Algoritmi trattati a lezione

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```
\begin{cases} \textbf{proc InsertionSort}\left(A\right) \\ \textbf{for } (j=2 \textbf{ to } A.length) \\ \left\{ key = A[j] \\ i = j - 1 \\ \textbf{while } ((i>0) \textbf{ and } (A[i] > key)) \\ \left\{ A[i+1] = A[i] \\ i = i - 1 \\ A[i+1] = key \\ \end{cases} \right.
```

```
\begin{array}{l} \textbf{proc RecursiveBinarySearch} \left(A,low,high,k\right) \\ \textbf{ if } \left(low>high\right) \\ \textbf{ then return nil} \\ mid = (high+low)/2 \\ \textbf{ if } \left(A[mid]=k\right) \\ \textbf{ then return } mid \\ \textbf{ if } \left(A[mid]<k\right) \\ \textbf{ then return } RecursiveBinarySearch(A,mid+1,high,k) \\ \textbf{ if } \left(A[mid]>k\right) \\ \textbf{ then return } RecursiveBinarySearch(A,low,mid-1,k) \\ \end{array}
```

```
\begin{cases} \textbf{proc SelectionSort}\left(A\right) \\ \textbf{for } (j=1 \textbf{ to } A.length-1) \\ \begin{cases} min=j \\ \textbf{for } (i=j+1 \textbf{ to } A.length) \\ \textbf{if } (A[i] < A[min]) \\ \textbf{then } min=i \\ SwapValue(A,min,j) \end{cases} \end{cases}
```

```
\begin{cases} & \text{proc } \textit{Merge} \left( A, p, q, r \right) \\ & n_1 = q - p + 1 \\ & n_2 = r - q \\ & \text{let } L[1, \dots, n_1] \text{ and } R[1, \dots, n_2] \text{ be new array} \\ & \text{for } (i = 1 \text{ to } n_1) \ L[i] = A[p + i - 1] \\ & \text{for } (j = 1 \text{ to } n_2) \ R[j] = A[q + j] \\ & i = 1 \\ & j = 1 \\ & \text{for } (k = p \text{ to } r) \\ & \left\{ \begin{array}{l} \text{if } (i \leq n_1) \\ \text{then} \\ \text{of } & \text{if } (L[i] \leq R[j]) \\ \text{then } & \text{copyFrom} L(i) \\ \text{else } CopyFrom} L(i) \\ \text{else } C
```

```
\begin{aligned} & \textbf{proc MergeSort}\left(A,p,r\right) \\ & \textbf{ (if } (p < r) \\ & \textbf{ then } \\ & \left\{ \begin{aligned} q &= [(p+r)/2] \\ MergeSort(A,p,q) \\ MergeSort(A,q+1,r) \\ Merge(A,p,q,r) \end{aligned} \right. \end{aligned}
```

```
\begin{aligned} & \operatorname{proc} \operatorname{Partition}\left(A,p,r\right) \\ & \begin{cases} x = A[r] \\ i = p-1 \\ & \text{for } (j = p \text{ to } r-1) \\ & \text{ if } (A[j] \leq x) \\ & \text{then} \\ & \begin{cases} i = i+1 \\ SwapValue(A,i,j) \\ SwapValue(A,i+1,r) \\ & \text{return } i+1 \\ \end{aligned} \end{aligned}
```

```
 \begin{aligned} & \textbf{proc } \textit{QuickSort} \left(A, p, r\right) \\ & \textbf{ if } \left(p < r\right) \\ & \textbf{ then } \\ & \left\{q = Partition(A, p, r) \\ & QuickSort(A, p, q - 1) \\ & QuickSort(A, q + 1, r) \\ \end{aligned} \right.
```

```
\begin{aligned} & \textbf{proc } \textit{RandomizedPartition } (A, p, r) \\ & \begin{cases} s = p \leq Random() \leq r \\ SwapValue(A, s, r) \end{cases} \\ & x = A[r] \\ & i = p - 1 \\ & \textbf{for } j = p \textbf{ to } r - 1 \\ & \begin{cases} \textbf{if } (A[j] \leq x) \\ \textbf{then} \\ & \begin{cases} i = i + 1 \\ SwapValue(A, i, j) \end{cases} \\ SwapValue(A, i + 1, r) \\ \textbf{return } i + 1 \end{cases} \end{aligned}
```

```
\begin{cases} & \text{proc } \textit{RandomizedQuickSort}\left(A,p,r\right) \\ & \text{if } (p < r) \\ & \text{then} \\ & \begin{cases} q = RandomizedPartition(A,p,r) \\ RandQuickSort(A,p,q-1) \\ RandQuickSort(A,q+1,r) \end{cases} \end{cases}
```

```
\begin{cases} \textbf{if} \ (Empty(S)) \\ \textbf{if} \ (Empty(S)) \\ \textbf{then return} \ ''underflow'' \\ S.top = S.top - 1 \\ \textbf{return} \ S[S.top + 1] \end{cases}
```

```
\begin{aligned} & \textbf{proc } \textit{Enqueue} \left(Q,x\right) \\ & \textbf{if } \left(Q.dim = Q.length\right) \\ & \textbf{then } \textbf{return''} over flow'' \\ & Q[Q.tail] = x \\ & \textbf{if } \left(Q.tail = Q.length\right) \\ & \textbf{then } Q.tail = 1 \\ & \textbf{else } Q.tail = Q.tail + 1 \\ & Q.dim = Q.dim + 1 \end{aligned}
```

```
\begin{aligned} & \textbf{proc Dequeue} \left(Q\right) \\ & \textbf{if } \left(Q.dim = 0\right) \\ & \textbf{then return }''underflow'' \\ & x = Q[Q.head] \\ & \textbf{if } \left(Q.head = Q.length\right) \\ & \textbf{then } \left(Q.head = 1\right) \\ & \textbf{else } \left(Q.head = Q.head + 1\right) \\ & Q.dim = Q.dim - 1 \\ & \textbf{return } x \end{aligned}
```

```
\mathbf{proc}\; \mathit{Empty}\; (S)
             if (S.head = nil)
                then return true
                                                      proc Push(S, x)
              return false
                                                       \{Insert(S, x)\}
                       proc Pop (S)
                         (if (Empty(S)))
then return "underflow"
                          x = S.head
                         Delete(S, x)
                         return x.key
       \operatorname{proc} \operatorname{{\it Empty}}(Q)
         if (Q.head = nil)
           then return true
                                                      proc Enqueue (Q, x)
          return false
                                                       \{Insert(Q, x)\}
                       \mathbf{proc}\; \mathbf{\textit{Dequeue}}\;(Q)
                         if (Empty(Q))
then return "underflow"
                          x = Q.tail
                          Q.tail = x.prev
                          Delete(Q, x)
                         return x.key
                      proc Enqueue (A, i, priority)
                       (if (i > A.length)
                          then return "overflow"
                         A[i] = priority
             \operatorname{proc} \operatorname{\textit{DecreaseKey}}\left(A, i, priority\right)
             (if ((A[i] < priority) or (A[i].empty = 1))
                 then return "error
               A[i] = priority
     proc ExtractMin (A)
       MinIndex = 0
       MinPriority = \infty
       for (i = 1 \text{ to } A.length)
         (if ((A[i] < MinPriority) and (A[i].empty = 0))
             then
           \int MinPriority = A[i]
           MinIndex = i
       if (MinIndex = 0)
         then return "underflow"
        A[MinIndex].empty = 1
       return MinIndex
proc\ Parent\ (i)
                                \operatorname{proc}\ \operatorname{Left}\ (i)
                                                              \mathbf{proc}\; \mathbf{\mathit{Right}}\,(i)
\left\{ \text{return } \left\lfloor \frac{i}{2} \right\rfloor \right\}
                                 { return 2 \cdot i
                                                              {return 2 \cdot i + 1
       proc MinHeapify (H, i)
         l = Left(i)
         r = Right(i)
         smallest=i
         if ((l \le H.heapsize) and (H[l] < H[i]))
           then smallest = l
         if ((r \leq H.heapsize) and (H[r] < H[smallest]))
           then smallest = r
         \textbf{if} \ smallest \neq i
           then
           \int SwapValue(H, i, smallest)
          MinHeapify(H, smallest)
      proc BuildMinHeap (H)
        \begin{cases} H.heapsize = H.length \\ \textbf{for} \; (i = \lfloor \frac{H.length}{2} \rfloor \; \textbf{downto} \; 1) \; MinHeapify(H,i) \end{cases} 
                proc HeapSort (H)
                   BuildMaxHeap(H)
                  \quad \text{for } (i=H.length \  \, \text{downto} \  \, 2)
                     SwapValue(H, i, 1)
```

H.heap size = H.heap size - 1

MaxHeapify(H, 1)

```
\mathbf{proc}\; \textit{Enqueue}\; (Q, priority)
           if (Q.heapsize = Q.length)
             then return "overflow"
           Q.heap size = Q.hep size + 1
           Q[heap size] = \infty
          DecreaseKey(Q,Q.heapsize,priority)
      proc\ \textit{DecreaseKey}\ (Q, i, priority)
         if (priority > Q[i])
then return "error"
         Q[i]=priority
         while ((i > 1) \text{ and } (Q[Parent(i)] > Q[i]))
          \int SwapValue(Q, i, Parent(i))
          i = Parent(i)
              \operatorname{proc} \operatorname{\textit{ExtractMin}}(Q)
                if (Q.heapsize < 1)
                   then return "under flow"
                 \begin{aligned} & \min = Q[1] \\ & Q[1] = Q[Q.heap size] \end{aligned}
                 Q.heap size = Q.heap size - 1 \\
                 MinHeapify(Q,1)
                 return min
                proc Enqueue (Q, i, priority)
                  \begin{cases} \text{if } (i > Q.length) \\ \text{then return "overflow"} \end{cases}
                   Q[i] = priority
       proc DecreaseKey (Q, i, priority)
        \begin{cases} \textbf{if } ((Q[i] < priority) \textbf{ or } (Q[i].empty = 1)) \\ \textbf{then return } "error" \end{cases}
         Q[i] = priority
\operatorname{proc} \operatorname{\textit{ExtractMin}}(Q)
   MinIndex = 0
   MinPriority = \infty
   for (i = 1 \text{ to } Q.length)
    (if ((Q[i] < MinPriority) and (Q[i].empty = 0))
      \int MinPriority = Q[i]
     \int MinIndex = i
   if (MinIndex = 0)
then return "underflow"
   Q[MinIndex].empty = 1
   return MinIndex
    proc CountingSort (A, B, k)
      let C[0,\ldots,k] new array
       for (i = 0 \text{ to } k) C[i] = 0
       for (j = 1 \text{ to } A.length) C[A[j]] = C[A[j]] + 1
       for (i = 1 \text{ to } k) C[i] = C[i] + C[i-1]
       for (j = A.length downto 1)
       \int B[C[A[j]]] = A[j]
       \left\{C[A[j]] = C[A[j]] - 1\right\}
    proc RadixSort (A, d)
    \{ for (i = 1 to d) \ AnyStableSort(A) on digit i \}
                   proc ListInsert (L, x)
                     x.next = L.head
                     if (L.head \neq nil)
                        then L.head.prev = x
                     L.head = x
                    x.prev = nil
    proc ListSearch (L, k)
      x = L.head
       while (x \neq \text{nil}\ ) and (x.key \neq k)\ x = x.next
      return x
               proc\ \textit{ListDelete}\ (L,x)
                  if (x.prev \neq nil)
                    then x.prev.next = x.next
                    else L.head = x.next
                  if (x.next \neq nil)
                   then x.next.prev = x.prev
```

```
proc Union (x, y)
                                         S_1 = x.head
                                         S_2 = y.head
                                         if (S_1 \neq S_2)
                                           then
                                           (S_1.tail.next = S_2.head)
                                            z = S_2.head
\textbf{proc MakeSet}\left(\mathcal{S}, S, x, i\right)
                                           while (z \neq \text{nil })
  \mathcal{S}[i].set = x
                                            \int z.head = S_1
   S.head=x
                                             z = z.next
  S.tail = x
                                            \hat{S}_1.tail = S_2.tail
                           proc\ FindSet\ (x)
                           {return x.head.head
                                     proc Union (x, y)
                                       S_1 = x.head
                                        S_2 = y.head
                                        if (S_1 \neq S_2)
                                         then
                                         (if (S_2.rank > S_1.rank)
                                            then
                                            S_{temp} = S_1
                                             S_1 = S_2
                                          \begin{cases} S_1 = S_2 \\ S_2 = S_{temp} \\ S_1.tail.next = S_2.head \end{cases}
                                          z = S_2.head
                                          while (z \neq \text{nil})
\textbf{proc MakeSet}\left(\mathcal{S}, S, x, i\right)
                                           \int z.head = S_1
  S[i].set = x

S.head = x
                                          \begin{cases} z = z.next \\ S_1.tail = S_2.tail \end{cases}
  S.tail = x
 S.rank = 0
                                         S_1.rank = S_1.rank + S_2.rank
                                  proc Union(x, y)
                                     x = Findset(x)
                                    y = Findset(y)
                                    if (x.rank > y.rank)
                                      then y.p = x
                                    if (x.rank \le y.rank)
                                      then
     proc MakeSet (x)
                                       x.p = u
                                       if (x.rank = y.rank)
      \int x \cdot p = x
                                        then y.rank = y.rank + 1
      x.rank = 0
                   proc\ FindSet\ (x)
                     if (x \neq x.p)
                        then x.p = FindSet(x.p)
                      return x.p
             proc TreeInOrderTreeWalk(x)
              (if (x \neq \text{nil})
                 then
                 TreeInOrderTreeWalk(x.left)
                  Print(x.key)
                 TreeInOrderTreeWalk(x.right)
             \textbf{proc} \,\, \textit{TreePreOrderTreeWalk} \, (x)
              (if (x \neq \text{nil})
                 then
                 Print(x.key)
                  TreeInOrderTreeWalk(x.left)
                TreeInOrderTreeWalk(x.right)
             proc TreePostOrderTreeWalk (x)
              (if (x \neq \text{nil})
                 TreeInOrderTreeWalk(x.left)
                  TreeInOrderTreeWalk(x.right)
                  Print(x.key)
```

 $\mathbf{proc}\; \textit{BSTTreeSearch}\; (x,k)$

then return x

if $(k \le x.key)$

(if ((x = nil) or (x.key = k))

then return BSTTreeSearch(x.left, k)

else return BSTTreeSearch(x.right, k)

```
proc BSTTreeMinimum(x)
        if ((x.left = nil)
          then return x
        \mathbf{return}\ BSTTreeMinimum(x.left)
proc BSTTreeSuccessor(x)
  if (x.right \neq nil)
    then return BSTTreeMinimum(x.right)
  while (y \neq \text{nil}) and (x = y.right)
  \int x = y
  y = y.p
  return y
        proc BSTTreeInsert(T, z)
          y = \mathbf{nil}
          x = T.root
          while (x \neq \text{nil })
           \begin{cases} y = x \\ \text{if } (z.key \le x.key) \end{cases}
             then x = x.left
           lelse x = x.right
           z.p = y
          if (y = \mathbf{nil})
           then T.root = z
          if ((y \neq \text{nil}) \text{ and } (z.key \leq y.key)
           then y.left = z
          if ((y \neq nil) and (z.key > y.key)
           then y.right = z
 proc BSTTreeDelete(T, z)
   (if (z.left = nil)
     then BSTTransplant(T, z, z.right)
    if ((z.left \neq \mathbf{nil}\ ) and (z.right = \mathbf{nil}\ ))
     then Transplant(T, z, z.left)
    if ((z.left \neq \mathbf{nil}\ ) and (z.right \neq \mathbf{nil}\ ))
     then
      Y = BSTTreeMinimum(z.right)
      if (y.p \neq z)
        then
       (BSTTransplant(T, y, y.right))
         y.right = z.right
       \int_{0}^{\infty} y.right.p = y
      BSTTransplant(T, z, y)
      y.left = z.left
     y.left.p = y
   \textbf{proc } \textit{BSTTreeTransplant} \ (T,u,v)
      if (u.p = nil)
       then T.root = v
      if ((u.p \neq nil) and (u = u.p.left)
       then u.p.left = v
      if ((u.p \neq nil) and (u = u.p.right)
       then u.p.right = v
      if (v \neq \hat{\mathbf{nil}})
       then v.p = u.p
   proc BSTTreeLeftRotate(T, x)
    y = x.right
     x.right = y.left

if (y.left \neq T.nil )
      then y.left.p = x
```

```
\begin{aligned} & \textbf{proc BSTTreeLeftRotate} \left(T,x\right) \\ & y = x.right \\ & x.right = y.left \\ & \textbf{if} \left(y.left \neq T.\textbf{nil}\right) \\ & \textbf{then } y.left.p = x \\ & y.p = x.p \\ & \textbf{if} \left(x.p = T.\textbf{nil}\right) \\ & \textbf{then } T.root = y \\ & \textbf{if} \left(\left(x.p \neq T.\textbf{nil}\right)\right) \\ & \textbf{then } x.p.left = y \\ & \textbf{if} \left(\left(x.p \neq T.\textbf{nil}\right)\right) \\ & \textbf{then } x.p.right = y \\ & y.left = x \\ & x.p = y \end{aligned}
```

```
 \begin{aligned} & \operatorname{proc} \textit{RBTreeInsert}\left(T,z\right) \\ & \left\{y = T.\operatorname{nil} \right. \\ & \left\{x = T.root \right. \\ & \left. \operatorname{while}\left(x \neq T.\operatorname{nil}\right) \right. \\ & \left\{y = x \right. \\ & \left. \operatorname{if}\left(z.key < x.key\right) \right. \\ & \left. \operatorname{then} x = x.left \\ & \left. \operatorname{else} \ x = x.right \right. \\ & \left. z.p = y \right. \\ & \left. \operatorname{if}\left(y = T.\operatorname{nil}\right) \right. \\ & \left. \operatorname{then} T.root = z \right. \\ & \left. \operatorname{if}\left((y \neq T.\operatorname{nil}\right) \right. \right. \\ & \left. \operatorname{nand}\left(z.key < y.key\right) \right. \\ & \left. \operatorname{then} y.left = z \right. \\ & \left. \operatorname{if}\left((y \neq T.\operatorname{nil}\right) \right. \\ & \left. \operatorname{nand}\left(z.key \geq y.key\right) \right. \\ & \left. \operatorname{then} y.right = z \right. \\ & \left. z.left = T.\operatorname{nil} \right. \\ & \left. z.right = T.\operatorname{nil} \right. \\ & \left. z.color = RED \right. \\ & \left. RBTreeInsertFixup(T,z) \right. \end{aligned}
```

```
 \begin{cases} \textbf{proc } \textit{RBTreeInsertFixup} \ (T,z) \\ \textbf{while} \ (z.p.color = RED) \\ \textbf{if} \ (z.p = z.p.p.left) \\ \textbf{then} \ RBTreeInsertFixUpLeft}(T,z) \\ \textbf{else} \ RBTreeInsertFixUpRight}(T,z) \\ T.root.color = BLACK \end{cases}
```

```
proc RBTreeInsertFixupLeft (T, z)
 y = z.p.p.right
 if (y.color = RED)
   then
   (z.p.color = BLACK)
   y.color = BLACK
   z.p.p.color = RED
  z = z.p.p
   else
   (if (z = z.p.right)
     then
    \int z = z.p
    \int LeftRotate(T,z)
    z.p.color = BLACK
    z.p.p.color = RED
   TreeRightRotate(T, z.p.p)
```

```
\mathbf{proc}\; \textit{RBTreeInsertFixupRight}\,(T,z)
  y = z.p.p.left
  if (y.color = RED)
    then
   (z.p.color = BLACK)
    y.color = BLACK
    z.p.p.color = RED \\
   z = z.p.p
    else
   (if (z = z.p.left)
      then
     \int z = z.v
     \left( Tree \hat{Right} Rotate(T,z) \right.
    z.p.color = BLACK
     z.p.p.color = RED
    TreeLeftRotate(T, z.p.p)
```

```
 \begin{aligned} & \operatorname{proc} \textit{BTreeSearch}\left(x,k\right) \\ & \begin{cases} i=1 \\ & \text{while} \; ((i \leq x.n) \; \operatorname{and} \; (k > x.key_i)) \; i=i+1 \\ & \text{if} \; ((i \leq x.n) \; \operatorname{and} \; (k = x.key_i)) \\ & \text{then return} \; (x,i) \\ & \text{if} \; (x.leaf = \operatorname{true}) \\ & \text{then return nil} \\ & DiskRead(x.c_i) \\ & \text{return} \; BTreeSearch(x.c_i,k) \end{aligned}
```

```
 \begin{aligned} & \textbf{proc BTreeCreate} \left( T \right) \\ & \begin{cases} x = Allocate() \\ x.leaf = \textbf{true} \end{cases} \\ & \begin{cases} x.n = 0 \\ DiskWrite(x) \\ T.root = x \end{aligned} \end{aligned}
```

```
\begin{aligned} & \operatorname{proc} \textit{BTreeSplitChild}\left(x,i\right) \\ & \left\{ \begin{aligned} z &= Allocate() \\ y &= x.c_i \\ z.leaf &= y.leaf \end{aligned} \right. \\ & \left\{ \begin{aligned} \text{for } j &= 1 \text{ to } t - 1 \text{ } z.key_j &= y.key_{j+t} \end{aligned} \right. \\ & \text{if } \left(y.leaf &= \text{false}\right) \\ & \text{then for } j &= 1 \text{ to } t \text{ } z.c_j &= y.c_{j+t} \\ y.n &= t - 1 \\ & \text{for } j &= x.n + 1 \text{ downto } i + 1 \text{ } x.c_{j+1} &= x.c_j \\ x.c_{i+1} &= z \\ & \text{for } j &= x.n \text{ downto } i \text{ } x.key_{j+1} &= x.key_j \\ x.key_i &= y.key_t \\ x.n &= x.n + 1 \\ DiskWrite(y) \\ DiskWrite(z) \\ DiskWrite(x) \end{aligned}
```

```
\begin{aligned} & \textbf{proc BTreeInsert}\left(T,k\right) \\ & r = T.root \\ & \textbf{if}\left(r.n = 2 \cdot t - 1\right) \\ & \textbf{then} \\ & \left\{ \begin{aligned} s &= Allocate() \\ T.root &= s \\ s.leaf &= \textbf{false} \end{aligned} \right. \\ s.n &= 0 \\ s.c_1 &= r \\ BTreeSplitChild(s,1) \\ BTreeInsertNonFull(s,k) \\ \textbf{else} & BTreeInsertNonFull(r,k) \end{aligned} \end{aligned}
```

```
proc BTreeInsertNonFull(x, k)
 i = x.n
  if (x.leaf = true)
   then
   (while ((i > 1)) and (k < x.key_i))
    \int x.key_{i+1} = x.key_i
    i = i - 1
   x.key_{i+1} = k
   x.n = x.n + 1
   DiskWrite(x)
   else
   while ((i \ge 1) and (k < x.key_i)) i = i - 1
    i = i + 1
    DiskRead(x.c_i)
   if (x.c_i.n = 2 \cdot t - 1)
     then
     BTreeSplitChild(x, i)
      if (k > x.key_i)
       then i = i + 1
    \hat{B}TreeInsertNonFull(x.c_i, k)
```

```
\begin{cases} \text{let } x \text{ be a new node with key } k \\ i = h(k) \\ ListInsert(T[i], x) \end{cases}
```

```
\begin{cases} i = h(k) \\ \text{return } ListSearch(T[i], k) \end{cases}
```

```
\begin{cases} \textit{proc HashDelete} \left( k \right) \\ i = h(k) \\ x = ListSearch(T[i], k) \\ ListDelete(T[i], x) \end{cases}
```

```
\begin{aligned} & \mathbf{proc} \; \textit{OaHashInsert} \left(T, k\right) \\ & \begin{cases} i = 0 \\ \mathbf{repeat} \\ & \begin{cases} j = h(k, i) \\ \mathbf{if} \; \left(T[j] = \mathbf{nil} \; \right) \\ \end{cases} \\ & \mathbf{then} \\ & \begin{cases} T[j] = k \\ \mathbf{return} \; j \\ \mathbf{else} \; \; i = i+1 \\ \mathbf{until} \; \; (i = m) \\ \mathbf{return} \; \; "overflow" \end{aligned}
```

```
\begin{aligned} & \text{proc } \textit{OaHashSearch} \left(T,k\right) \\ & \begin{cases} i=0 \\ \text{repeat} \\ & \\ \text{if } \left(T[j]=k\right) \\ & \text{then return } j \\ & i=i+1 \\ & \text{until } \left(\left(T[j]=\text{nil }\right) \text{ or } (i=m)\right) \\ & \text{return nil} \\ \end{aligned}
```

```
\begin{cases} \textbf{proc } \textit{HashComputeModulo} \ (w,B,m) \\ \textbf{let} \ \ d = |w| \\ z_0 = 0 \\ \textbf{for} \ (i=1 \ \textbf{to} \ d) \ z_{i+1} = ((z_i \cdot B) + a_i) \ mod \ m \\ \textbf{return} \ z_d + 1 \end{cases}
```

```
proc BreadthFirstSearch (G, s)
  for (u \in G.V \setminus \{s\})

\int u.color = WHITE
    u.d=\infty
   u.\pi = nil
  s.color = GREY
  s.d = 0
  s.\pi = \mathbf{nil}
  Q = \emptyset
  Enqueue(Q, s)
  while (Q \neq \emptyset)
    u = Dequeue(Q)
    for (v \in G.Adj[u])
     (if (v.color = WHITE)
        then
        v.color = GRAY
         v.d = u.d + 1
         v.\pi = u
       Enqueue(Q, v)
    u.color = BLACK
```

```
\begin{aligned} & \operatorname{proc} \operatorname{\textit{DepthFirstSearch}}\left(G\right) \\ & \left\{ \begin{aligned} & \operatorname{for}\left(u \in G.V\right) \\ & \left\{ u.color = WHITE \\ u.\pi = \operatorname{nil} \\ \end{aligned} \right. \\ & \left\{ time = 0 \\ & \operatorname{for}\left(u \in G.V\right) \\ & \left\{ \begin{aligned} & \operatorname{if}\left(u.color = WHITE\right) \\ & \operatorname{then} \operatorname{\textit{DepthVisit}}(G,u) \end{aligned} \right. \end{aligned} \end{aligned} \end{aligned}
```

```
\begin{aligned} & \textbf{proc Depth Visit} \left(G, u\right) \\ & \begin{cases} time = time + 1 \\ u.d = time \\ u.color = GREY \end{cases} \\ & \textbf{for} \left(v \in G.Adj[u]\right) \\ & \textbf{if} \left(v.color = WHITE\right) \\ & \textbf{then} \\ & \begin{cases} v.\pi = u \\ DepthVisit(G, v) \\ u.color = BLACK \\ time = time + 1 \\ u.f = time \end{cases} \end{aligned}
```

```
\begin{aligned} & \mathbf{proc} \ \mathbf{CycleDet} \ (G) \\ & \begin{cases} cycle = False \\ \mathbf{for} \ (u \in G.V) \ u.color = WHITE \\ \mathbf{for} \ (u \in G.V) \\ \end{cases} \\ & \begin{cases} \mathbf{if} \ (u.color = WHITE) \\ \mathbf{then} \ DepthVisitCycle} \\ \mathbf{return} \ cycle \end{cases} \end{aligned}
```

```
 \begin{cases} \textbf{proc DepthVisitCycle} \left(G,u\right) \\ \left\{u.color = GREY \\ \textbf{for} \left(v \in G.Adj[u]\right) \\ \textbf{fif} \left(v.color = WHITE\right) \\ \textbf{then } DepthVisitCycle(G,v) \\ \textbf{if} \left(v.color = GREY\right) \\ \textbf{then } cycle = True \\ u.color = BLACK \end{cases}
```

```
\begin{aligned} & \textbf{proc TopologicalSort}\left(G\right) \\ & \textbf{for } (u \in G.V) \ u.color = WHITE \\ & L = \emptyset \\ & time = 0 \\ & \textbf{for } (u \in G.V) \\ & \textbf{jif } (u.color = WHITE) \\ & \textbf{then } DepthVisitTS(G,u) \\ & \textbf{return } L \end{aligned}
```

```
\begin{aligned} & \textbf{proc Depth VisitTS}\left(G,u\right) \\ & \begin{cases} time = time + 1 \\ u.d = time \\ u.color = GREY \\ & \textbf{for}\left(v \in G.Adj[u]\right) \\ \end{cases} \\ & \begin{cases} \textbf{if}\left(v.color = WHITE\right) \\ \textbf{then } Depth VisitTS(G,v) \\ u.color = BLACK \\ time = time + 1 \\ u.f = time \\ ListInsert(L,u) \end{cases} \end{aligned}
```

```
proc StronglyConnectedComponents (G)
 for (u \in G.V)
  \int u.color = WHITE
  u.\pi = \text{nil}
  time = 0
  for (u \in G.V)
  \int \mathbf{if} \ (u.color = WHITE)
    then DepthVisit(G, u)
  for (u \in G.V)
  \int u.color = WHITE
  u.\pi = \text{nil}
  time = 0
  L = \emptyset
  for (u \in G^T.V \text{ in rev. finish time order})
  (if (u.color = WHITE)
     then DepthVisit(G^{T'}, u)
  ListInsert(L, u)
  rèturn L
```

```
\begin{aligned} & \operatorname{proc} \operatorname{\textit{MST-Prim}} (G, w, r) \\ & \operatorname{\textit{for}} (v \in G.V) \\ & \operatorname{\textit{do}} \\ & \left\{ v.key = \infty \right. \\ & \left\{ v.\pi = \operatorname{\textit{nil}} \right. \\ & \left. v.key = 0 \right. \\ & \left\{ Q = G.V \right. \\ & \operatorname{\textit{while}} (Q \neq \emptyset) \\ & \operatorname{\textit{do}} \\ & \left\{ u = \operatorname{\textit{ExtractMin}} (Q) \\ & \operatorname{\textit{for}} (v \in G.Adj[u]) \\ & \operatorname{\textit{do}} \\ & \left\{ if \left( (v \in Q) \text{ and } (W(u,v) < v.key) \right) \\ & \operatorname{\textit{do}} \\ & \left\{ v.\pi = u \\ v.key = W(u,v) \right. \end{aligned} \end{aligned} \right.
```

```
 \begin{aligned} & \operatorname{proc} \operatorname{\mathit{MST-Kruskal}} \left( G, w \right) \\ & \begin{cases} T = \emptyset \\ & \operatorname{for} \ (v \in G.V) \\ & \operatorname{\mathit{do}} \ MakeSet(v) \\ & \operatorname{\mathit{SortNoDecreasing}} (G.E) \\ & \operatorname{for} \ ((u,v) \in G.E - in \ order) \\ & \operatorname{\mathit{do}} \\ & \operatorname{\mathit{do}} \\ & \operatorname{\mathit{dif}} \ (FindSet(u) \neq FindSet(v)) \\ & \operatorname{\mathit{then}} \\ & \begin{cases} T = T \cup \{(u,v)\} \\ Union(u,v) \\ \end{array} \end{aligned}
```

```
\begin{aligned} & \text{proc } \textit{Union } (x,y) \\ & \begin{cases} S_1 = FindSet(x) \\ S_2 = FindSet(y) \\ & \text{if } (S_1 \neq S_2) \\ & \text{then} \\ & \begin{cases} S_1.tail.next = S_2.head \\ z = S_2.head \\ & \text{while } (z \neq \text{nil }) \\ \begin{cases} z.head = S_1.head \\ z = z.next \end{cases} \end{aligned}
```

```
\begin{aligned} & \operatorname{proc} \operatorname{Union} \left( x, y \right) \\ & \left\{ \begin{array}{l} S_1 = \operatorname{FindSet}(x) \\ S_2 = \operatorname{FindSet}(y) \\ & \operatorname{if} \left( |S_1| > |S_2| \right) \\ & \operatorname{then} \\ & \left\{ \begin{array}{l} S_2 = \operatorname{FindSet}(x) \\ S_1 = \operatorname{FindSet}(y) \\ & \operatorname{if} \left( S_1 \neq S_2 \right) \\ & \operatorname{then} \\ & \left\{ \begin{array}{l} S_1 + \operatorname{Isal} \cdot \operatorname{lnext} = S_2 \cdot \operatorname{head} \\ z = S_2 \cdot \operatorname{head} \\ & \operatorname{while} \left( z \neq \operatorname{nil} \right) \\ & \left\{ \begin{array}{l} z \cdot \operatorname{head} = S_1 \cdot \operatorname{head} \\ z = z \cdot \operatorname{next} \\ & \operatorname{return} S_1 \\ \end{array} \right. \end{aligned}
```

```
 \begin{cases} \textbf{proc InizializeSingleSource} \left(G,s\right) \\ \begin{cases} \textbf{for} \ \left(v \in G.V\right) \\ \begin{cases} v.d = \infty \\ v.\pi = \textbf{nil} \\ s.d = 0 \end{cases} \end{cases}
```

```
\begin{cases} & \operatorname{proc} \operatorname{\textit{Relax}}\left(u,v,w\right) \\ & \operatorname{if}\left(v.d > u.d + W(u,v)\right) \\ & \operatorname{then} \\ & \left\{v.d = u.d + W(u,v) \\ v.\pi = u \right. \end{cases}
```

```
 \begin{array}{l} \textbf{proc Bellman-Ford} \left(G, w, s\right) \\ \begin{cases} InizializeSingleSource(G, s) \\ \textbf{for } i = 1 \textbf{ to } |G.V| - 1 \\ \textbf{ for } \left((u, v) \in G.E) \ Relax(u, v, w) \\ \textbf{ for } \left((u, v) \in G.E\right) \\ \end{cases} \\ \begin{cases} \textbf{if } \left(v.d > u.d + W(u, v)\right) \\ \textbf{ then return false} \\ \textbf{ return true} \\ \end{cases}
```

```
\begin{aligned} & \text{proc } \textit{Dijkstra} \ (G, w, s) \\ & \begin{cases} InizializeSingleSource(G, s) \\ S = \emptyset \\ Q = G.V \\ & \text{while} \ (Q \neq \emptyset) \\ \begin{cases} u = ExtractMin(Q) \\ S = S \cup \{u\} \\ & \text{for} \ (v \in G.Adj[u]) \\ \begin{cases} & \text{if} \ (v \in Q) \\ & \text{then } Relax(u, v, w) \end{cases} \end{aligned}
```

```
\begin{aligned} & \operatorname{proc} \textit{ExtendShortestPaths} \left(L, W\right) \\ & \begin{cases} n = L.rows \\ \operatorname{let} L' \text{ be a new matrix} \\ \text{for } i = 1 \text{ to } n \end{cases} \\ & \begin{cases} \begin{cases} for \ j = 1 \text{ to } n \\ L'_{ij} = \infty \\ \text{for } k = 1 \text{ to } n \end{cases} \\ \begin{cases} L'_{ij} = \min\{L'_{ij}, L_{ik} + W_{kj}\} \\ \text{return } i + 1 \end{cases} \end{aligned} \end{aligned}
```

```
 \begin{cases} \operatorname{proc} \operatorname{SlowAllPairsMatrix}\left(W\right) \\ n = L.\operatorname{rows} \\ L^{1} = W \\ \operatorname{for} \ m = 2 \ \operatorname{to} \ n - 1 \\ L^{m} = \operatorname{ExtendShortestPaths}(L^{m-1}, W) \\ \operatorname{return} \ L^{n-1} \end{cases}
```

```
\begin{aligned} & \operatorname{proc} \operatorname{\textit{FastAllPairsMatrix}}\left(W\right) \\ & \left\{ \begin{aligned} & n = L.rows \\ & L^1 = W \\ & m = 1 \end{aligned} \right. \\ & \operatorname{\textbf{while}}\left(m < n - 1\right) \\ & \operatorname{\textbf{do}} \\ & \left\{ \begin{aligned} & L^{2 \cdot m} = ExtendShortestPaths(L^m, L^m) \\ & m = 2 \cdot m \end{aligned} \right. \end{aligned}
```

```
\begin{aligned} & \operatorname{proc} \operatorname{Floyd-Warshall}\left(W\right) \\ & n = W.rows \\ & D^0 = W \\ & \operatorname{for} \ k = 1 \ \operatorname{to} \ n \\ & \left\{ \begin{aligned} & \operatorname{let} \ D^k \ \operatorname{be} \ \operatorname{a} \ \operatorname{new} \ \operatorname{matrix} \\ & \operatorname{for} \ i = 1 \ \operatorname{to} \ n \\ & \left\{ \begin{aligned} & \operatorname{for} \ j = 1 \ \operatorname{to} \ n \\ & D_{ij}^k = \min\{D_{ij}^{k-1}, D_{ik}^{k-1} + D_{kj}^{k-1} \} \end{aligned} \right. \end{aligned} \end{aligned}
```