

PROJECT

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PROBLEM: SEPARATING POINTS BY AXIS PARALLEL LINES

GIVEN:

Input → set of n points in 2D plane

Output → Set S of combinations of vertical and horizontal lines separating all points.

Assumptions: No two points have same x or y coordinates.

MEASURE: minimise S .

GREEDY HEURISTIC

1. PSEUDOCODE:

READ_FILES():

Read all files from "Input" folder.

Read all points.

return points

CREATE_PAIRS(points):

$U = []$ // union of all possible pairs of all points.

for i in range(0, length(points)):

for j in ($i+1$, length(points)):

$U.append(points[i], points[j])$

return U // returns Union^{set} of all pairs

SEPARATE_POINTS (U, n):

(2)

Here n is the total no. of points.

Let $V_{Axis} \leftarrow$ set of all vertical axes,

and $H_{Axis} \leftarrow$ set of all horizontal axes.

$S = [] \rightarrow$ solution SET

while U is not Empty:

Let $countV \leftarrow$ no. of axis cut by a
vertical axis. (declared as dict)

and $countH \leftarrow$ no. of axis cut by a
horizontal axis. (declared as dict).

Count the total ~~line~~ pairs from U
separated by ~~a~~ vertical axes and store
in $countV$ as KEY-Value pair.

Do the same for ~~all~~ horizontal axes and
store in $countH$ as KEY-Value pair

Select ~~an~~ an axis that cuts the maximum
pair out of both vertical and horizontal
axis.

$S \leftarrow$ store that axis in solution set.

$U \leftarrow U - S^{(pairs)}$ // Remove those pairs from
Union set that are cut by S .
// $S(pairs) \rightarrow$ pairs cut by S .

return S // S is our solution set

2. ANALYSIS OF RUNTIME:

1 • Reading Files takes $\sim O(n)$ ^① (depending on how large n is). Here n can be no. of lines in each all the files.

2 • CREATE_PAIRS() function creates pair of a point with all other possible points. This runs in $T(n) = n * (n-1)$. Hence runtime complexity of creating pairs is $\sim O(n^2)$ ^②.

3 • SEPARATE_POINTS() function has following assignments and comparisons:-

* P.S → Here the small nos. in circles above the complexity is the ordering of the sequence in which code is running and not powers.

3.1 - Creating all vertical and horizontal axes takes $\rightarrow \sim O(n)$. Vertical and Horizontal lines are stored in a List. Addition to List is $\sim O(n)$ ^③

3.2 - There's a first while loop. The while loop has following assignments and comparison:- $\rightarrow \sim O(U - \text{scpairs}) \rightarrow \sim O(n/2)$ ^④

3.2.1 • First for loop for iterating in vertical or Horizontal axes List. This is $\rightarrow O(\text{size of each List}) \rightarrow \sim O(n)$ ^⑤
This for loop has subloop to iterate in U (Union set)

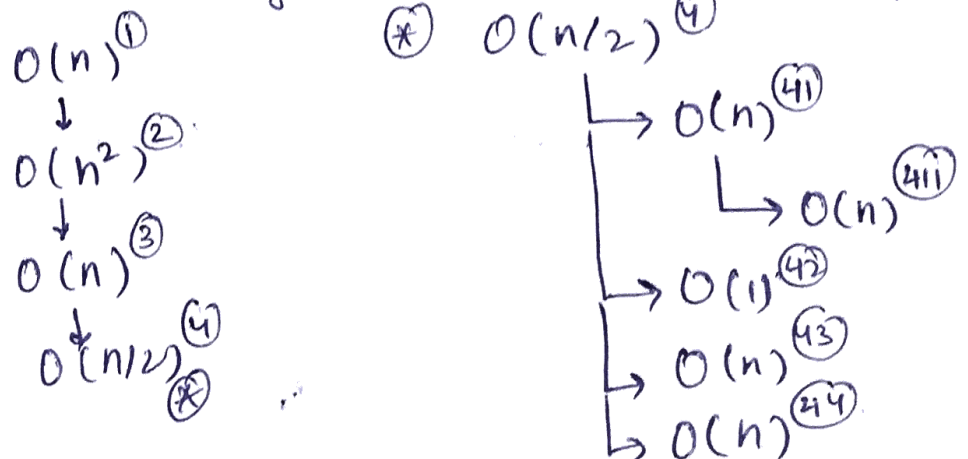
3.2.1.1.1 \triangleright This sub for loop iterates in U and has an if condition to check for those U in U that satisfies the condition. Hence this for loop has $(n+1)$ assignments and n comparisons in U .
 This is $\rightarrow O(n+1) + O(n) \rightarrow \sim O(n)$

3.2.2.2 \triangleright Here comparing for maximum between two axis takes $2 * O(1), \sim O(1)$.

3.2.3 \triangleright Here addition of axes to solution set (List in Python) takes a runtime of $O(\text{min. no. of axes}) \rightarrow \sim O(n)$.

3.2.2 \circ There's another for loop at the last to remove pairs ~~for~~ from U , which runs in $O(U\text{-size}) \rightarrow \sim O(n)$.

So our runtime is something as follows (in hierarchy):



We can see from the diagram, that

⑤

Adding all the runtime we get:-

$$\begin{aligned}\text{Total Runtime} &= O(n)^{(1)} + O(n^2)^{(2)} + O(n)^{(3)} + (O(n/2))^4 \times \\ &\quad O(n)^{(4)} \times O(n)^{(4)} + (O(n/2) \times O(1)) \\ &\quad + (O(n/2) \times O(n)) + (O(n/2) \times O(n)) \\ &= 2O(n) + O(n^2) + O(n^3/2) + O(n/2) + O(n^2/2)\end{aligned}$$

Here $O(n^3/2)$ has the highest degree.

So we can avoid all lower ordered terms.

$$\therefore \boxed{\text{Total Runtime (T(n))} \approx \underline{\underline{O(n^3)}}}$$

3. One instance where my greedy solution fail to return optimum solution is when all the points are in linearly increasing order or decreasing order.

For example, instance = $[(1,1), (2,2), (3,3), (4,4), (5,5)]$

The output was = 4

h 3.5

h 2.5

h 4.5

h 1.5

Hence all the points were separated using maximum @ parallel axes. Hence this type of instance will not be optimal.

And there is no better solution for this kind of ⑥
~~problem~~ instance.

~~Another instance~~

Another instance where ~~was~~ I did not get optimal solution was:

instance06.txt from website.

output: my output
15 axes

h 15.5
v 15.5
h 23.5
h 9.5
v 8.5
v 21.5
h 18.5
v 4.5
v 28.5
h 26.5
h 11.5
h 5.5
h 22.5
h 17.5
h 1.5

(Better Solution)
original output

14 axes.

v 15.5
h 14.5
h 22.5
h 5.5
v 22.5
v 7.5
v 26.5
v 11.5
h 10.5
v 4.5
v 13.5
v 16.5
v 20.5
v 27.5