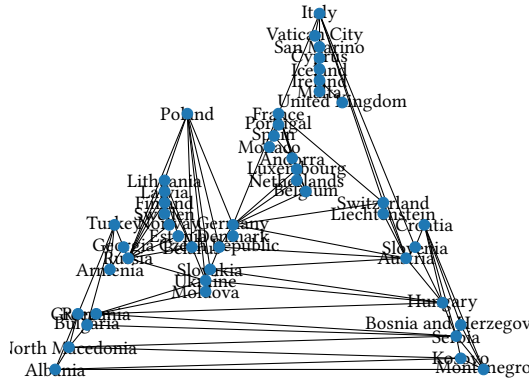


- The graph¹ of Europe $\mathcal{G}^* = \langle V, E \rangle$ is defined as follows: each vertex $v \in V$ is a Europe country²; two vertices are adjacent ($\{u, v\} \in E$) if the corresponding countries share a land border. Let \mathcal{G} be the largest connected component of \mathcal{G}^* .

- Draw \mathcal{G}^* with the minimum number of intersecting edges.



- Find $|V|$, $|E|$, $\delta(\mathcal{G})$, $\Delta(\mathcal{G})$, $\text{rad}(\mathcal{G})$, $\text{diam}(\mathcal{G})$, $\text{girth}(\mathcal{G})$, $\text{center}(\mathcal{G})$, $\kappa(\mathcal{G})$, $\lambda(\mathcal{G})$.
 - Find the minimum vertex coloring $Z : V \rightarrow \mathbb{N}$ of \mathcal{G} .
 - Find the minimum edge coloring $X : E \rightarrow \mathbb{N}$ of \mathcal{G} .
 - Find the maximum clique $Q \subseteq V$ of \mathcal{G} .
 - Find the maximum stable set $S \subseteq V$ of \mathcal{G} .
 - Find the maximum matching $M \subseteq E$ of \mathcal{G} .
 - Find the minimum vertex cover $R \subseteq V$ of \mathcal{G} .
 - Find the minimum edge cover $F \subseteq E$ of \mathcal{G} .
 - Find the shortest closed walk W that visits every vertex of \mathcal{G} .
 - Find the shortest closed walk U that visits every edge of \mathcal{G} .
 - Find all biconnected components (blocks) and draw the block-cut tree of \mathcal{G}^* .
 - Find all 2-edge-connected components of \mathcal{G}^* .
 - Construct an SPQR tree of the largest biconnected component of \mathcal{G}^* .
 - Add the weight function $w : E \rightarrow \mathbb{R}$ denoting the distance between capitals. Find the minimum (w.r.t. the total weight of edges) spanning tree T for the largest connected component of the weighted Europe graph $\mathcal{G}_w^* = (V, E, w)$.
 - Find $\text{centroid}(T)$ (w.r.t. the edge weight function w).
 - Construct the Prüfer code for T .
- Prove *rigorously* the following theorems:

Theorem 1 (TRIANGLE INEQUALITY). For any connected graph $G = \langle V, E \rangle$:

$$\forall x, y, z \in V : \text{dist}(x, y) + \text{dist}(y, z) \geq \text{dist}(x, z)$$

Theorem 2 (TREE). A connected graph $G = \langle V, E \rangle$ is a tree (i.e. acyclic graph) iff $|E| = |V| - 1$.

Theorem 3 (WHITNEY). For any graph $G : \kappa(G) \leq \lambda(G) \leq \delta(G)$.

¹ Hereinafter, “graphs” are “simple, finite, undirected and unweighted”, unless stated otherwise.

² https://simple.wikipedia.org/wiki/List_of_European_countries used as reference.

Theorem 4 (CHARTRAND). For a connected graph $G = \langle V, E \rangle$: if $\delta(G) \geq \lfloor \frac{|V|}{2} \rfloor$, then $\lambda(G) = \delta(G)$.

Theorem 5 (Menger). For any pair of non-adjacent vertices u and v in an undirected graph, the size of the minimum *vertex cut* is equal to the maximum number of pairwise *internally vertex-disjoint paths* from u to v .

Theorem 6 (HARARY). Every block of a block graph³ is a clique.

³ A block graph $H = B(G)$ is an intersection graph of all blocks (biconnected components) of G , *i.e.* each vertex $v \in V(H)$ corresponds to a block of G , and there is an edge $\{v, u\} \in E(H)$ iff “blocks” v and u share a cut vertex.