# Diophantine Equations and the Euclidean Algorithm

## Reading

- Section 5.2
- Section 5.3

#### **Practice Problems**

5.2 2-7, 11, 13
Challenge 5.2 (Optional) 14
5.3 1, 2-5, 12, 15
Challenge 5.3 (Optional) 25

#### **Notes**

### Case of c=gcd

The Euclidean Algorithm allows us to find a solution to the equation

$$ax + by = c$$

where  $c = d = \gcd(a, b)$ . We can then use this to find a solution for any c divisible by  $d = \gcd(a, b)$ .

Let us take a look at the first few steps in the Euclidean Algorithm:

$$a = q_1 b + r_1$$

$$b = q_2 r_1 + r_2$$

$$r_1 = q_3 r_2 + r_3$$

We will show that all  $r_1$ ,  $r_2$ ,  $r_3$  are explicit linear combinations of a and b. For  $r_1$  this is clear, as:

$$r_1 = a - q_1 b$$

For  $r_2$  we would use this fact, so we would have:

$$r_2 = b - q_2 r_1 = b - q_2 (a - q_1 b) = (1 + q_1 q_2) b - q_2 a$$

Similarly we can write  $r_3$  as an explicit compination of a and b by using its equation along with the fact that we already have a way to write  $r_1$  and  $r_2$ .

This continues:

In the Euclidean Division algorithm, each remainder  $r_n$  is an *explicit* integer linear combination of a and b.

Since the last step is the gcd(a, b), we now have a way of writing gcd(a, b) in an explicit way as an integer linear combination of a and b.

#### Case of other c's

We already know two things:

- How to find an explicit solution to ax + by = gcd(a, b)
- That the only way that ax + by = c has a solution is if gcd(a, b)|c

We now want to find an explicit solution in the case where gcd(a, b)|c. To do that:

- Write  $c = d \gcd(a, b)$ , where d is an integer.
- Write  $ax + by = \gcd(a, b)$  for some integers x, y.
- Then we have  $a(xd) + b(yd) = \gcd(a,b)d = c$  and we have our solution.

So every solution to the  $c = \gcd$  case can scale up to a solution of the c case for all those c for which there is a solution in the first place.