

1. Page Size = $4096 = 2^{12}$ bytes

\therefore Page Table Size = $2^{32-12} = 2^{20}$ bytes.

least significant 12 bits = $(456)_{16}$.

most significant 20 bits = $(11123)_{16}$.

- Using hardware, we can generate physical address.
- since only 2^{16} bytes of physical memory, frame size is 6 bits.
 - 6-bit frame address combined with 12-bit page displacement to generate physical address.

2. Effective access time = $(1-P)ma + P(\text{page fault service time})$.

mem access time = 100ns

Page Fault empty/nonempty = 8ms

Page Fault modify = 20ms

Fault frequency = $\frac{7}{10}$.

\therefore effective access time:

$$200\text{ns} = (1-p)100\text{ns} + P[(3/10)8\text{ms} + (7/10)20\text{ms}]$$

$$= 100\text{ns} - p(100\text{ns}) + P(12/5\text{ms} + 14\text{ms})$$

$$100\text{ns} = P\left(\frac{82}{5}\text{ms}\right)\left(\frac{10 \times 10^6\text{ns}}{\text{ms}}\right) - 100\text{ns}(p)$$

$$= P(16399900\text{ns})$$

$$\therefore P \approx 6.0976 \times 10^{-6}$$

$\Rightarrow \therefore$ max acceptable page fault rate for effective time of no more than 200ns