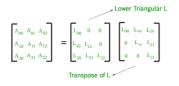
Algorithms and Data Structures

Matrix Decomposition Cholesky Decomposition





Learning goals

- Cholesky decomposition
- Properties of Cholesky decomposition

Aim: Solve LES of the form $\mathbf{A}\mathbf{x} = \mathbf{b}$

with $\mathbf{A} \in \mathbb{R}^{n \times n}$, \mathbf{A} positive-definite

- Write **A** as $\mathbf{A} = \mathbf{L} \mathbf{L}^{\top}$
- **2** Solve $\mathbf{L}\mathbf{y} = \mathbf{b}$ by forward substitution
- **3** Solve $\mathbf{L}^{\top}\mathbf{x} = \mathbf{y}$ by back substitution



Example: Let $\mathbf{A}\mathbf{x} = \mathbf{b}$ be a LES

$$\begin{pmatrix} 4 & 2 & 2 & 2 \\ 2 & 5 & 3 & 3 \\ 2 & 3 & 11 & 5 \\ 2 & 3 & 5 & 19 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} 22 \\ 33 \\ 61 \\ 99 \end{pmatrix}$$



• Write **A** as $\mathbf{A} = \mathbf{L} \mathbf{L}^{\top}$

$$\begin{pmatrix} l_{11} & \mathbf{0} & \mathbf{0} & \mathbf{0} \\ l_{21} & l_{22} & \mathbf{0} & \mathbf{0} \\ l_{31} & l_{32} & l_{33} & \mathbf{0} \\ l_{41} & l_{42} & l_{43} & l_{44} \end{pmatrix} \begin{pmatrix} l_{11} & l_{21} & l_{31} & l_{41} \\ \mathbf{0} & l_{22} & l_{32} & l_{42} \\ \mathbf{0} & \mathbf{0} & l_{33} & l_{43} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & l_{44} \end{pmatrix} = \begin{pmatrix} 4 & 2 & 2 & 2 \\ 2 & 5 & 3 & 3 \\ 2 & 3 & 11 & 5 \\ 2 & 3 & 5 & 19 \end{pmatrix}$$

$$I_{11}^2 = a_{11} \quad \rightarrow \quad I_{11} = \sqrt{a_{11}} = \sqrt{4} = 2$$



1 Write **A** as $\mathbf{A} = \mathbf{L} \mathbf{L}^{\top}$

$$\begin{pmatrix} l_{11} & 0 & 0 & 0 \\ l_{21} & l_{22} & \mathbf{0} & \mathbf{0} \\ l_{31} & l_{32} & l_{33} & 0 \\ l_{41} & l_{42} & l_{43} & l_{44} \end{pmatrix} \begin{pmatrix} l_{11} & l_{21} & l_{31} & l_{41} \\ \mathbf{0} & l_{22} & l_{32} & l_{42} \\ \mathbf{0} & 0 & l_{33} & l_{43} \\ \mathbf{0} & 0 & 0 & l_{44} \end{pmatrix} = \begin{pmatrix} 4 & 2 & 2 & 2 \\ 2 & 5 & 3 & 3 \\ 2 & 3 & 11 & 5 \\ 2 & 3 & 5 & 19 \end{pmatrix}$$

$$l_{11}^2 = a_{11} \rightarrow l_{11} = \sqrt{a_{11}} = \sqrt{4} = 2$$

 $l_{21} \cdot l_{11} = a_{21} \rightarrow l_{21} = \frac{a_{21}}{l_{11}} = \frac{2}{2} = 2$



• Write **A** as $\mathbf{A} = \mathbf{L} \mathbf{L}^{\top}$

$$\begin{pmatrix} l_{11} & 0 & 0 & 0 \\ l_{21} & l_{22} & \mathbf{0} & \mathbf{0} \\ l_{31} & l_{32} & l_{33} & 0 \\ l_{41} & l_{42} & l_{43} & l_{44} \end{pmatrix} \begin{pmatrix} l_{11} & l_{21} & l_{31} & l_{41} \\ 0 & l_{22} & l_{32} & l_{42} \\ 0 & 0 & l_{33} & l_{43} \\ 0 & 0 & 0 & l_{44} \end{pmatrix} = \begin{pmatrix} 4 & 2 & 2 & 2 \\ 2 & 5 & 3 & 3 \\ 2 & 3 & 11 & 5 \\ 2 & 3 & 5 & 19 \end{pmatrix}$$

$$\begin{aligned} & l_{11}^2 = a_{11} & \to & l_{11} = \sqrt{a_{11}} = \sqrt{4} = 2 \\ & l_{21} \cdot l_{11} = a_{21} & \to & l_{21} = \frac{a_{21}}{l_{11}} = \frac{2}{2} = 2 \\ & l_{22}^2 + l_{21}^2 = a_{22} & \to & l_{22} = \sqrt{a_{22} - l_{21}^2} = \sqrt{5 - 1^2} = 2 \end{aligned}$$



• Write **A** as $\mathbf{A} = \mathbf{L}\mathbf{L}^{\top}$

$$\begin{pmatrix} l_{11} & 0 & 0 & 0 \\ l_{21} & l_{22} & 0 & 0 \\ l_{31} & l_{32} & l_{33} & 0 \\ l_{41} & l_{42} & l_{43} & l_{44} \end{pmatrix} \begin{pmatrix} l_{11} & l_{21} & l_{31} & l_{41} \\ 0 & l_{22} & l_{32} & l_{42} \\ 0 & 0 & l_{33} & l_{43} \\ 0 & 0 & 0 & l_{44} \end{pmatrix} = \begin{pmatrix} 4 & 2 & 2 & 2 \\ 2 & 5 & 3 & 3 \\ 2 & 3 & 11 & 5 \\ 2 & 3 & 5 & 19 \end{pmatrix}$$

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General formula:
$$l_{jj} = \left(a_{jj} - \sum_{k=1}^{j-1} l_{jk}^2\right)^{\frac{1}{2}} \quad l_{ij} = \frac{1}{l_{jj}} \left(a_{ij} - \sum_{k=1}^{j-1} l_{ik} l_{jk}\right)$$



2 Solve Ly = b by forward substitution

$$\begin{pmatrix} 2 & 0 & 0 & 0 \\ 1 & 2 & 0 & 0 \\ 1 & 1 & 3 & 0 \\ 1 & 1 & 1 & 4 \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \\ y_3 \\ y_4 \end{pmatrix} = \begin{pmatrix} 22 \\ 33 \\ 61 \\ 99 \end{pmatrix}$$

$$\begin{pmatrix} 2y_1 \\ y_1 + 2y_2 \\ y_1 + y_2 + 3y_3 \\ y_1 + y_2 + y_2 + 4y_4 \end{pmatrix} = \begin{pmatrix} 22 \\ 33 \\ 61 \\ 99 \end{pmatrix}$$

$$\Rightarrow$$
 $v_1 = 11, v_2 = 11, v_3 = 13, v_4 = 16$



3 Solve $\mathbf{L}^{\top}\mathbf{x} = \mathbf{y}$ by back substitution

$$\begin{pmatrix} 2 & 1 & 1 & 1 \\ 0 & 2 & 1 & 1 \\ 0 & 0 & 3 & 1 \\ 0 & 0 & 0 & 4 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} 11 \\ 11 \\ 13 \\ 16 \end{pmatrix}$$

$$\Rightarrow x_4 = 4, x_3 = 3, x_2 = 2, x_1 = 1$$

Calculation of the lower triangular matrix (L):

$$\begin{pmatrix} l_{11} & 0 & 0 & 0 \\ l_{21} & l_{22} & 0 & 0 \\ l_{31} & l_{32} & l_{33} & 0 \\ l_{41} & l_{42} & l_{43} & l_{44} \end{pmatrix} \begin{pmatrix} l_{11} & l_{21} & l_{31} & l_{41} \\ 0 & l_{22} & l_{32} & l_{42} \\ 0 & 0 & l_{33} & l_{43} \\ 0 & 0 & 0 & l_{44} \end{pmatrix} = \begin{pmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \\ a_{41} & a_{42} & a_{43} & a_{44} \end{pmatrix}$$



Thus the entries of L (j rows, i columns) result from

$$I_{ij} = \begin{cases} 0 & \text{for } i < j \\ (a_{jj} - \sum_{k=1}^{j-1} I_{jk}^2)^{\frac{1}{2}} & \text{for } i = j \\ \frac{1}{I_{ji}} (a_{ij} - \sum_{k=1}^{j-1} I_{ik} I_{jk}) & \text{for } i > j \end{cases}$$

Important: Order of calculation (row by row) matters!

$$\rightarrow$$
 I_{11} , I_{21} , I_{22} , I_{31} , I_{32} , I_{33} ,..., I_{nn}

Algorithm Cholesky decomposition

- 1: **for** j = 1 to n **do**
- 2: $l_{jj} = \left(a_{jj} \sum_{k=1}^{j-1} l_{jk}^2\right)^{\frac{1}{2}}$
- 3: **for** i = j + 1 to n **do**
- 4: $I_{ij} = \frac{1}{I_{ji}} \left(a_{ij} \sum_{k=1}^{j-1} I_{ik} I_{jk} \right)$
- 5: end for
- 6: end for

If we consider only the (dominant) multiplications, we count in each step of the outer loop

- For diagonal elements: (j-1) multiplications
- For non-diagonal elements: (n-j)(j-1) multiplications



In total, we estimate the computational effort with

$$\sum_{j=1}^{n} [(j-1) + (n-j)(j-1)]$$

$$= \sum_{j=1}^{n} [j-1+nj-n-j^2+j] = \sum_{j=1}^{n} [(n+2)j-1-j^2]$$

$$= n \frac{(n+2)(n+1)}{2} - n - n \frac{(n+1)(2n+1)}{6}$$

$$= n \cdot \frac{3(n+2)(n+1) - 6 - (n+1)(2n+1)}{6}$$

$$= n \cdot \frac{3n^2 + 9n + 6 - 6 - 2n^2 - 2n - n - 1}{6}$$

$$\approx \frac{1}{6}n^3 + \mathcal{O}(n^2) \text{ for large } n$$



PROPERTIES OF CHOLESKY DECOMPOSITION

- Most important procedure for positive-definite matrices
- Algorithm is always stable (no pivoting necessary)
- Existence and uniqueness: The Cholesky decomposition exists and is unique for a positive-definite matrix A
- Runtime behavior:
 - Decomposition of the matrix: $\frac{n^3}{6} + \mathcal{O}(n^2)$ multiplications
 - Forward and back substitution: n²



PROPERTIES OF CHOLESKY DECOMPOSITION / 2

```
cholesky = function(a) {
 n = nrow(a)
  l = matrix(0, nrow = n, ncol = n)
  for (j in 1:n) {
   l[j, j] = (a[j, j] - sum(l[j, 1:(j - 1)]^2))^0.5
   if (j < n) {
     for (i in (j + 1):n) {
       1[i, j] = (a[i, j] -
         sum(1[i, 1:(j-1)] * 1[j, 1:(j-1)])) / 1[j, j]
 return(1)
```



PROPERTIES OF CHOLESKY DECOMPOSITION / 3

```
A = crossprod(matrix(runif(16), 4, 4))
cholesky(A)

t(chol(A))

A = crossprod(matrix(runif(1e+06), 1e+03, 1e+03))
system.time(cholesky(A))
system.time(chol(A))
```



APPLICATION EX.: MULTIVARIATE GAUSSIAN

Target: Efficient evaluation of the density of a normal distribution.

The density of the *d*-dimensional multivariate normal distribution is

$$f(\mathbf{x}) = \frac{1}{(2\pi)^{\frac{d}{2}} |\Sigma|^{\frac{1}{2}}} \exp\{-\frac{1}{2} (\mathbf{x} - \mu)^{\top} \Sigma^{-1} (\mathbf{x} - \mu)\}$$

with $\mathbf{x} \in \mathbb{R}^d$, $\mathsf{Cov}(\mathbf{x}) = \mathbf{\Sigma}, \mathbf{\Sigma}$ positive-definite.

With $\mathbf{z} = \mathbf{x} - \boldsymbol{\mu}, \mathbf{z} \in \mathbb{R}^d$ we obtain:

$$(\mathbf{x} - \boldsymbol{\mu})^{ op} \mathbf{\Sigma}^{-1} (\mathbf{x} - \boldsymbol{\mu}) = \mathbf{z}^{ op} \mathbf{\Sigma}^{-1} \mathbf{z}$$

Problem: Calculation of Σ^{-1} is numerically unstable and requires a long time.



APPLICATION EX.: MULTIVARIATE GAUSSIAN / 2

Solution: Use Cholesky decomposition to avoid inverting Σ^{-1} . Write Σ as $\Sigma = \mathbf{L}\mathbf{L}^{\top}$, $\mathit{rank}(\mathbf{L}) = d$.

Thus it holds:

$$\mathbf{z}^{\top} \mathbf{\Sigma}^{-1} \mathbf{z} = \mathbf{z}^{\top} (\mathbf{L} \mathbf{L}^{\top})^{-1} \mathbf{z}$$

 $= \mathbf{z}^{\top} (\mathbf{L}^{\top})^{-1} \mathbf{L}^{-1} \mathbf{z}$
 $= (\mathbf{L}^{-1} \mathbf{z})^{\top} \mathbf{L}^{-1} \mathbf{z}$
 $= \mathbf{v}^{\top} \mathbf{v}$

with $\mathbf{v} = \mathbf{L}^{-1}\mathbf{z}$, $\mathbf{v} \in \mathbb{R}^d$.

To avoid inverting ${\bf L}$ we can calculate ${\bf v}$ as a solution of the LES

$$\mathbf{L}\mathbf{v}=\mathbf{z}$$

Then we can calculate $\mathbf{v}^T \mathbf{v}$ as a scalar product of two *d*-dimensional vectors.

