Assumption: used 2 level page tables to implement this

- 1. Instruction fetch (00000,FFC)
- 2. Fault(00000,FFC, instruction) // first instruction fetch; PDE[0] and PTE[0] not present
- 3. Page allocate = 00001 // allocate new page table for low 4 MB region (PDE 0)
- 4. 00000[0].page = 00001
- 5. Page allocate = 00002 // allocate physical frame for program pg1
- 6. 00001[0].page = 00002
- 7. read block(/bin/program, 0, 00002)
- 8. Return from fault
- 9. Instruction fetch (00000,FFC)
- 10. attempt(PUSH #10)
- 11. store(08000,FFC)
- 12. fault(08000,FFC, data store) // stack region not yet mapped
- 13. Page allocate = 00003 // allocate new page table for stack region (PDE 32)
- 14. 00000[32].page = 00003
- 15. Page allocate = 00004 // allocate actual stack frame (physical page)
- 16. 00003[0].page = 00004
- 17. Return from fault
- 18. store(08000,FFC)
- 19. Success
- 20. Instruction fetch (00001,000)
- 21. fault(00001,000, instruction) // program page 1 not yet loaded
- 22. Page allocate = 00005 // allocate new frame for program pg2
- 23. 00001[1].page = 00005
- 24. read block(/bin/program, 1, 00005)
- 25. Return from fault
- 26. Instruction fetch (00001,000)
- 27. attempt(CALL 2,000)
- 28. store(08000,FF8)
- 29. Success
- 30. Instruction fetch (00002,000)
- 31. fault(00002,000, instruction) // program page 2 not yet loaded
- 32. Page allocate = 00006 // allocate new frame for program pg3
- 33. 00001[2].page = 00006
- 34. read block(/bin/program, 2, 00006)
- 35. Return from fault
- 36. Instruction fetch (00002,000)
- 37. attempt(MOV EAX -> \*(00010,000))

- 38. store(00010,000)
- 39. fault(00010,000, data store) // first access to data region no PTE[16] entry
- 40. Page allocate = 00007 // allocate physical frame for data page
- 41. 00001[16].page = 00007
- 42. Return\_from\_fault
- 43. attempt(MOV EAX -> \*(00010,000))
- 44. Success
- 45. Instruction fetch (00002,004)
- 46. attempt(HALT)
- 47. Success