

Solution 1:

The code for this programming exercise can be found in the `hw_shapley_py.sol.ipynb` or `hw_shapley_vanilla_py.sol.ipynb` files for Python and in the `hw_shapley_R.sol.Rmd` or `hw_shapley_vanilla_R.sol.Rmd` files for R.

a)

$$\begin{aligned}\text{payoff}(S) &= 10t + 10m + 10s + 2j + 20(t \wedge m) + 20(t \wedge m \wedge s) - 30((t \vee m \vee s) \wedge j) \\ \text{payoff}(\{t, m\}) &= 10 + 10 + 20 = 40 \\ \text{payoff}(\{t, j, s\}) &= 10 + 10 + 2 - 30 = -8\end{aligned}$$

b) Pseudocode of `payoff(coalition)`:

Algorithm 1 `payoff()`

Require: `coalition`: Coalition vector (`set` or `list` or ...)

```

1: t ← boolean if 't' is in coalition
2: s ← boolean if 's' is in coalition
3: m ← boolean if 'm' is in coalition
4: j ← boolean if 'j' is in coalition
5: l ← boolean if 'l' is in coalition
6: return  $10 * t + 10 * m + 2 * j + 20 * (t \text{ and } m) + 20 * (t \text{ and } m \text{ and } s) - 30 * ((t \text{ or } m \text{ or } s) \text{ and } j)$ 
```

$$\begin{aligned}\text{payoff}(\{\emptyset\}) &= 0 \\ \text{payoff}(\{t, m, s, j, l\}) &= 10 + 10 + 10 + 2 + 20 + 20 - 30 = 42\end{aligned}$$

Concerning `all_unique_subsets(population)`, both R and Python provide built-in functions that return the power set, i.e. the set of all subsets of a given set. Alternatively, one possible pseudocode of `all_unique_subsets(population)` would be:

Algorithm 2 `all_unique_subsets()`

Require: `population`: Vector / list / set containing all available players

```

1: if population =  $\emptyset$  then subsets  $\leftarrow \emptyset$ 
2: else if population  $\neq \emptyset$  then
3:   first_member  $\leftarrow$  pick any element from population
4:   population_wo_member  $\leftarrow$  population \ first_member
5:   subsets_wo_member  $\leftarrow$  all_unique_subsets(population_wo_member)
6:   subsets_w_member  $\leftarrow$  list of all sets in subsets_wo_member each with first_member added to it
7:   subsets  $\leftarrow$  Union(subsets_w_member, subsets_wo_member)
8: end if
9: return subsets
```

Pseudocode of `shapley_set(member, population, v_function)`:

Algorithm 3 `shapley_set()`

Require: `member`: individual player, or feature of interest
Require: `population`: vector containing all available players
Require: `v_function`: Some value function

```
1: remainder ← everyone from the population but member (or: population \ member)
2: all_sets_wo_member ← all_unique_subsets(remainder)
3: F ← length of population
4: result ← 0
5: for coalition in all_sets_wo_member do
6:   S ← length of coalition
7:   diff ← v_function(Union(coalition, member)) - v_function(coalition)
8:   factor ← S! * (F - S - 1)! / F!
9:   result ← result + factor * diff
10: end for
11: return result
```

- c) A function to generate all permutations can be implemented similar to the function `all_unique_subsets()`, also by first generating the first element and then calling the function on the set of remaining elements.

Pseudocode of `shapley_perm(member, population, v_function)`:

Algorithm 4 `shapley_perm()`

Require: `member`: individual player, or feature of interest
Require: `population`: vector containing all available players
Require: `v_function`: Some value function

```
1: all_perms ← All permutations of population
2: F ← length of population
3: result ← 0
4: for perm in all_perms do
5:   member_ix ← index of member in perm
6:   coalition ← coalition of perm before member_ix
7:   diff ← v_function(Union(coalition, member)) - v_function(coalition)
8:   result ← result + diff
9: end for
10: return result / F!
```

- d) Pseudocode of `shapley_perm_approx(member, population, v_function, num_iter)`:

Algorithm 5 `shapley_perm_approx()`

Require: `member`: individual player, or feature of interest
Require: `population`: vector containing all available players
Require: `v_function`: value function
Require: `num_iter`: number of iterations

```
1: vals ← Empty vector or list
2: for i in num_iter do
3:   perm ← draw a random permutation of population
4:   member_ix ← index of member in perm
5:   coalition ← coalition of perm before member_ix
6:   vals[i] ← v_function(Union(coalition, member)) - v_function(coalition)
7: end for
8: return average of vals
```

- e) (i) Pseudocode of `symmetry_check()`:

Algorithm 6 symmetry_check()

Require: j : first feature index
Require: k : second feature index
Require: population: vector containing all available players
Require: v_function: value function
Require: shapley_func: function for computing Shapley values

- 1: remainder \leftarrow everyone from the population but j, k
- 2: all_S \leftarrow all_unique_subsets(remainder)
- 3: **for** S in all_S **do**
- 4: surplus_j \leftarrow v_function(Union(coalition, j)) - v_function(coalition)
- 5: surplus_k \leftarrow v_function(Union(coalition, k)) - v_function(coalition)
- 6: save surplus_j and surplus_k in vectors surplus_j and surplus_k, respectively, for every iteration
- 7: **end for**
- 8: **if** surplus_j equal surplus_k **then**
- 9: **print** *equal surplus*
- 10: val_j \leftarrow shapley_func(j , population, v_function)
- 11: val_k \leftarrow shapley_func(k , population, v_function)
- 12: **return** val_j == val_k
- 13: **end if**
- 14: **return** TRUE

(ii) Pseudocode of dummy_check():

Algorithm 7 dummy_check()

Require: j : feature index
Require: population: vector containing all available players
Require: v_function: value function
Require: shapley_func: function for computing Shapley values

- 1: remainder \leftarrow everyone from the population but j
- 2: all_S \leftarrow all_unique_subsets(remainder)
- 3: **for** S in all_S **do**
- 4: surplus_j \leftarrow difference of v_function of S with j minus v_function of S
- 5: save surplus_j in vector surplus_j for every iteration
- 6: **end for**
- 7: **if** sum of $|surplus_j|$ == 0 **then**
- 8: **print** *has contribution*
- 9: val_j \leftarrow shapley_func(j , population, v_function)
- 10: **return** val_j == 0
- 11: **end if**
- 12: **return** TRUE

(iii) Pseudocode of additivity_check():

Algorithm 8 additivity_check()

Require: population: vector containing all available players
Require: v_function1: value function 1
Require: v_function2: value function 2
Require: shapley_func: function for computing Shapley values

- 1: combined \leftarrow addition of v_function1 and v_function2
- 2: vals1 \leftarrow Shapley values for all features using v_function1
- 3: vals2 \leftarrow Shapley values for all features using v_function2
- 4: vals_comb \leftarrow Shapley values for all features using combined
- 5: vals_additive \leftarrow vals1 + vals2
- 6: **return** vals_comb == vals_additive

(iv) Pseudocode of efficiency_check():

Algorithm 9 efficiency_check()

Require: population: vector containing all available players
Require: v_function: value function
Require: shapley_func: function for computing Shapley values

- 1: payoff_total \leftarrow v_function of population
- 2: shapley_vals \leftarrow Shapley values for all features using v_function
- 3: total_shapley_vals \leftarrow sum of shapley_vals
- 4: **return** payoff_total == total_shapley_vals

Solution 2:

Setup. Let $P = \{1, \dots, p\}$ be the set of $p := |P|$ players, v the value function and ϕ_j the Shapley value of player j .

Recall the two equivalent definitions of Shapley values:

Permutation-Based Definition of Shapley Values. Let $\text{Pred}_\pi(j)$ be the set of players that appear before player j in the permutation π . Then for any j , the Shapley value of player j is

$$\phi_j = \frac{1}{p!} \sum_{\pi \in \mathfrak{S}_P} \left(v(\text{Pred}_\pi(j) \cup \{j\}) - v(\text{Pred}_\pi(j)) \right),$$

where \mathfrak{S}_P is the set of all $p!$ permutations of P . In words, we look at the “marginal contribution” of j each time it arrives in a permutation (depending on whichever players arrived before it), and then average over all permutations.

Set-Based Definition of Shapley Values. For any j , we consider all subsets S of P which do not contain j , then the Shapley value of player j is

$$\phi_j = \sum_{S \subseteq P \setminus \{j\}} \frac{|S|! (p - |S| - 1)!}{p!} \left(v(S \cup \{j\}) - v(S) \right).$$

In essence, both definitions are almost the same, just that those terms from the first definition that yield the same subsets S (although from different permutations), are summed up and counted in the second definition, so that the second definition has less terms in the sum.

(a) **Proof of the first 3 Axioms:** The Dummy and Additivity properties are relatively easy to proof.

Theorem 1 (Null Player (Dummy)). *Let $\{\phi_j\}_{j \in P}$ be the Shapley values induced by v . For any player $j \in P$ it holds:*

$$\text{If } v(S \cup \{j\}) = v(S) \text{ for all } S \subseteq P \setminus \{j\}, \text{ then } \phi_j = 0.$$

Proof. Assume $p \geq 1$, $j \in P$ and $v(S \cup \{j\}) = v(S)$ for all subsets $S \subseteq P \setminus \{j\}$. We need to prove that $\phi_j = 0$.

This follows directly from either of the two definitions: The assumption states that the difference $v(S \cup \{j\}) - v(S)$ inside the sum is 0, for every $S \subseteq P \setminus \{j\}$, in other words for every summand. Therefore, also the whole sum is 0, so $\phi_j = 0$.

□

Theorem 2 (Additivity). *Let $v, v_1, v_2 : 2^P \rightarrow \mathbb{R}$ be value functions with $v = v_1 + v_2$, which means that for every subset $S \subseteq P$ we have $v(S) = v_1(S) + v_2(S)$. Let $\{\phi_{j,v}\}_{j \in P}$ be the Shapley values induced by v and $\{\phi_{j,v_1}\}_{j \in P}$ and $\{\phi_{j,v_2}\}_{j \in P}$ those for v_1 and v_2 respectively. Then for any player $j \in P$ it holds:*

$$\phi_{j,v} = \phi_{j,v_1+v_2} = \phi_{j,v_1} + \phi_{j,v_2}.$$

Proof. Assume $p \geq 1$, $j \in P$ and that v_1, v_2 are two value functions. We need to prove that $\phi_{j,v_1+v_2} = \phi_{j,v_1} + \phi_{j,v_2}$.

For the marginal contributions of player j on any subset $S \subseteq P$ we have: $v(S \cup \{j\}) - v(S) = v_1(S \cup \{j\}) - v_1(S) + v_2(S \cup \{j\}) - v_2(S)$. Plugging this into either of the two definitions yields the result:

$$\begin{aligned}\phi_{j,v} &= \phi_{j,v_1+v_2} = \sum_{S \subseteq P \setminus \{j\}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{j\}) - v(S)) \\ &= \sum_{S \subseteq P \setminus \{j,k\}} \frac{|S|! (p - |S| - 1)!}{p!} (v_1(S \cup \{j\}) - v_1(S) + v_2(S \cup \{j\}) - v_2(S)) \\ &= \sum_{S \subseteq P \setminus \{j,k\}} \frac{|S|! (p - |S| - 1)!}{p!} (v_1(S \cup \{j\}) - v_1(S)) + \sum_{S \subseteq P \setminus \{j,k\}} \frac{|S|! (p - |S| - 1)!}{p!} (v_2(S \cup \{j\}) - v_2(S)) \\ &= \phi_{j,v_1} + \phi_{j,v_2}.\end{aligned}$$

□

The symmetry axiom requires a little more calculation.

Theorem 3 (Symmetry). *Let $\{\phi_j\}_{j \in P}$ be the Shapley values induced by v . For any two players $j, k \in P$ it holds:*

$$\text{If } v(S \cup \{j\}) = v(S \cup \{k\}) \text{ for all } S \subseteq P \setminus \{j, k\}, \text{ then } \phi_j = \phi_k.$$

Proof. Assume $p \geq 2$ and $j, k \in P$ and $v(S \cup \{j\}) = v(S \cup \{k\})$ for all subsets $S \subseteq P \setminus \{j, k\}$. We need to prove that $\phi_j = \phi_k$.

We will use the set-based definition of Shapley values, and will show the claim by straightforward calculation and rearranging the sum. We first split the sum, which is over all sets $S \subseteq P \setminus \{j\}$, into two parts, depending on whether player k is in S or not:

$$\begin{aligned}\phi_j &= \sum_{S \subseteq P \setminus \{j\}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{j\}) - v(S)) \\ &= \sum_{S \subseteq P \setminus \{j,k\}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{j\}) - v(S)) + \sum_{\substack{S \subseteq P \setminus \{j\}, \\ k \in S}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{j\}) - v(S)).\end{aligned}$$

Now, in the first part of the sum we can directly use our assumption:

$$\sum_{S \subseteq P \setminus \{j,k\}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{j\}) - v(S)) = \sum_{S \subseteq P \setminus \{j,k\}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{k\}) - v(S))$$

For the second part, first note that both sums have the same number of summands, since there is a one-to-one correspondence (a bijection) between the subsets of $P \setminus \{j, k\}$ and those subsets of $P \setminus \{j\}$ that contain k , by either adding k to a given subset $S \subseteq P \setminus \{j, k\}$ or not. We can use this bijection to define an “index transformation” for the second sum: We replace the set S , which we know must contain k , with $S = \tilde{S} \cup \{k\}$, in other words we define $\tilde{S} := S \setminus \{k\}$. We can then use our assumption on the set \tilde{S} as well. This yields:

$$\begin{aligned}&\sum_{\substack{S \subseteq P \setminus \{j\}, \\ k \in S}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{j\}) - v(S)) \\ &= \sum_{\tilde{S} \subseteq P \setminus \{j,k\}} \frac{(|\tilde{S}| + 1)! (p - (|\tilde{S}| + 1) - 1)!}{p!} (v(\tilde{S} \cup \{j, k\}) - v(\tilde{S} \cup \{k\})) \\ &= \sum_{\tilde{S} \subseteq P \setminus \{j,k\}} \frac{(|\tilde{S}| + 1)! (p - |\tilde{S}| - 2)!}{p!} (v(\tilde{S} \cup \{j, k\}) - v(\tilde{S} \cup \{j\})) \\ &= \sum_{\substack{S \subseteq P \setminus \{k\}, \\ j \in S}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{k\}) - v(S)).\end{aligned}$$

In the last step, we used this transformation backwards, so we add / remove j instead of k this time, in other words $S = \tilde{S} \cup \{j\}$ and $\tilde{S} := S \setminus \{j\}$. This is possible since \tilde{S} contains neither j nor k .

We now use the first step from the beginning backwards and arrive at the desired result:

$$\begin{aligned}\phi_j &= \sum_{S \subseteq P \setminus \{j,k\}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{k\}) - v(S)) + \sum_{\substack{S \subseteq P \setminus \{k\}, \\ j \in S}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{k\}) - v(S)) \\ &= \sum_{S \subseteq P \setminus \{k\}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{k\}) - v(S)) = \phi_k.\end{aligned}$$

□

(b) Bonus: Proof of the Efficiency Axiom

Efficiency requires a little more effort than the others:

Theorem 4 (Efficiency). *Let $\{\phi_j\}_{j \in P}$ be the Shapley values induced by v . Then*

$$\sum_{j=1}^p \phi_j = v(P).$$

The proof idea is as follows: Because we sum up the Shapley values over all players, the values $v(S)$ of each coalition not equal to P , so $S \subsetneq P$, appear with both minus and plus signs that exactly cancel each other out. Apart from that, the term $v(P)$ occurs for every player, and one only needs to calculate that its weights sum up to 1.

We will use the order-based definition of Shapley values here, since with it the proof is simpler.

Proof. We first plug in the definition and swap the two sums:

$$\begin{aligned}\sum_{j=1}^p \phi_j &= \sum_{j=1}^p \frac{1}{p!} \sum_{\pi \in \mathfrak{S}_P} (v(\text{Pred}_\pi(j) \cup \{j\}) - v(\text{Pred}_\pi(j))) \\ &= \frac{1}{p!} \sum_{\pi \in \mathfrak{S}_P} \sum_{j=1}^p (v(\text{Pred}_\pi(j) \cup \{j\}) - v(\text{Pred}_\pi(j))).\end{aligned}$$

Now the inner sum corresponds to a single permutation, and sums over all players in this permutation.

Fix a particular permutation π . We consider the order of the players given by this permutation: $(\pi(1), \pi(2), \dots, \pi(p))$. Each marginal contribution inside the inner sum corresponds to one step in this ordering. We reorder the inner sum according to this ordering given by the permutation:

$$\begin{aligned}\sum_{j=1}^p (v(\text{Pred}_\pi(j) \cup \{j\}) - v(\text{Pred}_\pi(j))) &= \sum_{i=1}^p (v(\{\pi(1), \dots, \pi(i)\}) - v(\{\pi(1), \dots, \pi(i-1)\})) \\ &= \underbrace{v(\{\pi(1)\}) - v(\emptyset)}_{i=1 \text{ or } j=\pi(1)} + \underbrace{v(\{\pi(1), \pi(2)\}) - v(\{\pi(1)\})}_{i=2 \text{ or } j=\pi(2)} + \underbrace{v(\{\pi(1), \pi(2), \pi(3)\}) - v(\{\pi(1), \pi(2)\})}_{i=3 \text{ or } j=\pi(3)} \\ &\quad + \dots + \underbrace{v(\{\pi(1), \dots, \pi(p)\}) - v(\{\pi(1), \dots, \pi(p-1)\})}_{i=p \text{ or } j=\pi(p)} = v(\{\pi(1), \dots, \pi(p)\}) - v(\emptyset) = v(P).\end{aligned}$$

In other words, the inner sum is a telescope sum, where all the terms except the first and the last one cancel. Plugging this into the outer sum yields the desired result:

$$\begin{aligned}\sum_{j=1}^p \phi_j &= \frac{1}{p!} \sum_{\pi \in \mathfrak{S}_P} \left(\sum_{j=1}^p (v(\text{Pred}_\pi(j) \cup \{j\}) - v(\text{Pred}_\pi(j))) \right) \\ &= \frac{1}{p!} \sum_{\pi \in \mathfrak{S}_P} v(P) = \frac{1}{p!} \cdot (p! \cdot v(P)) = v(P).\end{aligned}$$

□

When using the set-based definition, one can show that the term $v(P)$ appears exactly once per player (hence p times in total) for the coalition $S = P \setminus \{j\}$ with the weight $\frac{|P-1|! (|P|-|P-1|-1)!}{|P|!} = \frac{1}{p}$ each, and then one just has to show that for all the other terms, their number of appearance together with their weights cause them to cancel out. See also here.

(c) Bonus: Proof of Linearity

This is as straightforward as the additivity:

Theorem 5 (Linearity). *Let $\alpha, \beta \in \mathbb{R}$ and $v, v_1, v_2 : 2^P \rightarrow \mathbb{R}$ be value functions with $v = \alpha v_1 + \beta v_2$, which means that for every subset $S \subseteq P$ we have $v(S) = \alpha \cdot v_1(S) + \beta \cdot v_2(S)$. Let $\{\phi_{j,v}\}_{j \in P}$ be the Shapley values induced by v and $\{\phi_{j,v_1}\}_{j \in P}$ and $\{\phi_{j,v_2}\}_{j \in P}$ those for v_1 and v_2 respectively. Then for any player $j \in P$ it holds:*

$$\phi_{j,v} = \phi_{j,\alpha v_1 + \beta v_2} = \alpha \phi_{j,v_1} + \beta \phi_{j,v_2}.$$

Proof. Since additivity was already proven in part (a), we only have to prove homogeneity, that means that $\phi_{j,\alpha v_1} = \alpha \phi_{j,v_1}$ for any $\alpha \in \mathbb{R}$.

So let $p \geq 1$, $j \in P$, $\alpha \in \mathbb{R}$ and v_1 be any value function. Denote $v = \alpha v_1$. For the marginal contributions of player j on any subset $S \subseteq P$ we have: $v(S \cup \{j\}) - v(S) = \alpha v_1(S \cup \{j\}) - \alpha v_1(S) = \alpha(v_1(S \cup \{j\}) - v_1(S))$. As for the additivity, plugging this into either of the two definitions yields the results:

$$\begin{aligned} \phi_{j,v} &= \sum_{S \subseteq P \setminus \{j\}} \frac{|S|! (p - |S| - 1)!}{p!} (v(S \cup \{j\}) - v(S)) = \sum_{S \subseteq P \setminus \{j\}} \frac{|S|! (p - |S| - 1)!}{p!} \alpha (v_1(S \cup \{j\}) - v_1(S)) \\ &= \alpha \sum_{S \subseteq P \setminus \{j\}} \frac{|S|! (p - |S| - 1)!}{p!} (v_1(S \cup \{j\}) - v_1(S)) = \alpha \phi_{j,v_1}. \end{aligned}$$

□