

THE UNIVERSITY OF BRITISH COLUMBIA  
CPSC 121: MIDTERM 1 – Individual – February 6, 2017

**Important notes about this examination**

1. **You have 70 minutes to complete this exam.**
2. You are allowed up to three textbooks and (the equivalent of) a 3" 3-ring binder of notes as references. Otherwise, no notes or aids are allowed. No electronic equipment is allowed.
3. Good luck!
4. There are a total of 70 marks available.

Full Name: \_\_\_\_\_

Signature: \_\_\_\_\_

UBC Student #: \_\_\_\_\_

Exam ID: \_\_\_\_\_



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## **Student Conduct during Examinations**

1. Each examination candidate must be prepared to produce, upon the request of the invigilator or examiner, his or her UBCcard for identification.
2. Examination candidates are not permitted to ask questions of the examiners or invigilators, except in cases of supposed errors or ambiguities in examination questions, illegible or missing material, or the like.
3. No examination candidate shall be permitted to enter the examination room after the expiration of one-half hour from the scheduled starting time, or to leave during the first half hour of the examination. Should the examination run forty-five (45) minutes or less, no examination candidate shall be permitted to enter the examination room once the examination has begun.
4. Examination candidates must conduct themselves honestly and in accordance with established rules for a given examination, which will be articulated by the examiner or invigilator prior to the examination commencing. Should dishonest behaviour be observed by the examiner(s) or invigilator(s), pleas of accident or forgetfulness shall not be received.
5. Examination candidates suspected of any of the following, or any other similar practices, may be immediately dismissed from the examination by the examiner/invigilator, and may be subject to disciplinary action:
  - i. speaking or communicating with other examination candidates, unless otherwise authorized;
  - ii. purposely exposing written papers to the view of other examination candidates or imaging devices;
  - iii. purposely viewing the written papers of other examination candidates;
  - iv. using or having visible at the place of writing any books, papers or other memory aid devices other than those authorized by the examiner(s); and,
  - v. using or operating electronic devices including but not limited to telephones, calculators, computers, or similar devices other than those authorized by the examiner(s)—  
(electronic devices other than those authorized by the examiner(s) must be completely powered down if present at the place of writing).
6. Examination candidates must not destroy or damage any examination material, must hand in all examination papers, and must not take any examination material from the examination room without permission of the examiner or invigilator.
7. Notwithstanding the above, for any mode of examination that does not fall into the traditional, paper-based method, examination candidates shall adhere to any special rules for conduct as established and articulated by the examiner.
8. Examination candidates must follow any additional examination rules or directions communicated by the examiner(s) or invigilator(s).

[8] 1. **Logical Equivalence Proof**

Using a sequence of logical equivalence rules from “Dave’s Awesome Handout”, simplify:

$$((q \vee (p \wedge p)) \wedge (q \rightarrow p)) \vee \sim(p \rightarrow (r \vee \sim s))$$

Please write the name of the law(s) you applied at each step. *Hint:* your final simplified form should be **very** simple.

[8] 2. **Tautologies, Contingencies, and Contradictions:**

Determine whether the following statements are **tautologies** (true for all assignments of truth values to their variables, i.e., logically equivalent to  $T$ ), **contradictions** (false for all assignments of truth values to their variables, i.e., logically equivalent to  $F$ ), or contingencies (true or false depending on the values of their variables). Indicate your answer by circling *one* of TAUTOLOGY, CONTRADICTION, and CONTINGENCY for each. You do **not** need to justify your answer.

[2] a.  $d$

TAUTOLOGY	CONTRADICTION	CONTINGENCY
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[2] b.  $(r \wedge s) \rightarrow r$

TAUTOLOGY	CONTRADICTION	CONTINGENCY
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[2] c.  $(p \oplus q) \leftrightarrow (q \vee p)$

TAUTOLOGY	CONTRADICTION	CONTINGENCY
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[2] d.  $(p \vee q) \wedge (p \vee \sim q) \wedge (\sim p \vee q) \wedge (\sim p \vee \sim q)$

TAUTOLOGY	CONTRADICTION	CONTINGENCY
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[23] 3. **Number Representation**

**THIS TEXT EXACTLY REPEATS THE PRE-READING YOU WERE GIVEN, UP TO “END OF PRE-READING” BELOW.**

In our discussion of number representation, we usually discuss **fixed-width** representations. That is, a number might be represented by 8 bits or 32 bits, but not by “however many bits it takes”.

However, when people write out numbers, we usually use a variable-width representation. For instance, the grade you get for this course might be a 3-digit number (100), a 2-digit number (10–99), or (hopefully not) a 1-digit number (0–9). If grades were “fixed-width”, then, we’d have to write a grade like 83 as 083 instead. An apparent advantage of our variable-width representation is that we can write numbers as large as we like just by taking up more space.

For this problem, we’ll consider an approach that uses bits to achieve this sort of variable-width representation of unsigned binary values. We call our proposed representation “flag bit” numbers: a number is composed of one or more 2-bit “blocks”. The first bit in each block is a normal base-2 digit. The second bit is the “flag bit”. It is 0 if this is the last block in the number and 1 if there is another block. To determine the value of the variable-width number, we collect up all the normal digits, **reverse them**, and interpret the result as an unsigned binary number.

For clarity, we’ll write “flag bit” numbers with commas separating the blocks, although those wouldn’t be included in the stored representation of the numbers.

For example, the numbers 0 and 1 would be 00 and 10 (respectively) as “flag bit” numbers.

The numbers 2 and 3 would be 01,10 and 11,10 (respectively).

The number 01,11,01,11,10 is the decimal number 26 because: the first four blocks’ flag bits indicate that the number continues up to the fifth block’s flag bit, which indicates that the number ends there; the binary digits in those five blocks, read from left to right, are 01011; we reverse these bits to get 11010; and, that unsigned binary number is the decimal number 26.

**END OF PRE-READING**

- [2] a. Translate the “flag bit” number 01,11,10 to a decimal number.
- [2] b. Express the decimal number 13 as a “flag bit” number. (Write commas between your blocks in your answer so we can read it easily!)
- [2] c. An 8 bit unsigned binary number can represent  $2^8 = 256$  distinct values. How many distinct values can we represent with a single “flag bit” number of 8 bits (or fewer)?

- [2] d. “Flag bit” numbers have the interesting property that we can “smash” them together and still tell where one ends and the next begins. For example, 01,10,10 must be two numbers (the decimal numbers 2 and then 1) rather than just one because the second block’s 0 flag bit ends the first number and the third block’s 0 flag bit ends the second number.

**Circle each separate “flag bit” number** in this sequence of blocks: 10,01,10,11,11,10

- [15] e. In this part, you will design a circuit that takes 4 bits of input  $f_1, f_2, f_3, f_4$  that constitute two blocks like  $f_1f_2, f_3f_4$ . You’ll first design a circuit to test whether the input contains at least one whole flag bit number. Then, you’ll design a circuit that assumes the input has at least one whole flag bit number and produces its value as a 2-bit unsigned binary number.

You **must** show clear design steps, **including truth table(s), logical expression(s), and a correct circuit**. Some credit is reserved for the elegance of your solution, but most can be earned for any clear, correct answer.

- i. **Using only  $f_2$  and  $f_4$ , produce an output  $v$  that is 1 (true) exactly when the input contains at least one whole flag bit number.** For example, all of these should produce  $v = 1$ : 01,10 is a single whole flag bit number, 00,10 is two whole flag bit numbers, 10,11 contains a whole flag bit number (10) even though it ends with an incomplete flag bit number. **However**, 11,11 should produce  $v = 0$  because it is not a whole flag bit number.

- ii. **Using only  $f_1$ ,  $f_2$ , and  $f_3$  (and assuming that  $f_4 = 0$ ), convert the *first whole flag bit number* in the input into the corresponding 2-bit unsigned binary number  $o = o_1o_2$ .** For example, 01,10 should produce  $o = 10$  (i.e.,  $o_1 = 1$  and  $o_2 = 0$ , representing the decimal number 2) because the whole input is the single flag bit number representing the decimal number 2. 00,10 should produce  $o = 00$  because the first whole flag bit number in 00,10 is 00, which represents the decimal value 0. Similarly, 10,11 should produce  $o = 01$ . Finally, even though 11,11 does not represent a whole flag bit number, we assume for this part that  $f_4$  is actually 0 instead and produce  $o = 11$  because the decimal value of the flag bit number 11,10 is 3.

[17] 4. **Propositional Logic Proof**

Prove/disprove the argument in the following two questions. The premises in both parts a and b are the same.

- [9] a. Prove that the following argument is **valid**. Justify each step using logical equivalences or rules of inference on Dave's Awesome Handout. Do not re-write the premises.

1.  $\sim(r \vee \sim s)$
2.  $(\sim u \wedge t) \rightarrow r$
3.  $\sim x \oplus \sim q$
4.  $\sim s \vee t$
5.  $t \rightarrow \sim q$

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$\therefore x$

- [8] b. Prove that the following argument is **invalid**.

1.  $\sim(r \vee \sim s)$
2.  $(\sim u \wedge t) \rightarrow r$
3.  $\sim x \oplus \sim q$
4.  $\sim s \vee t$
5.  $t \rightarrow \sim q$

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$\therefore \sim u \wedge t$



[7] 5. Translating From English to Predicate Logic

**THIS TEXT EXACTLY REPEATS THE PRE-READING YOU WERE GIVEN, UP TO “END OF PRE-READING” BELOW.**

Consider the following definitions:

- $A$ : all (non-human) animals near the barn.
- $P$ : all people near the barn.
- $L$ : all locations near the barn.
- $Cow(x)$ : animal  $x$  is a cow.
- $Pig(x)$ : animal  $x$  is a pig.
- $Friends(x, y)$ :  $x$  and  $y$  are friends.
- $Noise(x, y)$ :  $x$  has made an animal noise at  $y$ .
- $Enjoy(x, y)$ :  $x$  enjoys spending time at location  $y$ .

For example, we can translate the English statement “Alice enjoys spending time anywhere near the barn” into predicate logic as  $\forall x \in L, Enjoy(Alice, x)$ .

**END OF PRE-READING**

Translate each of the following English statements into predicate logic.

[1] a. There is a cow who is friends with Ryan.

[3] b. There is at least one cow and at least one pig who have made animal noises at Steve.

[3] c. Every cow enjoys spending time in the pasture (a particular location in  $L$ ), and is friends with at least one pig.

[4] 6. **Translating From Predicate Logic to English** Using the same definitions as in the previous question, translate the following predicate logic statement into English:

[2] a. Translate:

$$\exists x \in A, \text{Cow}(x) \wedge (\forall y \in A, \text{Pig}(y) \rightarrow \text{Friends}(x, y))$$

[2] b. Translate:

$$\exists x \in L, (\exists y \in A, \text{Cow}(y) \wedge \text{Enjoy}(y, x)) \wedge (\exists y \in A, \text{Pig}(y) \wedge \text{Enjoy}(y, x))$$

[3] 7. **A Little Set and Function Theory**

[2] a. Which of the following sets is equal to  $S = \{2, 3, 5, 7\}$ ? *Clearly* indicate **all** that apply.

- i.  $\{2, 3, 5, 7\}$
- ii.  $\{3, 2, 5, 7\}$
- iii.  $\{\{2, 3\}, \{5\}, \{7\}\}$
- iv.  $\{2, 3, 3, 7, 7, 5\}$
- v.  $\{\{2, 2, 3\}, \{5, 5\}, \{\}, \{7\}\}$
- vi.  $\{2, 3, 5, 7, \{\}, \{\}\}$

[1] b. Let  $P$  be a predicate on the domain  $D$ . We can conceive of  $P$  as a function mapping elements from  $D$  to true or false:  $\{T, F\}$ . Write a predicate logic expression that expresses the following fact about  $P$ :  $P$  never maps any element of  $D$  to  $F$ .

DO NOT DETACH; if you need extra space, please use this page, clearly labeling what problem your work goes with and noting at that problem's page that you're using this space