1.	Consider an alphabet with five letters, $\{a,b,c,d,e\}$, and suppose we know the frequencies $f_a=0.32$, $f_b=0.25$, $f_c=0.2$, $f_d=0.18$, and $f_e=0.05$. What is the expected number of bits used by Huffman's coding scheme to encode a 1000-letter document?		
		2230	
	正确 For e	example, $a=0$ 0, $b=0$ 1, $c=1$ 0, $d=1$ 10, $e=1$ 11.	
		2400	
		3000	
		3450	
2.	Under a Huffman encoding of n symbols, how long (in terms of number of bits) can a codeword possibly be?		
		n-1	
		the lower bound, take frequencies proportional to powers of 2. For the upper nd, note that the total number of merges is exactly $n-1.$	
		$\log_2 n$	
		$\ln n$	
		n	
3.	Which	of the following statements holds for Huffman's coding scheme? If the most frequent letter has frequency less than 0.33 , then all letters will be	
		coded with at least two bits.	
	invo itera rem	a letter will endure a merge in at least two iterations: the last one (which lives all letters), and at least one previous iteration. In the penultimate ation, if the letter has not yet endured a merge, at least one of the two other aining subtrees has cumulative frequency at least $(133)/2>.33$, so the er will get merged in this iteration.	
		A letter with frequency at least 0.5 might get encoded with two or more bits.	
		If a letter's frequency is at least 0.4 , then the letter will certainly be coded with	
		only one bit. If the most frequent letter has frequency less than 0.5 , then all letters will be	
		coded with more than one bit.	
4.	Which of the following is true for our dynamic programming algorithm for computing a maximum-weight independent set of a path graph? (Assume there are no ties.)		
		If a vertex is excluded from the optimal solution of a subproblem, then it is excluded from the optimal solutions of all bigger subproblems.	
		As long as the input graph has at least two vertices, the algorithm never selects the minimum-weight vertex.	
		The algorithm always selects the maximum-weight vertex.	
		If a vertex is excluded from the optimal solution of two consecutive subproblems, then it is excluded from the optimal solutions of all bigger	
		subproblems.	
		nduction, since the optimal solution to a subproblem depends only on the	
	Solu	tions of the previous two subproblems.	
5.	indeperation of the second of	our dynamic programming algorithm for computing the maximum-weight endent set of a path graph. Consider the following proposed extension to more all graphs. Consider an undirected graph with positive vertex weights. For a vertex win the graph $G'(v)$ by deleting v and its incident edges from G , and obtain the $G''(v)$ from G by deleting v , its neighbors, and all of the corresponding incident from G . Let $OPT(H)$ denote the value of a maximum-weight independent set of the G consider the formula $G(G) = \max\{OPT(G'(v)), w_v + OPT(G''(v))\}$, where v is an arbitrary vertex of weight w_v . Which of the following statements is true?	
		The formula is always correct in trees, and it leads to an efficient dynamic programming algorithm.	
	正确 Inde	ed. What running time can you get?	
		The formula is correct in path graphs but is not always correct in trees.	
		The formula is always correct in trees, but does not lead to an efficient dynamic programming algorithm.	
		The formula is always correct in general graphs, and it leads to an efficient dynamic programming algorithm.	