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Lecture Pattern Analysis

Part 19: HMMs Algorithms 1 and 2

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Algorithm 1: How Well Matches an Input Sequence the HMM?

- Calculate $p(\mathbf{x}_1, \dots, \mathbf{x}_T | \lambda)$ by marginalizing over all state sequences z_1, \dots, z_T
- Naive marginalization is exponential in the sequence length,

$$p(\mathbf{x}_1, \dots, \mathbf{x}_T) = \sum_{z_1=S_1}^{S_N} \sum_{z_2=S_1}^{S_N} \dots \sum_{z_T=S_1}^{S_N} \pi_{z_1} b_{z_1}(\mathbf{x}_1) a_{z_1 z_2} b_{z_2}(\mathbf{x}_2) \dots a_{z_{T-1} z_T} b_{z_T}(\mathbf{x}_T) , \quad (1)$$

with a computational cost of $O(N^T)$

- However, the chain structure of dependencies admits an efficient dynamic programming solution

$$p(\mathbf{x}_1, \dots, \mathbf{x}_T) = \sum_{z_1=S_1}^{S_N} \pi_{z_1} b_{z_1}(\mathbf{x}_1) \sum_{z_2=S_1}^{S_N} a_{z_1 z_2} b_{z_2}(\mathbf{x}_2) \dots \sum_{z_T=S_1}^{S_N} a_{z_{T-1} z_T} b_{z_T}(\mathbf{x}_T) \quad (2)$$

with caching of the N accumulated state probabilities at time t .

Each state transition requires $O(N^2)$ evaluations, the total cost is $O(N^2 T)$

The Forward Algorithm and the Backward Algorithm

- The forward algorithm directly implements Eqn. 2:

- Initialization (here and below, S_i implicitly creates a list of N entries):

$$t = 1: \quad \alpha_1(z_1 = S_i) = \pi_{z_1} b_{z_1}(\mathbf{x}_1)$$

- Iteration:

$$2 \leq t \leq T: \alpha_t(z_t = S_i) = \left(\sum_{z_{t-1}=S_1}^{S_N} \alpha_{t-1}(z_{t-1}) \cdot a_{z_{t-1}z_t} \right) b_{z_t}(\mathbf{x}_t)$$

- Summation over all end points:

$$t = T: \quad p(\mathbf{x}_1, \dots, \mathbf{x}_T) = \sum_{z_T=S_1}^{S_N} \alpha_T(z_T)$$

- Training requires the (almost identical) backward alg., starting at time T :

- Initialization:

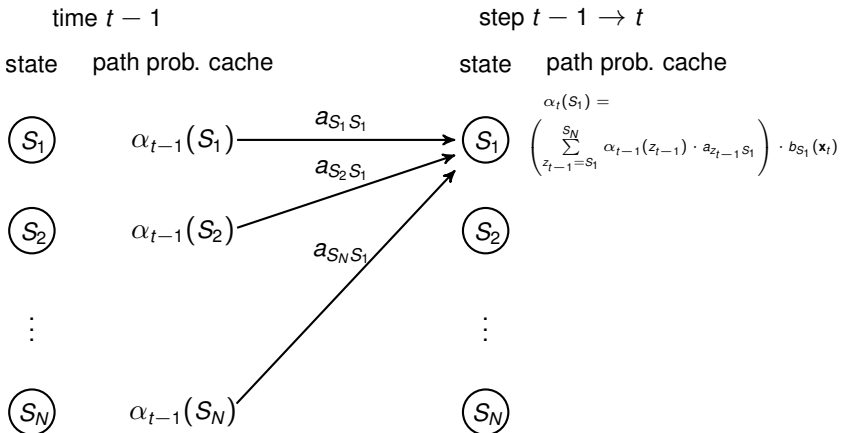
$$t = T: \quad \beta_T(z_1 = S_i) = 1$$

- Iteration: $1 \leq t \leq T: \beta_t(z_t = S_i) = \sum_{z_{t+1}=S_1}^{S_N} \beta_{t+1}(z_{t+1}) b_{z_{t+1}}(\mathbf{x}_{t+1}) a_{z_t z_{t+1}}$

- The backward iteration does not include \mathbf{x}_1 , which is OK for its training use

Forward Algorithm: State-Time Diagram

- This diagram over states (y-axis) and time (x-axis) illustrates the caching:



Algorithm 2: What is the Most Likely State Sequence?

- Finding $\operatorname{argmax}_{z_t=S_j} p(\mathbf{x}_1, \dots, \mathbf{x}_T, z_1, \dots, z_T | \lambda)$ is almost identical to forward alg.
 - Variable $\delta_t(z_i)$ caches the likelihood of the most likely path to z_i within time $1..t$
 - Variable $\psi_t(z_i)$ caches the predecessor state to reconstruct the path
- Viterbi algorithm:
 - Initialization:

$$t = 1: \quad \delta_1(z_1 = S_i) = \pi_{z_1} b_{z_1}(\mathbf{x}_1)$$

$$\psi_1(z_1 = S_i) = 0$$
 - Iteration:

$$1 \leq t \leq T: \delta_t(z_t = S_i) = \max_{S_1 \leq z_{t-1} \leq S_N} (\delta_{t-1}(z_{t-1}) \cdot a_{z_{t-1} z_t}) b_{z_t}(\mathbf{x}_t)$$

$$\psi_t(z_t) = \operatorname{argmax}_{S_1 \leq z_{t-1} \leq S_N} \delta_{t-1}(z_{t-1}) \cdot a_{z_{t-1} z_t}$$
 - Maximum over all end points:

$$t = T: \quad p^* = \max_{S_1 \leq z_T \leq S_N} \delta_T(z_T)$$

$$z_T^* = \operatorname{argmax}_{S_1 \leq z_T \leq S_N} \delta_T(z_T)$$
 - Path reconstruction: $z_t^* = \psi_{t+1}(z_{t+1}^*)$

Remarks

- The computational complexity of the Viterbi algorithm is also $O(N^2 T)$
- Interpretation of the caches $\alpha_t(S_i)$, $\beta_t(S_i)$, $\delta_t(S_i)$:
 - $\alpha_t(S_i)$ is the sum of probabilities of all paths from time 1 to time t that are at time t in state S_i
 - $\beta_t(S_i)$ is the sum of probabilities of all paths from time T to time t that are at time t in state S_i
 - $\delta_t(S_i)$ is the probability of the single, most likely path from time 1 to time t that is at time t in state S_i

We will use these interpretations to explain the training procedure

- The Viterbi algorithm has also many other applications, e.g., as a decoder for partially corrupted digital watermarks (see lecture “Multimedia Security”)