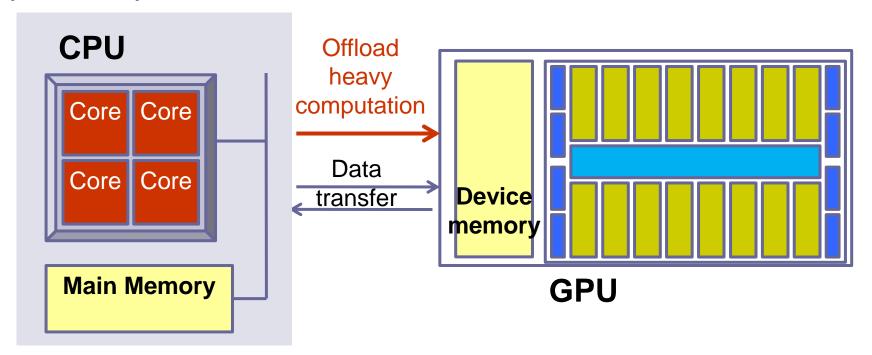




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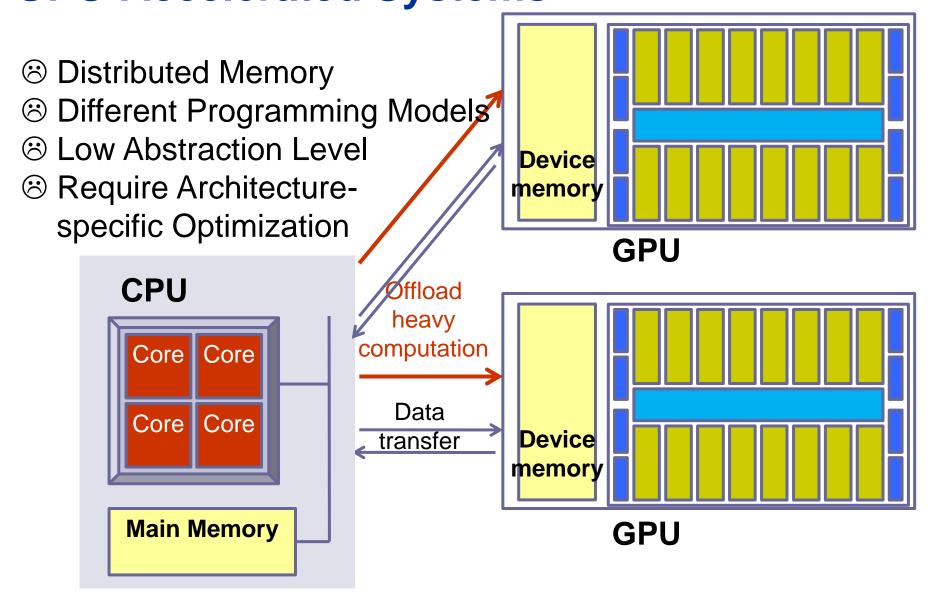
Example: GPU-Accelerated Systems

- Distributed Memory
- Different Programming Models
- Low Abstraction Level
- Require Architecture specific Optimization

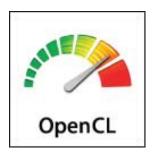




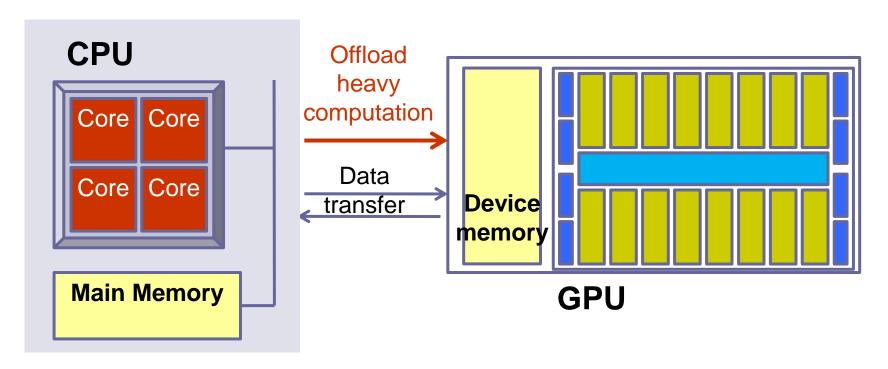
Example: GPU-Accelerated Systems



Programming of GPU-accelerated Systems ... with OpenCL™



- Code portability ©
- Programmability (low level)
- Performance portability (requires reoptimization)





Programming of Clusters ... with MPI

- Code portability ©
- Programmability (low level)
- Performance portability (8)





Example: 1D filter in C

```
float filter (float a, b, c) { return 0.25*a + 0.5*b + 0.25*c; }
void main ( int argc, char *argv[] )
 float *array ( n+2 );
 float *tmp ( n+2 );
 while ( globalerr > 0.1 ) {
                                                          array
      for (i=1; i<=n; i++)
         tmp[i] = filter( array[i-1], array[i], array[i+1] );
      globalerr = 0.0;
      for (i=1; i<=n; i++)
          globalerr = fmax (globalerr, fabs( array[i] – tmp[i] ));
      for (i=1; i<=n; i++)
         array[i] = tmp[i];
```

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Example: 1D filter in C + MPI

```
void main ( int argc, char *argv[] ) {
        MPI Comm com = MPI COMM WORLD;
        MPI Init (&argc, &argv);
        MPI Comm size (com, &np);
        MPI_Comm_rank ( com, &me );
         localsize = (int) ceil ( (float) n / np );
         FloatArray local( localsize + 2);
        while (globalers > 0.34
                                                                                                                                                                                                   Teft_neighbor, 10, com);
                    if (me>0)
                                                                     NPL Send (local+1, 1
                               Ti=1: <=localsite
                             tmp[i = file / // cani-1], localii, localii (p
                   if (me
                                                                                                                                                 NP_FLOAT, right_poi_por, 10, com,
                                                                                                                                                                                                          F CAT, left_neighbor
                   if (me>0)
                                                                     MPI_Recv (thep-localsize+1, 1
                    tmp[1] = filter( local[0], local[1]_local[2]
                    tmp[localsize] = filter( local[localsize], local[localsize]
                    localerr = 0.0:
                                                                                                                                                      Tabs ( loc [i]-tmp[i] ));

The second of the
                   for (i=1; i<=localsize; i++) localerr= for ax(localerr, fabs (localerr, fabs (localerr, fabs (localerr));
                    MPI Allreduce ( & localerr, & granderr,
```

Our goal:

Raise the level of abstraction

- * towards portable, future-proof programs for heterogeneous parallel systems
- * Try to make it as easy as sequential programming

Techniques:

- Skeleton Programming
- Portable Data Abstractions (Data-Containers)





Observation

- The same (platform-specific) technique for identifying/managing parallelism, communication, synchronization... is re-applicable for all occurrences of the same specific structure (pattern) of computation
 - Elementwise operations on arrays
 - Stencil computations
 - Reductions
 - Scan (Prefix-op)
 - Divide-and-Conquer
 - Farming independent tasks
 - Pipelining
 - ...
- Most of these have both sequential and parallel implementations
- □ Idea: **Reusable** (customizable) generic constructs (*skeletons*)



Data Parallelism

Given:

- Operand data-containers $\boldsymbol{a}, \boldsymbol{b}, \ldots$ with \boldsymbol{n} elements each, e.g., array(s) $\boldsymbol{a} = (a_1, \ldots, a_n), \ \boldsymbol{b} = (b_1, \ldots, b_n), \ldots$
- A side-effect-free operation f on elements of a, b, ...
 - Any arity: 0, 1, 2 or more input operands, 1 or more output operands

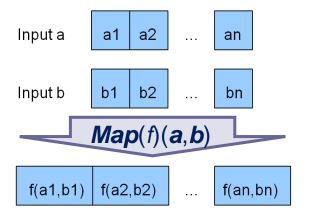
Compute: $(f(a_1,b_1,...), ..., f(a_n,b_n,...))$

Parallelizability:

- All elementwise computations are independent
- Easily partitioned into independent tasks of arbitrary granularity

Notation with higher-order function:

$$y = Map(f)(a, b, ...)$$



Result



Data Parallelism with Stencils

A generalization of *Map*:

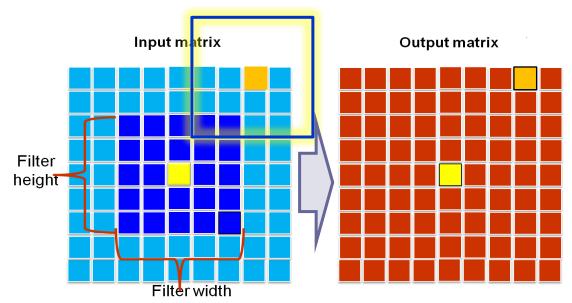
For calculating **y**[i], can access elements in **a** within a small (constant-sized) halo **region** around element index i

□ Access window on \mathbf{a} , e.g., for 1D region: $\mathbf{a}[i+K]$... $\mathbf{a}[i+K]$ for given K>0

For stencil computations (filters, convolutions, PDE solving, ...)

Typically, 2D, 3D or 4D operands and regions

y = MapOverlap (filterfunction) (a, ...)



Further customization:

Special treatment of elements at domain boundaries (edges)

- constant padding
- duplication padding
- cyclic / toroidal padding
- ..

Update schema

- Jacobi style
- In-place, Red-Black
- ..



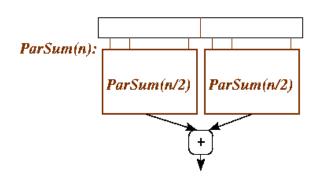
Data-Parallel Reduction

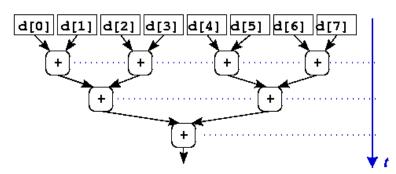
Given:

- A data container **d** with *n* elements, e.g. array $\mathbf{d} = (d_1, ..., d_n)$
- A <u>binary</u>, <u>associative</u> operation *op* on individual elements of *d* (e.g. *add*, *max*, *bitwise-or*, ...)

Compute: $s = OP_{i=1...n} d = d_1 op d_2 op ... op d_n$

Parallelizability: Exploit associativity of op





Notation with higher-order function:

$$s = Reduce(op)(a)$$

Further customization:

For higher-dimensional data-containers, e.g. matrix: total / row-wise / column-wise



MapReduce (pattern)

Common combination:

$$a, \dots \longrightarrow Map(f, a, \dots) \xrightarrow{y} Reduce(g, y, \dots) \longrightarrow s$$

$$s = MapReduce(f, g)(a, ...)$$

Example:

Dot product of two vectors \mathbf{a} , \mathbf{b} : $s = \sum_i a_i * b_i$

f =scalar multiplication,

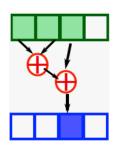
g = scalar addition



Further Patterns ...

More Data-parallel patterns

- Scan (prefix-sums)



Task-parallel patterns

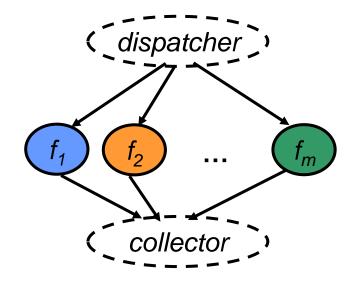
- Task farming
- Parallel divide-and-conquer
- ---

Stream-parallel patterns

- Pipelining
- Stream farming
- ...

Domain-specific patterns

o ...

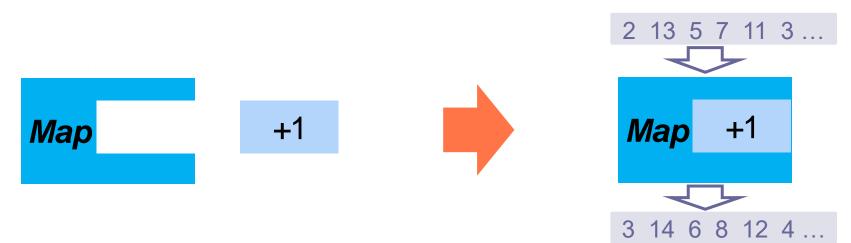




Skeleton / Pattern Programming

[Cole'89]

- A skeleton is a pre-defined generic program building block based on a higher-order function that
 - models a common computation / dependence pattern map, reduce, scan, stencil, farm, pipe, ...
 - can be parameterized with sequential user code
 - provides a sequential interface
 - can underneath have optimized implementations



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Example revisited: 1D filter in C++ Skeletons

```
float filter (Region1D<float> r, Vec<float> w) { return w[0]*r[-1] + w[1] * r[0] + w[2] * r[1]; }
float fabsdiff (float a, b) { return fabs (a - b); }
float idem ( float a ) { return a; }
                                                         Generic data-containers
void main ( int argc, char *argv[] )

    used as operands

  Vector<float> array(n);

    encapsulate metadata

  Vector<float> tmp ( n );
                                                           (e.g., size)
  Vector<float> weights { 0.25, 0.5, 0.25 };
  // skeleton instantiations:
  auto stencil = MapOverlap(filter); stencil.setOverlap(1);
  auto maxchange = MapReduce( fabsdiff, fmax );
  auto copy = Map( idem );
  while ( globalerr > 0.1 ) {
     stencil( tmp, array, weights, n);
     globalerr = maxchange ( array, tmp );
     copy( array, tmp);
```



High-Level Parallel Programming with Skeletons

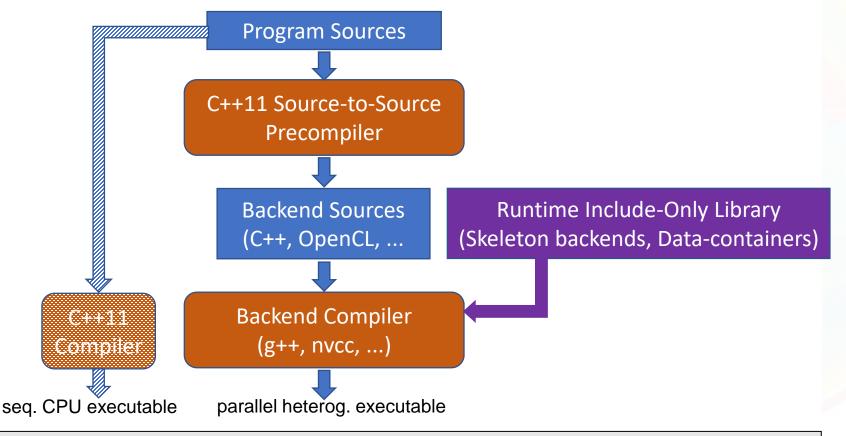
Skeletons + Data-containers implement (parallel) algorithmic patterns

- Abstraction, hiding complexity (parallelism and low-level programming)
- Sequential(-looking) programming interface Parallelization for free
- Easier to analyze and transform
- Portability
- Enforces code (re)structuring
- Available skeleton set does not always fit
- May lose some efficiency compared to manual parallelization
- Idea developed since the late 1980s [Cole'89]
- Many (esp., academic) frameworks exist, mostly as libraries atop C++
 - actively researched: FastFlow, SkePU, GrPPI, LiFT, Musket, ...
- Industry has adopted skeletons, too
 - map, reduce, scan in many modern parallel programming APIs
 - e.g., Intel *TBB*: parallel for, parallel reduce, pipe
 - Thrust, MapReduce, Spark etc.



SkePU

- C++11 based
 - Every SkePU program is a valid C++11 program
 - Variadic template metaprogramming include library and a source-to-source pre-compiler



A. Ernstsson, L. Li, C. Kessler: **SkePU 2: Flexible and type-safe skeleton programming for heterogeneous parallel systems.** *International Journal of Parallel Programming* 46(1):62, Jan. 2018



SkePU 3

- C++11 based
 - Every SkePU program is a valid C++11 program
 - Variadic template metaprogramming include library and a source-to-source pre-compiler
- Data-parallel skeletons
 - Map, Reduce, MapReduce, MapOverlap, MapPairs, MapPairsReduce, Scan, ...
 - Fully variadic in elementwise, random-access and scalar operands
- STL-like data containers wrapping operand arrays
 - Vector<..>, Matrix<..>, Tensor3<..>, Tensor4<..>
 - Container windows for different operand access patterns in user functions: e.g., elementwise; **Region** (stencils); **Vec**, **Mat** (random access); **MatRow**; ...
- Multiple back-ends:
 - C, OpenMP, OpenCL, CUDA
 - New hybrid stand-alone backend [Öhberg, Ernstsson, K. 2018]
 - Multi-GPU support
 - Cluster backend for StarPU-MPI
- ◆ Tunable for time [Dastgeer et al.'11,'13], energy [Li, K. J Supercomp.'16]
- Open source: https://skepu.github.io



SkePU Example: Dot Product

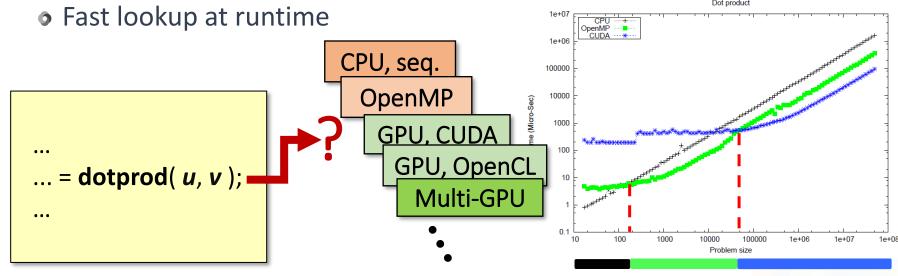
```
#include <skepu>
                                                 User function
                                            (platform-independent,
float add (float a, float b)
                                                side-effect free)
  return a + b;
float mult ( float a, float b )
                                             Data-containers
  return a * b;
                                              for array-based
                                          skeleton-call operands
int main()
  ... arrays A, B ...
                                              Skeleton instantiation:
  Vector<float> u(A, N);
                                        generates a multi-backend function
  Vector<float> v(B, N);
  auto dotprod = MapReduce<2>(mult, add);
  \dots = dotprod(u, v);
                                    Call the skeleton instance
                                                                  H2020, Project No. 801015
```



Automatic Backend Selection for a Call

• For a **skeleton-instance call**, find *automatically* the expected **best backend** and settings depending on execution context (operand data size / location)

Off-line training of performance models



- Or pre-select a backend manually
 - global default or call-specific



New in SkePU 3: Multi-Variant User Functions

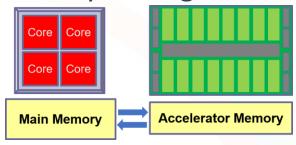
- In *addition* to the default (universal) user function, further **platform-specific user function variants** can be provided e.g. by a platform expert programmer.
 - Can also apply to multiple elements at a time
 - SkePU generates additional backend variants from these, selectable only if the feature is available on the target system
- Example: Custom SIMDization, special instructions





Smart Data-Containers

- Extended generic STL-like containers (Vector<>, Matrix<>, Tensor3<>, Tensor4<>)
- Software caching of operand element data in device memories
- Runtime optimization of data transfers, memory management
 - Lazy data transfer [Enmyren, K.'10]
- Iterators
 - → Coherence at arbitrary *sub*array level [Dastgeer, K. *IJPP*'16]
 - Copy plan locate closest overlapping valid copies
 - Reuse invalidated allocations of device memory
 - Speedup (up to 10³ x over 'dumb' containers on iterative applications)



```
skepu::Vector<float> a, w, x; ... skel1 ( w, a ); OCPU? GPU1? ... Write w Read w CPU? GPU1? ...
```



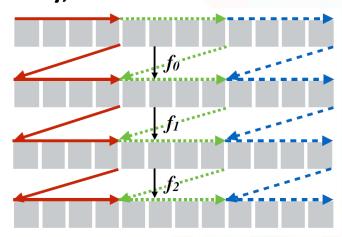
Lazy Execution of Transformation-Style (Map, MapOverlap) Skeleton

- Lazy execution of Map* calls: build run-time dependence graph
 - Lineage of data-containers cf. Spark RDDs
 - Not limited to static scope e.g., entire iterative stencil loop
- Lineage is evaluated when needed (e.g., at a reduction/scan) by **tiling** the postponed **Map*** computations
 • Efficient cache utilization on CPU, local memory on accelerator

 - Tile size is a tuning parameter

Standard (eager) execution Map* call 0 Map* call 1 Map* call 2

Lazy, tiled execution





Cluster Execution Support in SkePU

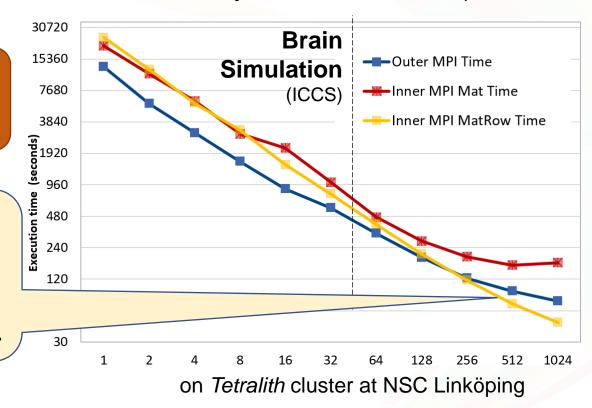
2 "modes" of cluster support in SkePU code:

- Outer MPI (user-written MPI + SkePU at node level)
- Inner MPI (SkePU uses MPI internally cluster backend)

SPMD exec. abstracted away.
No syntactic difference from single-node SkePU code

Poor scaling if using the most general matrix access window (Mat<>, full random access).

MatRow<> scales much better.



A. Ernstsson, J. Ahlqvist, S. Zouzoula, C. Kessler: **SkePU 3: Portable high-level programming of heterogeneous systems and HPC clusters.** Accepted for *Int. J. of Parallel Programming*, 2021



Programmability: One Code to Rule Them All

```
Sequential C++:
                          int main()
                            double s:
                            ... arravs A. B ...
OpenMP:
                            s = 0.0:
                            for (i=0; i<N; i++)
  #include <omp.h>
                               s += A[i] * B[i];
  int main()
    double s:
   ... arrays A, B ...
   s = 0.0:
  #pragma omp parallel shared(s, A, B)
  #pragma omp for reduction ( s: + )
    for (i=0; i<N; i++)
       s += A[i] * B[i];
```

MPI:

```
#include <mpi.h>
 MPI_Scatter( A, N, MPI_double, myA,myN,
 for (i=0; i<myN; i++)
    mys += myA[i] * myB[i];
 MPI_Reduce( &mys, 1, MPI_double,
              &s. 1. MPI double. 0.
```

```
Valid C++(11+) code.
  No need to write
platform-specific code
```

Hybrid:

OpenCL:

SkePU source code:

```
return a + b:
float mult (float a, float b)
 return a * b:
int main()
  ... arrays A, B ...
 Vector<float> u(A, N);
 Vector<float> v(B, N):
 auto dotprod = MapReduce<2>
 \dots = dotprod(u, v);
```

float add (float a, float b)

CUDA:

```
#include <cuda.h>
                   global void dot (int *a, int *b, int *s)
                     shared int temp[THREADS PER BLOCK]
                    int index = threadldx.x + blockldx.x * blockDin
                   temp[threadldx.x] = a[index] * b[index];
                    syncthreads():
                   if (threadIdx.x == 0) {
                     int sum = 0:
                     for (int i = 0; i < N; i++) {
                       sum += temp[i];
                     atomicAdd(s, sum);
                 int main()
                    ... allocate arrays A, B.
No complex code base
  cluttered with #ifdefs
   for various platform-
                                                   cpyHos
                                                   cpvHos
    specific code versions
                                                    ovHost
                                              >> >(dev A. d
                                    s, sizeof(int), cudaMemor
                    cudaFree(s):
```



SkePU Evolution Milestones











	SkePU 1 (2010)	SkePU 2 (2016)	SkePU 3 (2020)
API based on	C, C++ (pre-2011), C preprocessor	C++11, Precompiler (clang)	C++11, Precompiler
Skeletons	Map, Reduce, Scan, MapReduce, MapArray, MapOverlap, Generate	Map, Reduce, Scan, MapReduce, MapOverlap, Call	Map, Reduce, MapReduce, MapPairs/MapPairsReduce, Scan, MapOverlap№, Call
User functions as	C preprocessor macros Not type-safe	C++ functions Multi-way variadic Type-safe	As SkePU 2, + multi-variant user func.s + multi-value return
Data- containers, Access proxy	Vector<>, Matrix<>	Vector<>, Matrix<>	Vector<>, Matrix<>, Tensor3<>, Tensor4<>; MatRow<>, RegionND<>,
Platforms supported	CPU (C, OpenMP), GPU (CUDA, OpenCL); Myriad	CPU (C++, OpenMP), GPU (CUDA, OpenCL)	CPU, GPU, hybrid CPU/GPU, Cluster (StarPU-MPI)
Memory model	Sequential consistency	Seq. consistency	Weak consistency (default), alternatively sequential c.



Summary, so far

- Goal: Simplify the use of heterogeneous parallel systems and enable writing future-proof, portable, high-level code
- Skeletons implement parallelizable patterns
- Single-source, sequential-looking code
- SkePU skeleton programming framework
 - API: Multi-variadic, customizations, type-safe, compatibility with C++11
 - Smart data-containers
 - Data-container access windows model access patterns
 - Multi-backend, tunable selection
- Open source, https://skepu.github.io
- Used in research and teaching
- Lots of ideas for the future
- Cooperations welcome!

Now let's take a closer look at SkePU ...

Thanks!

skepu.github.io





Glossary

Performance Portability

... is the ability of a program to automatically adapt to a new execution platform to achieve an automated best-effort optimization of performance on the new target system, without manual rewriting / reoptimization.

[Algorithmic] Skeleton

is a pre-defined, generic software construct for high-level programming that implements a specific *pattern* of control and data flow, that can be *parameterized* by problem-specific code to instantiate a problem-specific function, and whose implementation internally *encapsulates* all platform-specific details such as parallelism, heterogeneity, communication and synchronization.



References on SkePU

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August Ernstsson: **Designing a Modern Skeleton Programming Framework for Parallel and Heterogeneous Systems**. Licentiate Thesis no. 1886, Nov. 2020, Linköping University. DOI: 10.3384/lic.diva-170194 https://doi.org/10.3384/lic.diva-170194

August Ernstsson, Johan Ahlqvist, Stavroula Zouzoula, Christoph Kessler: **SkePU 3: Portable high-level programming of heterogeneous systems and HPC clusters.** *Int. J. of Parallel Programming,* May 2021. https://link.springer.com/article/10.1007/s10766-021-00704-3

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SkePU documentation and download: https://skepu.github.io