

1. Prove the following using the original definitions of O , Ω , θ , and ω .

1.1. $15n^2 + 5n + 8 = O(n^2)$

$g(n)$ is $O(f(n))$ if for positive constants c and N , $0 \leq g(n) \leq cf(n)$ for all $n \geq N$

$$g(n) = 15n^2 + 5n + 8$$

$$f(n) = n^2$$

$$0 \leq 15n^2 + 5n + 8 \leq cn^2$$

$$0 \leq 15 + 5/n + 8/n^2 \leq c \text{ (Divide both sides by } n^2 \text{)}$$

$$\text{For } N = 1, c = 28, 0 \leq 28 \leq 28$$

Expression is satisfied.

1.2. $n^3 + 2n + 5 = \Omega(n^3)$

$g(n)$ is $\Omega(f(n))$ if for positive constants c and N , $0 \leq cf(n) \leq g(n)$ for all $n \geq N$

$$g(n) = n^3 + 2n + 5$$

$$f(n) = n^3$$

$$0 \leq cn^3 \leq n^3 + 2n + 5$$

$$0 \leq c \leq 1 + 2/n^2 + 5/n^3 \text{ (Divide both sides by } n^3 \text{)}$$

$$\text{For } N = 1, c = 8, 0 \leq 8 \leq 8$$

Expression is satisfied.

1.3. $\sum_{i=1}^k a_i n^i = \Omega(n^k)$ for $k > 1$. Assume that $a_i > 0$ for all i .

$g(n)$ is $\Omega(f(n))$ if for positive constants c and N , $0 \leq cf(n) \leq g(n)$ for all $n \geq N$

$$g(n) = \sum_{i=1}^k a_i n^i$$

$$f(n) = n^k$$

$$0 \leq cn^k \leq \sum_{i=1}^k a_i n^i$$

$$0 \leq cn^k \leq a(1-n^k)/(1-n) \text{ (Geometric Series)}$$

$$0 \leq cn^k \leq a/(1-n) - an^k/(1-n) \text{ (Separate terms)}$$

$$0 \leq c \leq a/n^k(1-n) - a(1-n) \text{ (Divide both sides by } n^k \text{)}$$

$$\text{For } N = 2, k = 1, a = 1, c = 1/2, 0 \leq 1/2 \leq 1/2$$

Expression is satisfied.

1.4. $n^7 + 3n^2 = \theta(n^7)$

$g(n)$ is $\theta(f(n))$ if for positive constants c , d , and N , $0 \leq cf(n) \leq g(n) \leq df(n)$ for all $n \geq N$

$$0 \leq cn^7 \leq n^7 + 3n^2 \leq dn^7$$

$$0 \leq cn^7 \leq n^7 + 3n^2$$

$$0 \leq c \leq 1 + 3/n^5$$

$$\text{For } N = 1, c = 4, 0 \leq 4 \leq 4$$

Expression is satisfied.

$$n^7 + 3n^2 \leq dn^7$$

For $N = 1$, $d = 4$, $4 \leq 4$

Expression is satisfied.

2. Prove the following using limits.

$$\lim_{n \rightarrow \infty} f(n)/g(n) = \begin{cases} c & \text{then } f(n) = \theta(g(n)) \text{ if } c > 0 \\ 0 & \text{then } f(n) = o(g(n)) \\ \infty & \text{then } f(n) = \omega(g(n)) \end{cases}$$

2.1. $n^k = o(3^n)$ where $k \geq 0$

$$f(n) = n^k$$

$$g(n) = 3^n$$

$$\lim_{n \rightarrow \infty} n^k / 3^n$$

$$\lim_{n \rightarrow \infty} kn^{k-1} / \ln(3) * 3^n \text{ (L'Hopital's)}$$

$$\lim_{n \rightarrow \infty} kn^{k-1} / \ln(3) * 3^n = \lim_{n \rightarrow \infty} k(k-1)n^{k-2} / \ln^2(3) * 3^n = \dots \text{ (Continue taking derivative)}$$

$$\lim_{n \rightarrow \infty} k! / \ln^k(3) * 3^n = 0$$

2.2. $n \ln n^7 = o(n^3)$

$$f(n) = n \ln n^7$$

$$g(n) = n^3$$

$$\lim_{n \rightarrow \infty} n \ln n^7 / n^3$$

$$\lim_{n \rightarrow \infty} \ln n^7 / n^2$$

$$\lim_{n \rightarrow \infty} 7 / 2n^2 = 0 \text{ (L'Hopital's)}$$

3. Prove that $7^n - 1$ is divisible by 6 for $n = 1, 2, 3, \dots$ using induction

$$m \bmod n = m - n \text{ floor}(m / n)$$

Base: For $n = 1$

$$7^n - 1 = 7^1 - 1 = 6. 6 - 6 \text{ floor}(1) = 0. \text{ Satisfied.}$$

Hypothesis: Assume true for $(n+1)$.

Induction: $7^{n+1} - 1$ is divisible by 6.

$7 * 7^n - 1 - (7^n - 1)$. If we subtract a multiple of 6 from a value, we should get another multiple of 6.

$6 * 7^n$ is a multiple of 6 as it is multiplied by 6.

4. Analyze the worst-case time complexity in terms of Big Oh notation. Briefly, yet clearly explain your answer.

4.1. $O(\lg(n))$. By instruction count, we see that $\lg(n) + 1$ instructions are performed in the loop due to the statement $i=2*i$.

4.2. $O(n \lg(n))$. By instruction count, the outer for loop executes n times. The three function calls have time complexities of $O(1)$, $O(\lg(n))$, and $O(1)$, respectively. Since each of these are called n times, the highest order term is then $n \lg(n)$.

4.3. $O(2^s)$, where s is the number of bits for a value n . When the value computed for a factorial gets large enough, it will no longer fit in a word and multiplication will no longer be a constant time operation.