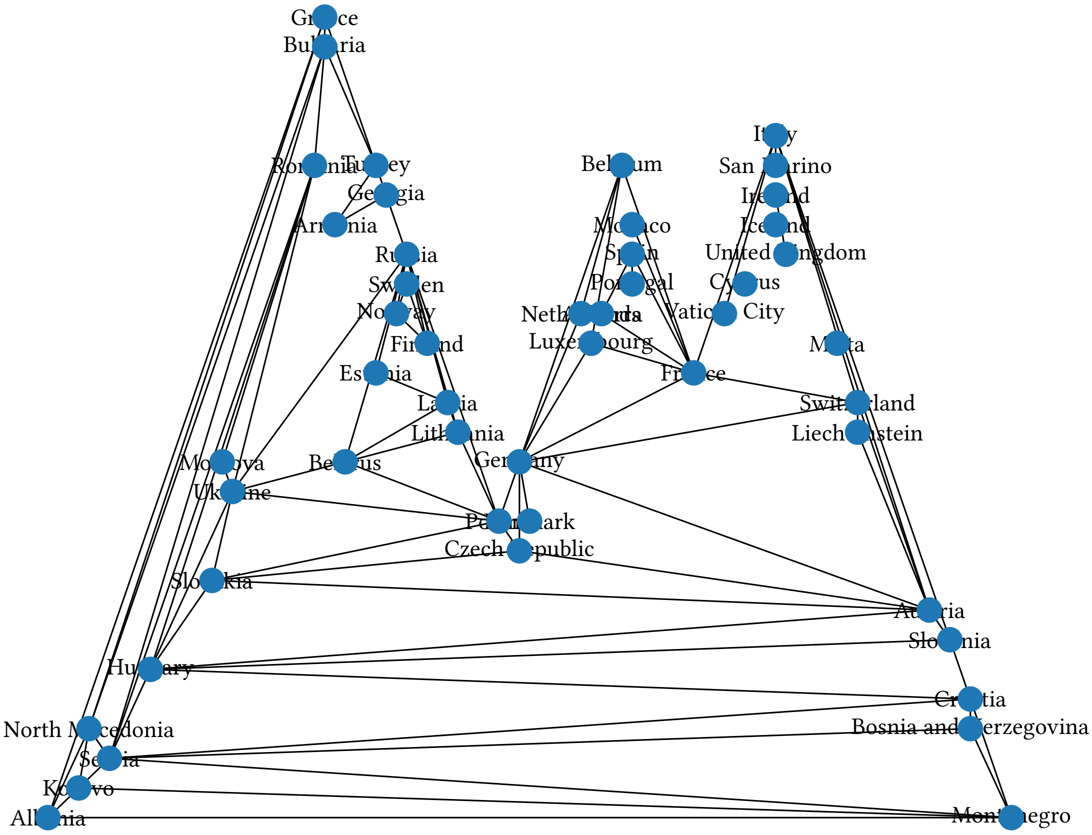
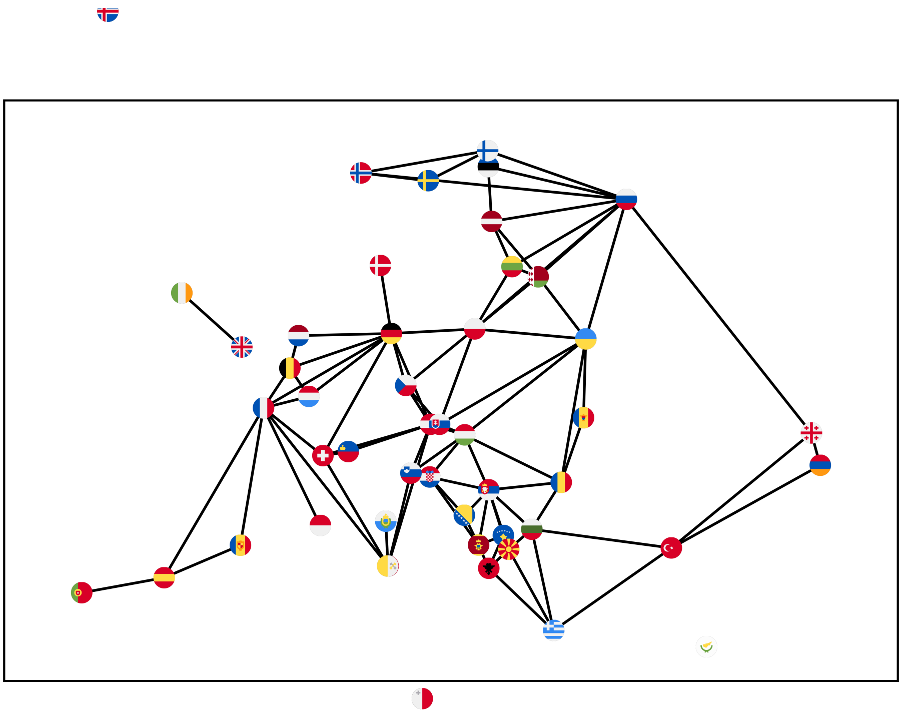
1. The graph[1] of Europe $\mathcal{G}^* = \langle V, E\rangle$ is defined as follows: each vertex $v \in V$ is a Europe country[2]; two vertices are adjacent ($\{u, v\} \in E$) if the corresponding countries share a land border. Let $\mathcal{G}$ be the largest connected component of $\mathcal{G}^*$.

   (a) Draw $\mathcal{G}^*$ with the minimum number of intersecting edges.
   The graph was created in Python. More about creation can be found here.
   A planar version of the graph (ugly):



[1] Hereinafter, "graphs" are "simple, finite, undirected and unweighted", unless stated otherwise.
[2] `https://simple.wikipedia.org/wiki/List_of_European_countries` used as reference.

A prettier version (but unfinished):



(b) Find $|V|$, $|E|$, $\delta(\mathcal{G})$, $\Delta(\mathcal{G})$, $\mathrm{rad}(\mathcal{G})$, $\mathrm{diam}(\mathcal{G})$, $\mathrm{girth}(\mathcal{G})$, $\mathrm{center}(\mathcal{G})$, $\varkappa(\mathcal{G})$, $\lambda(\mathcal{G})$.
$|V| = 49$, $|E| = 92$

(c) Find the minimum vertex coloring $Z : V \longrightarrow \mathbb{N}$ of $\mathcal{G}$.

(d) Find the minimum edge coloring $X : E \longrightarrow \mathbb{N}$ of $\mathcal{G}$.

(e) Find the maximum clique $Q \subseteq V$ of $\mathcal{G}$.

(f) Find the maximum stable set $S \subseteq V$ of $\mathcal{G}$.

(g) Find the maximum matching $M \subseteq E$ of $\mathcal{G}$.

(h) Find the minimum vertex cover $R \subseteq V$ of $\mathcal{G}$.

(i) Find the minimum edge cover $F \subseteq E$ of $\mathcal{G}$.

(j) Find the shortest closed walk W that visits every vertex of $\mathcal{G}$.

(k) Find the shortest closed walk U that visits every edge of $\mathcal{G}$.

(l) Find all biconnected components (blocks) and draw the block-cut tree of $\mathcal{G}^*$.

(m) Find all 2-edge-connected components of $\mathcal{G}^*$.

(n) Construct an SPQR tree of the largest biconected component of $\mathcal{G}^*$.

(o) Add the weight function $w : E \longrightarrow \mathbb{R}$ denoting the distance between capitals. Find the minimum (*w.r.t.* the total weight of edges) spanning tree T for the largest connected component of the weighted Europe graph $\mathcal{G}^*_w = (V, E, w)$.

(p) Find $\mathrm{centroid}(T)$ (*w.r.t.* the edge weight function w).

(q) Construct the Prüfer code for T.

2. Prove *rigorously* the following theorems:

Theorem 1 (Triangle Inequality). For any connected graph $G = \langle V, E \rangle$:

$$\forall x, y, z \in V : \mathrm{dist}(x, y) + \mathrm{dist}(y, z) \geq \mathrm{dist}(x, z)$$

Theorem 2 (Tree). A connected graph $G = \langle V, E \rangle$ is a tree (*i.e.* acyclic graph) *iff* $|E| = |V| - 1$.

Theorem 3 (Whitney). For any graph $G : \varkappa(G) \leq \lambda(G) \leq \delta(G)$.

Theorem 4 (Chartrand). For a connected graph $G = \langle V, E \rangle$: if $\delta(G) \geq \lfloor \frac{|V|}{2} \rfloor$, then $\lambda(G) = \delta(G)$.

Theorem 5 (Menger). For any pair of non-adjacent vertices u and v in an undirected graph, the size of the minimum *vertex cut* is equal to the maximum number of pairwise *internally vertex-disjoint paths* from u to v.

Theorem 6 (Harary). Every block of a block graph[3] is a clique.

[3] A block graph $H = \mathrm{B}(G)$ is an intersection graph of all blocks (biconnected components) of G, *i.e.* each vertex $v \in V(H)$ corresponds to a block of G, and there is an edge $\{v, u\} \in E(H)$ iff "blocks" v and u share a cut vertex.