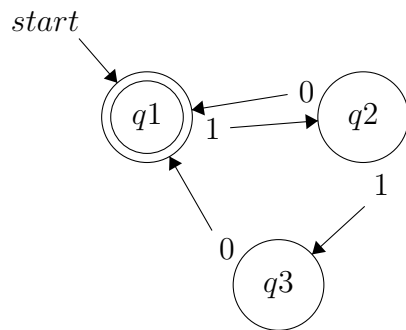


# CSC236, A2

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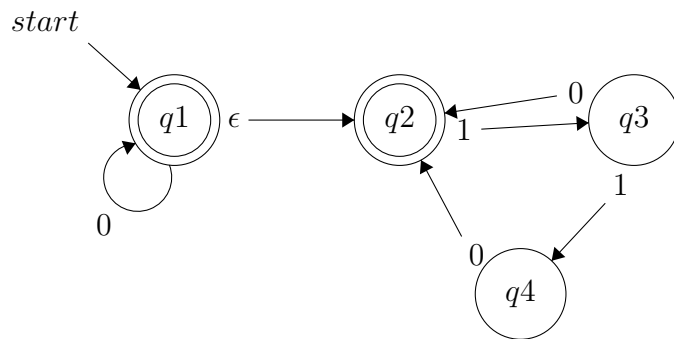
## Q1

(a) i.



Text

ii.



(b) i.  $(b + \epsilon)(ab)^*(a + \epsilon)$

ii.  $a^*ba^*ba^* + a^*ba^*ba^*ba^*$

## Q2

First, by the given fact, we have:

1.  $R^* \subseteq (R + S)^*$  (by the given fact)
2.  $S^* \subseteq (R + S)^*$  (by the given fact)

Then, by the given fact, we have:

$$\begin{aligned}
 R^* S^* &\subseteq (R + S)^* (R + S)^* && \text{(by the given fact)} \\
 &= (R + S)^* && \text{(by the property of the kleene star)} \\
 (R^* S^*)^* &\subseteq ((R + S)^*)^* && \text{(by the given fact)} \\
 (R^* S^*)^* &\subseteq (R + S)^* && \text{(by the given fact)}
 \end{aligned}$$

Secondly, we also have:

1.  $(R^*)^* \subseteq (R^* + S^*)^*$  (by the given fact)
- $R^* \subseteq (R^* + S^*)^*$  (by the given fact)
2.  $(S^*)^* \subseteq (R^* + S^*)^*$  (by the given fact)
- $S^* \subseteq (R^* + S^*)^*$  (by the given fact)

Then, we have:

$$\begin{aligned}
 R^* S^* &\subseteq (R^* + S^*)^* (R^* + S^*)^* && \text{(by the given fact)} \\
 &= (R^* + S^*)^* && \text{(by the property of the kleene star)} \\
 (R^* S^*)^* &\subseteq ((R^* + S^*)^*)^* && \text{(by the given fact)} \\
 (R^* S^*)^* &\subseteq (R^* + S^*)^* && \text{(by the given fact)}
 \end{aligned}$$

Now, note that  $R + S \subseteq R^* S^*$  since  $R \subseteq R^*$  and  $S \subseteq S^*$ .

Also we know that  $R^* + S^* \subseteq R^* S^*$  by the property of union and concatenation operations.

1.  $R + S \subseteq R^* S^*$
- $(R + S)^* \subseteq (R^* S^*)^*$  (by the given fact)
2.  $R^* + S^* \subseteq R^* S^*$
- $(R^* + S^*)^* \subseteq (R^* S^*)^*$  (by the given fact)

Altogether, we have:

$$\begin{aligned}
 (R^* S^*)^* &\subseteq (R + S)^* \\
 (R^* S^*)^* &\subseteq (R^* + S^*)^* \\
 (R + S)^* &\subseteq (R^* S^*)^* \\
 (R^* + S^*)^* &\subseteq (R^* S^*)^*
 \end{aligned}$$

We have  $(R^* S^*)^* \subseteq (R + S)^*$ ,  $(R^* + S^*)^* \subseteq (R^* S^*)^*$ . This means that  $(R + S)^*$  and  $(R^* + S^*)^*$  are equivalent. Therefore,  $(R + S)^* \equiv (R^* + S^*)^*$ . ■

### Q3

The answer is yes.  $Rev(L)$  is regular.

Let there be the NFA,  $M_1$ , for  $L$ .

Then,  $\mathcal{L}(M_1) = L$ .

Let  $M_1 = \langle Q_1, \Sigma, \delta_1, \mathcal{S}_1, \mathcal{F}_1 \rangle$

Let  $M_2$  be the NFA for  $Rev(L)$ .

Then,  $\mathcal{L}(M_2) = Rev(L)$ .

Let  $q$  be a new initial state for  $M_2$ .

Then, the new set of states for  $Rev(L)$  should include this new initial state  $q$  that is linked with an epsilon to the next state.

Since  $Rev(L)$  is reversing the original string in  $L$ , the alphabet of symbols used will still be the same for  $Rev(L)$ .

Let  $w \in \Sigma^*$ .  $w \in Rev(L)$  iff  $w^R \in L$ .

Let  $q', p' \in Q_1$ . In  $L$ , a string  $w$  will be read from a state  $q'$  to a state  $p'$ .

In other words,  $p' \in \delta_1^*(q', w)$ .

Then, for  $Rev(L)$ , we have  $q' \in \delta_r^*(p', w)$  since all the paths should be reversed.

Again, since  $Rev(L)$  is reversing the original string in  $L$ , all the initial states should now become the accepting states. Finally, we have the formal definition for  $Rev(L)$ .

That is,  $M_2 = \langle Q_1 \cup \{q\}, \Sigma, \delta_r, q, \mathcal{S}_1 \rangle$ .

### Q4

Assume for a contradiction that  $L$  is a regular.

Then  $L$  must satisfy PLC. That is, there exists  $p \in \mathbb{N}^+$  such that  $\forall w \in L$ , if  $|w| \geq p$ , then  $\exists x, y, z$  such that the conditions,  $w = xyz$ ,  $|y| \geq 1$ ,  $|xy| \leq p$ , and  $\forall i \geq 0, xy^iz \in L$  hold.

Let  $w = a^{(p+1)}b^p$ .

Then  $w \in L$  and  $|w| \geq p$ .

Assume  $w$  satisfies the first three conditions,  $w = xyz$ ,  $|y| \geq 1$ , and  $|xy| \leq p$ .

That is,  $w = xyz$  such that  $|y| \geq 1$  and  $|xy| \leq p$ .

Since  $w$  starts with  $(p+1)$   $a$ 's and  $|xy| \leq p$ ,  $x$  and  $y$  must only contain  $a$ 's.

Let  $x = a^s$ ,  $y = a^t$ , and  $z = a^r b^p$ , where  $s + t + r = p + 1$  and  $t \geq 1$  (by  $|y| \geq 1$ ).

Consider when  $i = 0$ .

Then  $xy^0z = xz = a^s a^r b^p$  and  $s + r \leq p$  because  $s + r = p - t + 1$  and  $t \geq 1$ .

This means that for  $i = 0$ ,  $xy^iz \notin L$ .

Since the length of  $a$ 's is not greater than the length of  $b$ 's,  $xy^0z \notin L$  and  $L$  does not satisfy PLC.

This is a contradiction. So  $L$  is non-regular. ■