

2 Global Optimization

2.1 An abstract description of the global optimization problem

Let \mathcal{F} be a finite family of segment sequences. The global optimization problem is to find

$$\max_{s \in \mathcal{F}} \max_{\Delta \in \{-1, 0, 1\}^{|s|}} \left(\min_{\gamma \in [-1, 1]^{|s|}} c_{s, \Delta}^T \cdot \gamma \quad \text{subject to} \quad A_{s, \Delta} \cdot \gamma \geq 0 \right)$$

Here, the inner minimization problem is a linear program, where $A_{s, \Delta}$ is a matrix of constraints that depends on s and Δ . Similarly, $c_{s, \Delta}$ is a vector of costs that depends on s and Δ .

2.2 Specifying the inner minimization problem

Here, we specify how $A_{s, \Delta}$ and $c_{s, \Delta}$ are defined. Let $s = s_1 \hookrightarrow \dots \hookrightarrow s_n$ be a sequence of segments with total number of transitions m . Let $\Delta \in \{-1, 0, 1\}^m$ be a vector of input perturbations.

We define some notation as follows:

- Let t_j^i denote the j th transition in segment i , and ε_j^i denote the noise added to the input before that transition.
- Let Δ_j^i denote the entry of Δ that corresponds to the input perturbation for the j th transition in segment i .
- Let γ_j^i denote the entry of $\gamma \in [-1, 1]^m$ that corresponds to the coupling shift for the j th transition in segment i – this is to be determined by the inner minimization problem.

Then, the minimization problem over γ is as follows:

$$\begin{aligned} \min_{\gamma \in [-1, 1]^m} \quad & \sum_{i=1}^n \sum_{j=1}^{|s_i|} |\gamma_j^i - \Delta_j^i| \varepsilon_i \\ \text{subject to} \quad & \gamma_k^i \leq \gamma_0^i && \text{if } t_k^i \text{ has guard } < \\ & \gamma_k^i \geq \gamma_0^i && \text{if } t_k^i \text{ has guard } \geq \\ & \gamma_0^i \leq \gamma_0^k && \text{if } s_k \hookrightarrow s_i \text{ and guard}(s_i) \text{ is } < \\ & \gamma_0^i \geq \gamma_0^k && \text{if } s_k \hookrightarrow s_i \text{ and guard}(s_i) \text{ is } \geq \\ & \gamma_k^i = 0 && \text{if } t_k^i \text{ outputs } \texttt{insample} \\ & \gamma_k^i = \Delta_k^i && \text{if } t_k^i \text{ belongs to a cycle} \end{aligned}$$

This can be rewritten as a linear program using standard techniques, producing a constraint matrix $A_{s, \Delta}$ and a cost vector $c_{s, \Delta}$.

2.3 Suspected ways to simplify the problem

- We might be able to determine the maximizing Δ in the second minimization problem in linear time given $s \in \mathcal{F}$.
 - I suspect this is true since I see that the maximizing Δ always has $\Delta_j^i = -1$ if t_j^i has guard \geq , and $\Delta_j^i = 1$ if t_j^i has guard $<$. In the case that t_j^i has guard *true* and is an assignment transition, the value of Δ_j^i seems to depend on the costs ε_j^i in the segment s_i .
 - If this is true, we need not check exponentially many Δ in the second maximization problem.
- It might be possible to solve a local minimization problem over segments instead of segment sequences, and then use the results to solve a global constraint system that is much smaller than the one described above.