

HASHING.

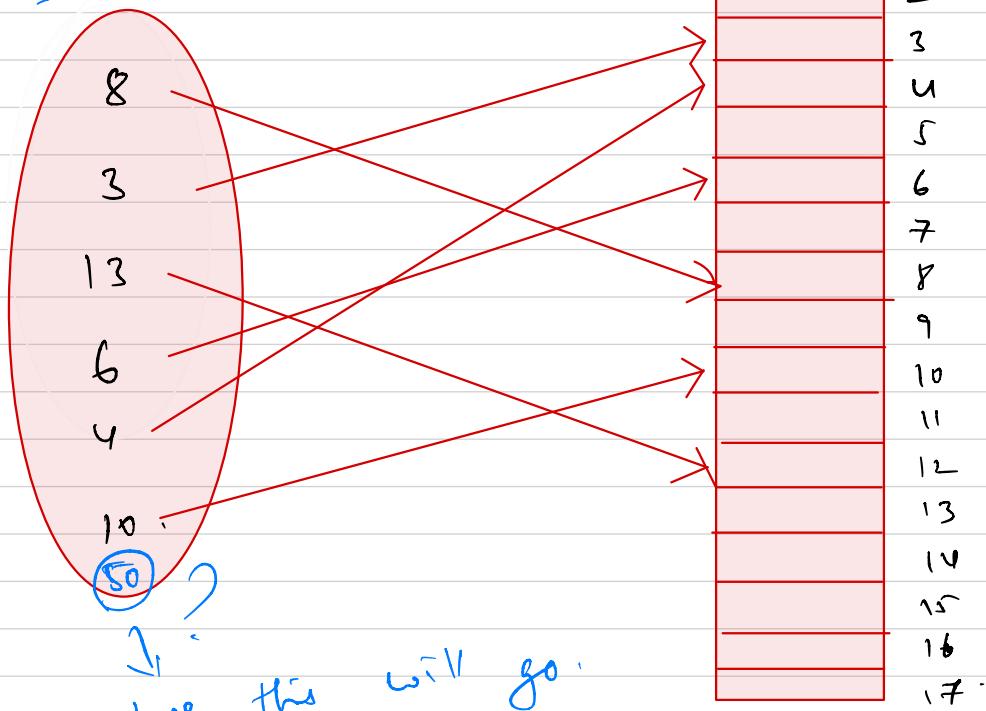
HASHING

Theory :-

- Hashing is a technique or process of mapping keys and values into hashtable by using hash function.
- It usually has worst case $O(1)$. *insert all these values in this table*

$$h(a) = x$$

Keys.



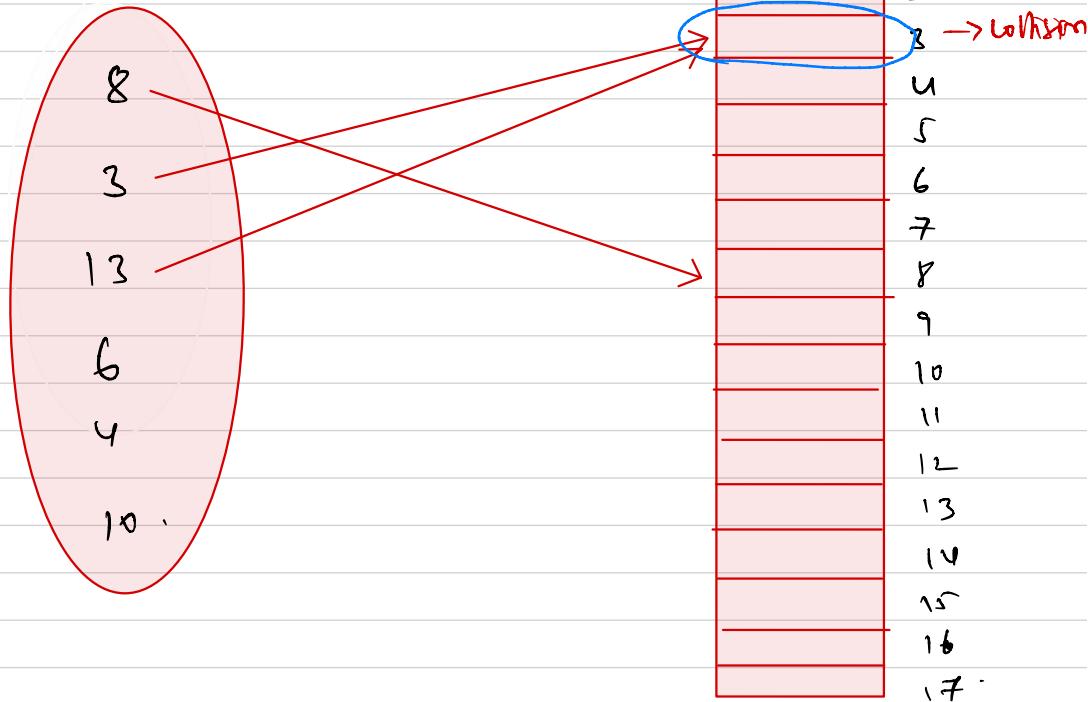
→ This can be resolved using proper hash functions.

Hash functions: There can be different types of Hash functions.

- ① $K \bmod 10$
 - ② $K \bmod n$
 - ③ Mid square.
 - ④ Folding method.
- **most important**.

→ Let us now consider Hash function $K \bmod 10$.

$$h(n) = F' \cdot 10$$



~~★ ★ *~~

→ When you change the hash function there is a possibility that two keys are now mapped to same index of hash table. This is called Collision.

Collision resolution Techniques:-

① Open Hashing.

(i) Chaining.

② Closed Hashing.

(i) Linear probing.

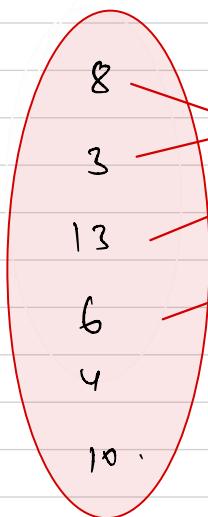
(ii) Quadratic probing.

① Open Hashing:- When there is a collision element will be added to a new node in a chain attached to the index. Chain is usually made of linked list.

D's advantages:- Now the search complexity of Hash table is not $O(1)$ because now we have to search over chains.



$$h(x) = x \% 10$$



Collision:

0
1
2
3
4
5
6
7
8
9
10
11
12
13
14
15
16
17

0

1

2

3

4

5

6

7

8

9

10

11

12

13

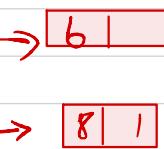
14

15

16

17

chains



- ② Linear probing :- Here if there is a collision then we search for next empty slot in hash table and store it there.

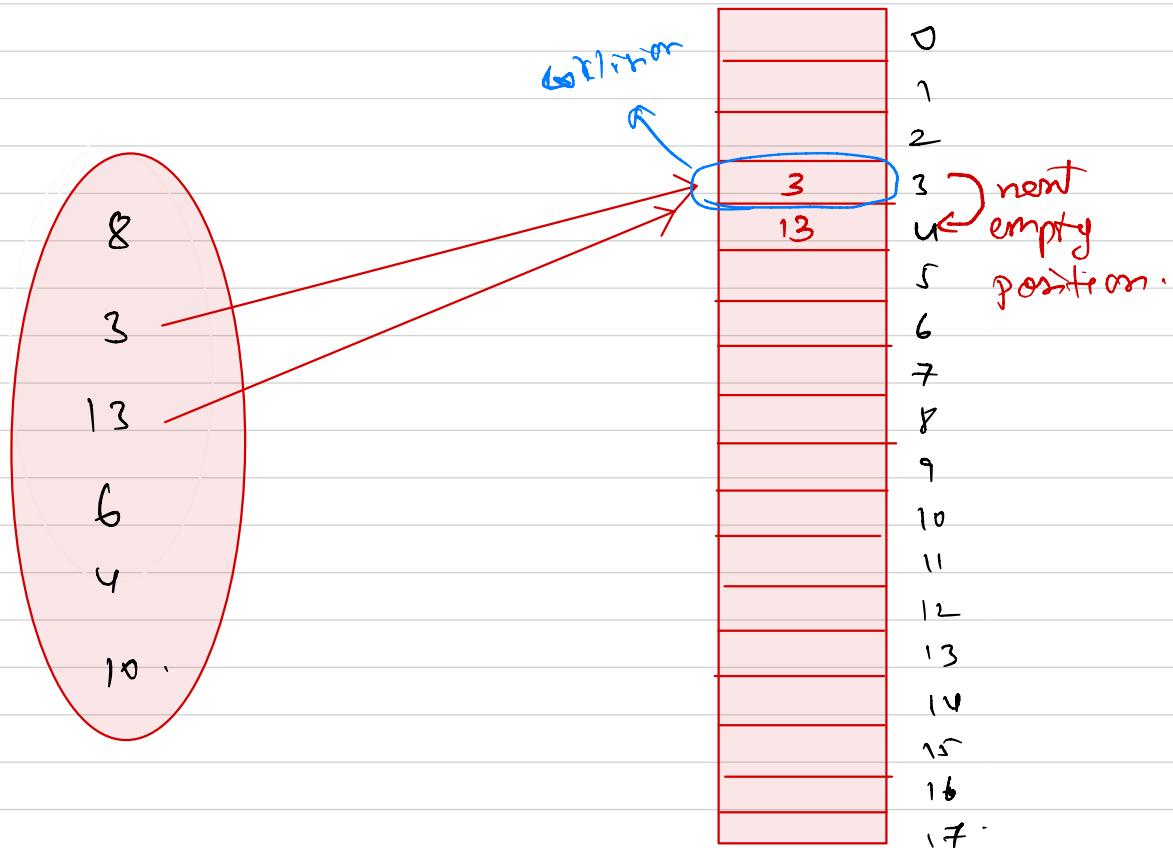
$$h(x) = x \% \text{size}$$

} new functions

$$h(x) = [h(x) + f(i)] \% \text{size}$$

$$f(i) = i$$

$$i = 1, 2, \dots$$



$$h'(13) = [h(13) + f(10)] \mod 10 \quad f(i) = 0, 1, 2, 3, \dots \\ = 3.$$

$$h'(13) = [h(13) + f(10)] \mod 10 \\ = 4. \rightarrow \text{next index}$$

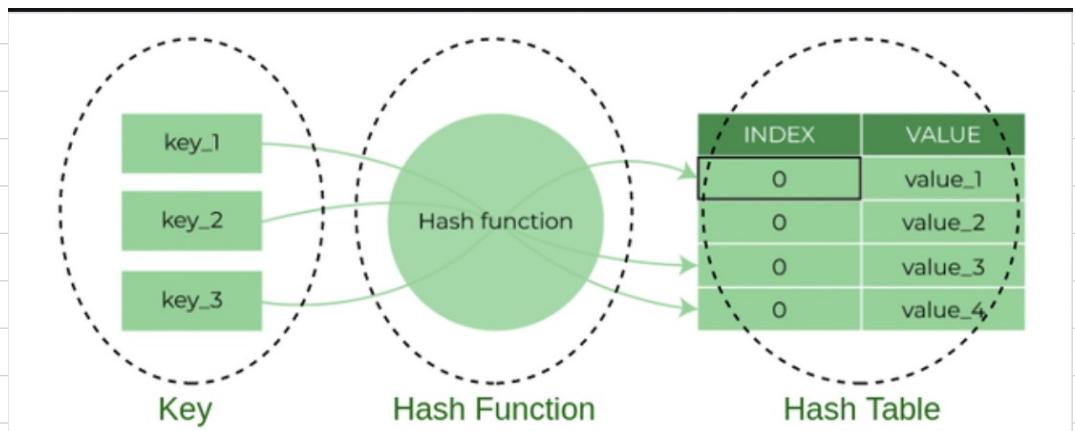
↳ keep doing until
empty index is found.

③ Quadratic Probing :- This is also similar to linear probing but slightly different hash function.

$$h'(x) = [h(x) + f(i)] \% \text{size}$$

$$f(i) = i^2$$

$$i=0, 1, 2, 3, \dots$$



Components of Hashing

Reference :- Books for books.

Problems:-

Two Sum:-

Problem 1 i:- find indices in array such that they add upto target.

Input :- $\text{nums} = [2, 7, 11, 15]$, $\text{target} = 9$.

Output - $[0, 1]$.

$$\text{num}[0] + \text{num}[1] == 9.$$

Soln:-

- The most important fact for this problem is "There is exactly one solution".
- We can calculate complement and check if that is present.

$$\text{Complement} = \text{target} - \text{current num}.$$



check if other number which adds target is present.

{

```
Map<int, Integer> map = new Map();
```

```
for (int i=0; i< nums.length; i++)
```

{

```
int complement = target - nums[i];
```

→ most important thing

```
if (map.containsKey(complement))
```

d

```
return new Int[] {i, map.get(complement)};
```

else {

```
map.put(nums[i], i);
```

y

}

0	1	2	3
1	2	7	10

1	2	7	10
---	---	---	----

↑ → ↑ → ↑ → look for complement(2).

Key
nums.

value
index

1	0
2	1

[2, 1]

= [2, 1]

Problem 2 :- Random note.

Given two strings ransomNote and magazine, return true if ransomNote can be constructed by using letters from magazine.

Input:- $\text{ransomNote} = \text{"aa"}$.
 $\text{magazine} = \text{"aab"}.$
= True;

most imp
App. \rightarrow fast.
Each letter can be used only once.

Soli-

- Main logic:- The number of characters in ransom note should be present in magazine
- Basically we have to calculate both and every character of magazine and see all characters of ransom note are present.

$\text{magazine} = \text{"aab"}$.

$\text{ransomNote} = \text{"aa"}$.

key	Count
a	2 1 0
b	1

a
a

↓
keep looking in both
task and reduce number.

- Every character should be greater than 0 in the table.