

Recuperare Protocoale criptografice

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Examen

Ex #1 Prototipul polinomial $x^3 + x^2 + 1$ este ireductibil peste \mathbb{F}_2 . Este ω o radacina a polinomului.

Calculati elementul $(\omega^2 + \omega + 1)^{-1}$ in $\mathbb{F}_8 = \mathbb{F}_2[\omega]$

Denumire

Este $f(x) = x^3 + x^2 + 1$. Observam ca f are acelasi punct de zero peste \mathbb{F}_2 deoarece

$$\begin{aligned} f(0) &= 0 + 0 + 1 \stackrel{2}{=} 1 \\ f(1) &= 1 + 1 + 1 \stackrel{2}{=} 1 \end{aligned} \quad \text{Astfel}$$

Dacă ar fi fost redusibil, ar fi găsit un factor de $\deg = 1$ și un factor de $\deg = 2$.

Dacă ω este radacina a lui f , regula de calcul din $\mathbb{F}_8 = \mathbb{F}_2[\omega]$ este $\omega^3 = \omega^2 + 1$, iar elementele din \mathbb{F}_8 sunt de forma $a + bw + cw^2$ cu $a, b, c \in \mathbb{F}_2$. Către totuntem un astfel de element să se satisfacă

$$(\omega^2 + \omega + 1)(a + bw + cw^2) = 1$$

Calculare

$$c\omega^4 + (b+c)\omega^3 + (a+b+c)\omega^2 + (a+b)\omega + a = 1 \quad (*)$$

Stimmt es $\omega^3 = \omega^2 + 1$, dann

$$\begin{aligned}\omega^4 &= \omega^3 + \omega \\ &= \omega^2 + 1 + \omega\end{aligned}$$

Dann $(*)$, ausmultiplizieren:

$$c(\omega^2 + 1 + \omega) + (b+c)(\omega^2 + 1) + (a+b+c)\omega^2 + (a+b)\omega + a = 1$$

$$\begin{aligned}\cancel{(c+b+c)} + \cancel{a+b+c} \omega^2 + (c+a+b)\omega + \\ + \cancel{(c+b+c)} + a &\stackrel{\text{(mod 2)}}{=} 1\end{aligned}$$

Ausarbeiten

$$\cancel{(a+c)}\omega^2 + \cancel{(a+b+c)}\omega + \cancel{(a+b)} = 1 + \underbrace{0 \cdot \omega}_{\text{mod 2}} + \underbrace{0 \cdot \omega^3}_{\text{mod 2}}$$

$$\Rightarrow \begin{cases} a+b = 1 & (1) \\ a+b+c = 0 & (2) \quad (\text{mod 2}) \\ a+c = 0 & (3) \end{cases}$$

$$(3) \Rightarrow a = c$$

$$(2) \Rightarrow 2a+b = 0 \Rightarrow b = 0$$

$$(1) \Rightarrow a+0 = 1 \Rightarrow a = 1 \Rightarrow c = 1$$

Prim zermore, $(a,b,c) = (1,0,1)$, dann

$$(\omega^2 + \omega + 1)^{-1} = 1 + \omega^2$$

Verificare:

$$\begin{aligned}
 (1+\omega+\omega^2)(1+\omega^2) &= 1+\cancel{\omega^2} + \omega + \omega^3 + \cancel{\omega^2} + \omega^4 \\
 &= 1 + \cancel{\omega} + \cancel{\omega^2} + 1 + \cancel{\omega^2} + 1 + \omega \\
 &= 1.
 \end{aligned}$$

□

Ex #2 Showur Secret Sharing. Fie $P \in \mathbb{Z}_{19}[x]$ un pol. de $\deg = 2$. Se consideră zdrobitorile perechi $(\alpha, P(\alpha))$ unde $\alpha \in \mathbb{Z}_{19}^*$ și $P(x) \in \mathbb{Z}_{19}$:

$(10, 16), (11, 0)$ și $(12, 5)$. Deduciți secretul pe care îl reprezintă $\Delta = P(0) \in \mathbb{Z}_{19}$.

Din
nu

Considerăm $P(x) = \Delta + ax + bx^2 \in \mathbb{Z}_{19}[x]$. Atunci avem

$$\begin{cases}
 \Delta + 10a + 100b \stackrel{19}{=} 16 \\
 \Delta + 11a + 121b \stackrel{19}{=} 0 \quad r=7 \\
 \Delta + 12a + 144b \stackrel{19}{=} 5
 \end{cases}$$

$$\begin{cases}
 \Delta + 10a + 5b = 16 \\
 \Delta + 11a + 4b = 0 \\
 \Delta + 12a + 11b = 5
 \end{cases}$$

$$\left[\begin{array}{ccc|c}
 1 & 10 & 5 & 16 \\
 1 & 11 & 4 & 0 \\
 1 & 12 & 11 & 5
 \end{array} \right] \xrightarrow[L_2 - L_1]{L_3 - L_1} \left[\begin{array}{ccc|c}
 1 & 10 & 5 & 16 \\
 0 & 1 & 2 & 3 \\
 0 & 2 & 6 & 8
 \end{array} \right] \xrightarrow[\frac{1}{2}L_3]{}$$

$$\left[\begin{array}{ccc|c} 1 & 10 & 5 & 16 \\ 0 & 1 & 2 & 3 \\ 0 & 1 & 3 & 4 \end{array} \right] \xrightarrow{L_3 - L_2} \left[\begin{array}{ccc|c} 1 & 10 & 5 & 16 \\ 0 & 1 & 2 & 3 \\ 0 & 0 & 1 & 1 \end{array} \right] \xrightarrow{\frac{L_2 - 2L_3}{L_1 - 5L_3}} \left[\begin{array}{ccc|c} 1 & 10 & 5 & 16 \\ 0 & 1 & 0 & 3 \\ 0 & 0 & 1 & 1 \end{array} \right]$$

$$\left[\begin{array}{ccc|c} 1 & 10 & 0 & 11 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{array} \right] \xrightarrow{L_1 - 10L_2} \left[\begin{array}{ccc|c} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{array} \right]$$

$$\Rightarrow \begin{cases} A = 1 \\ Q = 1 \\ b = 1 \end{cases}$$

□

RSA. O persoană encriptează mesajul $m \bmod 85$. folosind cheia publică $e=9$ și obține $c=10$. Decriptarea mesajului folosind fumetă

- a) $\varphi(n)$
- b) $\lambda(n)$

Deu

[a] Stim $n = 85 = 5 \cdot 17$
 $ed = 1 \bmod \varphi(n)$ $\textcircled{*}$

$$\varphi(n) = (p-1)(q-1) = (5-1)(17-1) = 4 \cdot 16$$

$$\Rightarrow \varphi(85) = 64$$

Dein $\textcircled{*} \Rightarrow d = 9^{-1} \bmod 64$

Calc. $g^{-1} \pmod{64}$ folosind Euclid extins.

$$64 = g \cdot 7 + 1 \Rightarrow 1 = 64 - g \cdot 7 \pmod{64}$$
$$\Rightarrow 1 = -g \cdot 7 \pmod{64}$$
$$\Rightarrow 1 = g \cdot 57 \pmod{64}$$

Deci $d = 57 \pmod{64}$.

Călcălare $a_n = c^d \pmod{n}$, ie

$$a_n = 10^{57} \pmod{85}$$

folosind alg. de exponentiere rapidă.

$$57 = 1 + 56$$

$$57 = 32 + 16 + 8 + 1$$

Aveam

$$\begin{aligned} 10^2 &= 10 \\ 10^4 &= 55 \\ 10^8 &= 50 \\ 10^{16} &= 35 \\ 10^{32} &= 35 \end{aligned} \quad \left. \begin{array}{l} \\ \\ \\ \end{array} \right\} \Rightarrow a_n = 10^{57} = 35 \cdot 35 \cdot 50 \cdot 10 \\ &= 75 \\ \Rightarrow \boxed{a_n = 75} \end{aligned}$$

□

b) Stiu $a_n = 85 = 5 \cdot 17$

$$\frac{e}{\lambda(n)} d = 1 \pmod{\lambda(n)}$$

$$\lambda(n) = \text{lcm}(5-1, 17-1) = \text{lcm}(4, 16) = 16$$

Calculation $d = g^{-1} \pmod{16}$ see Euclid extius.

$$\begin{array}{l} 16 = g \cdot 1 + 4 \\ g = 7 \cdot 1 + 2 \Rightarrow 1 = 4 - 2 \cdot 3 \\ 7 = 2 \cdot 3 + 1 \end{array} \quad \begin{aligned} &= 7 - (g-7) \cdot 3 \\ &= 7 - 4 - g \cdot 3 \\ &= (16-g) \cdot 4 - g \cdot 3 \end{aligned}$$

$$\Rightarrow l = 16 \cdot 4 - g \cdot 7 \pmod{16} \Rightarrow l = -g \cdot 7 \pmod{16}$$

$$\rightarrow d = -7 = 9 \pmod{16} \rightarrow \boxed{d=9}$$

Decipher: $a_u = c^d \pmod{n}$

$$\Rightarrow a_u = 10^9 \pmod{85}$$

Exponentiation modulo: $g = 8+1$

$$10^2 = 15$$

$$10^4 = 55 \pmod{85}$$

$$10^8 = 50$$

Azunder $a_u = 50 \cdot 10 = 45 \Rightarrow \boxed{u=45}$.

Algoritm de afcore a logoritmului discret

Baby Step - Giant Step

OBS • Breake force \rightarrow It's poss calc. mod 19.

Vrem x cu $\boxed{\cancel{2^x = y \pmod{19}}}$

Algorithm

Stim y, g și x . Vrem $g^x = y \pmod{p}$

x sol. logaritmul discret a lui y
 $x = \log_g y \pmod{p}$

Făcem folial eò fiecare $x < p$ și poate scrie

$$x = \lceil \sqrt{p} \rceil x_1 + x_2, \text{ unde } 0 \leq x_1, x_2 \leq \lfloor \sqrt{p} \rfloor$$

$$\lceil 3,14 \rceil = 4 \quad \lfloor 3,14 \rfloor = 3$$

Aveam

$$y = g^x = g^{\lceil \sqrt{p} \rceil x_1 + x_2} = (g^{\lceil \sqrt{p} \rceil})^{x_1} \cdot g^{x_2}$$

$$\Leftrightarrow y \cdot (g^{-1})^{x_2} = (g^{\lceil \sqrt{p} \rceil})^{x_1}$$

Caculam $g^{-1} \pmod{p} = z$ și $g^{\lceil \sqrt{p} \rceil} \pmod{p} = w$

Scriem zistele

$$L_1 = \{(x_1, \omega^{x_1}) \mid x_1 = 0, 1, \dots, \lfloor \sqrt{p} \rfloor\}$$

$$L_2 = \{(x_2, yz^{x_2}) \mid x_2 = 0, 1, \dots, \lfloor \sqrt{p} \rfloor\}$$

Se canta o cotejamento de tipos $(x_1, \delta) \in L_1$ e $(x_2, \delta) \in L_2$. Aí temos $\alpha = \lceil \sqrt{p} \rceil x_1 + x_2$ é o
segundo cortejo.

$$\boxed{2^x \equiv 6 \pmod{19}}$$

$$p = 19 \quad y = 6 \\ g = 2$$

$$\underline{\text{OBS: }} \sqrt{19} \approx 4,35$$

$$\lceil \sqrt{19} \rceil = 5 \text{ e } \lfloor \sqrt{19} \rfloor = 4$$

Considerar $\alpha < p$ em $\alpha = 5x_1 + x_2$, sendo

$$0 \leq x_1, x_2 \leq 4$$

Calcular $g^{-1} \pmod{p} = 2^{-1} \pmod{19}$ (Euclides)

$$\Rightarrow \dots \Rightarrow 2^{-1} \equiv 10 \pmod{19} \Rightarrow \boxed{z = 10}$$

Calcular $g^{\lceil \sqrt{p} \rceil} \pmod{p} = 2^5 \pmod{19}$ (expo. rápida)

$$\Rightarrow \dots \Rightarrow 2^5 \equiv 13 \pmod{19} \Rightarrow \boxed{\omega = 13}$$

Construir 2 istele

$$L_1 = \{(0, 1); (1, 13); \cancel{(2, 17)}, (3, 12), (4, 1)\}$$

$$L_2 = \{(0, 6); (1, 3); (2, 11); (3, 15); \cancel{(4, 17)}\}$$

$$\begin{aligned}\omega^0 &= 13^0 = 1 \\ \omega^1 &= 13^1 = 13 \\ \omega^2 &= 13^2 = 17 \\ \omega^3 &= 13^3 = 12 \\ \omega^4 &= 13^4 = 4\end{aligned}$$

$$\begin{aligned}yz^0 &= 6 \cdot 10^0 = 6 \\ yz^1 &= 6 \cdot 10^1 = 3 \\ yz^2 &= 6 \cdot 10^2 = 11 \\ yz^3 &= 6 \cdot 10^3 = 15 \\ yz^4 &= 6 \cdot 10^4 = 17\end{aligned}$$

Gösim çözümleri

- $(x_1, y) = (2, 17) \in L_1$
- $(x_2, y) = (4, 17) \in L_2$

Şimdi calculation $x = 5 \cdot x_1 + x_2 \pmod{19}$

$$x = 5 \cdot 2 + 4 = 14 \pmod{19}$$

□

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