STDIN\_FILENO
File number of stdin. It is 0.
STDOUT\_FILENO
File number of stdout. It is 1.
STDERR\_FILENO
File number of stderr. It is 2.

fchtl.h

Type Definitions

The size\_t, ssize\_t, uid\_t, gid\_t, off\_t and pid\_t types are defined as described in  $\leq sus/types.h \geq$ .

The **useconds\_t** type is defined as described in <a href="mailto:sys/types.h"><a href="mailto:sys/types.h">sys/types.h</a>.</a>.

The intptr\_t type is defined as described in  $\leq inttypes.h \geq$ .

### Declarations

The following are declared as functions and may also be defined as macros. Function prototypes must be provided for use with an ISO C compiler.

```
vaccess(const char *, int);
unsigned int alarm(unsigned int);
             brk(void *);
int
             chder(const char *);
             chcoot(const char *); (LEGACY)
           invo(const char *, uid_t, gid_t);
             configuration, char *, size t);
SIZE 1
            *crypt(const char *, const char *);
char
char
            *stermed(char *);
            *cuserid(char *s); (LEGACY)
             dup(int);
int
             duo2(int, int);
             encrypt(char[64], int);
             execl(const char *, const char *, ...);
int
             execle(const char *, const char *, ...);
int
             execlp(const char *, const char *, ...);
execv(const char *, char *const []);
int
int
             exerve(const char *, char *const [], char *const []);
              execup(const char *, char *const []);
int
              exit(int);
            fcnown(int, uid_t, gid_t);
              fendir(int);
              fastasync(int);
              fork(void);
pid_t
long int
              fpashconf(int, int);
              fsync(int);
              ftruccate(int, off_t);
             *gettwd(char *, size_t);
             getdtablesize(void); (LEGACY)
gid_t
             setruld(void);
           vestsis(void);
              Setscoups(int, gid_t []);
long
              gethostid(void),
             *getlogin(void);
             getlogin r(char *, size_t);
getopt(int, char * const [], const char *);
              getpagesize(void); (LEGACY)
             *getpass(const char *); (LEGACY)
             getpgid(pid_t);
```

```
_pid_t
                getppid(void);
  pid t
                getsid(pid t);
  uid t
                getuid(void);
  char
               *getwd(char *);
  int
                isatty(int);
                lchown(const char *, uid_t, gid_t);
link(const char *, const char *);
  int
  int
                lockf(int, int, off_t);
  off_t
                lseek(int, off_t, int);
 /int
                nice(int);
   long int
                pathconf(const char *, int);
 Int
              pause(void);
   int
                pipe(int [2]);
   ssize_t
                pread(int, void *, size_t, off_t);
  int
                pthread atfork(void (*)(void), void (*)(void),
                    void(*)(void));
  ssize_t
                pwrite(int, const void *, size_t, off_t);
Ssize_t
               read(int, void *, size t);
  int
                readlink(const char *, char *, size_t);
  int
                rmdir(const char *);
  void
                *sbrk(intptr t);
-Int
                setgid(gid t);
 Int
                setpgid(pid_t, pid_t);
  pid_t
                setpgrp(void);
  int
                setregid(gid_t, gid_t);
  int
                setreuid(uid_t, uid_t);
  pid t
                setsid(void);
  int
                setuid(uid t);
unsigned int sleep(unsigned int);
  void
                swab(const void *, void *, ssize_t);
  int
                symlink(const char *, const char *);
  void
                sync(void);
  long int
                 sysconf(int);
  pid t
                regetpgrp(int);
  int
                tcsetpgrp(int, pid_t);
  int
                 truncate(const char *, off_t);
                *ttyname(int);
                ttyname r(int, char *, size_t);
  useconds t
                ualarm(useconds_t, useconds_t);
  int
                unlink(const char *);
  int
                usleep(useconds t);
  pid t
                 vfork(void);
 -ssize t
               write(int, const void *, size_t);
```

The following external variables are declared:

```
extern char *optarg;
extern int optind, opterr, optopt;
```

#### APPLICATION USAGE

None.

### **FUTURE DIRECTIONS**

None.

### SEE ALSO

(access(), alarm(), chdir(), chown(), close(), crypt(), ctermid(), dup(), encrypt(), environ(), exec, exit(), f./hdir(), fchown(), fcutl(), fork(), fpathconf(), fsync(), ftruncate(), getcwd(), getegid(), geteg

- It ends a process and setums a value to its parent. status is an integer between 0 to 255. This number is returned to the parent via went () as the exit status ext() is a library function that calls the system call -exit(), exit() cleans up the standard I/o streams (tempfile()) before Calling \_exit(). Calling - exit() instead of exit() will bypass this cleanup procedure. wait() seturms the exit status multiplied by 256 (Stored in upper 8 bits). The status value is shifted signt 8 bit (divided by 256) to obtain the correct value.

int chown (coast char \*path, wid owner, gid group);

- It sets the owner ID and grup ID of the file specified by path to owner by grup respectively.
- If owner or group ID is specified as -1 or -1 respectively.

  The corresponding ID of file is unchanged.
- It deturns o upon successful completion, otherwise 1 is beturned.
- It feils if searching permission is denied.

int fehown (int fd, vid, gid)!
-fehown () has the same effect as chown () except that the file path and file description parameter.

pid & fork();

Per child's PD

C ← O insufficient
-1 -if Error ⇒ No Resources, memory and CHIDMAX

pid + getpid (); Return pid of calling process. No setum value is reserved to indicate an error.

pid + getppid(): feturn parent process ID of calling Process

-) If called in patent process, it returns shell's procession.

int nice (int incr): feturn new inice value, - 1 in error

- adds the value of incr to the rice value of the Calling process.
  - I n'ice value is a non-negative number.
- minimum nice value of O (3ero) are imposed by system
- No effect on SCHED\_FIFO BSCHED\_RR
- If the process is multi-threeded, the nice value affects all system scope threads in the process.
- more positive value results an a lower cpu priority

unsigned int sleep (unsigned int seconds);

- Returns the seguisted time has elapsed, the value setward will be O. - No errors are defined.

& Pid value will always be virgue. Once the process finish execution its pid value is returned to the kernel which my

→ Pid value Can't be changed when arigned during execution.

- \* fork() duplicates the variousle which are before fork(). It also duplicated even the Global variable (declared before mains)
- \* fork() simply shares the files instead of duplicates.
- \* globally opened file tases the same discriptor. All processes use that one descriptor.

int pause (void):

It suspends the execution of the calling process untill it receives a signal.

- There is no successful completion setum value. 1 is defuned in case of error.
- It Sail if a signal is caught by the calling process and control is returned from the rignal-catching function.

ent sead (fd, buffer pointer, int nyte)!
int write (fd, buffer pointer, int nyte)!

- fd identifies the I/o channel.
- buffer pointer points to the area in memory where the data is stored for a read () or where the data is taken for a write .
- nbyte defines the maximum number of characters or bytes transfered between the file and the buffer.
- seadub writer setum the number of bytes transfered.
- There is no limit on royse rize or depends on SSIZE-MAX. (implementation-dependent).
- A neyte of 1 is used to toquefer a byte at a time
- nost efficient value for nbyte is the size of the largest phyrical second the I/o dannel is likely to have to have

- read() D write() setum a non-negative integer |
  Indicating the number of byte actually read or we cre
   In case of failure, setum 1 b set earno to ingrant
  exros.
  - int close(int fd);
  - It will deallocate the fd. The deallocated fd will be availed. For further use for open () b creat (). All the autobanding becord locks owned by the process on the file associated with fd will also be semoved.
- It setums an successful completion. Otherwise, it setums I and set earns to indicate the error.
  - a Int access (const char's path, int amode);
- It checks the file pointed by path argument is accessible according to the amode, wing realwer ID and real group ID.
- The amode constants are defined in systfile b

F-OK existence X-OK execute W-OK write R-OK read

Ored to check more than one access permession.

- It returns 0, if requested access is permitted. Otherwise returns -1 to set the error.
- access() only answer the question "do I have this permission." It can't answer the question "what permission do I have

pid. I setsid (void)

regroup leader. Upon return the calling process will be the servicen leader of the new session, will be the process your leader of a new process group, and will have no controlling-terminal. The process group ID of the calling process will be set equal to the process ID of the colling process.

The calling process will be the only process in the new process grup and the only process in the new session.

-) It octums the value of the process group ID of the calling process on successful completion. Otherwise setuons - 1 & set ermo to indicate the error.

- It fails if the calling process is already a process group leader.

int setpgid (pid.t pid, pid.t pgid)

- It is used either to join an existing process grup or coeate a new process group within the session of the calling process. The process group ID of a session leader will not change upon successful completion, the process group ID of the process with a process ID that matches pid will be set to paid.

As a special case, if pid is 0, the process ID of the calling process will be used. Also, if poid is 0, the process group ID of the

indicated process will be used,

upon successful completion, it returns 0, otherwise - 1 is setumed and eomo is set, set pgsp() is exhivalent to setpgid (0,0)

# pidt galpgid (pid t pid)

- It setums the process group ID of the process whose Is process ID is equal to pid.
- If pid is equal to O(3ers), it returns the process group to of the calling process.
- It is fails, it setures as set the exmo to indicate the error.
  - geologid ) fails if the process whose process ID is equal to pid is not in the same session as the calling process, and the implementation does not allow access to the process grupID of There is no process with a process ID equal to pid.

# Pid. + getpgrp(void)

- . It deturns the process grup ID of the calling process.
- It always successful and no return value is peserved to indicate an
- + getpgrk() is equivalent to getpgid(o).
  - Each process group is a member of a session and each process is a member of the serian of which its process grup is a member.

# Rid t setpgrp (void);

If the calling process is not already a sersion leader, setpgops) sets the process group ID of the calling process to the process ID of the calling process. If setpgrp() creates a new serion, then the new servion has no controlling terminal.

setpgrp() has no effect when the calling process is a servicen

No any error is defined

upon Completion, it returns the process group ID.

## ent kill (Pd-t pid, int rig)

- -> It will send a rignal to a process or a group of processes specified by pid. The signal to be sent is specified by sig and is either one from the list given in Lingual. h) or O.
- + If sog is O (the null signal), error cheeking is performed but no signal is actual sent. The null signal can be used to check the validity of pid.

+ If pid is greater than 0, mg will be sent to the process

whose process ID is equal to pid.

- If pid is 0, sog will be sent to all processes whose process group ID is equal to the process group ID of the sender, and for which the process has permission to send

a rignal.

The pid is -1, rig will be sent to all processes for which the process has permission to send that rignal.

- processed whose process group ID is equal to the absolute value of pid, and for which the process has permission; send a signal.
- the processes specified by pid.
- If kall () fails, no rignal will be sent, and returns -1, and set erro.
- unsupported signal number. The process does not have permission to send the signal to any receiving process.

  3 No process or process group can be found corresponding to that specified by pid.

### #include <unistd.h>

- > The arguments specified by a program with one of the exec functions are passed on to the new process image in the corresponding main() arguments.
- > The argument path points to a pathname that identifies the new process image file.
- The exec functions replace the current process image with a new process image. The new image is constructed from a regular, executable file called the new process image file. There is no return from a successful exec, because the calling process image is overlaid by the new process image.
- > When a C-language program is executed as a result of this call, it is entered as a C-language function call as follows:

```
int main (int argc, char *argv[]);
```

where argc is the argument count and argv is an array of character pointers to the arguments themselves.

In addition, the environ variable:

```
extern char **environ;
```

Penviron variable is initialized as a pointer to an array of character pointers to the environment strings.

The argv and environ arrays are each terminated by a <u>null pointer</u>. The <u>null pointer</u> terminating the argv array is not counted in argc.

Strigs

- The argument file is used to construct a pathname that identifies the new process image file. If the file argument contains a slash character, the file argument is used as the pathname for this file. Otherwise, the path prefix for this file is obtained by a search of the directories passed as the environment variable. If this environment variable is not present, the results of the search are implementation-dependent.
- ➤ If the process image file is not a valid executable object, execlp() and execvp() use the contents of that file as standard input to a command interpreter conforming to system(). In this case, the command interpreter becomes the new process image.
- These strings constitute the argument list available to the new process image. The list is terminated by a null pointer. The argument arg0 should point to a filename that is associated with the process being started by one of the exec functions.
- The argument argv is an array of character pointers to null-terminated strings. The last member of this array must be a null pointer. These strings constitute the argument list available to the new process image. The value in argv[0] should point to a filename that is associated with the process being started by one of the exec functions.
- The argument envp is an array of character pointers to null-terminated strings. These strings constitute the environment for the new process image. The envp array is terminated by a null pointer.
- For those forms not containing an envp pointer (eg. execl , execv(), execlp() and execvp()), the environment for the new process image is taken from the external variable environ in the calling process.
- The number of bytes available for the new process' combined argument and environment lists is {ARG\_MAX}. It is implementation-dependent whether null terminators, pointers, and/or any alignment bytes are included in this total.

- File descriptors open in the calling process image remain open in the new process image, except for those whose close-on-exec flag FD\_CLOEXEC is set. For those file descriptors that remain open, all attributes of the open file description, including file locks remain unchanged.
- Directory streams open in the calling process image are closed in the new process image.
- > RETURN VALUE: If one of the exec functions returns to the calling process image, an error has occurred; the return value is -1, and error is set to indicate the error.
- > The new process also inherits at least the following attributes from the calling process image:
  - ✓ nice value (see nice())
  - ✓ semadj values (see semop())
  - ✓ process ID
  - ✓ parent process ID
  - ✓ process group ID
  - ✓ session membership
  - ✓ real user ID
  - ✓ real group ID
  - √ supplementary group IDs
  - ✓ time left until an alarm clock signal (see alarm())
  - ✓ current working directory
  - ✓ root directory
  - ✓ file mode creation mask (see umask())
  - ✓ file size limit (see ulimit())
  - ✓ process signal mask (see sigprocmask())
  - ✓ pending signal (see sigpending())
  - ✓ tms\_utime, tms\_stime, tms\_cutime, and tms\_cstime (see times())
  - ✓ resource limits

- ✓ controlling terminal
- ✓ interval timers
- All other process attributes defined in this document will be the same in the new and old process images. The inheritance of process attributes not defined by this specification is implementation-dependent.

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <process.h>
int main(int argc, char* argv[]) {
   char* arr[] = {
       "ping",
       "google.com",
       NULL
   };
   char* env[] = {
       "TEST=environment variables",
       NULL
   };
    // Without env and PATH variable
   vexecv("C:\\Windows\\System32\\ping.exe",
         "C:\\Windows\\System32\\ping.exe", "google.com", NULL);
   execl("C:\\Windows\\System32\\ping", arr);
   // Without env but with PATH variable
  execvp("ping", "ping", "google.com", NULL);
 execlp("ping", arr);
   // With env and PATH variable
    execvpe("ping", "ping", "google.com", NULL, env);
    execlpe("ping", arr, env);
   return 0;
```

- \* pid-t waitpid (pid-t pid, int + stat-loc, int options)
  - > Arg. pid. + pid is greater than O, It specify the process ID of a ringle child process whose status is
- If pid is O, status is sequised for any child process whose process group ID is equal to that of the calling process.
- > If Pid is -1, Status is requested for any child process. Then waitpid () is equivalent to wait ().

Options: WCONTINUED (Ask status)
waitpoid() report the status of any continued child process. WNOHANG (No More weig)

waitpid() shall not suspend the execution of the calling thread if status is not immediately available for child process.

### WUNTRACED

child process is stopped but his status is not yet reported. These status seported to the sequesting process.

Zombie: - A process that has finished the execution but still has entry in the process table example - parent is in suspended or sleep state & child exits.

Orphan: A process whose parent process no more exists eg.

Finished or terminated without weiting for its

Child process to terminate. eg. farent finished but

Child still executing

sys/wait.h Synchronization wait b waitpid - wait for a child process to stop or It forces the parent to suspend execution until the child's fruished

wait setums the process ID of a child process that finished.

Statt-loc is a printer to an integer where value. seturned of wait is stored.

= wait () fails by two complitions:

1 The process has no children to wait for

2 stat-loc points to an invalid address.

I Ended by exit().

2nd Lowest byte is set to the argument by exit(). Lowest byk is set to Bens

2. Ended by signal.

Lowest byte is set to signal number that ended the process. 2nd Lowest byte is set to zero