# Let it Rot



# 目录

| 1 | 图论   |                                    | 1         |   | 3.5 AC 自动机 2                            | 25         |
|---|------|------------------------------------|-----------|---|---|------------|
|   | 1.1  | 二分图匹配   最小边覆盖                      | 1         |   | W. 1-41.14                              |            |
|   | 1.2  | 网络最大流   dinic                      | 4         | 4 | *************************************** | 26         |
|   | 1.3  | 最小费用流   dijkstra                   | 5         |   | 4.1 区间加区间求和树状数组 2                       |            |
|   | 1.4  | 最小费用流   spfa                       |           |   | 4.2 zkw 线段树                             |            |
|   | 1.5  | 最小树形图                              |           |   | 4.3 Link Cut Tree                       | 28         |
|   | 1.6  | 缩点   kasaraju                      | 10        | 5 | 计算几何 2                                  | 29         |
| 2 | 3.5  |                                    |           |   | 5.1 向量                                  | 29         |
|   | Math |                                    | 11        |   | 5.2 直线半平面                               |            |
|   | 2.1  | 万能欧几里得                             | 11        |   | 5.3 半平面交 3                              |            |
|   | 2.2  | 直线下点数   欧几里得                       | 12        |   |   |            |
|   | 2.3  | 扩展欧几里得                             | 13        |   | 5.4 线段                                  |            |
|   | 2.4  | 扩展中国剩余定理                           | 13        |   | 5.5 多边形                                 |            |
|   | 2.5  | Miller-Rabin                       | 14        |   | 5.6 Graham 凸包                           |            |
|   | 2.6  | Pollard-rho                        |           |   | 5.7 上凸壳                                 | 37         |
|   | 2.7  | Fast Fourier Transform             | 17        | 6 | Dirty Hacks 3                           | 8          |
|   | 2.8  | Number Theoretic Transform $\dots$ | 18        |   | 6.1 Pragma                              | 38         |
|   | 2.9  | Generating function                | 19        |   | 6.2 Barrett                             | 38         |
|   |      |                                    |           |   | 6.3 LCS                                 | 38         |
| 3 |      |                                    | <b>20</b> |   | 6.4 日期公式 3                              | 39         |
|   | 3.1  | 后缀自动机   SAM                        | 20        |   | 302 A/MAZY                              | , ,        |
|   | 3.2  | 回文自动机   PAM                        | 21        | 7 | 配置 4                                    | ł0         |
|   | 3.3  | 回文自动机 border 处理                    | 22        |   | 7.1 vimrc                               | 10         |
|   | 3.4  | 后缀数组   SA                          | 24        |   | 7.2 bash 命令                             | <b>1</b> C |

## 1 图论

## 1.1 二分图匹配 | 最小边覆盖

```
// 匈牙利, 左到右单向边, D(M|match|)
std::vector<int> edge[N];
inline bool dfs(int x, std::vector<int> & vis, std::vector<int> & match) {
   for(int y : edge[x]) if(!vis[y])
      if(vis[y] = 1, !match[y] || dfs(match[y], vis, match))
         return match[y] = x, 1;
   return 0;
}
inline std::vector<int> match(int nl, int nr) {
   std::vector<int> vis(nr + 1), match(nr + 1), ret(nl + 1);
   for(int i = 1;i <= nl;++i) if(dfs(i, vis, match))</pre>
      memset(vis.data(), 0, vis.size() << 2);</pre>
   for(int i = 1;i <= nr;++i) ret[match[i]] = i;</pre>
   return ret[0] = 0, ret;
}
/* 最小边覆盖,很可能没用,真别抄,抄这个不用抄 match */
inline std::pair<std::vector<int>, std::vector<int>> minedgecover(int nl, int nr) {
   std::vector<int> vis(nr + 1), match(nr + 1), ret(nl + 1);
   for(int i = 1;i <= nl;++i) if(dfs(i, vis, match))</pre>
      memset(vis.data(), 0, vis.size() << 2);</pre>
   for(int i = 1;i <= nr;++i) ret[match[i]] = i;</pre>
   ret[0] = 0;
   for(int i = 1;i <= nl;++i) if(!ret[i]) dfs(i, vis, match);</pre>
   std::vector<int> le, ri;
   for(int i = 1;i <= nl;++i) if(ret[i] && !vis[ret[i]]) le.push_back(i);</pre>
   for(int i = 1;i <= nr;++i) if(vis[i]) ri.push_back(i);</pre>
   return std::make_pair(le, ri);
}
// 匈牙利, 左到右单向边, bitset, O(n^2/w | match|)
using set = std::bitset<N>;
set edge[N];
inline bool dfs(int x, set & unvis, std::vector<int> & match) {
   for(set z = edge[x];;) {
      z &= unvis;
      int y = z._Find_first();
      if(y == N) return 0;
      if(unvis.reset(y), !match[y] || dfs(match[y], unvis, match))
         return match[y] = x, 1;
   }
}
```

```
inline std::vector<int> match(int nl, int nr) {
   set unvis; unvis.set();
   std::vector<int> match(nr + 1), ret(nl + 1);
   for(int i = 1;i <= nl;++i)</pre>
      if(dfs(i, unvis, match))
          unvis.set();
   for(int i = 1;i <= nr;++i) ret[match[i]] = i;</pre>
   return ret[0] = 0, ret;
}
// HK, 左到右单向边, O(M sqrt(|match|))
std::vector<int> edge[N];
std::vector<int> L, R, vis, matchl, matchr;
inline bool bfs(int nl, int nr) {
   L.assign(nl + 1, 0), R.assign(nr + 1, 0);
   static std::queue<int> Q;
   for(int i = 1;i <= nl;++i)</pre>
      if(!matchl[i]) L[i] = 1, Q.push(i);
   int succ = 0;
   for(;Q.size();) {
      int x = Q.front(); Q.pop();
      for(int i : edge[x]) if(!R[i]) {
          R[i] = L[x] + 1;
          if(int p = matchr[i]) L[p] = R[i] + 1, Q.push(p);
          else succ = 1;
      }
   }
   return succ;
inline bool dfs(int x) {
   for(int i : edge[x]) if(R[i] == L[x] + 1 && !vis[i]) {
      vis[i] = 1;
      if(!matchr[i] || dfs(matchr[i]))
          return matchl[x] = i, matchr[i] = x, 1;
   }
   return 0;
}
inline std::vector<int> match(int nl, int nr) {
   matchl.assign(nl + 1, 0);
   matchr.assign(nr + 1, 0);
   vis.resize(nr + 1);
   for(;bfs(nl, nr);) {
      memset(vis.data(), 0, vis.size() << 2);</pre>
      for(int i = 1;i <= nl;++i) if(!matchl[i]) dfs(i);</pre>
   }
```

北京大学-摆烂人

```
return matchl;
}

/* 最小边覆盖,很可能没用,真别抄 */
inline std::pair<std::vector<int>, std::vector<int>> minedgecover(int nl, int nr) {
   auto res = match(nl, nr);
   std::vector<int> ls, rs;
   for(int i = 1;i <= nl;++i) if(!L[i]) ls.push_back(i);
   for(int i = 1;i <= nr;++i) if(R[i]) rs.push_back(i);
   return std::make_pair(ls, rs);
}
```

PKU-Let it Rot 3

#### 1.2 网络最大流 | dinic

```
// S 编号最小, T 最大, 或者改一下清空
struct Dinic {
   struct T {
       int to, nxt, v;
   } e[N << 3];
   int h[N], head[N], num = 1;
   inline void link(int x, int y, int v) {
       e[++num] = {y, h[x], v}, h[x] = num;
       e[++num] = \{x, h[y], 0\}, h[y] = num; // !!!
   }
   int dis[N];
   bool bfs(int s, int t) {
       std::queue<int> Q;
       for(int i = s;i <= t;++i) dis[i] = -1, head[i] = h[i]; //如果编号不是[S,T], 只要改这里
       for(Q.push(s), dis[s] = 0;!Q.empty();) {
          int x = Q.front(); Q.pop();
          for(int i = h[x];i;i = e[i].nxt) if(e[i].v && dis[e[i].to] < 0) {</pre>
             dis[e[i].to] = dis[x] + 1, Q.push(e[i].to);
          }
       return dis[t] >= 0;
   }
   int dfs(int s, int t, int lim) {
       if(s == t || !lim) return lim;
       int ans = 0, mn;
       for(int & i = head[s];i;i = e[i].nxt) {
          if(dis[e[i].to] == dis[s] + 1 && (mn = dfs(e[i].to, t, std::min(lim, e[i].v)))) {
             e[i].v -= mn, e[i ^ 1].v += mn;
             ans += mn, lim -= mn;
             if(!lim) break;
          }
      return ans;
   }
   inline int flow(int s, int t) {
       int ans = 0;
       for(;bfs(s, t);) ans += dfs(s, t, 1e9);
       return ans;
   }
} G;
```

## 1.3 最小费用流 | dijkstra

```
// S 编号最小, T 最大, 或者改一下清空
namespace mcmf {
   using pr = std::pair<11, int>;
   const int N = 10005, M = 1e6 + 10;
   struct edge {
      int to, nxt, v, f;
   } e[M << 1];
   int h[N], num = 1;
   inline void link(int x, int y, int v, int f) {
       e[++num] = {y, h[x], v, f}, h[x] = num;
      e[++num] = \{x, h[y], 0, -f\}, h[y] = num;
   }
   11 d[N], dis[N];
   int vis[N], fr[N];
   inline void spfa(int s, int t) {
       std::queue<int> Q;
      std::fill(d + s, d + t + 1, 1e18);
      for(d[s] = 0, Q.push(s);!Q.empty();) {
          int x = Q.front(); Q.pop(); vis[x] = 0;
          for(int i = h[x];i;i = e[i].nxt)
             if(e[i].v \&\& d[e[i].to] > d[x] + e[i].f) {
                 d[e[i].to] = d[x] + e[i].f;
                 if(!vis[e[i].to]) vis[e[i].to] = 1, Q.push(e[i].to);
             }
      }
   inline bool dijkstra(int s, int t) {
       std::priority_queue<pr, std::vector<pr>, std::greater<pr>>> Q;
       std::fill(dis + s, dis + t + 1, 1e18);
       std::fill(vis + s, vis + t + 1, 0);
       for(Q.emplace(dis[s] = 0, s);!Q.empty();) {
          int x = Q.top().second; Q.pop();
          if(vis[x]) continue;
          vis[x] = 1;
          for(int i = h[x];i;i = e[i].nxt) {
              const ll v = e[i].f + d[x] - d[e[i].to];
             if(e[i].v && dis[e[i].to] > dis[x] + v) {
                 fr[e[i].to] = i;
                 Q.emplace(dis[e[i].to] = dis[x] + v, e[i].to);
             }
      }
      return dis[t] < 1e17;</pre>
   }
```

```
std::pair<11, 11> EK(int s, int t) {
    spfa(s, t);
    11 f = 0, c = 0;
    for(;dijkstra(s, t);) {
        11 fl = 1e18;
        for(int i = fr[t];i;i = fr[e[i ^ 1].to]) fl = std::min<1l>(e[i].v, fl);
        for(int i = s;i <= t;++i) d[i] += dis[i];
        f += fl, c += fl * d[t];
        for(int i = fr[t];i;i = fr[e[i ^ 1].to])
              e[i].v -= fl, e[i ^ 1].v += fl;
    }
    return std::make_pair(f, c);
}</pre>
```

## 1.4 最小费用流 | spfa

```
using ll = long long;
// S 编号最小, T 最大, 或者改一下清空
namespace mcmf {
   using pr = std::pair<11, int>;
   const int N = 10005, M = 1e6 + 10;
   struct edge {
       int to, nxt, v, f;
   } e[M << 1];
   int h[N], num = 1;
   inline void link(int x, int y, int v, int f) {
       e[++num] = {y, h[x], v, f}, h[x] = num;
       e[++num] = \{x, h[y], 0, -f\}, h[y] = num;
   }
   ll d[N];
   int vis[N], fr[N];
   inline bool spfa(int s, int t) {
       std::queue<int> Q;
       std::fill(d + s, d + t + 1, 1e18);
       for(d[s] = 0, Q.push(s);!Q.empty();) {
          int x = Q.front(); Q.pop(); vis[x] = 0;
          for(int i = h[x];i;i = e[i].nxt)
              if(e[i].v \&\& d[e[i].to] > d[x] + e[i].f) {
                 d[e[i].to] = d[x] + e[i].f, fr[e[i].to] = i;
                 if(!vis[e[i].to]) vis[e[i].to] = 1, Q.push(e[i].to);
              }
       return d[t] < 1e17;</pre>
   }
   inline std::pair<11, 11> EK(int s, int t) {
       11 f = 0, c = 0;
       for(;spfa(s, t);) {
          11 fl = 1e18;
          for(int i = fr[t];i;i = fr[e[i ^ 1].to]) fl = std::min<ll>(e[i].v, fl);
          f += fl, c += fl * d[t];
          for(int i = fr[t];i;i = fr[e[i ^ 1].to])
              e[i].v -= fl, e[i ^ 1].v += fl;
       }
       return std::make_pair(f, c);
   }
}
```

#### 1.5 最小树形图

抄罗大的,返回值是边的集合,如果没有最小树形图会返回-1,可以修改建图。

```
namespace DMST {
   struct edge {
       int u, v, id; ll w;
       inline bool operator < (const edge & y) const {</pre>
          return w < y.w;</pre>
       }
   } ent[N], val[M];
   int ls[M], rs[M], size[M], cc; ll tag[M];
   int fs[N], fw[N], rt[N];
   inline void put(int x, 11 v) {
       if(x) val[x].w += v, tag[x] += v;
   inline void pushdown(int x) {
       put(ls[x], tag[x]);
      put(rs[x], tag[x]);
       tag[x] = 0;
   }
   inline int merge(int x, int y) {
       if(!x || !y) return x | y;
       if(val[y] < val[x]) std::swap(x, y);</pre>
       pushdown(x), rs[x] = merge(rs[x], y);
       if(size[rs[x]] > size[ls[x]]) {
          std::swap(ls[x], rs[x]);
       size[x] += size[y];
       return x;
   inline void ins(int & x, const edge & z) {
       val[++cc] = z, size[cc] = 1;
       x = merge(x, cc);
   inline void pop(int & x) { x = merge(ls[x], rs[x]); }
   edge top(int x) { return val[x]; }
   inline int find(int x, int * anc) {
       return anc[x] == x ? x : anc[x] = find(anc[x], anc);
   }
   inline void link(int u, int v, int w, int id) {
       ins(rt[v], {u, v, id, w});
   }
   int pa[N * 2], tval[N * 2], up[N * 2], end_edge[M], cmt, baned[M];
   std::vector<int> solve(int r) {
       std::queue<int> roots;
```

}

```
for(int i = 1;i <= n;++i) {</pre>
       fs[i] = fw[i] = i, tval[i] = ++ cmt;
       if(i != r) roots.push(i);
   std::vector<edge> H;
   std::vector<int> ret;
   for(;!roots.empty();) {
       int k = roots.front(); roots.pop();
       if(!rt[k]) return ret;
       edge e = top(rt[k]); pop(rt[k]);
       int i = e.u, j = e.v;
       if(find(i, fs) == k) roots.push(k);
       else {
          H.push_back(e); end_edge[e.id] = tval[k];
          if(find(i, fw) != find(j, fw)) {
              fw[find(j, fw)] = i;
              ent[k] = e;
          } else {
              pa[tval[k]] = ++ cmt, up[tval[k]] = e.id;
              put(rt[k], -e.w);
              for(;(e = ent[find(e.u, fs)]).u;) {
                 int p = find(e.v, fs);
                 pa[tval[p]] = cmt;
                 up[tval[p]] = e.id;
                 put(rt[p], -e.w);
                 rt[k] = merge(rt[k], rt[p]);
                 fs[p] = k;
              }
              tval[k] = cmt;
              roots.push(k);
          }
       }
   reverse(H.begin(), H.end());
   for(edge i : H) if(!baned[i.id]) {
       ret.push_back(i.id);
       for(int j = i.v;j != end_edge[i.id];j = pa[j]) ++ baned[up[j]];
   }
   sort(ret.begin(), ret.end());
   return ret;
}
```

## 1.6 缩点 | kasaraju

时间复杂度  $O(\frac{n^2}{w})$ ,可以对于边修改不多的图快速计算。

```
using set = std::bitset<N>;
// re 是反向边,需要连好
set e[N], re[N], vis;
std::vector<int> sta;
inline void dfs0(int x, set * e) {
   vis.reset(x);
   for(;;) {
       int go = (e[x] & vis)._Find_first();
       if(go == N) break;
      dfs0(go, e);
   }
   sta.push_back(x);
inline std::vector<std::vector<int>> solve() {
   vis.set();
   for(int i = 1;i <= n;++i) if(vis.test(i)) dfs0(i, e);</pre>
   vis.set();
   auto s = sta;
   std::vector<std::vector<int>> ret;
   for(int i = n - 1;i >= 0;--i) if(vis.test(s[i])) {
       sta.clear(), dfs0(s[i], re), ret.push_back(sta);
   }
   return ret;
}
```

## 2 Math

## 2.1 万能欧几里得

```
// 万欧
// 前提 : r < q, r >= q 先提几个 U 出来再用
// 使用: Y * q <= X * p + r, 斜率 p/q, U表示向上, R表示到达一个顶点, 先一些 U 再一个 R
template<class T>
inline T power(T a, 11 k) {
   // 有效率需求可以改为半群乘法
   if(!k) return T();
   T res = a;
   for(--k;k;) {
      if(k & 1) res = res + a;
      if(k >>= 1) a = a + a;
   }
   return res;
}
template<class T>
T solve(ll p, ll q, ll r, ll l, T U, T R) {
   if (p >= q)
      return solve(p % q, q, r, 1, U, power(U, p / q) + R);
   11 m = ((_int128)p * 1 + r) / q;
   if (!m) return power(R, 1);
   11 \text{ cnt} = 1 - ((_int128)q * m - r - 1) / p;
   return power(R, (q - r - 1) / p) + U + solve(q, p, (q - r - 1) % p, m - 1, R, U) + power(R,
       cnt);
}
```

## 2.2 直线下点数 | 欧几里得

$$n < 2^{32}, 1 \le m < 2^{32}$$

$$result = \sum_{i=0}^{n-1} \lfloor \frac{ai+b}{m} \rfloor \pmod{2^{64}}$$

```
u64 floor_sum(u64 n, u64 m, u64 a, u64 b) {
       u64 \text{ ans} = 0;
       for(;;) {
              if(a >= m) {
                     ans += n * (n - 1) / 2 * (a / m);
                      a %= m;
              if(b >= m) {
                     ans += n * (b / m);
                     b %= m;
              }
              u64 ymax = a * n + b;
              if(ymax < m) break;</pre>
              n = ymax / m;
              b = ymax % m;
              std::swap(m, a);
       }
       return ans;
```

#### 2.3 扩展欧几里得

```
// exgcd
// result : -b < x < b AND -a < y <= a when a,b != 0
inline void exgcd(ll a, ll b, ll & x, ll & y) {
   if(!b) return x = 1, y = 0, void();
   exgcd(b, a % b, y, x), y -= a / b * x;
}</pre>
```

## 2.4 扩展中国剩余定理

```
11 exCRT(11 a1, 11 p1, 11 a2, 11 p2) {
        11 a, b, gcd = std::gcd(p1, p2);
        if((a1 - a2) % gcd)
            return -1;
        exgcd(p1, p2, a, b);
        11 k = i128((a2 - a1) % p2 + p2) * (a + p2) % p2;
        return p1 / gcd * k + a1;
}
```

#### 2.5 Miller-Rabin

```
using f64 = long double;
11 p;
f64 invp;
inline void setmod(ll x) {
       p = x, invp = (f64) 1 / x;
}
inline 11 mul(11 a, 11 b) {
       11 z = a * invp * b + 0.5;
       ll res = a * b - z * p;
       return res + (res >> 63 & p);
inline ll pow(ll a, ll x, ll res = 1) {
       for(;x;x >>= 1, a = mul(a, a))
              if(x & 1) res = mul(res, a);
       return res;
}
inline bool checkprime(ll p) {
       if(p == 1) return 0;
       setmod(p);
       ll d = \_builtin\_ctzll(p - 1), s = (p - 1) >> d;
       for(11 a : {2, 3, 5, 7, 11, 13, 82, 373}) {
              if(a \% p == 0)
                     continue;
              11 x = pow(a, s), y;
              for(int i = 0;i < d;++i, x = y) {</pre>
                     y = mul(x, x);
                     if(y == 1 && x != 1 && x != p - 1)
                            return 0;
              if(x != 1) return 0;
       }
       return 1;
}
```

#### 2.6 Pollard-rho

```
inline ll rho(ll n) {
       if(!(n & 1))
              return 2;
       static std::mt19937_64 gen((size_t)"hehezhou");
       ll c = gen() \% (n - 1) + 1, y = gen() \% (n - 1) + 1;
       auto f = [&](11 o) {
              o = mul(o, o) + c;
              return o >= n ? o - n : o;
       };
       setmod(p);
       for(int 1 = 1;;1 <<= 1) {</pre>
              11 x = y, g = 1;
              for(int i = 0;i < 1;++i) y = f(y);</pre>
              const int d = 512;
              for(int i = 0;i < 1;i += d) {</pre>
                     11 sy = y;
                     for(int j = 0; j < std::min(d, 1 - i);++j) {</pre>
                            y = f(y), g = mul(g, (y - x + n));
                     g = std::_gcd(n, g);
                     if(g == 1)
                            continue;
                     if(g == n)
                            for(g = 1, y = sy;g == 1;)
                                   y = f(y), g = std::_gcd(n, y - x + n);
                     return g;
              }
       }
inline std::vector<ll> factor(ll x) {
       std::queue<11> q; q.push(x);
       std::vector<ll> res;
       for(;q.size();) {
              11 x = q.front(); q.pop();
              if(x == 1) continue;
              if(checkprime(x)) {
                     res.push_back(x);
                     continue;
              }
              11 y = rho(x);
              q.push(y), q.push(x / y);
       }
       sort(res.begin(), res.end());
       return res;
```

}

#### 2.7 Fast Fourier Transform

```
using db = double;
using C = std::complex<db>;
// C::real, C::imag, std::conj, std::arg
const db pi = std::acos(-1);
int rev[N], lim, invlim;;
C wn[N];
void init(int len) {
   lim = 2 << std::__lg(len - 1);
   invlim = mod - (mod - 1) / lim;
   for(static int i = 1;i < lim;i += i) {</pre>
       for(int j = 0; j < i; ++ j) {</pre>
           wn[i + j] = std::polar(1., db(j) / i * pi);
       }
   }
   for(int i = 1;i < lim;++i) {</pre>
       rev[i] = rev[i >> 1] >> 1 | (i % 2u * lim / 2);
   }
}
void DFT(C * a) {
   for(int i = 0;i < lim;++i) {</pre>
       if(rev[i] < i) std::swap(a[rev[i]], a[i]);</pre>
   for(int i = 1;i < lim;i += i) {</pre>
       for(int j = 0; j < lim; j += i + i) {</pre>
           for(int k = 0; k < i; ++k) {
              C x = a[i + j + k] * wn[i + k];
              a[i + j + k] = a[k + j] - x;
              a[k + j] += x;
           }
       }
   }
}
void IDFT(C * a) {
   DFT(a), std::reverse(a + 1, a + lim);
   for(int i = 0;i < lim;++i)</pre>
       a[i] /= lim;
}
```

#### 2.8 Number Theoretic Transform

```
int rev[N], wn[N], lim, invlim;
int norm(int x) {
   return x \ge mod ? x - mod : x;
int pow(int a, int b, int ans = 1) {
   for(;b;b >>= 1, a = (u64) a * a % mod) if(b & 1)
       ans = (u64) ans * a % mod;
   return ans;
}
void init(int len) {
   lim = 2 << std::__lg(len - 1);
   invlim = mod - (mod - 1) / lim;
   for(static int i = 1;i < lim;i += i) {</pre>
       wn[i] = 1;
       const int w = pow(3, mod / i / 2);
       for(int j = 1; j < i; ++j) {
           wn[i + j] = (u64) wn[i + j - 1] * w % mod;
       }
   }
   for(int i = 1;i < lim;++i) {</pre>
       rev[i] = rev[i >> 1] >> 1 | (i % 2u * lim / 2);
   }
void DFT(int * a) {
   static u64 t[N];
   for(int i = 0;i < lim;++i) {</pre>
       t[i] = a[rev[i]];
   for(int i = 1;i < lim;i += i) {</pre>
       for(int j = 0; j < lim; j += i + i) {</pre>
           for(int k = 0; k < i; ++k) {
              const u64 x = t[i + j + k] * wn[i + k] % mod;
              t[i + j + k] = t[k + j] + mod - x, t[k + j] += x;
          }
       }
   }
   for(int i = 0;i < lim;++i) a[i] = t[i] % mod;</pre>
void IDFT(int * a) {
   DFT(a), std::reverse(a + 1, a + lim);
   for(int i = 0;i < lim;++i)</pre>
       a[i] = (u64) a[i] * invlim % mod;
}
```

## 2.9 Generating function

```
void cpy(int * a, int * b, int n) {
   if(a != b) memcpy(a, b, n << 2);
   memset(a + n, 0, (lim - n) << 2);
void inv(int * a, int * b, int n) { // mod x^n
   if(n == 1) return void(*b = pow(*a, mod - 2));
   static int c[N], d[N];
   int m = (n + 1) / 2;
   inv(a, b, m);
   init(n + m), cpy(c, b, m), cpy(d, a, n);
   DFT(c), DFT(d);
   for(int i = 0;i < lim;++i) c[i] = (u64) c[i] * c[i] % mod * d[i] % mod;</pre>
   IDFT(c);
   for(int i = m;i < n;++i) b[i] = norm(mod - c[i]);</pre>
}
void log(int * a, int * b, int n) {
   static int c[N], d[N];
   inv(a, c, n), init(n + n);
   for(int i = 1;i < n;++i) d[i - 1] = (u64) a[i] * i % mod;</pre>
   cpy(d, d, n - 1), cpy(c, c, n);
   DFT(c), DFT(d);
   for(int i = 0;i < lim;++i) c[i] = (u64) c[i] * d[i] % mod;</pre>
   IDFT(c), *b = 0;
   for(int i = 1;i < n;++i) b[i] = pow(i, mod - 2, c[i - 1]);</pre>
}
```

## 3 字符串

## 3.1 后缀自动机 | SAM

需要两倍点数量。

```
int c[N][26], mx[N], fail[N], tot = 1;
int append(int id, int w) {
   int p = id, now = ++ tot;
   //right[now] = id;
   for(mx[now] = mx[p] + 1;p && !c[p][w];p = fail[p])
       c[p][w] = now;
   if(!p) fail[now] = 1;
   else {
       int q = c[p][w];
       if(mx[q] == mx[p] + 1) fail[now] = q;
       else {
          int x = ++ tot; mx[x] = mx[p] + 1;
          memcpy(c[x], c[q], sizeof(c[0])), fail[x] = fail[q]; //right[x] = right[q];
          for(fail[q] = fail[now] = x;p && c[p][w] == q;p = fail[p])
              c[p][w] = x;
      }
   }
   return now;
}
void 后缀树() { // 倒着建
   for(int i = 2;i <= tot;++i)</pre>
       son[fa[i]][s[right[i] + mx[fa[i]]] - 'a'] = i;
}
```

## 3.2 回文自动机 | PAM

```
int c[N][26], fail[N], len[N], tot;
void init() {
   fail[0] = 1, len[++tot] = -1;
   // root is 1
int get_fail(int o, char * x) {
   for(;*x != x[-len[o] - 1];)
       o = fail[o];
   return o;
}
int append(int o, char * x) {
   o = get_fail(o, x);
   int & p = c[o][*x - 'a'];
   if(!p) {
       fail[++tot] = c[get_fail(fail[o], x)][*x - 'a'];
      len[p = tot] = len[o] + 2;
   return p;
}
```

#### 3.3 回文自动机 border 处理

```
int c[N][26], fail[N], len[N], tot;
void init() {
   fail[0] = 1, len[++tot] = -1;
   // root is 1
int get_fail(int o, char * x) {
   for(;*x != x[-len[o] - 1];)
      o = fail[o];
   return o;
}
int append(int o, char * x) {
   o = get_fail(o, x);
   int & p = c[o][*x - 'a'];
   if(!p) {
      fail[++tot] = c[get_fail(fail[o], x)][*x - 'a'];
      len[p = tot] = len[o] + 2;
   return p;
}
BorderPam
info atom[N];
int c[N][26], fail[N], len[N], tot;
int diff[N], bigfail[N];
/*----*/
// 严格 logn 插入, 但是很可能不要
void init() {
   bigfail[0] = fail[0] = 1, len[++tot] = -1;
   // root is 1
int get_fail(int o, char * x) {
   for(;*x != x[-len[o] - 1];) {
      o = (*x != x[-len[fail[o]] - 1] ? bigfail : fail)[o];
   }
   return o;
void init() {
   fail[0] = 1, len[++tot] = -1;
   // root is 1
}
int get_fail(int o, char * x) {
   for(;*x != x[-len[o] - 1];)
      o = fail[o];
   return o;
```

```
int append(int o, char * x) {
   o = get_fail(o, x);
   int & p = c[o][*x - 'a'];
   if(!p) {
       fail[++tot] = c[get_fail(fail[o], x)][*x - 'a'];
       len[p = tot] = len[o] + 2;
      diff[p] = len[p] - len[fail[p]];
      bigfail[p] = diff[p] == diff[fail[p]] ? bigfail[fail[p]] : fail[p];
   return p;
}
info node[N];
info query(int x, int i) {
   info z = \{\};
   for(;x;x = bigfail[x]) {
       const int go = bigfail[x], fa = fail[x];
       info & t = node[x] = atom[i - len[go] - diff[x]];
      if(fa != go) t = node[fa] + t;
      z = t + z;
   }
   return z;
}
```

## 3.4 后缀数组 | SA

```
char s[N];
int rank[N], sa[N], h[N], n, L;
bool cmp(int a, int b) {
   if(rank[a] != rank[b]) return rank[a] < rank[b];</pre>
   return b + L <= n && (a + L > n || rank[a + L] < rank[b + L]);</pre>
}
void SA() { // s[n + 1] need '0'
   static int a[N], t[N];
   for(int i = 1;i <= n;++i) rank[i] = s[i];</pre>
   std::iota(a, a + n + 1, 0);
   for(L = 1;;L <<= 1) {</pre>
       std::sort(a + 1, a + n + 1, cmp);
       for(int i = 1, r = 0;i \le n;++i) t[a[i]] = r += !i || cmp(a[i - 1], a[i]);
       memcpy(rank + 1, t + 1, n << 2);
       if(t[a[n]] == n) break;
   for(int i = 1;i <= n;++i) sa[rank[i]] = i;</pre>
   for(int i = 1, k = 0;i <= n;++i) if(rank[i] < n) {</pre>
       int j = sa[rank[i] + 1];
       for(k -= !!k;s[i + k] == s[j + k];++k);
       h[rank[i]] = k;
   }
}
```

#### 3.5 AC 自动机

```
const int sig = 26;
int son[N][sig], fail[N], cnt;
int ins(const char * c) {
   int x = 0;
   for(;*c;++c) {
       int & p = son[x][*c - 'a'];
      if(!p) p = ++ cnt;
      x = p;
   }
   return x;
void build_ac() {
   std::queue<int> Q;
   for(int i = 0;i < sig;++i) if(son[0][i]) Q.push(son[0][i]);</pre>
   for(;Q.size();) {
       int x = Q.front(); Q.pop();
       for(int i = 0;i < sig;++i)</pre>
          if(son[x][i]) fail[son[x][i]] = son[fail[x]][i], Q.push(son[x][i]);
          else son[x][i] = son[fail[x]][i];
   }
}
```

## 4 数据结构

## 4.1 区间加区间求和树状数组

```
// 后缀加, 前缀求和
struct BIT {
   ll a[N], b[N];
   inline void add(ll p, int v) {
       for(int i = p;i < N;i += i & -i)</pre>
          a[i] += v, b[i] += p * v;
   }
   inline ll qry(ll p) {
       11 \text{ res} = 0;
       for(int i = p;i;i &= i - 1) res += (p + 1) * a[i] - b[i];
       return res;
   inline void add(int 1, int r, int v) {
       add(1, v), add(r + 1, -v);
   }
   inline 11 qry(int 1, int r) {
       return qry(r) - qry(1 - 1);
   }
} bit;
```

#### 4.2 zkw 线段树

```
struct seg {
   11 o[1 << 20]; int L;</pre>
   void upt(int x) {
       o[x] = o[x << 1] + o[x << 1 | 1];
   void init(int n, int * w) {
       L = 2 \ll std::__lg(n + 1);
       for(int i = 1;i <= n;++i) o[i + L] = w[i];</pre>
       for(int i = L;i >= 1;--i) upt(i);
   void upt(int p, int v) {
       for(o[p += L] += v;p >>= 1;upt(p));
   11 qry(int 1, int r) {
       1 += L - 1, r += L + 1;
       11 \text{ ans} = 0;
       for(;l ^ r ^ 1;l >>= 1, r >>= 1) {
          if((1 \& 1) == 0) ans += o[1 ^ 1];
          if((r \& 1) == 1) ans += o[r ^ 1];
       return ans;
   }
   // if there is no I
   11 qry2(int 1, int r) {
       if(l == r) return o[l + L];
       11 le = o[1 + L], ri = o[r + L];
       1 += L, r += L;
       for(;1 ^ r ^ 1;1 >>= 1, r >>= 1) {
          if((1 & 1) == 0) le = le + o[l ^ 1];
          if((r & 1) == 1) ri = o[r ^ 1] + ri;
       }
       return le + ri;
   }
} sgt;
```

#### 4.3 Link Cut Tree

```
int son[N][2], fa[N], rev[N];
int get(int x, int p = 1) {
   return son[fa[x]][p] == x;
}
void update(int x) {
int is_root(int x) {
   return !(get(x) || get(x, 0));
void rotate(int x) {
   int y = fa[x], z = fa[y], b = get(x);
   if(!is_root(y)) son[z][get(y)] = x;
   son[y][b] = son[x][!b], son[x][!b] = y;
   fa[son[y][b]] = y, fa[y] = x, fa[x] = z;
   update(y);
void put(int x) {
   if(x) rev[x] ^= 1, std::swap(son[x][0], son[x][1]);
void down(int x) {
   if(rev[x]) {
      put(son[x][0]);
      put(son[x][1]);
      rev[x] = 0;
   }
}
void pushdown(int x) {
   if(!is_root(x)) pushdown(fa[x]);
   down(x);
}
void splay(int x) {
   for(pushdown(x);!is_root(x);rotate(x)) if(!is_root(fa[x]))
       rotate(get(x) ^ get(fa[x]) ? x : fa[x]);
   update(x);
}
void access(int x) {
   for(int t = 0;x;son[x][1] = t, t = x, x= fa[x])
       splay(x);
}
void makeroot(int x) {
   access(x), splay(x), put(x);
}
```

## 5 计算几何

#### 5.1 向量

```
using db = long double;
const db eps = 1e-10;
inline int add(int a, int b, int p) {
   return a += b, a >= p ? a - p : a;
}
inline db sgn(db x) {
   return x < -eps ? -1 : x > eps;
}
struct vec2 {
   db x, y;
   inline vec2() { }
   inline vec2(db a, db b) : x(a), y(b) { }
   inline db norm() const {
       return x * x + y * y;
   }
   inline db abs() const {
       return std::sqrt(x * x + y * y);
   }
};
inline vec2 operator + (const vec2 & x, const vec2 & y) {
   return vec2(x.x + y.x, x.y + y.y);
}
inline vec2 operator - (const vec2 & x, const vec2 & y) {
   return vec2(x.x - y.x, x.y - y.y);
}
inline vec2 operator / (const vec2 & x, db y) {
   return vec2(x.x / y, x.y / y);
inline db operator * (const vec2 & x, const vec2 & y) {
   return x.x * y.y - x.y * y.x;
inline vec2 operator * (const db & x, const vec2 & y) {
   return vec2(x * y.x, x * y.y);
inline vec2 operator * (const vec2 & y, const db & x) {
   return vec2(x * y.x, x * y.y);
inline db operator % (const vec2 & x, const vec2 & y) {
   return x.x * y.x + x.y * y.y;
}
```

```
inline db dist(const vec2 & x, const vec2 & y) {
    return (x - y).abs();
}
// 逆时针极角排序
inline int half(const vec2 & x) {
    return x.y < 0 || (x.y == 0 && x.x <= 0);
}
inline bool cmp(const vec2 & A, const vec2 & B) {
    if(half(A) != half(B)) return half(B);
    return A * B > 0;
}
```

#### 5.2 直线半平面

```
struct line {
        db a, b, c;
        // a * x + b * y + c (= or >) 0
        inline line(db A, db B, db C) : a(A), b(B), c(C) {}
        inline line(const vec2 & x, const vec2 & y) : a(x.y - y.y), b(y.x - x.x), c(x * y) { }
        // 左侧 > 0
        inline db operator ()(const vec2 & x) const { return a * x.x + b * x.y + c; }
        inline line perp() const { return line(b, -a, 0); }
        inline line para(const vec2 & o) { return line(a, b, c - (*this)(o)); }
        inline vec2 normVec() const { return vec2(a, b); }
        inline db norm() const { return normVec().norm(); }
};
inline vec2 operator & (const line & x, const line & y) {
        return vec2(vec2(x.c, x.b) * vec2(y.c, y.b), vec2(x.a, x.c) * vec2(y.a, y.c)) / -(vec2(x.a, x.c)) / -(ve
                 x.b) * vec2(y.a, y.b));
}
inline vec2 proj(const vec2 & x, const line & 1) {
        return x - 1.normVec() * (1(x) / 1.norm());
}
inline vec2 refl(const vec2 & x, const line & 1) {
        return x - 1.normVec() * (1(x) / 1.norm()) * 2;
}
inline bool is_para(const line & x, const line & y) {
        return std::fabs(x.normVec() * y.normVec()) < eps;</pre>
inline bool is_perp(const line & x, const line & y) {
        return std::fabs(x.normVec() % y.normVec()) < eps;</pre>
}
inline bool online(const vec2 & x, const line & 1) {
        return std::fabs(l(x)) < eps;</pre>
}
inline int ccw(vec2 a, vec2 b, vec2 c) {
        int sign = sgn((b - a) * (c - a));
        if(sign == 0) {
                if(sgn((b - a) % (c - a)) == -1) return 2;
                if((c - a).norm() > (b - a).norm() + eps) return -2;
        }
        return sign;
}
// 三条直线关系
db det(line a, line b, line c) {
        vec2 A = vec2(a), B = vec2(b), C = vec2(c);
        return c.c * (A * B) + a.c * (B * C) + b.c * (C * A);
```

```
}
db check(line a, line b, line c) { // sgn same as c(a & b), 0 if error
    return sgn(det(a, b, c)) * sgn(vec2(a) * vec2(b));
}
```

#### 5.3 半平面交

```
std::vector<vec2> HalfPlaneI(std::vector<line> vs) {
   auto cmp = [](line a, line b) {
       if(paraS(a, b)) {
          return dist(a) < dist(b);</pre>
       return ::cmp(vec2(a), vec2(b));
   };
   sort(vs.begin(), vs.end(), cmp);
   int ah = 0, at = 0, n = size(vs);
   std::vector<line> deq(n + 1);
   std::vector<vec2> ans(n);
   deq[0] = vs[0];
   for(int i = 1;i <= n;++i) {</pre>
       line o = i < n ? vs[i] : deq[ah];
       if(paraS(vs[i - 1], o)) {
           continue;
       for(;ah < at && check(deq[at - 1], deq[at], o) < 0;) -- at;</pre>
       if(i != n)
       for(;ah < at && check(deq[ah], deq[ah + 1], o) < 0;) ++ ah;</pre>
       if(!is_para(o, deq[at])) {
           ans[at] = o & deq[at];
           deq[++at] = o;
       }
   }
   if(at - ah <= 2) return {};</pre>
   return {ans.begin() + ah, ans.begin() + at};
}
```

#### 5.4 线段

```
struct seg {
   vec2 x, y;
   inline seg() {}
   inline seg(const vec2 & A, const vec2 & B) : x(A), y(B) {}
   inline bool onseg(const vec2 & o) const {
       return (o - x) % (o - y) < eps && std::fabs((o - x) * (o - y)) < eps;</pre>
   inline line to_l() const {
      return line(x, y);
   }
};
inline bool is_isc(const seg & x, const seg & y) {
   return
       ccw(x.x, x.y, y.x) * ccw(x.x, x.y, y.y) \le eps &&
       ccw(y.x, y.y, x.x) * ccw(y.x, y.y, x.y) \le eps;
}
inline db dist(const seg & o, const vec2 & x) {
   vec2 z = proj(x, o.to_1());
   if(o.onseg(z)) {
       return dist(x, z);
   } else {
       return std::min(dist(o.x, x), dist(o.y, x));
   }
}
inline db dist(const seg & x, const seg & y) {
   if(is_isc(x, y)) return 0;
   return std::min({ dist(y, x.x), dist(y, x.y), dist(x, y.x), dist(x, y.y), });
}
```

#### 5.5 多边形

```
using polygon = std::vector<vec2>;
// counter-clockwise
inline db area(const polygon & x) {
   db res = 0;
   for(int i = 2;i < (int) x.size();++i) {</pre>
       res += (x[i - 1] - x[0]) * (x[i] - x[0]);
   return res / 2;
}
inline bool is_convex(const polygon & x, bool strict = 1) {
   // warning, maybe wrong
   const db z = strict ? eps : -eps;
   for(int i = 2;i < (int) x.size() + 2;++i) {</pre>
       if((x[(i-1) \% x.size()] - x[i-2]) * (x[i \% x.size()] - x[i-2]) < z) return 0;
   }
   return 1;
}
inline int inpoly(const vec2 & x, const polygon & o) {
   bool in = false;
   for(int i = 0;i < (int) o.size();++i) {</pre>
       if(seg(o[i], o[add(i, 1, size(o))]).onseg(x)) {
          return 1;
       }
       auto a = o[i] - x, b = o[add(i, 1, size(o))];
   }
}
```

#### 5.6 Graham 凸包

没有判断左下角重点。结果为逆时针。

```
bool eql(db x, db y) {
       return fabs(x - y) < eps;</pre>
}
db cross(vec2 x, vec2 y, vec2 z) {
       return (y.x - x.x) * (z.y - x.y) - (y.y - x.y) * (z.x - x.x);
std::vector<vec2> gethull(std::vector<vec2> o) {
       iter_swap(o.begin(), min_element(o.begin(), o.end(), [](vec2 x, vec2 y) {
              return eql(x.x, y.x) ? x.y < y.y : x.x < y.x;</pre>
       }));
       sort(o.begin() + 1, o.end(), [&](vec2 x, vec2 y) {
              db c = cross(o[0], x, y);
              if(fabs(c) > eps) {
                     return c > 0;
              return (x - o[0]).norm() < (y - o[0]).norm();</pre>
       });
       std::vector<vec2> stack;
       for(vec2 x : o) {
              for(;stack.size() >= 2 && cross(stack.rbegin()[1], stack.back(), x) <= eps;) {</pre>
                     stack.pop_back();
              }
              stack.push_back(x);
       for(;stack.size() > 2 && cross(stack.rbegin()[1], stack.back(), o[0]) <= eps;) {</pre>
              stack.pop_back();
       return stack;
}
```

## 5.7 上凸壳

结果显然为顺时针。

```
std::vector<vec2> gethull(std::vector<vec2> o) {
       sort(o.begin(), o.end(), [](vec2 x, vec2 y) {
              if(x.x == y.x) {
                     return x.y > y.y; // gt => 1t
              } else {
                     return x.x < y.x;</pre>
       });
       std::vector<vec2> stack;
       for(vec2 x : o) {
              if(stack.size() && stack.back().x == x.x) {
                     continue;
              for(;stack.size() \ge 2 \&\& cross(stack.rbegin()[1], stack.back(), x) \ge 0;) { // gt}
                  => 1t
                     stack.pop_back();
              }
              stack.push_back(x);
       return stack;
}
```

## 6 Dirty Hacks

#### 6.1 Pragma

```
#pragma GCC optimize("Ofast")
#pragma GCC optimize("unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse4,popcnt,abm,mmx,avx,avx2")
```

#### 6.2 Barrett

```
struct DIV {
     u64 x;
     void init(u64 v) { x = -1ull / v + 1; }
};
// works while x*(y-1)<2^64
u64 operator / (const u64 & x, const DIV & y) {
     return (u128) x * y.x >> 64;
}
```

#### 6.3 LCS

```
int lim;
struct bitset {
    static const int B = 63;
    u64 a[N / B + 1];
    inline void set(int p) { a[p / B] |= 1ull << (p % B); }
    inline bool test(int p) { return a[p / B] >> (p % B) & 1; }
    inline void run(const bitset & o) {
        u64 c = 1;
        for(int i = 0; i < lim; ++i) {
            u64 x = a[i], y = x | o.a[i];
            x += x + c + (-y & (1ull << 63) - 1);
            a[i] = x & y, c = x >> 63;
        }
    }
} dp;
```

## 6.4 日期公式

```
// Mon = 0, ... % 7
// days since 1/1/1
int getday(int y, int m, int d) {
    if(m < 3) -- y, m += 12;
    return (365 * y + y / 4 - y / 100 + y / 400 + (153 * (m - 3) + 2) / 5 + d - 307);
}
void date(int n, int & y, int & m, int & d) {
    n += 429 + ((4 * n + 1227) / 146097 + 1) * 3 / 4;
    y = (4 * n - 489) / 1461;
    n -= y * 1461 / 4;
    m = (5 * n - 1) / 153;
    d = n - m * 153 / 5;
    if (--m > 12) m -= 12, ++y;
}
```

## 7 配置

#### 7.1 vimrc

set si ci ts=4 sw=4 nu cino=j1 backup undofile syntax on

map<F9> <ESC>: !g++ % -o %< -Wall -g<CR>

map<F10> <ESC>:!%<<CR>
map<F5> <ESC>:!gdb %<<CR>>

## 7.2 bash 命令

ulimit -s 1048576 ulimit -m 1048576