# Inverse Problems — Summary

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# Contents

	Generalised Solutions				
	1.1	Generalised inverses		3	
	1.2	Compa	act operators	5	
1.1 1.2 2 Cl 2.1	Cla	lassical regularisation theory			
	2.1	Param	eter choice rules	9	
		2.1.1	A priori parameter choice rules	9	
		2.1.2	A posteriori parameter choice rules	9	
		2.1.3	Heuristic parameter choice rules	10	
	2.2		al regularisation		

A direct problem is a problem where given an object or cause, we must determine the data or effect. In an inverse problem, we observe data and wish to recover the object.

A problem is called *well-posed* if a unique solution exists that depends continuously on the data. Most inverse problems are, unfortunately, ill-posed.

# 1 Generalised Solutions

**Recap 1.1.** 1. An linear operator  $A: \mathcal{X} \to \mathcal{Y}$  is called *bounded* if

$$||A||_{\mathcal{B}(\mathcal{X},\mathcal{Y})} \coloneqq \sup_{u \neq 0} \frac{||Au||_{\mathcal{Y}}}{||u||_{\mathcal{X}}} = \sup_{||u||_{\mathcal{X}} \leq 1} < \infty.$$

It is known that a linear operator between normed spaces is continuous if and only if it is bounded. The set of bounded linear operators from  $\mathcal{X}$  to  $\mathcal{Y}$  is denotes  $\mathcal{B}(\mathcal{X}, \mathcal{Y})$ .

- 2. We let  $\mathcal{D}(A)$ ,  $\mathcal{N}(A)$  and  $\mathcal{R}(A)$  denote the domain, null space, and range of A respectively.
- 3. We will assume  $\mathcal{X}$  and  $\mathcal{Y}$  are Hilbert spaces, so there is an inner product  $\langle \cdot, \cdot \rangle$  and any bounded operator  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$  has a unique adjoint  $A^* \in \mathcal{B}(\mathcal{Y}, \mathcal{X})$  which satisfies

$$\langle Au, v \rangle_{\mathcal{Y}} = \langle u, A^*v \rangle_{\mathcal{X}}$$
 for all  $u \in \mathcal{X}, v \in \mathcal{Y}$ .

4. For any  $\mathcal{X}' \subseteq \mathcal{X}$  we define the *orthogonal complement* of  $\mathcal{X}'$  as

$$(\mathcal{X}')^{\perp} := \{ u \in \mathcal{X} \mid \langle u, v \rangle_{\mathcal{X}} = 0 \ \forall v \in \mathcal{X}' \}.$$

It is known that  $(\mathcal{X}')^{\perp}$  is a closed subspace of  $\mathcal{X}$  and that  $\mathcal{X}' \subseteq ((\mathcal{X}')^{\perp})^{\perp}$ , where equality holds if and only if  $\mathcal{X}'$  is a closed subspace of  $\mathcal{X}$ . For a non-closed subspace  $\mathcal{X}'$  we have  $((\mathcal{X}')^{\perp})^{\perp} = \overline{\mathcal{X}'}$ .

- 5. If  $\mathcal{X}'$  is a closed subspace of  $\mathcal{X}$ , then for any  $u \in \mathcal{X}$  there exist unique  $x_u \in \mathcal{X}'$ ,  $x_u^{\perp} \in (\mathcal{X}')^{\perp}$  such that  $u = x_u + x_u^{\perp}$ . The map  $u \mapsto x_u$  is denoted  $P_{\mathcal{X}'}$  and is called the *orthogonal* projection on  $\mathcal{X}'$ . Properties are:
  - (a)  $P_{\mathcal{X}'}$  is bounded and self-adjoint with norm 1;
  - (b)  $P_{X'} + P_{(X')^{\perp}} = I;$
  - (c)  $P_{\mathcal{X}'}u$  minimises the distance from u to  $\mathcal{X}'$ ;
  - (d)  $x = P_{\mathcal{X}'}u$  if and only if  $x \in \mathcal{X}'$  and  $u x \in (\mathcal{X}')^{\perp}$ .
- 6. For any  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$  we have

$$\mathcal{R}(A)^{\perp} = \mathcal{N}(A^*)$$
 and  $\mathcal{N}(A)^{\perp} = \overline{\mathcal{R}(A^*)}$ 

**Lemma 1.2.** For any  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$  we have  $\overline{\mathcal{R}(A^*A)} = \overline{\mathcal{R}(A^*)}$ .

*Proof.* It is trivial that  $\overline{\mathcal{R}(A^*A)} \subseteq \overline{\mathcal{R}(A^*)}$ .

Now, suppose  $u \in \overline{\mathcal{R}(A^*)}$  and let  $\varepsilon > 0$ . Then there exists  $v \in \mathcal{X}$  such that  $||A^*v - u|| < \varepsilon/2$ . Writing v = e + f with  $e \in \mathcal{N}(A^*)$ ,  $f \in \mathcal{N}(A^*)^{\perp} = \overline{\mathcal{R}(A)}$ , we see that  $||A^*f - u|| < \varepsilon/2$ .

Since  $f \in \overline{\mathcal{R}(A)}$ , there exists  $x \in \mathcal{X}$  such that  $||Ax - f|| < \varepsilon/(2||A||)$ . We now compute

$$\|A^*Ax-u\|\leq \|A^*Ax-A^*f\|+\|A^*f-u\|<\|A^*\|\frac{\varepsilon}{2\|A\|}+\frac{\varepsilon}{2}=\varepsilon,$$

and conclude that  $u \in \overline{\mathcal{R}(A^*A)}$ . This shows that  $\overline{\mathcal{R}(A)} \subseteq \overline{\mathcal{R}(A^*A)}$ .

# 1.1 Generalised inverses

We consider the equation

$$Au = f, (1)$$

where  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$  and f are known, and we wish to find u.

**Definition 1.3.** An element  $u \in \mathcal{X}$  is called a *least-squares solution* of eq. (1) if u is a minimiser of the function  $v \mapsto ||Av - f||_{\mathcal{Y}}$ . It is called a *minimal-norm solution* of eq. (1) if it has minimal norm among all least-squares solutions.

Note that a least-squares solution may not exist. If a least-squares solution u exists, then the affine subspace of all least-squares solutions is given by  $u + \mathcal{N}(A)$ . By writing  $u = u^{\dagger} + v$  for  $u^{\dagger} \in \mathcal{N}(A)^{\perp}$ ,  $v \in \mathcal{N}(A)$ , we find that the space of least-squares solutions is given by  $u^{\dagger} + \mathcal{N}(A)$ , and it is now clear that  $u^{\dagger}$  is the unique minimum-norm solution.

**Theorem 1.4.** Let  $f \in \mathcal{Y}$  and  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$ . Then the following are equivalent:

- 1.  $u \in \mathcal{X}$  satisfies  $Au = P_{\overline{\mathcal{R}(A)}}f$ ;
- 2. u is a least-squares solution of eq. (1):
- 3. u solves the normal equation

$$A^*f = A^*Au. (2)$$

*Proof.* " $(1) \implies (2)$ ": We have

$$||Au - f||_{\mathcal{Y}} = \left||P_{\overline{\mathcal{R}(A)}}f - f\right|| = \inf_{g \in \overline{\mathcal{R}(A)}} ||g - f|| \le \inf_{g \in \overline{\mathcal{R}(A)}} ||g - f|| = \inf_{u \in \mathcal{X}} ||Au - f||.$$

"(2)  $\Longrightarrow$  (3)": Let  $u \in \mathcal{X}$  be a least-squares solution and  $v \in \mathcal{X}$  arbitrary. Define the quadratic polynomial

$$F: \mathbb{R} \to \mathbb{R}: \lambda \mapsto ||A(u + \lambda v) - f||^2$$

$$= \langle Au + \lambda Av - f, Au + \lambda Av - f \rangle$$

$$= \lambda^2 ||Av||^2 - 2\lambda \langle Av, f - Au \rangle + ||f - Au||^2.$$

As u is a least-squares solution, we know that F attains a minimum in  $\lambda = 0$  and therefore that

$$0 = F'(0) = 2\langle Av, f - Au \rangle = 2\langle v, A^*(f - Au) \rangle.$$

Since v is arbitrary, we must have  $A^*(f - Au) = 0$ , so u satisfies eq. (2).

"(3)  $\implies$  (1)": From the normal equation we know that  $A^*(f - Au) = 0$ . For any  $x \in \mathcal{X}$ , we have

$$\langle Ax, f - Au \rangle = \langle x, A^*(f - Au) \rangle = \langle x, 0 \rangle = 0,$$

so  $f - Au \in \mathcal{R}(A)^{\perp}$ .

So we have  $Au \in \overline{\mathcal{R}(A)}$  and  $f - Au \in \mathcal{R}(A)^{\perp} = \overline{\mathcal{R}(A)}^{\perp}$ , from which it follows that  $Au = P_{\overline{\mathcal{R}(A)}}f$ .  $\square$ 

The following lemma gives a precise condition for when a least-squares solution exists:

**Lemma 1.5.** Equation (1) has a least-squares solution if and only if  $f \in \mathcal{R}(A) \oplus \mathcal{R}(A)^{\perp}$ .

*Proof.* " $\Longrightarrow$ " Suppose u is a least-squares solution. Then  $f - Au \in \mathcal{R}(A)^{\perp}$ , so  $f = Au + (f - Au) \in \mathcal{R}(A) \oplus \mathcal{R}(A)^{\perp}$ .

"  $\Leftarrow$ " Suppose f = Au + g for some  $u \in \mathcal{X}$ ,  $g \in \mathcal{R}(A)^{\perp} = \overline{\mathcal{R}(A)}^{\perp}$ . Then by the previous theorem,  $Au = P_{\overline{\mathcal{R}(A)}}f$ , so u is a least-squares solution.

Corollary 1.6. If  $\mathcal{R}(A)$  is closed, then eq. (1) always has a least-squares solution.

In particular, this holds if  $\mathcal{R}(A)$  is finite-dimensional. Therefore, if either  $\mathcal{X}$  or  $\mathcal{Y}$  is finite-dimensional, eq. (1) has a least-squares solution for any A.

We have already seen that if a least-squares solution u exists, then the affine subspace of all least-squares solutions is  $u + \mathcal{N}(A)$ , and the unique minimum-norm solution is the projection of 0 onto this affine subspace, which is the unique element of  $u + \mathcal{N}(A)$  that lies in  $\mathcal{N}(A)^{\perp}$ .

**Definition 1.7.** Let  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$ , and define

$$\tilde{A} := A \upharpoonright_{\mathcal{N}(A)^{\perp}} : \mathcal{N}(A)^{\perp} \to \mathcal{R}(A).$$

Clearly  $\tilde{A}$  is bijective and we define the Moore-Penrose inverse

$$A^{\dagger} \colon \mathcal{R}(A) \oplus \mathcal{R}(A)^{\perp} \to \mathcal{N}(A)^{\perp} \colon f \mapsto \tilde{A}^{-1}P_{\overline{\mathcal{R}(A)}}f.$$

Remark. Note that  $\overline{\mathcal{R}(A) \oplus \mathcal{R}(A)^{\perp}} = \overline{\mathcal{R}(A)} \oplus \mathcal{R}(A)^{\perp} = \overline{\mathcal{R}(A)} \oplus \overline{\mathcal{R}(A)}^{\perp} = \mathcal{Y}$ , and therefore the operator  $\tilde{A}$  is densely defined, and it is defined on all of  $\mathcal{Y}$  if and only if  $\mathcal{R}(A)$  is closed.

We will not prove the following theorem, but it is interesting:

**Theorem 1.8.** The Moore-Penrose inverse  $A^{\dagger}$  is continuous if and only if  $\mathcal{R}(A)$  is closed.

The following characterises all important facts about the Moore-Penrose inverse:

**Theorem 1.9** (Moore-Penrose equations). The operator  $A^{\dagger}$  satisfies the following equations:

- (1)  $A^{\dagger}A = P_{\mathcal{N}(A)^{\perp}};$
- (2)  $AA^{\dagger} = P_{\overline{\mathcal{R}(A)}} \upharpoonright_{\mathcal{D}(A^{\dagger})};$
- (3)  $AA^{\dagger}A = A$ ;
- (4)  $A^{\dagger}AA^{\dagger} = A^{\dagger}$ .

Conversely, if any linear operator  $B \colon \mathcal{R}(A) \oplus \mathcal{R}(A)^{\perp} \to \mathcal{N}(A)^{\perp}$  satisfies  $BA = P_{\mathcal{N}(A)^{\perp}}$  and  $AB = P_{\overline{\mathcal{R}(A)}} \upharpoonright_{\mathcal{D}(A^{\dagger})}$  then  $B = A^{\dagger}$ .

*Proof.* We have

$$A^{\dagger}Au = \tilde{A}^{-1}AP_{\mathcal{N}(A)^{\perp}}u = P_{\mathcal{N}(A)^{\perp}}u,$$

which proves (1). Furthermore, we have

$$AA^{\dagger}f = A\tilde{A}^{-1}P_{\overline{\mathcal{R}(A)}}f = P_{\overline{\mathcal{R}(A)}}f,$$

which proves (2). Finally, (3) follows from (1) and (4) follows from (2).

Now, suppose B satisfies (1) and (2). First we show that  $B|_{\mathcal{R}(A)} = \tilde{A}^{-1}$ , then we show that  $B|_{\mathcal{R}(A)^{\perp}} = 0$ . This shows that  $B = A^{\dagger}$ . Let  $f = Au \in \mathcal{R}(A)$  with  $u \in \mathcal{N}(A)^{\perp}$ , then

$$Bf = BAu = P_{\mathcal{N}(A)^{\perp}}u = u = \tilde{A}^{-1}f, \text{ so } B \upharpoonright_{\mathcal{R}(A)} = \tilde{A}^{-1}.$$

Finally, let  $f \in \mathcal{R}(A)^{\perp}$ , then  $ABf = P_{\overline{\mathcal{R}(A)}}f = 0$ , and since  $Bf \in \mathcal{N}(A)^{\perp}$  this implies Bf = 0. We conclude that  $B \upharpoonright_{\mathcal{R}(A)^{\perp}} = 0$ , and this concludes the proof.

The Moore-Penrose inverse has the important property that it maps every f in its domain to the corresponding minimum-norm least-squares solution:

**Theorem 1.10.** For every  $f \in \mathcal{D}(A^{\dagger})$ , the minimum-norm solution  $u^{\dagger}$  to eq. (1) is given by  $u^{\dagger} = A^{\dagger}f$ .

*Proof.* Since  $f \in \mathcal{R}(A) \oplus \mathcal{R}(A)^{\perp}$ , we know that there exists a unique minimum-norm solution  $u^{\dagger} \in \mathcal{N}(A)^{\dagger}$ . We write

$$u^\dagger = P_{\mathcal{N}(A)^\perp}(u^\dagger) = A^\dagger A u^\dagger = A^\dagger P_{\overline{\mathcal{R}(A)}} f = A^\dagger A A^\dagger f = A^\dagger f.$$

Remark. We can also consider the normal equation  $A^*f = A^*Au$  as a least-squares problem, whose minimum-norm solution is  $(A^*A)^{\dagger}A^*f$ . It is clear that this expression must equal the minimum-norm solution  $u^{\dagger}$  from eq. (1).

## 1.2 Compact operators

**Definition 1.11.** Let  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$ . Then A is called *compact* if for any bounded  $B \subseteq \mathcal{X}$ , the image A(B) is precompact in  $\mathcal{Y}$ . The set of compact operators in  $\mathcal{B}(\mathcal{X}, \mathcal{Y})$  is denoted  $\mathcal{K}(\mathcal{X}, \mathcal{Y})$ .

**Lemma 1.12.** Let  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$ . Then A is compact if and only if, for every bounded sequence  $(x_n) \subseteq X$ , the sequence  $(Ax_n) \subseteq Y$  has a convergent subsequence.

**Theorem 1.13.** Let  $A \in \mathcal{K}(\mathcal{X}, \mathcal{Y})$  with  $\dim(\mathcal{R}(A)) = \infty$ . Then  $A^{\dagger}$  is discontinuous.

*Proof.* If dim  $\mathcal{R}(A) = \infty$ , then  $\mathcal{X}$  and  $\mathcal{N}(A)^{\perp}$  are infinite-dimensional as well. Chose an orthonormal sequence  $(x_n) \subseteq \mathcal{N}(A)^{\perp}$ , then after taking a subsequence if necessary, we can assume that  $f_n := Ax_n$  converges. However, we have

$$\|A^{\dagger}(f_n - f_m)\|^2 = \|A^{\dagger}A(x_n - x_m)\|^2 = \|P_{\mathcal{N}(A)^{\perp}}(x_n - x_m)\|^2 = \|x_n - x_m\|^2 = 2$$

and in particular the sequence  $(A^{\dagger}f_n)$  does not converge. This shows that  $A^{\dagger}$  is discontinuous.

In particular, combining this with theorem 1.8 shows that the range of a compact operator is always open, and that not every element in  $\mathcal{Y}$  has a least-squares solution.

We will need the following theorem, an infinite-dimensional analogue of the spectral theorem:

**Theorem 1.14** (Eigenvalue decomposition of self-adjoint compact operators). Let  $\mathcal{X}$  be a Hilbert space, and  $A \in \mathcal{K}(\mathcal{X}, \mathcal{X})$  self-adjoint. Then there exists an orthonormal basis  $(x_j)$  of  $\overline{\mathcal{R}(A)}$  and a sequence of eigenvalues  $|\lambda_1| \geq |\lambda_2| \geq \cdots > 0$  such that for all  $u \in \mathcal{X}$  we have

$$Au = \sum_{j=1}^{\infty} \lambda_j \langle u, x_j \rangle x_j.$$

The sequence  $(\lambda_i)$  is either finite or converges to 0.

The previous theorem gives rise to an infinite-dimensional analogue of the SVD:

**Theorem 1.15.** Let  $A \in \mathcal{K}(\mathcal{X}, \mathcal{Y})$ . Then there exists a (not necessarily infinite) sequence  $\sigma_1 \geq \sigma_2 \geq \cdots > 0$  converging to 0, and orthonormal bases  $(x_j)$ ,  $(y_j)$  of  $\mathcal{N}(A)^{\perp}$  and  $\overline{\mathcal{R}(A)}$  respectively, such that

$$Ax_j = \sigma_j y_j, \quad A^* y_j = \sigma_j x_j \quad \text{for all } j \in \mathbb{N},$$

and such that for all  $u \in \mathcal{X}$  and  $f \in \mathcal{Y}$  we have

$$Au = \sum_{j=1}^{\infty} \sigma_j \langle u, x_j \rangle y_j, \quad A^*f = \sum_{j=1}^{\infty} \sigma_j \langle f, y_j \rangle x_j.$$

The sequence  $\{(\sigma_j, x_j, y_j)\}$  is called the <u>singular value decomposition</u> (SVD) of A.

*Proof.* Define  $B := A^*A$  and  $C := AA^*$ , which are both compact, self-adjoint, and positive semi-definite operators. By the previous theorem, we can write

$$Cf = \sum_{j=1}^{\infty} \sigma_j^2 \langle f, y_j \rangle y_j,$$

where  $(y_j)$  is a basis of  $\overline{\mathcal{R}(AA^*)} = \overline{\mathcal{R}(A)}$  and  $(\sigma_j)$  is a positive decreasing sequence converging to 0. Note that

$$BA^*y_j = A^*AAy_j = A^*Cy_j = A^*\sigma^2y_j = \sigma_j^2A^*y_j,$$

so  $A^*y_j$  is an eigenvector of B with eigenvector  $\sigma_i^2$ .

We show that  $\left(\frac{A^*y_j}{\sigma_j}\right)$  is an orthonormal basis of  $\mathcal{R}(A)^{\perp}$ . is an orthonormal basis of  $\mathcal{N}(A)^{\perp}$ : their inner product is given by

$$\left\langle \frac{A^* y_j}{\sigma_j}, \frac{A^* y_k}{\sigma_k} \right\rangle = \frac{1}{\sigma_j \sigma_k} \langle y_j, C y_k \rangle = \frac{\sigma_k}{\sigma_j} \langle y_j, y_k \rangle = 0,$$

and since the  $(y_j)$  are a basis of  $\overline{\mathcal{R}(A)} = \mathcal{N}(A^*)^{\perp}$  it is clear that the span of  $(A^*y_j)$  is dense in  $\overline{\mathcal{R}(A^*)} = \mathcal{N}(A)^{\perp}$ .

If we choose  $x_j = \frac{A^* y_j}{\sigma_i}$ , we find by construction that  $A^* y_j = \sigma_j x_j$  and

$$Ax_j = \frac{AA^*y_j}{\sigma_i} = \frac{Cy_j}{\sigma_i} = \sigma_j y_j.$$

Finally, we see that

$$Au = \sum_{j=1}^{\infty} \langle u, x_j \rangle Ax_j = \sum_{j=1}^{\infty} \sigma_j \langle u, x_j \rangle y_j \quad \text{and} \quad A^*f = \sum_{j=1}^{\infty} \langle f, y_j \rangle A^*y_j = \sum_{j=1}^{\infty} \sigma_j \langle f, y_j \rangle x_j.$$

**Theorem 1.16.** Let  $A \in \mathcal{K}(\mathcal{X}, \mathcal{Y})$  with  $SVD \{(\sigma_i, x_i, y_i)\}$  and let  $f \in \mathcal{D}(A^{\dagger})$ . Then

$$A^{\dagger} f = \sum_{j=1}^{\infty} \frac{1}{\sigma_j} \langle f, y_j \rangle x_j.$$

Remark. Note that this is comparable to  $A^*f = \sum_{j=1}^{\infty} \sigma_j \langle f, y_j \rangle x_j$ , except that  $A^*$  is a smoothing operator (since  $\sigma_j \to 0$ ), while  $A^{\dagger}$  does the opposite. Furthermore,  $A^{\dagger}$  amplifies the right singular vectors corresponding to small singular values the most — intuitively, the corresponding left singular vectors are vectors where A doesn't "see much".

*Proof.* Define  $Bf = \sum_{j=1}^{\infty} \frac{1}{\sigma_j} \langle f, y_j \rangle x_j$ . Then by theorem 1.9, we must check that  $BA = P_{\mathcal{N}(A)^{\perp}}$  and  $AB = P_{\overline{\mathcal{R}(A)}} \upharpoonright_{\mathcal{D}(A^{\perp})}$ .

For the first equation, we compute

$$BAu = \sum_{j=1}^{\infty} \frac{1}{\sigma_j} \left\langle \sum_{i=1}^{\infty} \sigma_i \langle u, x_i \rangle y_i, y_j \right\rangle x_j = \sum_{j=1}^{\infty} \sum_{i=1}^{\infty} \frac{\sigma_i}{\sigma_j} \langle u, x_i \rangle \langle y_i, y_j \rangle x_j = \sum_{j=1}^{\infty} \langle u, x_j \rangle x_j.$$

Since  $(x_j)$  is a basis of  $\mathcal{N}(A)^{\perp}$ , this proves that  $BA = P_{\mathcal{N}(A)^{\perp}}$ .

For the second equation, an analogous computation gives  $ABf = \sum_{i=1}^{\infty} \langle f, y_i \rangle y_i$ , and since  $(y_i)$  is a basis of  $\overline{\mathcal{R}(A)}$ , this proves that  $AB = P_{\overline{\mathcal{R}(A)}} \upharpoonright_{\mathcal{D}(A^{\dagger})}$ .

**Definition 1.17.** Let  $A \in \mathcal{K}(\mathcal{X}, \mathcal{Y})$  have SVD  $\{(\sigma_j, x_j, y_j)\}$ . We say that  $f \in \mathcal{Y}$  satisfies the *Picard criterion* if

$$\sum_{j} \frac{\left| \langle f, y_j \rangle \right|^2}{\sigma_j^2} < \infty.$$

Note that the expression on the left corresponds to  $\|A^{\dagger}f\|^2$  if  $f \in \mathcal{D}(A^{\dagger})$ .

**Theorem 1.18.** Let  $f \in \overline{\mathcal{R}(A)}$ . Then  $f \in \mathcal{R}(A)$  if and only if f satisfies the Picard criterion.

*Proof.* ' $\Longrightarrow$ ' Write f = Au, then

$$\sum_{j} \frac{\left| \langle f, y_{j} \rangle \right|^{2}}{\sigma_{j}^{2}} = \sum_{j} \frac{\left| \langle Au, y_{j} \rangle \right|^{2}}{\sigma_{j}^{2}} = \sum_{j} \frac{\left| \langle u, A^{*}y_{j} \rangle \right|^{2}}{\sigma_{j}^{2}} = \sum_{j} \left| \langle u, x_{j} \rangle \right|^{2} < \infty.$$

'  $\longleftarrow$  ' Define  $u := \sum_{j=1}^{\infty} \frac{1}{\sigma_i} \langle f, y_j \rangle x_j$  (note that by assumption this sum converges). Then

$$Au = A\sum_{j=1}^{\infty} \frac{1}{\sigma_j} \langle f, y_j \rangle x_j = \sum_{j=1}^{\infty} \frac{1}{\sigma_j} \langle f, y_j \rangle Ax_j = \sum_{j=1}^{\infty} \langle f, y_j \rangle y_j = P_{\overline{\mathcal{R}(A)}} f = f,$$

so Au = f which implies  $f \in \mathcal{R}(A)$ .

We have seen that the stability of  $A^{\dagger}$  depends on the speed of decay of the singular values  $(\sigma_j)$ . We formalise this:

**Definition 1.19.** Let  $A \in \mathcal{K}(\mathcal{X}, \mathcal{Y})$  have singular values  $(\sigma_j)$ . Then the ill-posed inverse problem Au = f is called *mildly ill-posed* if the  $\sigma_j$  decay polynomially (i.e.,  $\frac{1}{\sigma_n} \leq Cn^{\gamma}$  for some  $C, \gamma$ ) and severely ill-posed otherwise.

Example 1.20. Consider the heat equation with initial conditions and boundary values:

$$\begin{cases} v_t - v_{xx} = 0 & (x,t) \in (0,\pi) \times \mathbb{R}_{>0}, \\ v(0,t) = v(\pi,t) = 0 & t \ge 0, \\ v(x,0) = u(x) & x \in (0,\pi), \\ v(x,T) = f(x) & x \in (0,\pi). \end{cases}$$

Then the forward problem is to determine f given u, while the inverse problem is to determine u given f. The solution for the foward problem is given by

$$f = Au := \sum_{j=1}^{\infty} e^{-j^2 T} \langle u, \sin(jx) \rangle \sin(jx),$$

and the eigenvalues are therefore  $\sigma_j = e^{-j^2T}$ . Since these clearly decay exponentially, this problem is severely ill-posed.

# 2 Classical regularisation theory

Let  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$  such that  $\mathcal{R}(A)$  is not closed (this happens for example when A is compact and does not have finite rank), and consider the inverse problem Au = f. Suppose we measure not f, but noisy data  $f_{\delta}$  such that  $||f_{\delta} - f|| \leq \delta$ . Then since  $A^{\dagger}$  is discontinuous, we cannot expect that  $A^{\dagger}f_{\delta} \to A^{\dagger}f$  as  $\delta \to 0$ . Therefore, we must replace  $A^{\dagger}$  by operators that approximate it.

**Definition 2.1.** Let  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$ . A family  $(R_{\alpha})_{\alpha>0}$  of continuous operators is called a *regularisation* of  $A^{\dagger}$  if

$$\lim_{\alpha \to 0} R_{\alpha} f = A^{\dagger} f \quad \text{for all } f \in \mathcal{D}(A^{\dagger}).$$

If all  $R_{\alpha}$  are linear (TODO: and bounded?), then we speak of a linear regularisation of  $A^{\dagger}$ .

**Theorem 2.2** (Banach-Steinhaus). Let  $\mathcal{X}, \mathcal{Y}$  be Hilbert spaces and  $\{A_{\alpha}\} \subseteq \mathcal{B}(\mathcal{X}, \mathcal{Y})$  a family of pointwise bounded operators. Then  $\{A_{\alpha}\}$  is bounded in norm.

Corollary 2.3. Let  $\mathcal{X}, \mathcal{Y}$  be Hilbert spaces and  $(A_j) \subseteq \mathcal{B}(\mathcal{X}, \mathcal{Y})$ . Then  $(A_j)$  converges pointwise to some  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$  if and only if  $\{A_j\}$  is norm-bounded and converges pointwise on some dense subset  $\mathcal{X}' \subseteq \mathcal{X}$ .

**Theorem 2.4.** Let  $\mathcal{X}, \mathcal{Y}$  be Hilbert spaces,  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$  and  $(R_{\alpha})_{\alpha>0}$  a linear regularisation. If  $A^{\dagger}$  is not continuous,  $(R_{\alpha})$  is not norm-bounded. In particular, there exists  $f \in \mathcal{Y}$  with  $||R_{\alpha}f|| \to \infty$ .

*Proof.* Suppose  $(R_{\alpha})$  is norm-bounded. Let  $\alpha_j \to 0$ , then we know that  $R_{\alpha_j} \to A^{\dagger}$  pointwise on  $\mathcal{D}(A^{\dagger})$ . Since  $\mathcal{D}(A^{\dagger})$  is dense in  $\mathcal{Y}$ , corollary 2.3 then tells us that  $A^{\dagger}$  is bounded and therefore continuous, a contradiction.

By the Banach-Steinhaus theorem, if  $(R_{\alpha})$  is not norm-bounded, it is not pointwise bounded, so there must exist  $f \in \mathcal{Y}$  such that  $\{\|R_{\alpha}f\|\}$  is not bounded.

**Recap 2.5.** Recall that any bounded sequence in a Hilbert space has a weakly convergent subsequence.

**Theorem 2.6.** Let  $A \in \mathcal{B}(\mathcal{X}, \mathcal{Y})$  and  $(R_{\alpha})$  a linear regularisation of  $A^{\dagger}$ . If  $\{\|AR_{\alpha}\|\}_{\alpha>0}$  is bounded, then  $\|R_{\alpha}f\| \to \infty$  as  $\alpha \to 0$  for every  $f \notin \mathcal{D}(A^{\dagger})$ .

*Proof.* Define  $u_{\alpha} := R_{\alpha} f$  for  $f \notin \mathcal{D}(A^{\dagger})$ , and assume there exists a sequence  $\alpha_k \to 0$  such that  $\{\|u_{\alpha_k}\|\}$  is bounded. After taking a subsequence if necessary, we may assume that  $u_{\alpha_k} \to u$  for some  $u \in \mathcal{X}$ , and therefore we also have  $Au_{\alpha_k} \to Au$ .

We also have  $\lim_{\alpha\to 0} AR_{\alpha}f = AA^{\dagger}f = P_{\overline{\mathcal{R}(A)}}f$  for  $f\in\mathcal{D}(A^{\dagger})$ , and since we assumed  $\{AR_{\alpha}\}$  was norm-bounded, by corollary 2.3 we have  $\lim_{\alpha\to 0} AR_{\alpha}f = P_{\overline{\mathcal{R}(A)}}f$  for all  $f\in\mathcal{Y}$ .

Since  $Au_{\alpha_k}$  is convergent and has weak limit Au, it must also have limit Au, so we find  $Au = P_{\overline{\mathcal{R}(A)}}f$  so  $f \in \mathcal{D}(A^{\dagger})$ , a contradiction.

We need some process to choose a parameter. To this end, note that we have

$$\left\| R_{\alpha} f_{\delta} - A^{\dagger} f \right\| \le \left\| R_{\alpha} (f_{\delta} - f) \right\| + \left\| (R_{\alpha} - A^{\dagger}) f \right\| \le \delta \|R_{\alpha}\| + \left\| (R_{\alpha} - A^{\dagger}) f \right\|. \tag{3}$$

The first term is called the *data error* and is unbounded for  $\alpha \to 0$ , and the second term is called the approximation error which does vanish for  $\alpha \to 0$ . Therefore, we want to choose  $\alpha$  small enough to have a low approximation error, while keeping the data error at bay.

## 2.1 Parameter choice rules

**Definition 2.7.** A function  $\alpha: \mathbb{R}_{>0} \times \mathcal{Y} \to \mathbb{R}_{>0}: (\delta, f_{\delta}) \mapsto \alpha(\delta, f_{\delta})$  is called a *parameter choice rule* (PCR). We distinguish three types:

- 1. An a priori PCR depends only on  $\delta$ ;
- 2. An a posteriori PCR depends on both  $\delta$  and  $f_{\delta}$ ;
- 3. A heuristic PCR depends only on  $f_{\delta}$ .

**Definition 2.8.** Let  $(R_{\alpha})_{\alpha>0}$  be a regularisation of  $A^{\dagger}$  and  $\alpha$  a parameter choice rule. We call  $(R_{\alpha}, \alpha)$  a convergent regularisation if

$$\lim_{\delta \to 0} \sup_{f_{\delta}: \|f - f_{\delta}\| \le \delta} \|R_{\alpha} f_{\delta} - A^{\dagger} f\| = 0$$

and

$$\lim_{\delta \to 0} \sup_{f_{\delta}: \|f - f_{\delta}\| \le \delta} \alpha(\delta, f_{\delta}) = 0.$$
(4)

#### 2.1.1 A priori parameter choice rules

We will not prove the following theorem, which guarantees the existence of a priori PCRs:

**Theorem 2.9.** Let  $(R_{\alpha})_{\alpha>0}$  be a regularisation of  $A^{\dagger}$ . Then there exists an a priori PCR  $\alpha=\alpha(\delta)$  such that  $(R_{\alpha}, \alpha)$  is convergent.

We can characterise PCRs in the following way:

**Theorem 2.10.** Let  $(R_{\alpha})_{\alpha>0}$  be a linear regularisation of  $A^{\dagger}$ , and  $\alpha=\alpha(\delta)$  an a priori PCR. Then  $(R_{\alpha}, \alpha)$  is convergent if and only if

$$\lim_{\delta \to 0} \delta \| R_{\alpha(\delta)} \| = 0 \quad and \quad \lim_{\delta \to 0} \alpha(\delta) = 0.$$

Proof. " $\Longrightarrow$ " Suppose  $(R_{\alpha}, \alpha)$  is convergent. It is clear that  $\lim_{\delta \to 0} \alpha(\delta) = 0$  by eq. (4). Suppose  $\lim_{\delta \to 0} \delta \|R_{\alpha(\delta)}\| \neq 0$ . Then there exists a sequence  $(\delta_k) \to 0$  and a constant C > 0 such that  $\delta_k \|R_{\alpha(\delta_k)}\| \geq C$  for all k. This implies we can find a sequence  $(g_k) \subseteq \mathcal{Y}$  with  $\|g_k\| = 1$  and  $\delta_k \|R_{\alpha(\delta_k)}g_k\| \geq C$  for all k.

Now let  $f \in \mathcal{D}(A^{\dagger})$  and define  $f_k := f + \delta_k g_k$ , then clearly we have  $f_k \to f$ , but also

$$C \le \left\| R_{\alpha(\delta_k)}(\delta_k g_k) \right\| = \left\| R_{\alpha(\delta_k)}(f_{\delta_k} - f) \right\| \le \left\| R_{\alpha(\delta_k)} f_{\delta_k} - A^{\dagger} f \right\| + \left\| (R_{\alpha(\delta_k)} - A^{\dagger}) f \right\|.$$

In particular we find that  $\|(R_{\alpha(\delta_k)} - A^{\dagger})f\| \ge C$ , so clearly  $R_{\alpha}$  is not convergent. " $\Leftarrow$ " This follows immediately from eq. (3).

A problem with a priori PCRs is that they are scale-invariant: if  $\alpha = \alpha(\delta)$  gives a convergent regularisation, then  $\hat{\alpha} = \alpha(k\delta)$  also gives a convergent regularisation for any k. In practice, it is not always clear which scale should be chosen.

#### 2.1.2 A posteriori parameter choice rules

Let  $f \in \mathcal{D}(A^{\dagger})$  and  $f_{\delta}$  s.t.  $||f - f_{\delta}|| \leq \delta$ . Letting  $u^{\dagger}$  denote the minimum-norm solution of the problem Au = f, and defining  $\mu := ||Au^{\dagger} - f|| = \inf_{u \in \mathcal{X}} ||Au - f||$ , we see that

$$||Au^{\dagger} - f_{\delta}|| \le ||Au^{\dagger} - f|| + ||f - f_{\delta}|| \le \mu + \delta.$$

Therefore, it is not useful to choose  $\alpha(\delta, f_{\delta})$  with  $||Au_{\alpha} - f_{\delta}|| < \mu + \delta$ : if this is the case, we are most likely overfitting.

This motivates *Morozov's discrepancy principle*:

**Definition 2.11.** Let  $(R_{\alpha})$  be a (TODO: linear?) regularisation of  $A^{\dagger}$  and assume  $\mathcal{R}(A)$  is dense in  $\mathcal{Y}$ . Fix  $\eta > 1$ , and define

$$\alpha(\delta, f_{\delta}) = \sup \{ \alpha > 0 : ||AR_{\alpha}f_{\delta} - f_{\delta}|| \le \eta \delta \}.$$

Then  $\alpha(\delta, f_{\delta})$  is said to satisfy Morozov's discrepancy principle.

It can be shown that the above  $\alpha$  indeed gives a convergent regularisation.

### 2.1.3 Heuristic parameter choice rules

Heuristic parameter choice rules unfortunately only work if the original problem was well-posed:

**Theorem 2.12** (Bakushinskii). Let  $(R_{\alpha})$  be a regularisation of  $A^{\dagger}$  and suppose there exists a heuristic parameter choice rule  $\alpha$  such that  $(R_{\alpha}, \alpha)$  is convergent. Then  $A^{\dagger}$  is continuous from  $\mathcal{Y}$  to  $\mathcal{X}$ .

# 2.2 Spectral regularisation

We will now start with specific examples of regularisations. Spectral regularisations are derived from the spectral decomposition

$$A^{\dagger} f = \sum_{j=1}^{\infty} \sigma_j^{-1} \langle f, y_j \rangle x_j.$$

We construct a regularisation by replacing  $\sigma_j^{-1}$  by some function  $g_{\alpha}(\sigma_j)$ , i.e.,

$$R_{\alpha}f = \sum_{j=1}^{\infty} g_{\alpha}(\sigma_j) \langle f, y_j \rangle x_j.$$
 (5)

Let us explore which conditions  $g_{\alpha}$  must satisfy:

**Theorem 2.13.** Let, for  $\alpha > 0$ , the function  $g_{\alpha} : \mathbb{R}_{>0} \to \mathbb{R}_{>0}$  satisfy

- 1.  $\lim_{\alpha \to 0} g_{\alpha}(\sigma) = \frac{1}{\sigma}$  for all  $\sigma > 0$ ;
- 2.  $g_{\alpha}(\sigma) \leq C_{\alpha}$  for some  $C_{\alpha} > 0$ ;
- 3.  $\sup_{\alpha,\sigma} \sigma g_{\alpha}(\sigma) \leq \gamma \text{ for some } \gamma > 0.$

Then collection  $(R_{\alpha})$  defined by eq. (5) is a linear regularisation of  $A^{\dagger}$ , and in particular, we have  $||R_{\alpha}|| \leq C_{\alpha}$ .

*Proof.* From condition 2 it follows that all  $R_{\alpha}$  are bounded. Since

$$\langle f, y_j \rangle = \langle P_{\overline{\mathcal{R}(A)}} f, y_j \rangle = \langle AA^{\dagger} f, y_j \rangle = \langle A^{\dagger} f, A^* y_j \rangle = \sigma_j \langle u^{\dagger}, x_j \rangle,$$

we compute

$$(R_{\alpha} - A^{\dagger})f = \sum_{j} (g_{\alpha}(\sigma_{j}) - \sigma_{j})^{-1} \langle f, y_{j} \rangle x_{j} = \sum_{j} (\sigma_{j}g_{\alpha}(\sigma_{j}) - 1) \langle u^{\dagger}, x_{j} \rangle x_{j},$$

and since  $\sigma g_{\alpha}(\sigma_i) \leq \gamma$ , we have  $(\sigma_i g_{\alpha}(\sigma_i) - 1)^2 \leq 1 + \gamma^2$ , so that

$$\|(R_{\alpha} - A^{\dagger})f\|^{2} \le (1 + \gamma^{2})\|u^{\dagger}\|^{2} < \infty.$$

Since  $\|(R_{\alpha} - A^{\dagger})f\|$  is finite, we may apply the reverse Fatou lemma to the sum and obtain

$$\limsup_{\alpha \to 0} \left\| (R_{\alpha} - A^{\dagger}) f \right\|^2 \le \sum_{j} \left( \sigma_{j} \limsup_{\alpha \to 0} g_{\alpha}(\sigma_{j}) - 1 \right)^2 \langle u^{\dagger}, x_{j} \rangle^2 = 0,$$

and therefore  $R_{\alpha}f \to A^{\dagger}f$  as  $\alpha \to 0$ .

**Example 2.14.** The first, very simple example is the truncated SVD: we simply define

$$g_{\alpha}(\sigma) = \begin{cases} 1/\sigma & \sigma \ge \alpha, \\ 0 & \sigma < \alpha. \end{cases}$$

It is easy to check that  $g_{\alpha}$  satisfies the conditions of theorem 2.13, and that all  $R_{\alpha}$  are finite-rank operators with  $||R_{\alpha}|| \leq \frac{1}{\alpha}$ . Therefore, if we choose  $\alpha = \alpha(\delta)$  such that  $\delta/\alpha(\delta) \to 0$ , then we obtain a convergent regularisation.

This also highlights the problem with this method: as  $\delta$  gets smaller, we need more and more singular vectors which are generally expensive to compute.

**Example 2.15.** The second example is *Tikhonov regularisation*. Here, we define  $g_{\alpha}(\sigma) = \frac{\sigma}{\sigma^2 + \alpha}$ , and again it is easily checked that the conditions of theorem 2.13 are satisfied, noting that

$$\frac{\sigma}{\sigma^2 + \alpha} \le \frac{\sigma}{2\sigma\sqrt{\alpha}} = \frac{1}{2\sqrt{\alpha}} =: C_{\alpha}.$$

Therefore, if  $\delta/\sqrt{\alpha(\delta)} \to 0$ , the regularisation is convergent.

This method does not require computing the SVD of A: it is easily shown that  $u_{\alpha} := R_{\alpha}f$  is the unique solution to the regularised normal equation

$$(A^*A + \alpha I)u_{\alpha} = A^*f.$$

While  $A^*A + \alpha I$  is always invertible, computing the inverse is expensive, so we usually use some approximation of the inverse.

Finally, it can also be shown that

$$u_{\alpha} = \min_{u \in \mathcal{X}} \|Au - f\|^2 + \alpha \|u\|^2,$$

so we can also view  $u_{\alpha}$  as the solution of an optimisation problem.