

Algorithm

Review: Data Structures

Instructor: Jae-Pil Heo (허재필)

Arrays

- An *array* is a data structure consisting of a collection of *values*, each identified by array *index*.

- In C/C++ programming language,
 - A consecutive set of memory location
 - Logical order is the same as the physical order
 - Creation of an array

```
Type d[10];
```

```
Type *d = new Type [ size ];
```

- Accessing an element by array index

```
d[5] = 2;
```

- Release the allocated memory

```
delete [] d;
```

Arrays

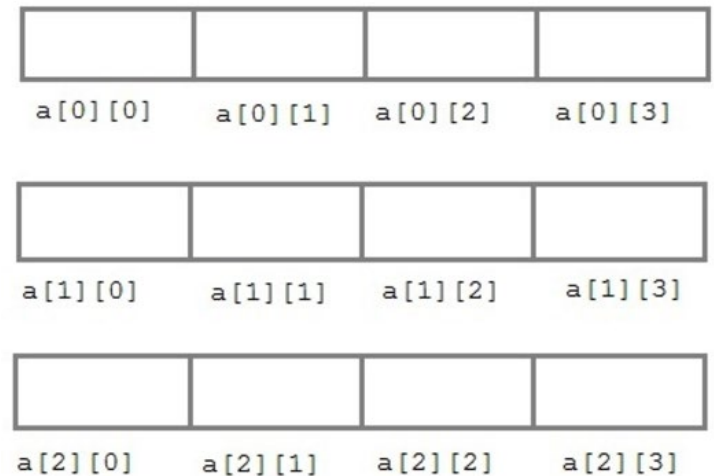
■ 1 dimensional array

- `int arr [10];`
- `int *arr = new int [10];`



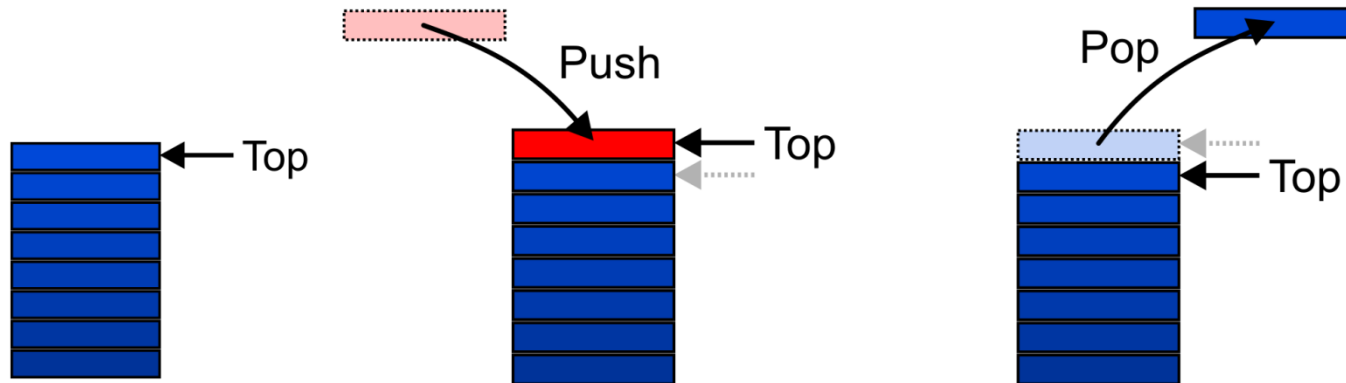
■ 2 dimensional array

- `int a[3][4];`
- `int **a = new int * [3];`
 `for(int i=0;i<3;i++)`
 `{ a[i] = new int [4]; }`



Stacks

- A stack has a *last-in-first-out (LIFO)* behavior
 - Insertion (push) and deletion (pop) are made at one position called 'top'.

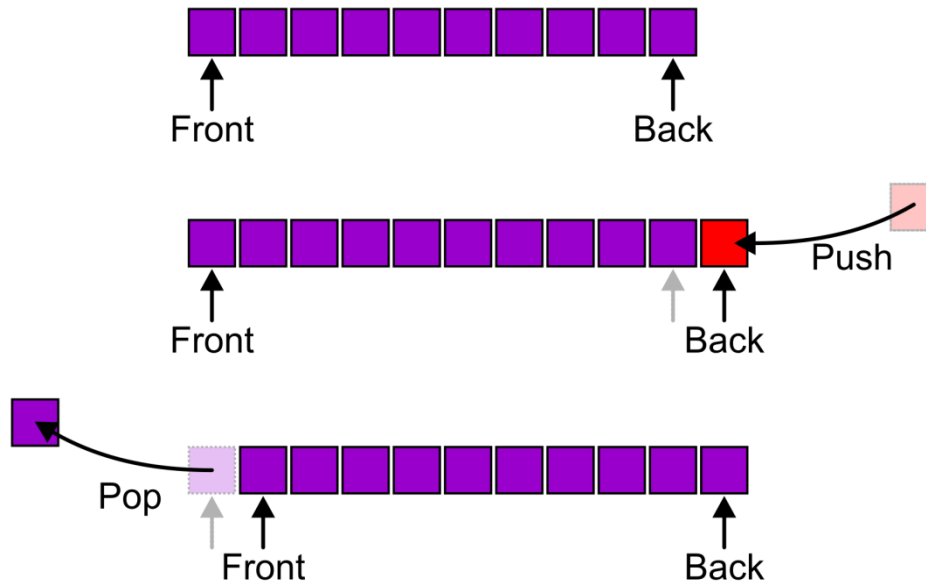


Applications of Stacks

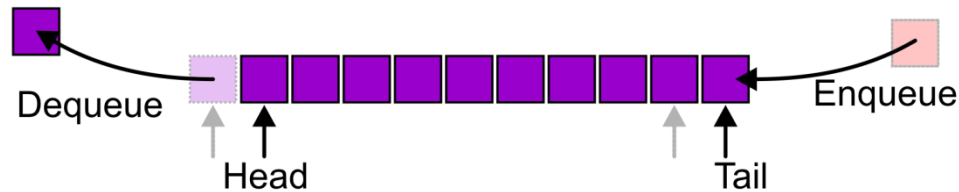
- Code parsing
 - Matching parenthesis
 - XML, HTML, ...
- Tracking functions calls
- Dealing with undo/redo operations
- Tree/Graph traversal

Queues

- A queue has a first-in-first-out (FIFO) behavior.
 - Push (enqueue) is made at the position called 'back (tail)'
 - Pop (dequeue) is made at the position called 'front (head)'

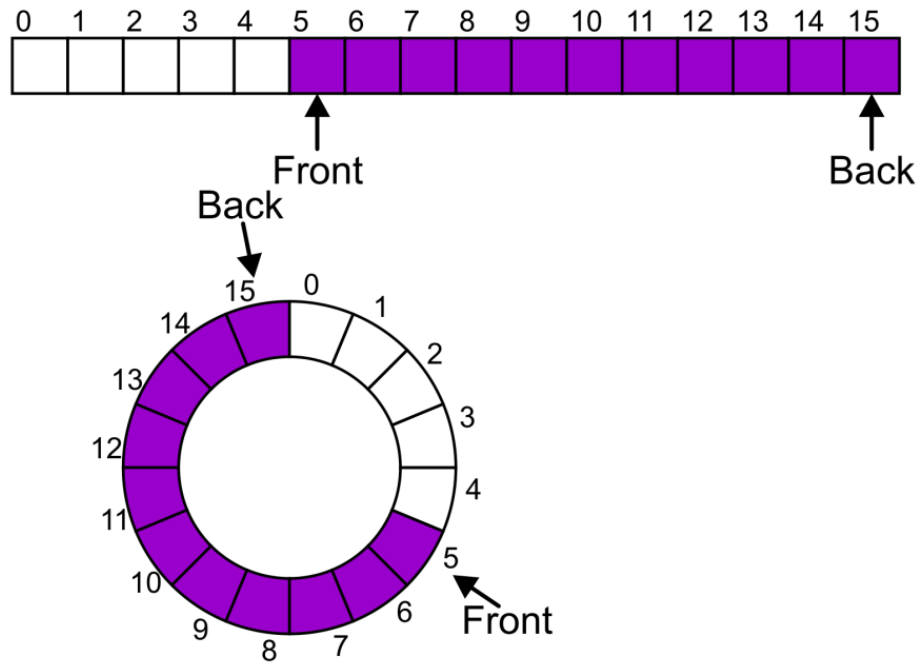


Alternative terms



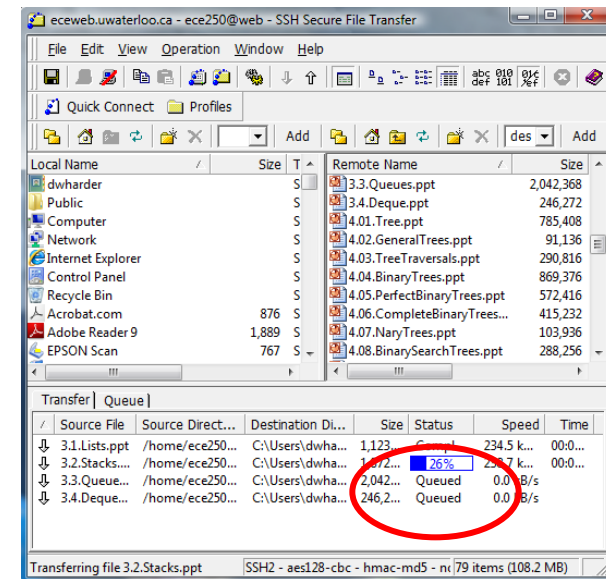
Circular Queue

- Consider the indices being cyclic:



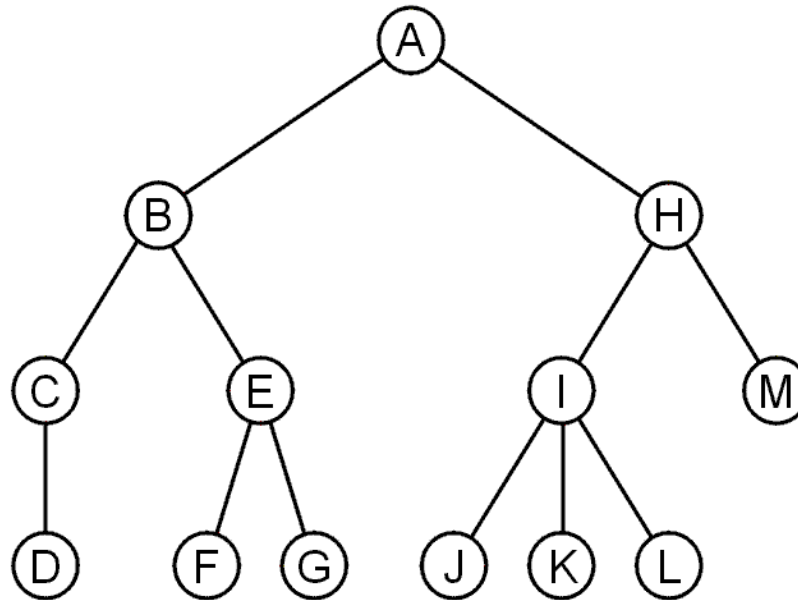
Applications of Queues

- Client-Server (client-casher) models
 - Multiple clients requesting services
 - Web, file ftp, database, mail, printer servers
- Breadth First Search (BFS)



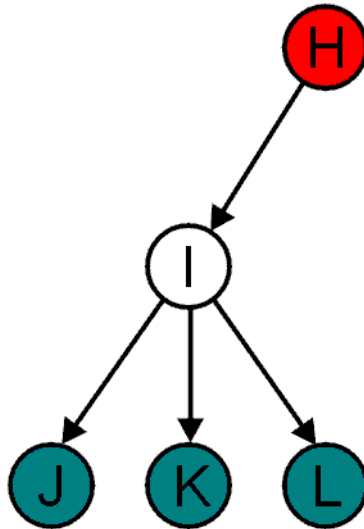
Trees

- A tree data structure stores information in nodes
 - There is a first node, or root
 - Each node has variable number of successors
 - Each node, other than the root, has exactly one node pointing to it



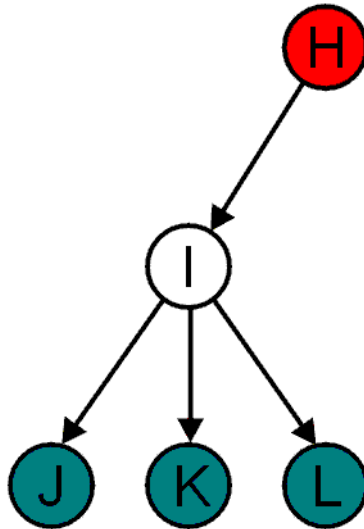
Trees

- All nodes will have zero or more child nodes or children
 - I has three children: J, K, and L
- For all nodes other than the root node, there is one parent node
 - H is the parent of I



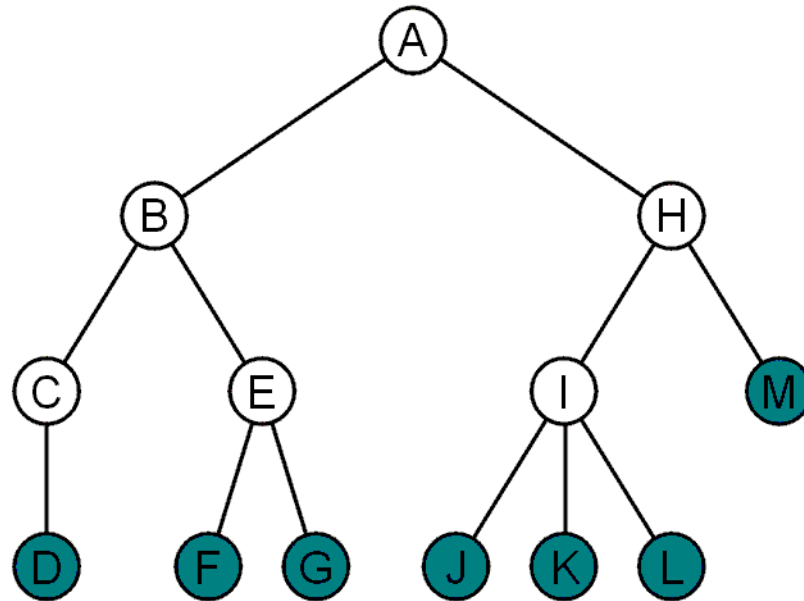
Trees

- The *degree* of a node is defined as the number of its children:
 - $\deg(I) = 3$
- Nodes with the same parent are *siblings*
 - J, K, and L are siblings



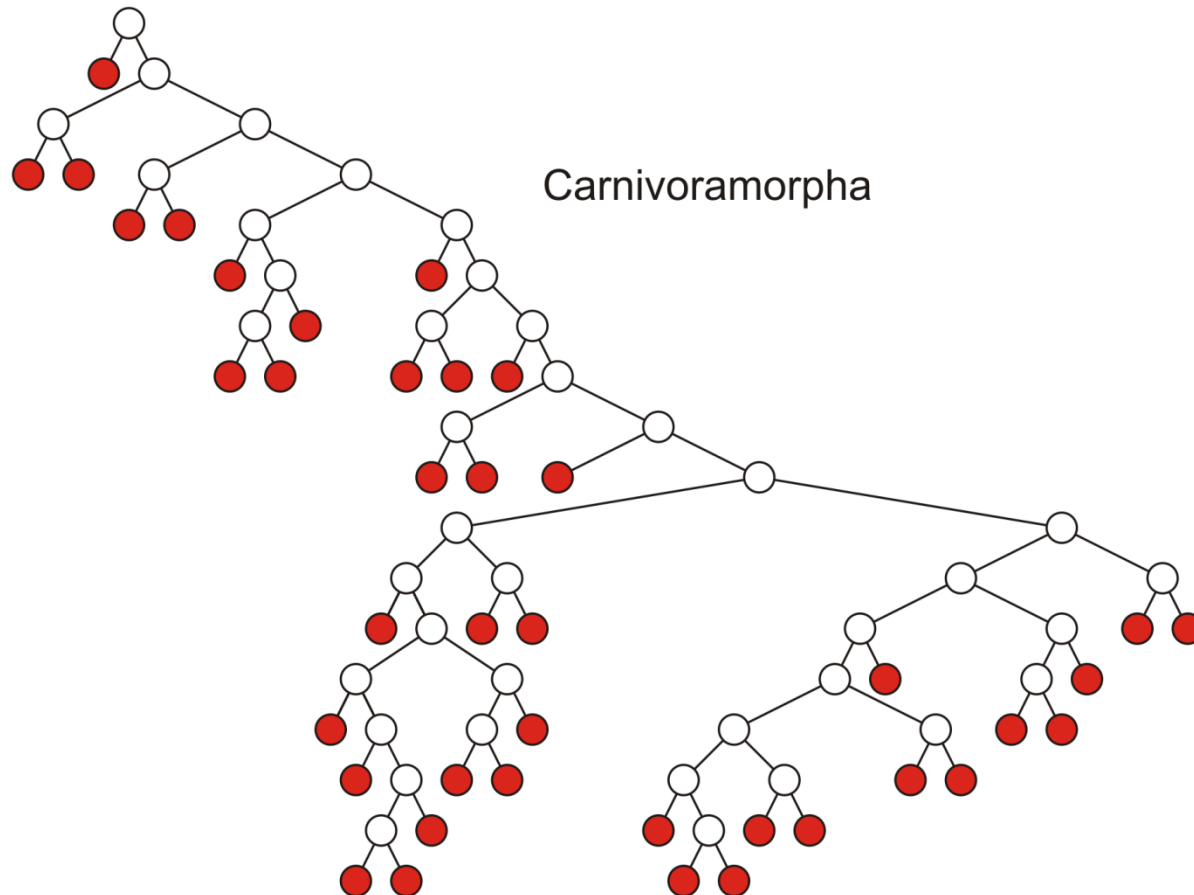
Trees (Leaf and Internal Nodes)

- Nodes with degree zero are also called *leaf nodes*.
- All other nodes are said to be *internal nodes*, that is, they are internal to the tree.



Trees (Leaf and Internal Nodes)

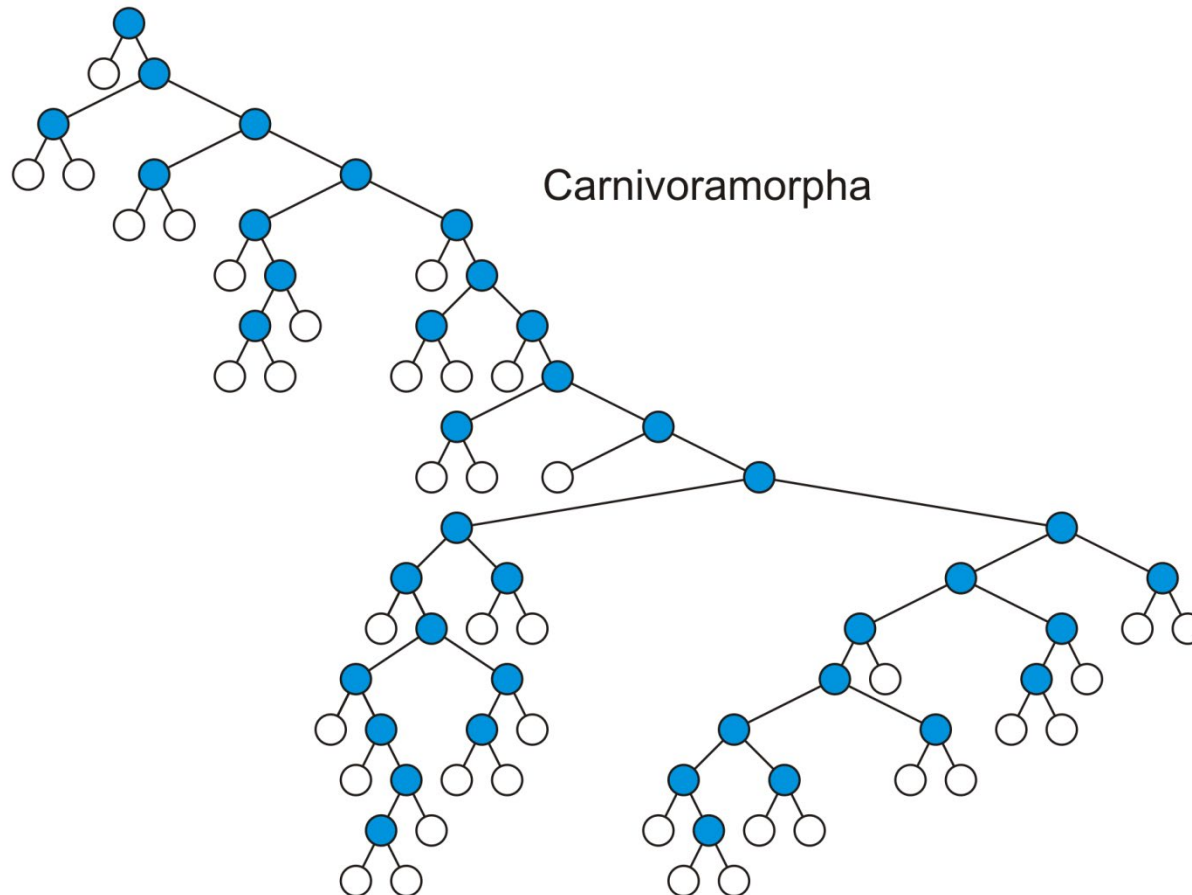
- Leaf nodes



Wesley-Hunt, G. D.; Flynn, J. J. "Phylogeny of the Carnivora: basal relationships among the Carnivoramorpha, and assessment of the position of 'Miacoidea'"

Trees (Leaf and Internal Nodes)

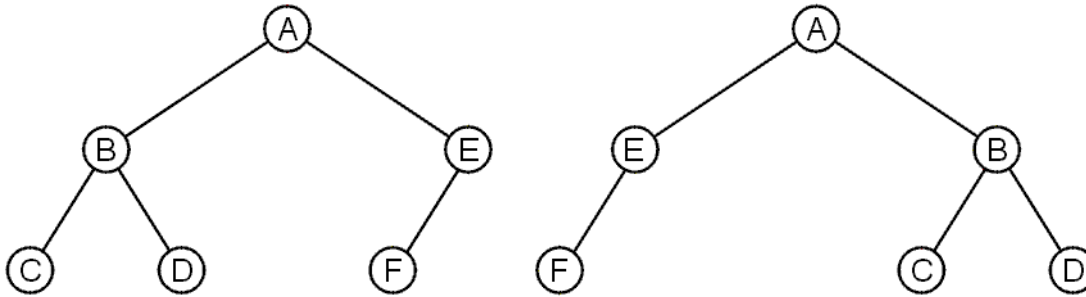
- Internal nodes



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Trees (Unordered & Ordered)

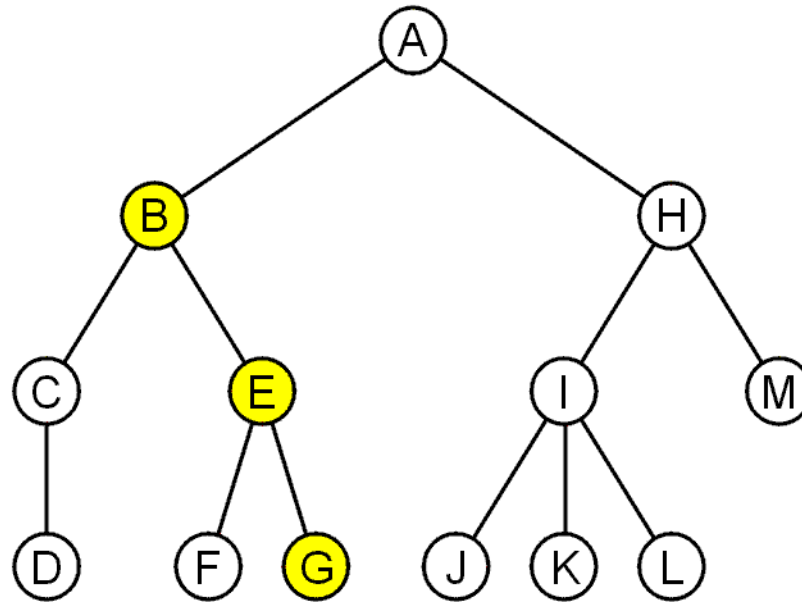
- These trees are equal if the order of the children is ignored
 - unordered tree



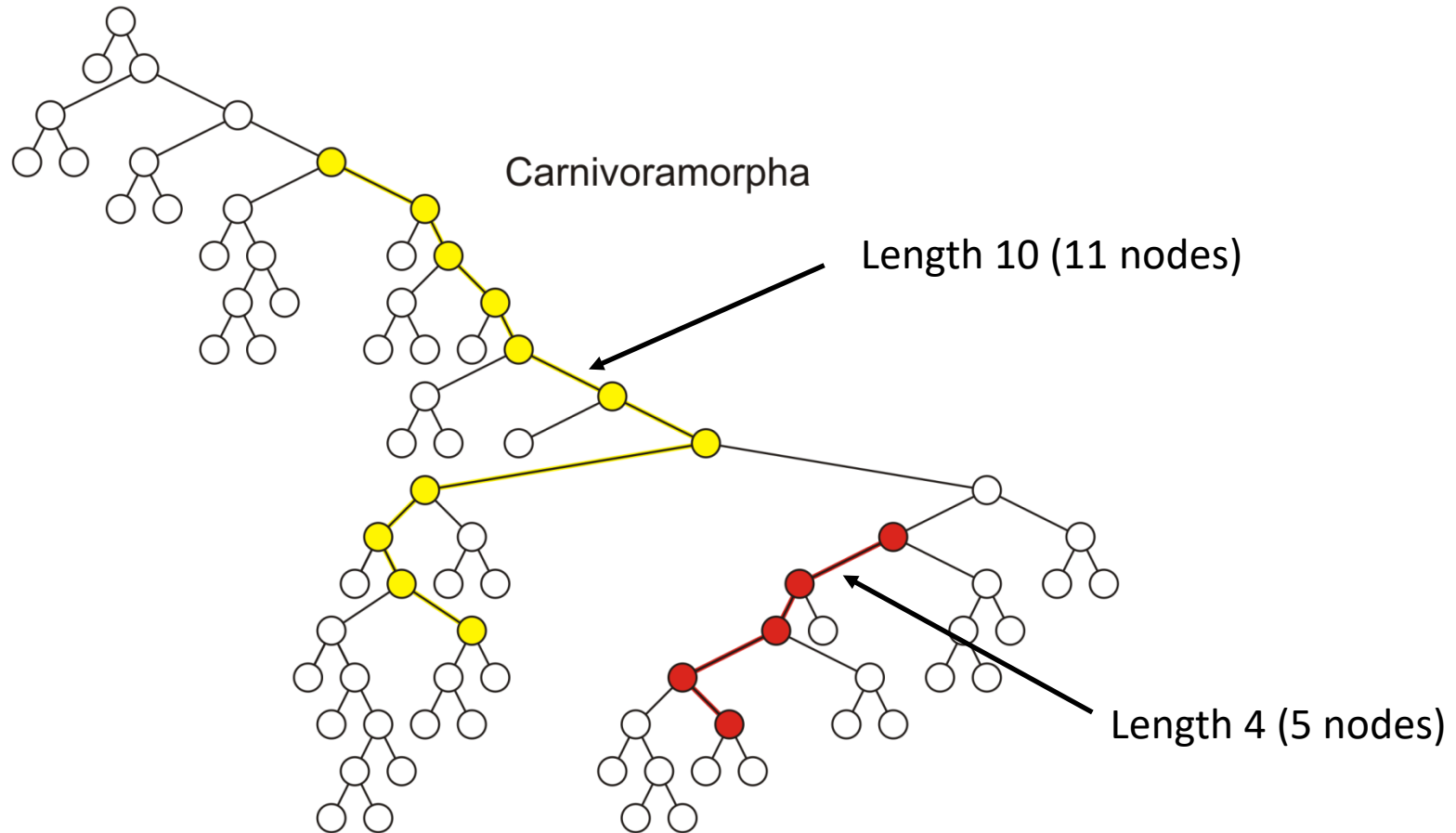
- They are different if order is relevant (ordered trees)

Trees - Paths

- A path is a sequence of nodes (a_0, a_1, \dots, a_n) , where a_{k+1} is a child of a_k .
 - The length of this path is n
 - The path (B, E, G) has length 2



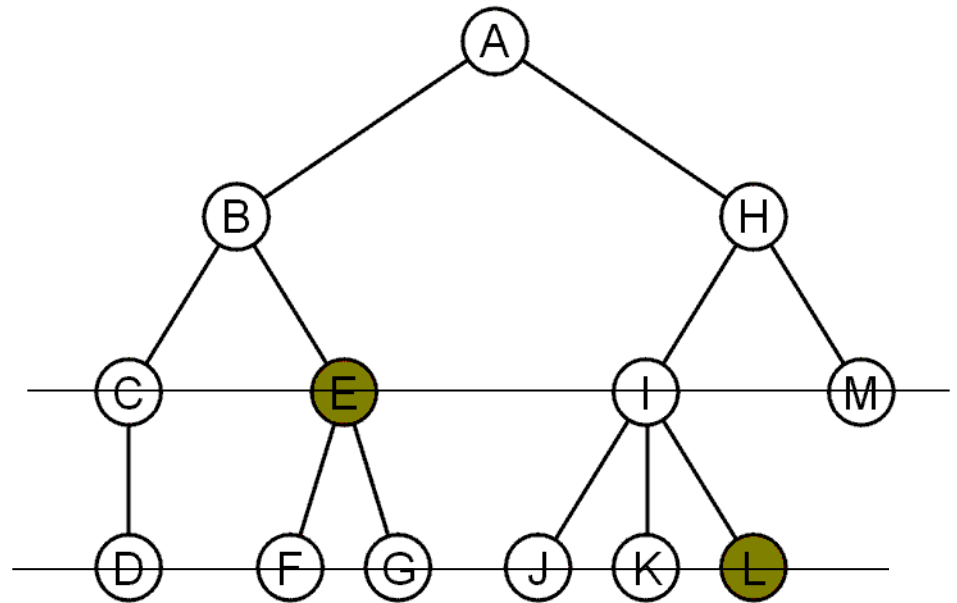
Trees - Paths



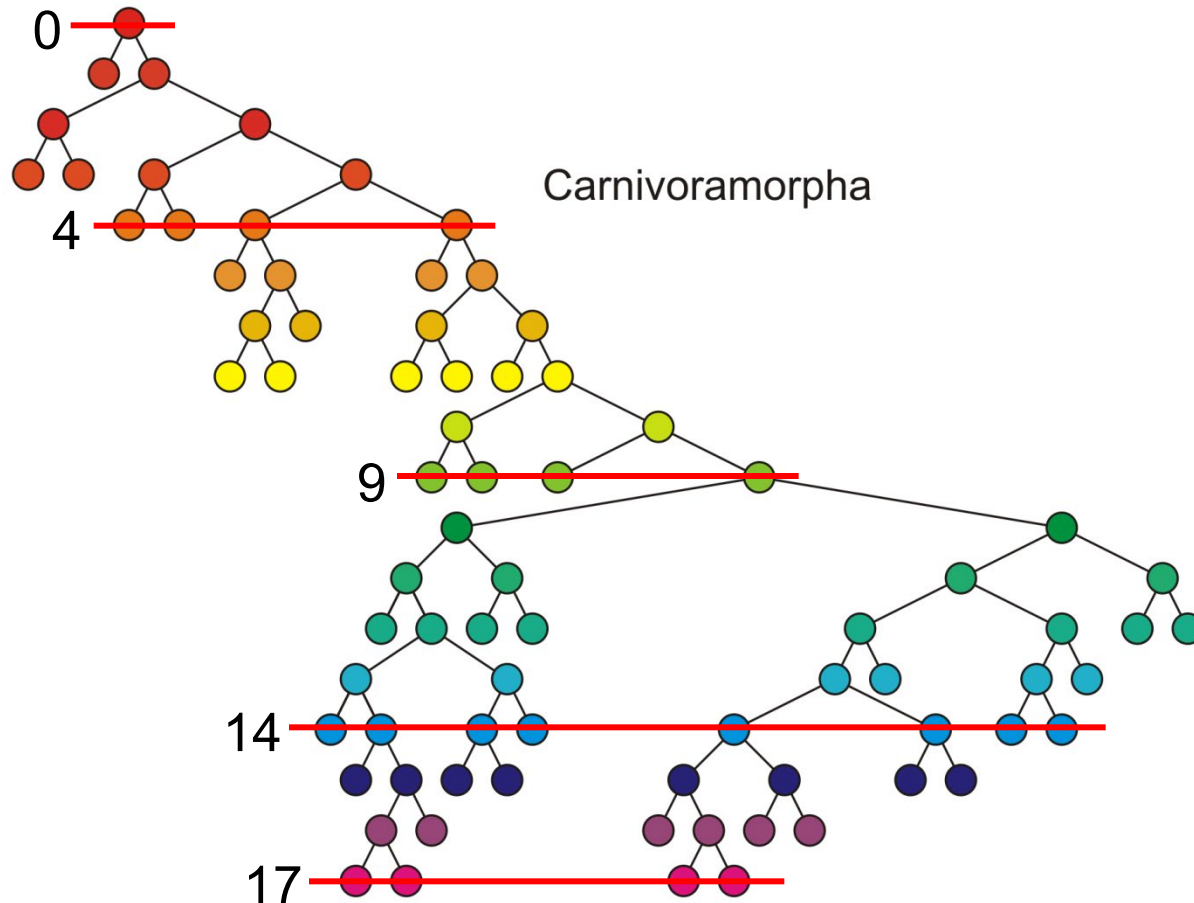
Wesley-Hunt, G. D.; Flynn, J. J. "Phylogeny of the Carnivora: basal relationships among the Carnivoramorpha, and assessment of the position of 'Miacoidea'"

Trees - Depth

- For each node in a tree, there exists a unique path from the root node to that node.
- The length of this path is *depth* of the node
 - E has depth 2
 - L has depth 3



Trees - Depth



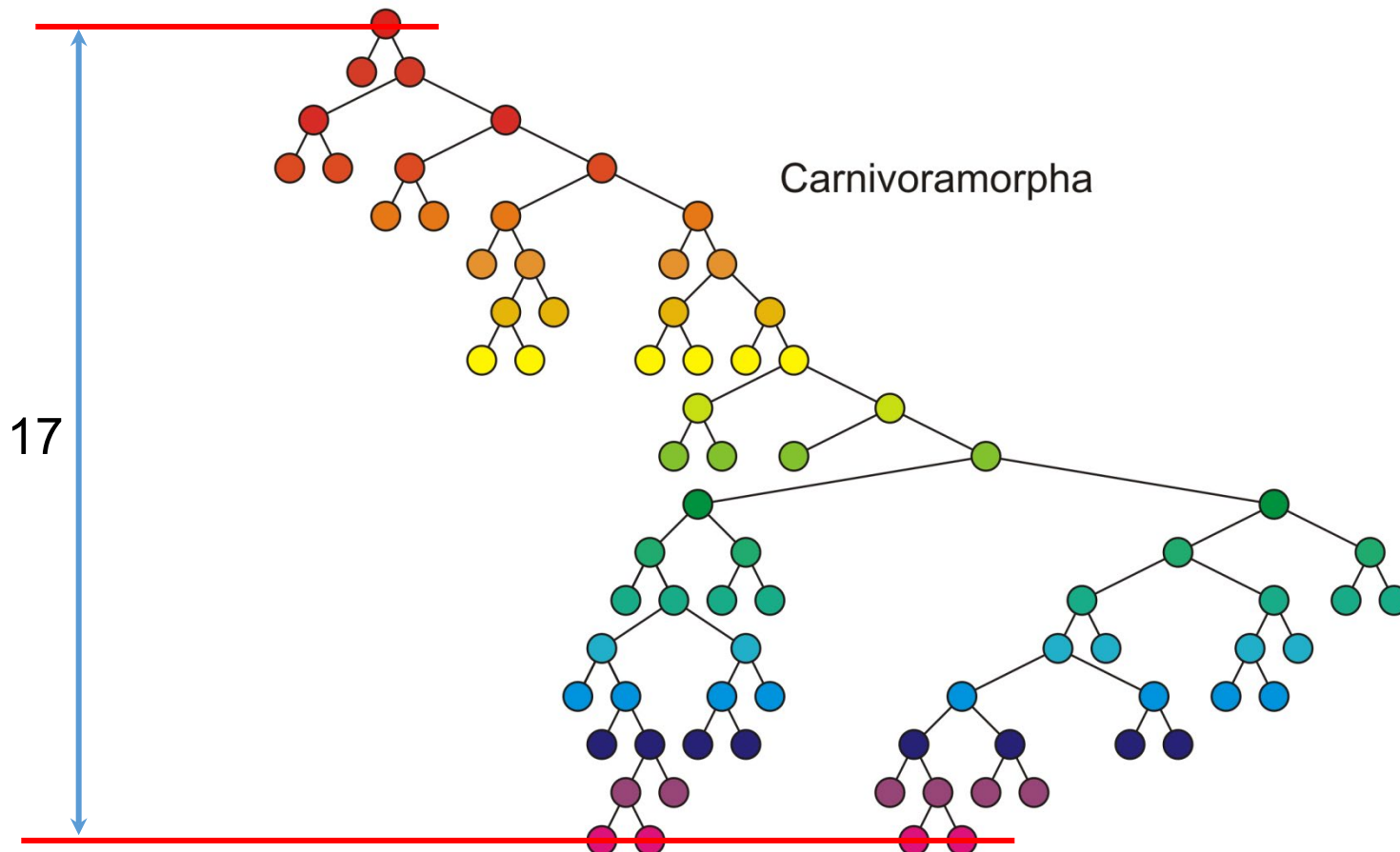
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Trees - Height

- The height of a tree is defined as the maximum depth of any node within the tree
 - The height of a tree with one node is 0
 - Just the root node
 - For convenience, we define the height of the empty tree to be -1

Trees - Height

- The height of this tree is 17



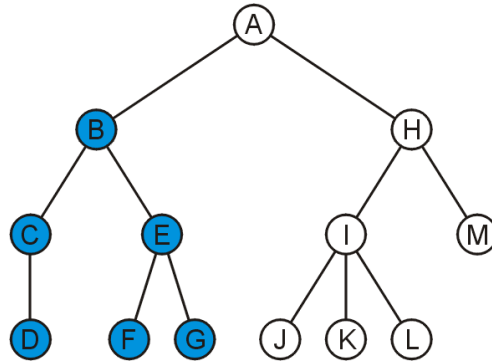
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Trees – Ancestor and Descendant

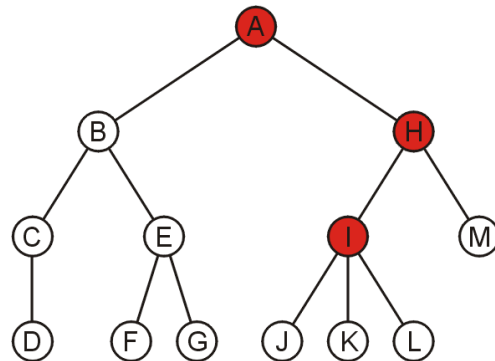
- If a path exists from node a to node b
 - a is an ancestor of b
 - b is a descendant of a
- Thus, a node is both an ancestor and a descendant of itself
 - We can add the adjective *strict* to exclude equality: a is a *strict descendant* of b if a is a descendant of b but $a \neq b$
- The root node is an ancestor of all nodes

Trees – Ancestor and Descendant

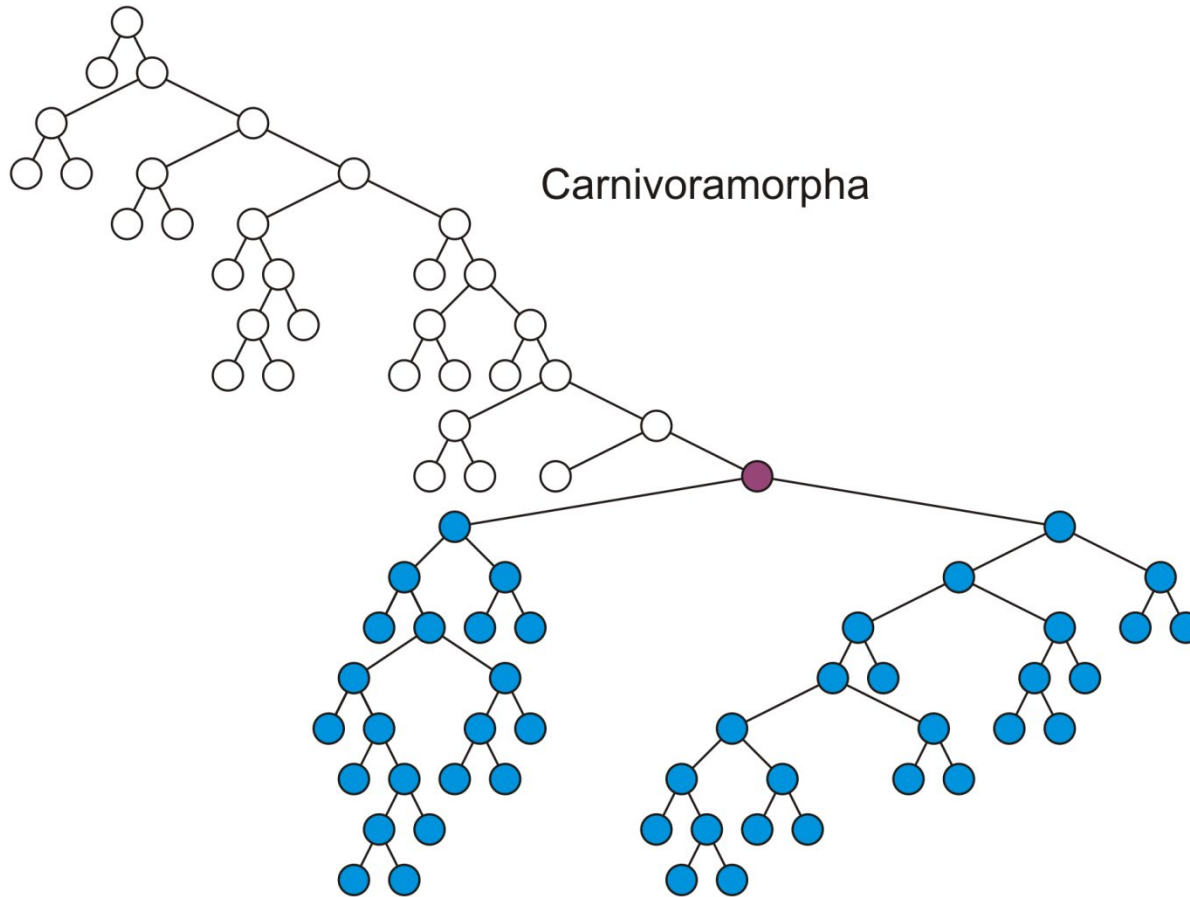
- The descendants of node B are B, C, D, E, F, and G:



- The ancestors of node I are I, H, and A:

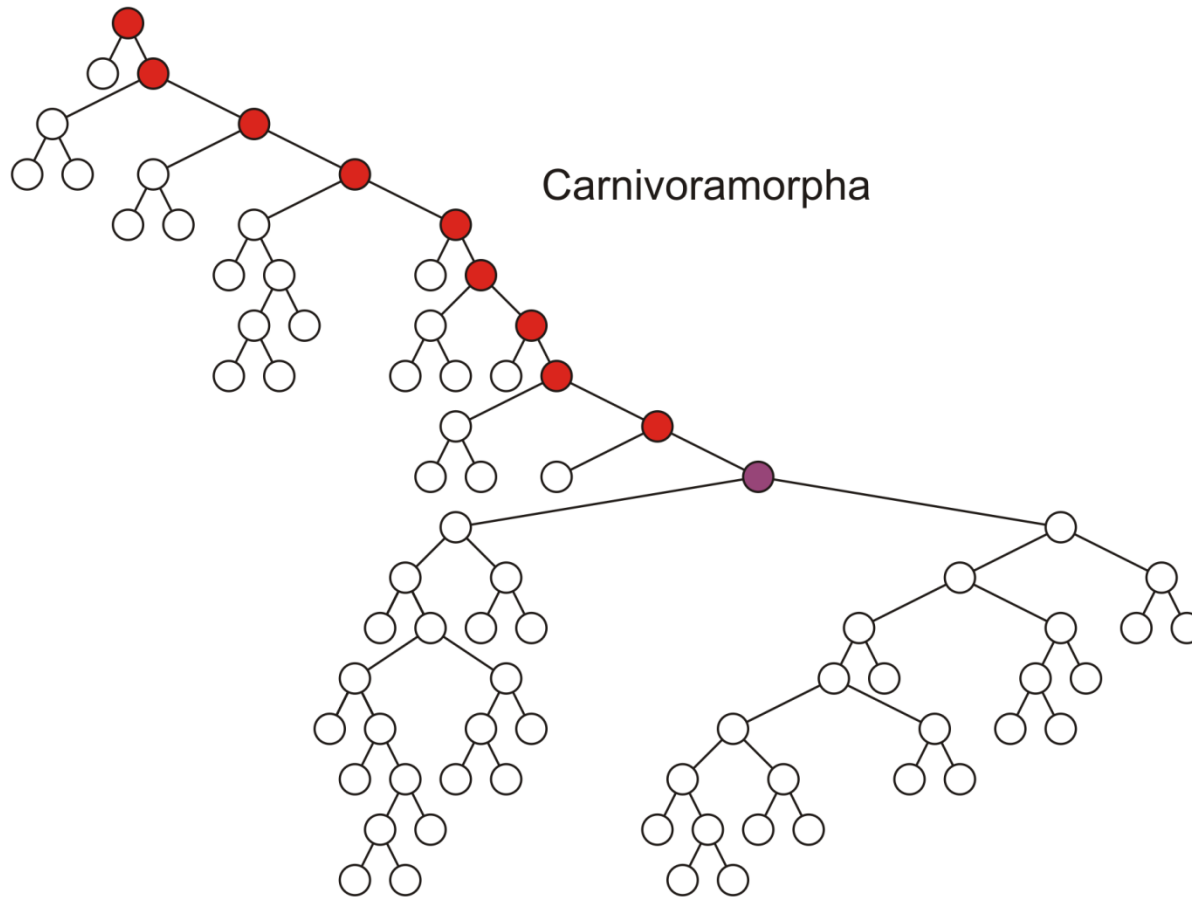


Trees - Descendants



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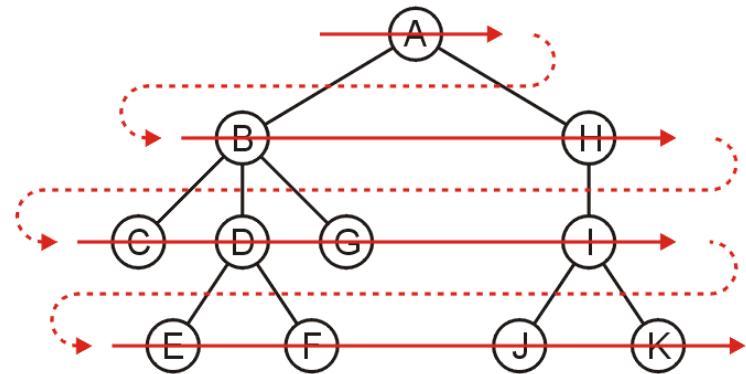
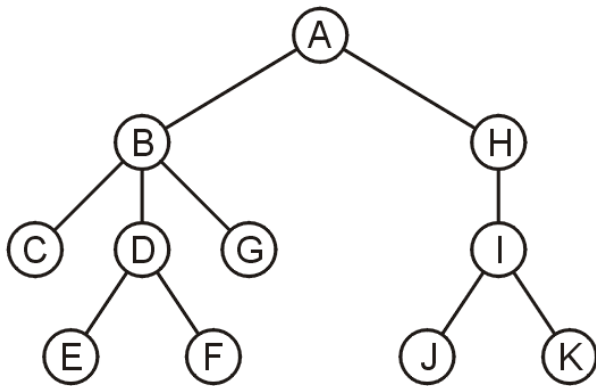
Trees - Ancestors



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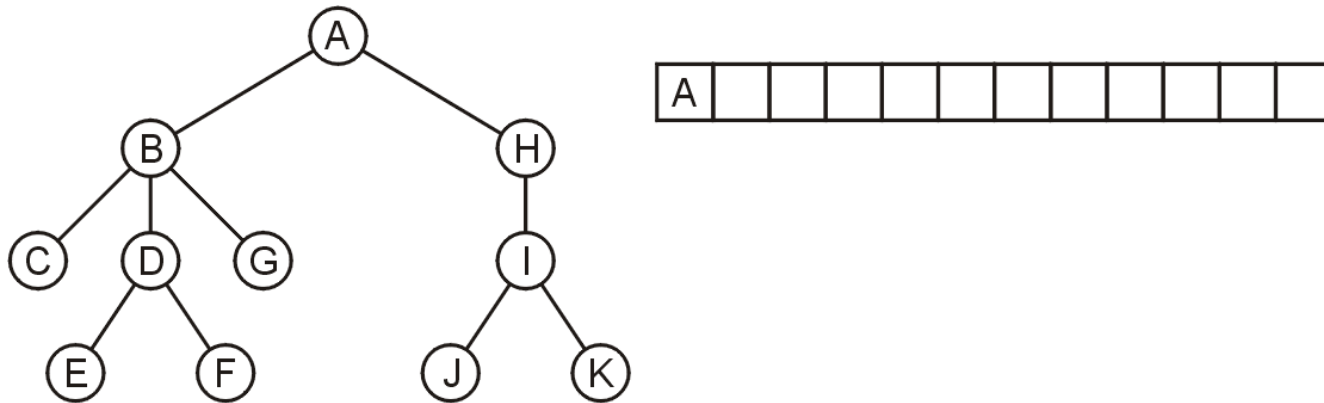
Breadth First Search (BFS)

- Search all sibling nodes at one level before descending a level



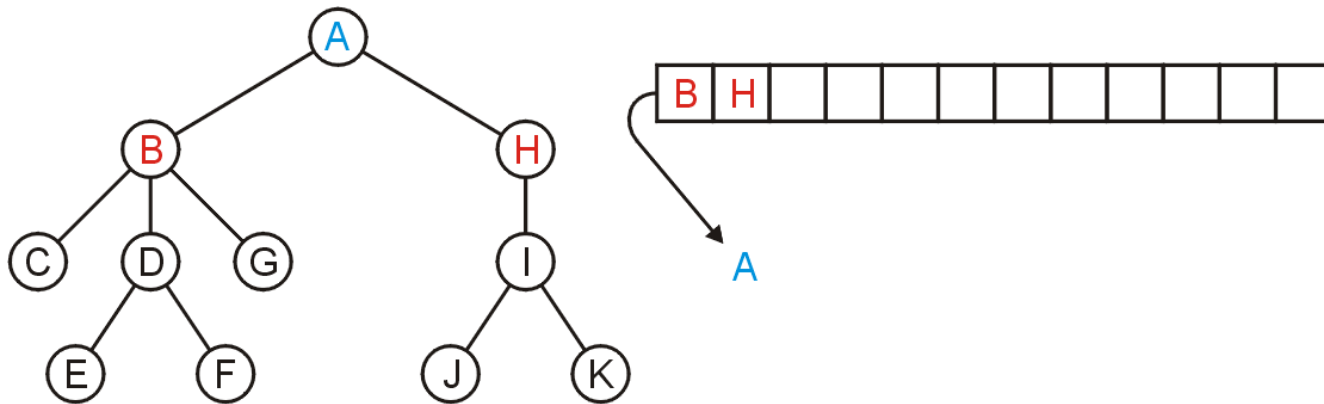
Breadth First Search using Queue

Push the root directory A



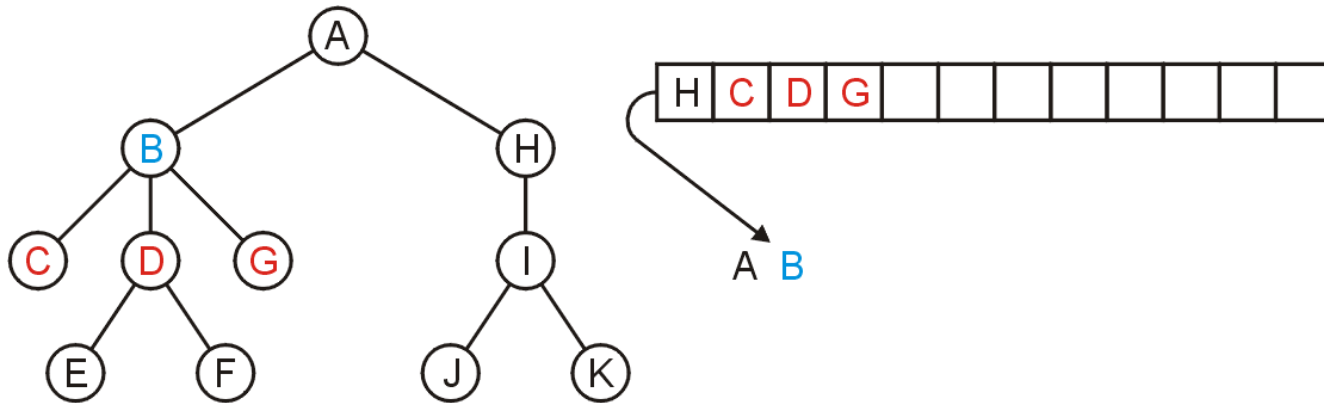
Breadth First Search using Queue

Pop A and push its two children: B and H



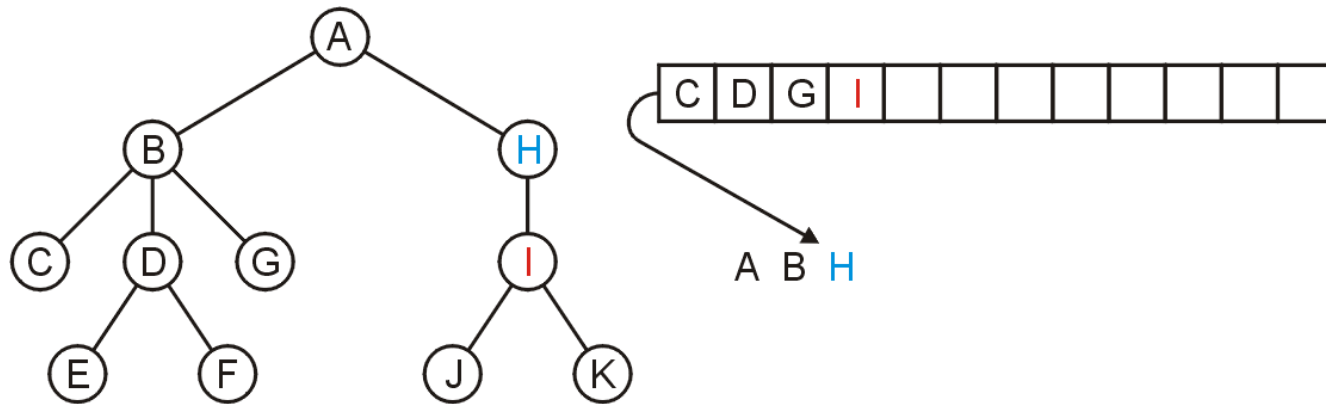
Breadth First Search using Queue

Pop B and push C, D, and G



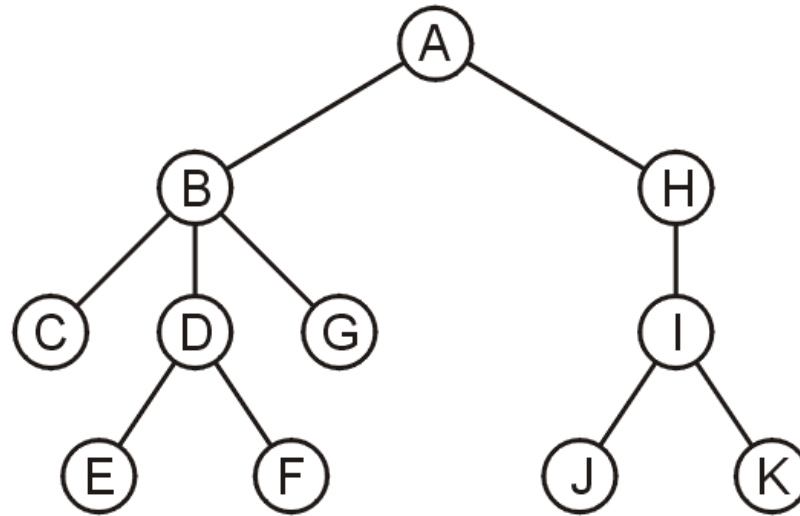
Breadth First Search using Queue

Pop H and push its one children I



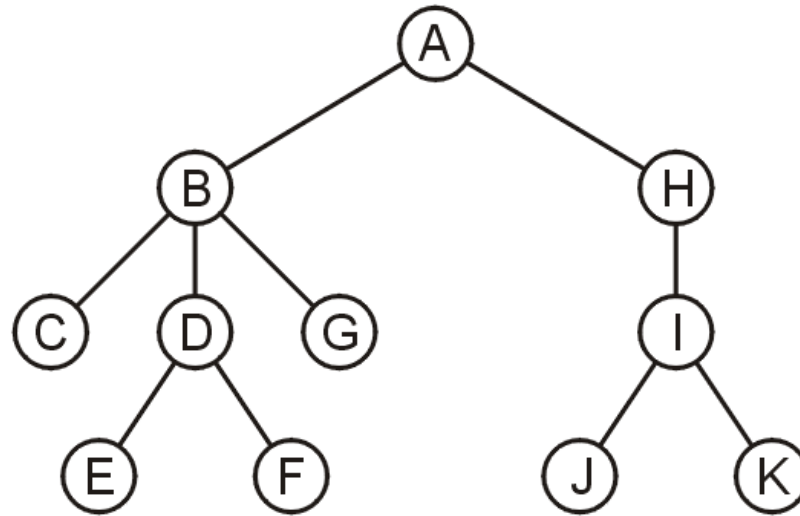
Breadth First Search using Queue

The resulting order: A B H C D G I E F J K is in breadth-first order



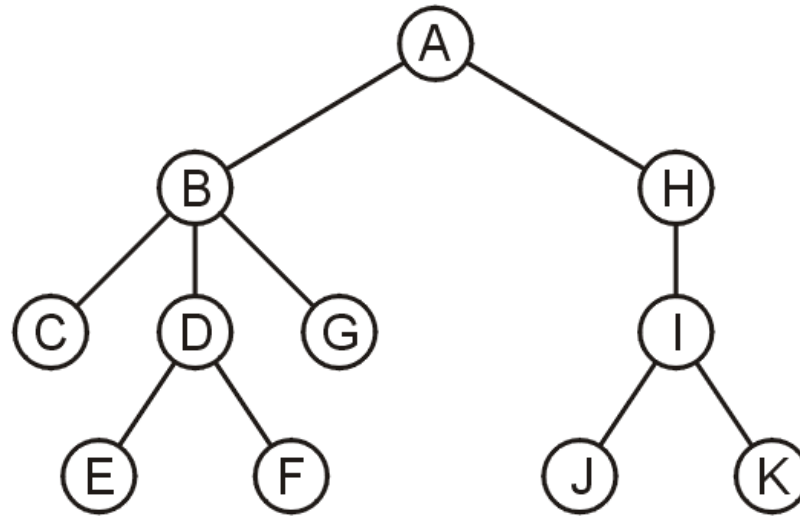
Breadth First Search using Queue

The resulting order: A B H C D G I E F J K is in breadth-first order



Depth First Search using Stack

The resulting order: A B C D E F G H I J K is in depth-first order



Any Question?