

- a.
  - i. The complexity of `pushright(x)` is  $O(n)$   
The complexity of `popleft()` is  $O(1)$   
The complexity of `element_at(i)` is  $O(i)$  which in general is  $O(n)$
  - ii. This does not satisfy the fundamental inequality of amortized cost. The aggregate true cost of  $n$  calls to `element_at(n)` is  $kn^2$  which is  $O(n^2)$  for some constant  $k$ . However the proposed aggregate amortized cost for the same actions would be  $O(n)$ . Therefore there always exists some  $n$  for which the true aggregate cost is greater than the proposed aggregate amortized cost, and so the proposal must be incorrect.
- b.
  - i.

```

class RandomAccessQueue:
    array = empty array of length 8
    head_pointer = 0
    tail_pointer = 0

    def pushright(x):
        array[tail_pointer] = x
        tail_pointer++
        if (tail_pointer == array.length):
            new_array = empty array of length (array.length * 2)
            for (int i=0; i<tail_pointer-head_pointer; i++) {
                new_array[i] = array[head_pointer+i]
            }
            array = new_array
            tail_pointer -= head_pointer
            head_pointer = 0

    def popleft():
        if (head_pointer >= tail_pointer):
            throw Exception("The queue is empty")
            # or return null
        else:
            x = array[head_pointer]
            head_pointer++
            return x

    def element_at(i):
        index = head_pointer+i
        if (index >= tail_pointer):
            throw Exception("Index out of bounds")
            # or return null
        else:
            return array[index]

```

- ii. The true cost of `element_at(i)` is  $O(1)$ . The true cost of `popleft()` is  $O(1)$ . For `pushright(x)`, the cost of copying the array into a larger one is  $kn$  where  $n$  is the initial size of the array and  $k$  is a constant. After  $n$  calls to `pushright`, the total cost is  $n$  (for the inserting) +  $k(1+2+4+8+\dots+2^{\lfloor \log_2(n-1) \rfloor})$  which is strictly less than  $k(2n-3)+n$  which is  $O(n)$ . Therefore the amortized cost of a single call to `pushright(x)` is  $O(1)$ . Therefore all of the function calls have amortized cost  $O(1)$ .