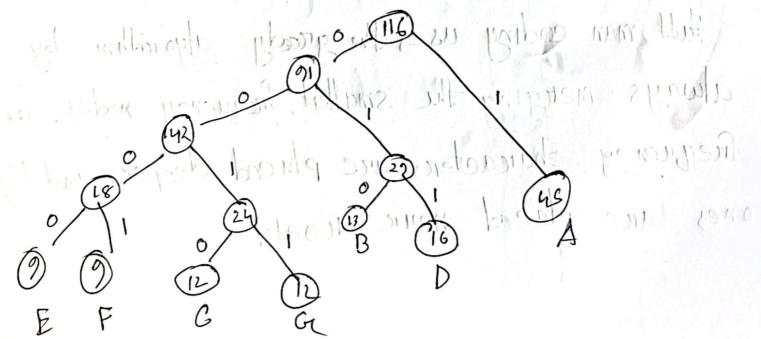
Med 81	- (01) Ans 1)	Plans	congs :
Chan	Freq	code	and an end
A	45	1203	X45=45
В	1.33 Joseph de	0/100 000	3×13 = 39
С	12	0010	4 *12= 48
D	16	0 11	3 x 16 = 4 8
E	2	0000	4×2=36
F	2	0001	489=376
or	12	0 0 11	4 × 12 = 48
	4116	190	£ 300 bits

3 bit encoding (116×3) - 348 bits



saved (348-300) = 48 bits : Space shos lans. Pans 11815=16 Ans:b for each character code 8 HA = 12 H 3=010 C=0010 D=011 E = 0000 F= 000 1 31131 G=0011

Ansi Que - (E > 311). Bribooms tid &

And man coding uses the grady algorithm by always menging the smallest frequency nodes low frequency characters are placed deeper and his ones are placed near root.

This approach ensures that the tree is built
bottom up priorrotizing the combination of the
least frequency symbol.
2 RXV 1 Ams: D
Both of them will generate the same output
and will generate same total length. Because
huffman coding optimally depends only on freque
Among them the first friends approach is the beg
Ans: 2
Merged characters chabaches
Chanader (iii) Frequency
2000 = 0
b 5
4 1000
Ç 13

ZINI Hind en sont It & Alle: 2 dis man Monon gan to charman frequisifforceden mothed 20 mp 10000 13x2 5 0.1 tighto ones who will generalle du same output actions of April 1 Same Some Sold forgin Because How I them the direct Intends approach is the bost. / Kins / Morelly chair 15222 Chandelen (iii) 2: end replicated a= 000 2=001

ans: 2(1)

Freoded mag: éb abdéb bádébé:

hall

0110001006101161 0 11000 100 10111000000 10 1101

ons: 2(v)

Fore encoded msg = 25 bits

for characters-(4x8) = 32 bits How he represents the hold

code: 9 bits.

Mail pet the 1/2 Place

charles .