

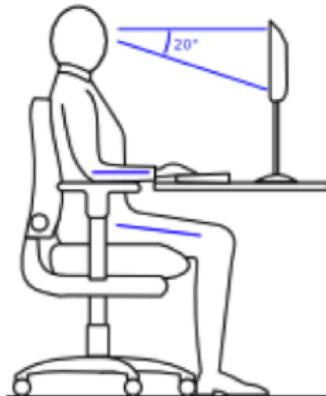
# An Introduction to High Performance Computing

Stuart Rankin

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Research Computing Services (<http://www.hpc.cam.ac.uk/> <http://www.csd3.cam.ac.uk/>)  
University Information Services (<http://www.uis.cam.ac.uk/>)

# Health and Safety



# Welcome

- ▶ Please sign in on the **attendance sheet**.
- ▶ Please fill in the **online feedback** at the end of the course:  
<http://feedback.training.cam.ac.uk/ucs/form.php>
- ▶ Keep your belongings with you.
- ▶ Please ask questions and let us know if you need assistance.

# Plan of the Course

Part 1: Introduction

Part 2: Cluster Basics

Part 3: Advanced HPC

09:30	WELCOME
10:00-11:00	Introduction and basics
11:00-11:30	Practical and break
12:00-12:30	Practical
12:30-13:30	LUNCH
14:00-14:30	Practical
14:45-15:15	Practical
15:30-CLOSE	Further discussion

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## **Part I: Introduction**

# Basics: Outline

Who are we?

Training Accounts

CSD 3 and Login node for today

Security

Basics: Pre-requisites

Connecting

Basic practicals

Connecting from other clients

Your trainers for today will be:

- ▶ Paul Sumption: Research Computing Technical Liaison
- ▶ Arjen Tumerus: Research Software Engineer

## Basics: Training accounts

- ▶ For our practical exercise's we will use HPC training accounts.
- ▶ You will find two pieces of paper on your desk.
- ▶ 1: A terms and conditions form for you to sign.
- ▶ 2: Your training account details.
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- ▶ The main difference is the login node names and slightly different submit scripts as CSD 3 has more node types.

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- ▶ Boring but very, very important...
- ▶ Cambridge IT is under constant attack by would-be intruders.
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2. Always choose strong passwords.
3. Your UIS password is used for multiple systems so keep it secure!
4. Keep the software on your laptops/tablets/PCs up to date this includes home computers especially if you are using the VPN to connect in.
5. Don't share accounts (this is against the rules anyway).

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# Pre-requisites

- ▶ A pre-requisite of this course:
- ▶ Basic Unix/Linux command line experience
- ▶ <https://www.training.cam.ac.uk/ucs/Course/ucs-unixintro1>  
Unix: Introduction to the Command Line Interface (Self-paced)
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## Basics: Connecting

- ▶ When connecting to the cluster will we use the SSH secure protocol only.
- ▶ We will use the Linux workstations during this course
- ▶ Please check your workstation is booted into Ubuntu Linux, ask if you need help with this.
- ▶ You are welcome to use your own laptop, however you may need to install some software in order to connect.
- ▶ Later in this section of the slides we will cover the software you may need to install on your own computer.

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Connection from home possible via the VPN service  
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or SSH tunnel through a departmental gateway.

# Connecting: Linux Clients

- ▶ ssh, scp, sftp, rsync  
*Installed (or installable), in Ubuntu we will use 'Terminal'.*
- ▶ X Windows, for using graphical applications remotely *This is already installed on your desktop.*

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# Connecting: Login

- ▶ From Linux/MacOSX/UNIX (or Cygwin):  
`ssh -Y abc123@login.hpc.cam.ac.uk`
- ▶ From graphical clients:  
Host: `login.hpc.cam.ac.uk`  
Username: `abc123` (your UCAM account name)
- ▶ `login.hpc` will map to a random login node  
i.e. one of `login-sand1`, `login-sand2`, ..., `login-sand8`  
NB Not `darwin.hpc` (the head node).
- ▶ Non-registered addresses will fail with “Connection refused”.
- ▶ Similarly for other systems (e.g. `cardio-login.hpc`,  
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## Connecting: First time login

- ▶ The first connection to a particular hostname produces the following:

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The authenticity of host 'login-sand2.hpc.cam.ac.uk (131.111.1.214)'  
can't be established.
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RSA key fingerprint is

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0b:ef:59:90:fb:13:4a:c9:56:82:7b:cd:4b:2b:e1:3b.
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Are you sure you want to continue connecting (yes/no)? yes

Warning: Permanently added 'login-sand2.hpc.cam.ac.uk' (RSA) to the  
list of known hosts.

- ▶ One should always check the fingerprint before typing “yes”.
- ▶ Graphical SSH clients *should* ask a similar question.
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# Connecting: First time login

- ▶ You may be presented with any of the following fingerprints (depending on your client):

MD5:0b:ef:59:90:fb:13:4a:c9:56:82:7b:cd:4b:2b:e1:3b

SHA256:sSkVfzpwpjwiFvxLcdPoDpN8IsN3kt0ZSywhDhPKZPAG

MD5:34:9b:f2:d2:c6:b3:5c:63:99:b7:27:da:5b:c8:16:fe

SHA256:HsiY10e0M8tS6JwR76PeQQA/VB7r8675BzG50YQ4h34

MD5:64:7c:7c:ff:05:9d:0e:dc:06:fe:f1:c2:10:37:7a:85

SHA256:wq91jBfPa71XXpQq+rk5JTBXLJ0/kXj0c5A7rp4ENzA

## Basics: A practical example before some theory

- ▶ Some simple exercises will help us gauge your experience.
- ▶ Exercise 1: Login with SSH.
- ▶ Exercise 2: Navigating the command line.
- ▶ Exercise 3: SFTP file transfer.

# Navigating your terminal

Useful commands for navigating your terminal.

- ▶ `cd <dirname >` - change into a directory
- ▶ `ls <dirname >` - list the contents of a directory
- ▶ `cd` or `cd ~` - change into your home folder
- ▶ `cd ..` - change back one folder
- ▶ `man ls` - will bring up the manual page for the ls command
- ▶ `pwd` - print working directory

## Exercise 1: Navigating your terminal

- ▶ Start a terminal by double clicking on the terminal icon
- ▶ Try the `ls` and `cd` commands in a terminal.
- ▶ Look at the man page for the `ls` command
- ▶ Close the terminal

## Exercise 2: Login

Using a Linux terminal you will login to the cluster with your HPC training account.

- ▶ Start the terminal by double clicking on the terminal icon
- ▶ In your terminal enter:
- ▶ `ssh -Y abc123@login.hpc.cam.ac.uk`  
Replace abc123 with your training account username
- ▶ Enter your password as supplied on the sheet
- ▶ Leave this terminal open, you will need it for exercise 3!

## Exercise 3: Transfer some files

You will need to transfer the exercise files to the cluster.

- ▶ Open a second Linux terminal on your training computer.
- ▶ Enter this command: `cd ~\Course_material`
- ▶ Check the file 'exercises.tar.gz' is in your directory listing
- ▶ Hint: `ls`

## Exercise 3: Transfer some files

Transfer the exercises.tar.gz to your HPC home folder.

- ▶ `sftp abc123@login.hpc.cam.ac.uk`  
Change abc123 to your training account username
- ▶ The command: `put exercises.tar.gz` will transfer the file from your local computer to the remote one
- ▶ In terminal logged into the cluster type: `ls`
- ▶ Use: `tar -xvf exercise.tar -` to unzip the file

## Connecting: File Transfer

In exercise 3 we used sftp. Another tool to be familiar with is rsync especially when you want to transfer lots of directories and files.

- ▶ rsync is a powerful tool and has many features.
  - \* You can rerun it to update or resume after interruption.
  - \* All transfers are checksummed.

# Connecting: File Transfer

- ▶ From Linux/MacOSX/UNIX (or Cygwin):

`rsync -av old_directory/ abc123@login.hpc.cam.ac.uk:scratch/new_directory`  
copies contents of old\_directory to `~/scratch/new_directory`.

`rsync -av old_directory abc123@login.hpc.cam.ac.uk:scratch/new_directory`  
copies old\_directory (and contents) to  
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## Basics: Connecting from other clients

- ▶ When using your own computer you may need to install some software in order to connect.
- ▶ There are quite a few choices of software packages for this, we will just cover a few.
- ▶ Our website  
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Installed (or installable) OS X has a native terminal package. This can be launched by clicking Apple >Go >Utilities and then clicking the Terminal icon.

- ▶ TurboVNC (for remote desktop, 3D optional)

<http://sourceforge.net/projects/turbovnc/files/>

- ▶ On MacOSX, install XQuartz to display remote graphical applications.

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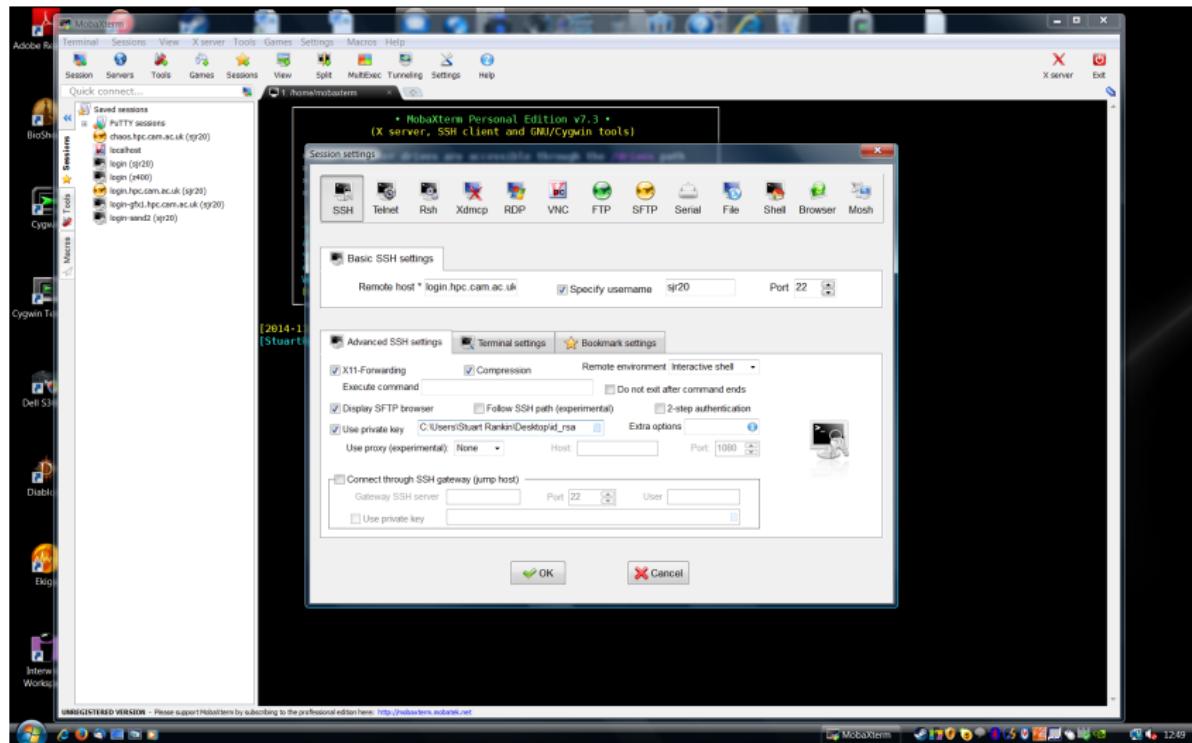
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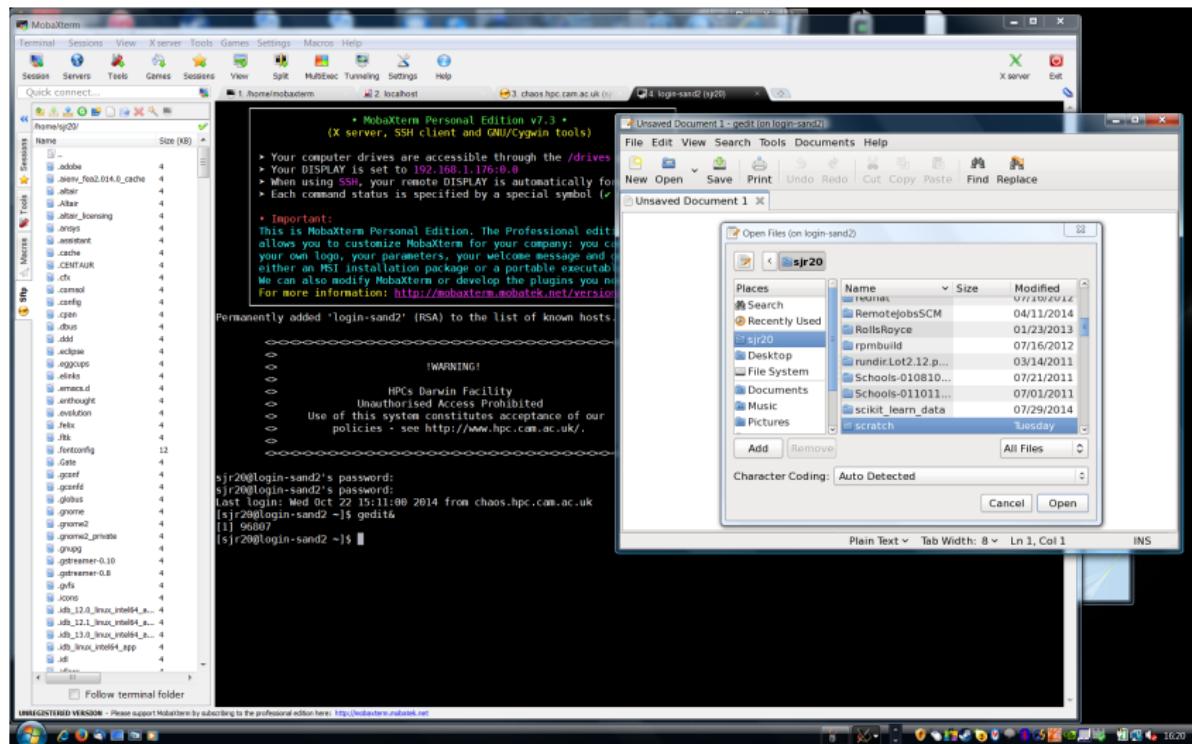
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# MobaXterm SSH (Windows)



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## **Part II: HPC Basics**

# HPC Basics: Our Clusters

- ▶ Our clusters are built using lots of commodity servers which then operate as a 'super computer'.
- ▶ We have CPU and GPU cluster nodes.
- ▶ A cluster has a scheduler which runs jobs from a queue.
- ▶ You submit jobs to the queue using a submission script.
- ▶ Jobs have service levels and QOS (quality of service) associated with them.
- ▶ There is a user environment this allows you to load or unload versions of software.
- ▶ We will look at each of these aspects in more detail during the course.

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# HPC Basics: Why use an HPC cluster?

There are many reasons why you may need to use a big computer or cluster.

*Limited local resources:* Your jobs can no longer run on your own computer or run very slowly.

*Storage Intensive:* Your data will exceed that of your local storage capacity.

*Software:* Many of the software packages you require have a complex installation process or require specialist support.

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# HPC Basics: Types of problem part 1.

In this section we will cover high throughput and compute intensive jobs. We will introduce the user and module environment and submit some High Throughput jobs to some CPU nodes.

*High Throughput:* Many unrelated problems to be executed in bulk.

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# HPC Basics: High Throughput

- ▶ Distribute **independent**, **multiple problems** across multiple CPUs to reduce the overall execution time.
- ▶ Workload is trivially (or *embarrassingly*) parallel:
  - \* Workload breaks up naturally into *independent* pieces.
  - \* Each piece is performed by a separate process/thread on a separate CPU (concurrently).
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# HPC Basics: Data Intensive Problems

- ▶ Distribute the **data** for a **single problem** across multiple CPUs to reduce the overall execution time.
- ▶ The *same* work may be done on each data segment.
- ▶ Rapid movement of data to and from disk is more important than inter-CPU communication.
- ▶ **Big Data** problems of great current interest -
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# HPC Basics: User Environment

- ▶ The user environment can be changed by loading modules.
- ▶ At login some default modules are loaded.
- ▶ When you start to compile software or write your own programs you will need to know more about your user environment and the underlying hardware.

# HPC Basics: User Environment

- ▶ Scientific Linux 6.8 ([Red Hat Enterprise Linux 6.8 rebuild](#))
  - ▶ bash
  - ▶ GNOME2 or XFCE4 desktop ([if you want](#))
- ▶ Lustre (patched), Mellanox OFED, CUDA
- ▶ But you don't need to know that.
- ▶ Upgrade to Scientific Linux/Red Hat Enterprise Linux 7 underway (Wilkes and new CSD 3 nodes are already on SL7).

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- ▶ Multi-petabytes split across multiple filesystems with tape.
- ▶ Lustre cluster filesystem:
  - \* Multiple RAID6 back-end disk volumes.
  - \* Multiple object storage servers.
  - \* Single metadata server.
  - \* Tape-backed HSM on newest filesystems.
  - \* 4 GB/sec overall read or write.
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# User Environment: Filesystems

When you apply for an HPC account a home directory is created for you.

- ▶ [/home/abc123](#)

- ▶ 40GB quota.
- ▶ Visible equally from all nodes.
- ▶ Single storage server.
- ▶ Hourly, daily, weekly snapshots copied to tape.
- ▶ Not intended for job outputs or large/many input files.

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# Our storage services

- ▶ We have several storage services for users that need to exceed 1TB.
- ▶ <http://www.uis.cam.ac.uk/initiatives/storage-strategy/storage-services>
- ▶ The most relevant services to HPC are RCS and RDS.
- ▶ RCS - Research Cold Store is for data that isn't changing, data goes to disk then two sets of tapes.
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Usage on /scratch (lfs quota -u abc123 /scratch):
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Disk quotas for user abc123 (uid 456):
      Filesystem  kbytes  quota  limit  grace  files  quota  limit  grace
      /scratch/abc123  9298708  1073741824  1181116006  -    165588      0      0      -
...
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- ▶ Aim to stay below the soft limit (*quota*).
- ▶ Once over the soft limit, you have 7 days grace to return below.
- ▶ When the grace period expires, or you reach the hard limit (*limit*), no more data can be written.
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# Filesystems: Backups

- ▶ Disk snapshots and tape (as of May 2017).
- ▶ They are not an undelete - take care when deleting.
- ▶ Successful restoration depends on:
  - ▶ The file having existed long enough to have been backed up at all.
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- ▶ Scratch files are not backed up.

# Filesystems: Automounter

- ▶ Directories under /scratch are **automounted**.
- ▶ They only appear under /scratch when explicitly referenced.
- ▶ Thus when browsing /scratch may appear too empty
  - use *ls* or *cd* to reference /scratch/abc123 explicitly.

# Filesystems: Permissions

- ▶ Be careful and if unsure, please ask [support@hpc.cam.ac.uk](mailto:support@hpc.cam.ac.uk).
  - ▶ Can lead to accidental destruction of your data or account compromise.
- ▶ Avoid changing the permissions on your home directory.
  - ▶ Files under /home are particularly security sensitive.
  - ▶ Easy to break passwordless communication between nodes.

# Using HPC: Software

- ▶ Free software accompanying Red Hat Enterprise is (or can be) provided.
- ▶ Other software (free and non-free) is available via modules.
- ▶ Some proprietary software may not be generally accessible.
- ▶ See <http://www.hpc.cam.ac.uk/using-clusters/software>.
- ▶ New software may be possible to provide on request.
- ▶ Self-installed software must be properly licensed.
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# User Environment: Environment Modules

- ▶ Modules load or unload additional software packages.
- ▶ Some are **required** and automatically loaded on login.
- ▶ Others are optional extras, or possible replacements for other modules.
- ▶ **Beware** unloading default modules in `~/.bashrc`.
- ▶ **Beware** overwriting environment variables such as `PATH` and `LD_LIBRARY_PATH` in `~/.bashrc`. If necessary append or prepend.

# User Environment: Environment Modules

- ▶ Currently loaded:

```
module list
Currently Loaded Modulefiles:
 1) dot                      6) intel/impi/4.1.3.045   11) default-impi
 2) scheduler                 7) global
 3) java/jdk1.7.0_60          8) intel/cce/12.1.10.319
 4) turbovnc/1.1              9) intel/fce/12.1.10.319
 5) vgl/2.3.1/64              10) intel/mkl/10.3.10.319
```

- ▶ Available:

```
module av
```

# User Environment: Environment Modules

## ► Show:

```
module show castep/impi/7.0.3
-----
/usr/local/Cluster-Config/modulefiles/castep/impi/7.0.3:
module-whatis    adds CASTEP 7.0.3 (Intel MPI) to your environment
Note that this software is restricted to registered users.
prepend-path      PATH /usr/local/Cluster-Apps/castep/impi/7.0.3/bin:/usr/local/...
-----
```

## ► Load:

```
module load castep/impi/7.0.3
```

## ► Unload:

```
module unload castep/impi/7.0.3
```

# User Environment: Environment Modules

- ▶ Matlab

```
module load matlab/r2015b
```

- ▶ Invoking matlab in batch mode:

```
matlab -nodisplay -nojvm -nosplash command
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where the file `command.m` contains your matlab code.

- ▶ The University site license contains the Parallel Computing Toolbox.

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# User Environment: Environment Modules

- ▶ Purge:

```
module purge
```

- ▶ Defaults:

```
module show default-impi
module unload default-impi
module load default-impi-LATEST
```

- ▶ If you have compiled software yourself your run time environment must match compile time environment!

## Excercise 4: Environment Modules

- ▶ Connect to the cluster using your training account: See excercise 2 if you have closed your terminal.
- ▶ Get a list of modules that are currently loaded

*Hints:* [module list](#)

- ▶ Get a list of available R modules

*Hints:* [module av R](#)

## Excercise 5: Run an Rscript

- ▶ Connect to the cluster using your training account: See exercise 2 if you have closed your terminal.
- ▶ In the exercises folder there is a file called test.r
- ▶ Run this script using: Rscript hello.r
- ▶ Load the module for: R/3.3.2

*Hints:* [module load R/3.3.2](#)

- ▶ Run the script again: Rscript hello.r

**Service Level 1** Paying, intended for large projects with long-term, consistent requirement.

**Service Level 2** Paying, intended for medium-sized projects with irregular requirement.

**Service Level 3** Non-paying, intended for interim or pump priming, small-scale use.

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- Guaranteed fraction of resources per quarter.

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Very low priority, very restricted, very limited. Best efforts continuation.

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# HPCS: How To Apply

- ▶ Submit the online application form:  
<https://www.hpc.cam.ac.uk/services/applying-for-resources/hpc-application>
- ▶ The PI should be someone senior enough to have funding.  
E.g. supervisor, head of research group.
- ▶ Funding is not necessary.
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- ▶ Further information can be found on the web site:  
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## Part III: Using HPC

# Advanced HPC: Types of problem.

In this section we will cover Memory Intensive and Data Intensive jobs.

*Memory Intensive:* A single problem requiring a large amount of memory.

*Data Intensive:* A single problem operating on a large amount of data.

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# Advanced: Compute Intensive Problems

- ▶ Distribute the **work** for a **single problem** across multiple CPUs to reduce the execution time as far as possible.
- ▶ Program workload must be *parallelised*:
  - Parallel programs split into copies (processes or threads).
  - Each process/thread performs a part of the work on its own CPU, concurrently with the others.
  - A well-parallelised program will fully exercise as many CPUs as there are processes/threads.
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# Advanced: Scaling & Amdahl's Law

- ▶ Using more CPUs is not necessarily faster.
- ▶ Typically parallel codes have a scaling limit.
- ▶ Partly due to the system overhead of managing more copies, but also to more basic constraints;
- ▶ Amdahl's Law (idealized):

$$S(N) = \frac{1}{(1 - p + \frac{p}{N})}$$

where

$S(N)$  is the fraction by which the program has sped up  
relative to  $N = 1$

$p$  is the fraction of the program which can be parallelized  
 $N$  is the number of CPUs.

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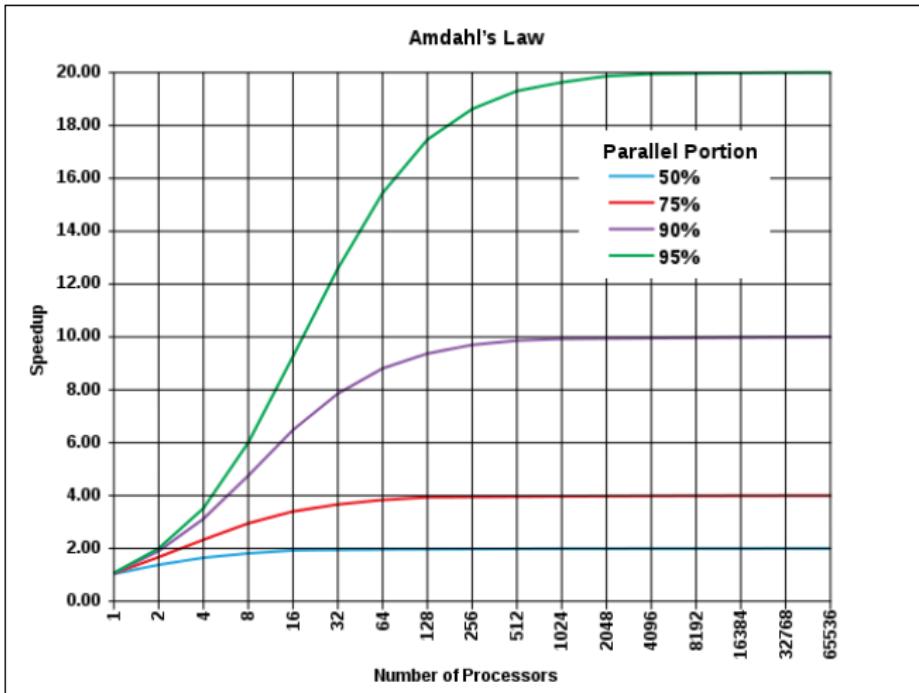
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<http://en.wikipedia.org/wiki/File:AmdahlsLaw.svg>

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- ▶ Parallelisation requires effort:
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# Advanced HPC: Memory Intensive Problems

- ▶ Require aggregation of large memory rather than multiple CPUs.  
NB Memory (fast, volatile) vs disk (slow, non-volatile).
- ▶ Technically more challenging to build machines (interconnecting memory and CPUs).
- ▶ Coding/porting easier (memory appears seamless, allowing a single system image).
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# Advanced HPC: Architecture

- ▶ So far we have abstracted you from the architecture and the concept of compiling code.
- ▶ When does the underlying architecture start to become important?
- ▶ When you are compiling code
- ▶ working with MPI/OpenMP
- ▶ large job optimisation

# Advanced HPC: Inside a Modern Computer

- ▶ Today's commodity servers already aggregate both CPUs and memory.
- ▶ Even small computers now have multiple CPU cores per socket
- ▶ Larger computers have multiple sockets (each with local memory): all CPU cores (unequally) share the node memory

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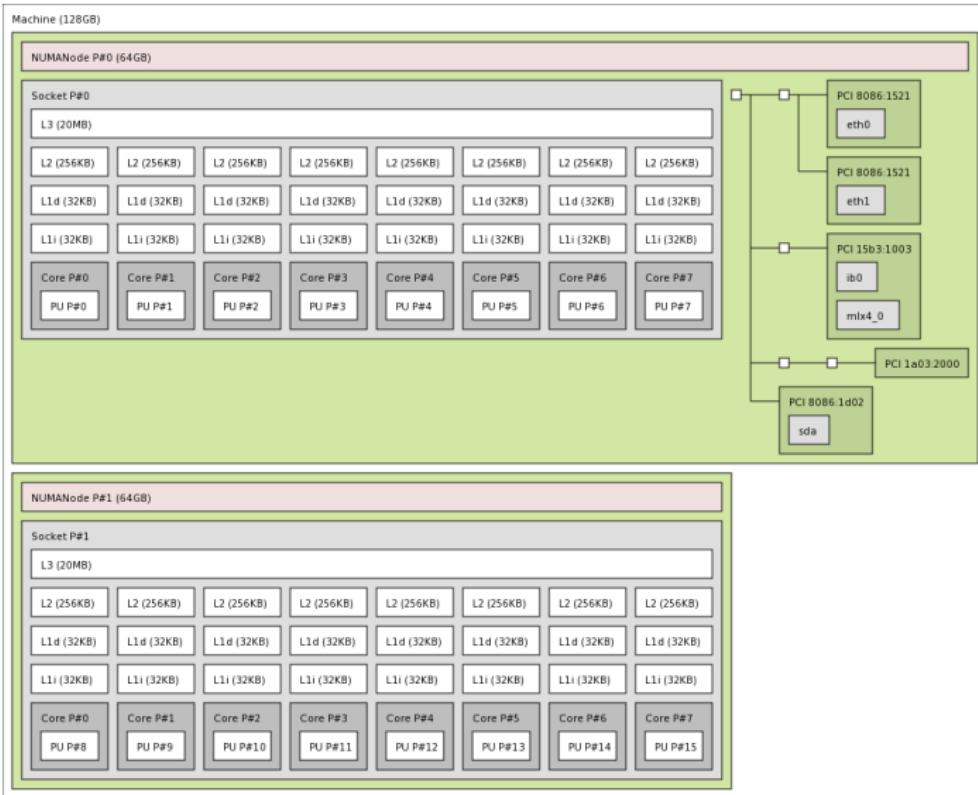
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    but users still see a single computer (**single system image**).

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## Advanced: How to Build a Supercomputer

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## 1. Take some (multicore) CPUs plus some memory.

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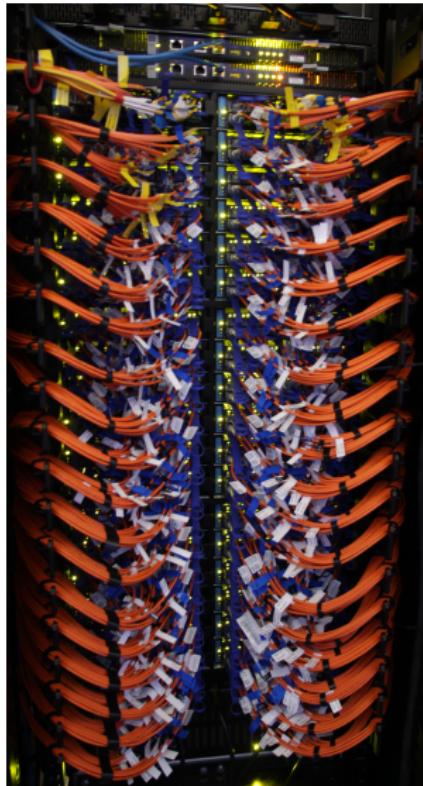
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Gbit Ethernet: 100 MB/sec

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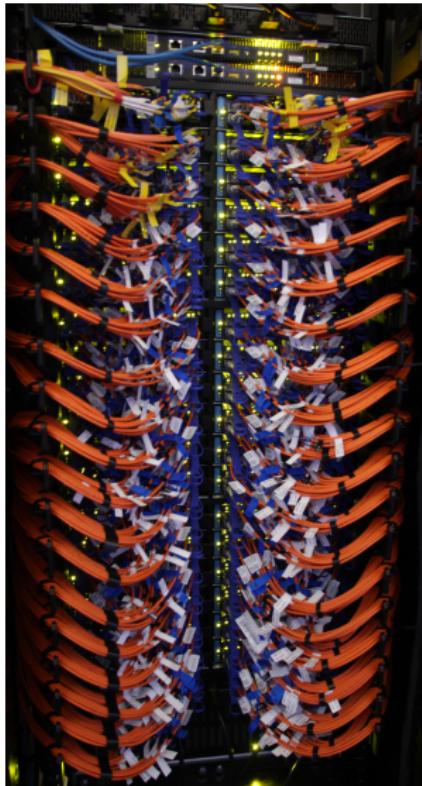
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## 3. Logically bind the nodes

- ▶ Clusters consist of distinct nodes (i.e. separate Linux computers) on common private network(s) and controlled centrally.
  - \* Private networks allow CPUs in different nodes to communicate.
  - \* Clusters are distributed memory machines:  
Each process/thread sees only its local node's CPUs and memory (without help).
  - \* Each process/thread must fit within a single node's memory.
- ▶ More expensive machines logically bind nodes into a single system i.e. CPUs and memory.
  - \* E.g. SGI UV (**COSMOS** system in DAMTP).
  - \* Private networks allow CPUs to see CPUs and memory in other nodes.
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## **Part V: The HPC Service**

# HPCS: A Brief History

**Created:** 1996 (as the HPCF).

**Mission:** Delivery and support of a large HPC resource for use by the University of Cambridge research community.

**Self-funding:** Paying and non-paying service levels.

**User base:** Includes external STFC & EPSRC plus industrial users.

**Plus:** Dedicated group nodes and research projects.

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# Darwin1 (2006–2012)



# Darwin3 (2012)(b) & Wilkes (2013)(f)



# Darwin: an Infiniband CPU Cluster

- ▶ Each compute node:
  - \* 2x8 cores, Intel Sandy Bridge 2.6 GHz.
  - \* 64 GB RAM (63900 MB usable).
  - \* 56 Gb/sec (4X FDR) Infiniband.
- ▶ 600 compute nodes (300 belong to Cambridge).
- ▶ 8 login nodes ([login.hpc.cam.ac.uk](http://login.hpc.cam.ac.uk)).

# Darwin: an Infiniband CPU Cluster

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  - \* 5 GB/sec Infiniband (for MPI and storage)
- ▶ 600 compute nodes (300 belong to Cambridge).
- ▶ 8 login nodes ([login.hpc.cam.ac.uk](http://login.hpc.cam.ac.uk)).
- ▶ 1 Petaflop upgrade availability October 2017

# Wilkes: a Dual-Rail Infiniband GPU Cluster

- ▶ Each compute node:
  - \* 2 × NVIDIA Tesla K20c GPU.
  - \* 2x6 cores, Intel Ivy Bridge 2.6 GHz.
  - \* 64 GB RAM (63900 MB usable).
  - \* 2 × 56 Gb/sec (4X FDR) Infiniband.
- ▶ 128 compute nodes.
- ▶ 8 login nodes ([login.hpc.cam.ac.uk](http://login.hpc.cam.ac.uk)).
- ▶ Environment shared with Darwin (same filesystems, user environment, scheduler).
- ▶ 1 Petaflop upgrade (Wilkes2) early access June 2017

# Wilkes: a Dual-Rail Infiniband GPU Cluster

- ▶ Each compute node:
  - \* 2 GPUs
  - \* 12 cores
  - \* 63900 MB
  - \*  $2 \times 5$  GB/sec Infiniband (for MPI and storage)
- ▶ 128 compute nodes.
- ▶ 8 login nodes ([login.hpc.cam.ac.uk](http://login.hpc.cam.ac.uk)).
- ▶ Environment shared with Darwin (same filesystems, user environment, scheduler).
- ▶ 1 Petaflop upgrade (Wilkes2) early access June 2017

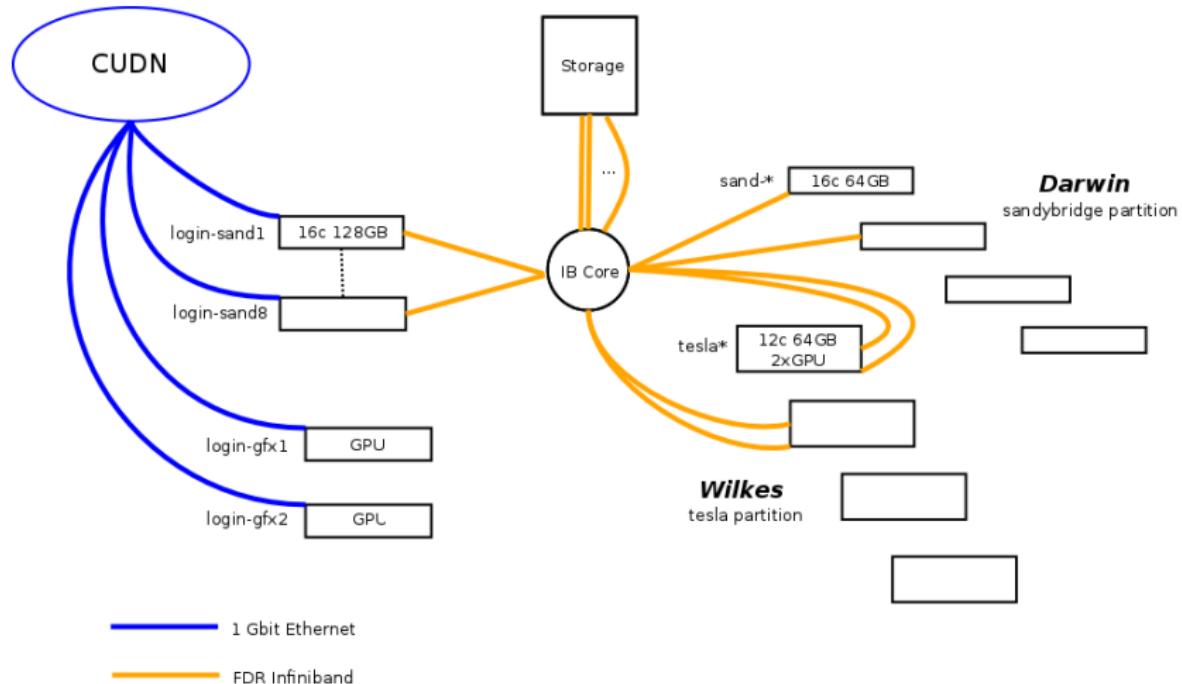
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- ▶ 1 Petaflop upgrade (Wilkes2) early access June 2017

# HPCS Production Cluster Schematic



# User Environment: Compilers

Intel: `icc`, `icpc`, `ifort` (recommended)

```
icc -O3 -xHOST -ip code.c -o prog  
mpicc -O3 -xHOST -ip mpi_code.c -o mpi_prog
```

GCC: `gcc`, `g++`, `gfortran`

```
gcc -O3 -mtune=native code.c -o prog  
mpicc -cc=gcc -O3 -mtune=native mpi_code.c -o mpi_prog
```

PGI: `pgcc`, `pgCC`, `pgf90`

```
pgcc -O3 -tp=sandybridge code.c -o prog  
mpicc -cc=pgcc -O3 -tp=sandybridge mpi_code.c -o mpi_prog
```

# User Environment: Compilers

Intel: `icc`, `icpc`, `ifort` (recommended)

```
icc -O3 -xHOST -ip code.c -o prog  
mpicc -O3 -xHOST -ip mpi_code.c -o mpi_prog
```

GCC: `gcc`, `g++`, `gfortran`

```
gcc -O3 -mtune=native code.c -o prog  
mpicc -cc=gcc -O3 -mtune=native mpi_code.c -o mpi_prog
```

PGI: `pgcc`, `pgCC`, `pgf90`

```
pgcc -O3 -tp=sandybridge code.c -o prog  
mpicc -cc=pgcc -O3 -tp=sandybridge mpi_code.c -o mpi_prog
```

## Exercise 6: Modules and Compilers

- ▶ Go to the `exercises` directory of your HPCS training account.
- ▶ Try to compile the `hello.c` program using the default `icc` compiler (it will fail because there is a deliberate bug).
- ▶ To fix the problem, open the `hello.c` file in the `gedit` editor.

## Exercise 6: Modules and Compilers

- ▶ Go to the `exercises` directory of your HPCS training account.

*Hints:* Firstly you may need to review Exercise 2 in order to reconnect to your HPCS training account. At the HPCS command prompt, change to the exercises directory (`cd ~/exercises`).

- ▶ Try to compile the `hello.c` program using the default `icc` compiler (it will fail because there is a deliberate bug).
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*Hints:* `icc hello.c -o hello`

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- ▶ Try to compile the `hello.c` program using the default `icc` compiler (it will fail because there is a deliberate bug).

*Hints:* `icc hello.c -o hello`

- ▶ To fix the problem, open the `hello.c` file in the `gedit` editor.

*Hints:* Launch gedit in the background by doing `gedit&`. A gedit window should appear. Remove the word `BUG` and save the file.

## Exercise 6: Modules and Compilers (ctd)

- ▶ Try again to compile the `hello.c` program using the default `icc` compiler, and run it. You should see “*node says: Hello, World!*”.
- ▶ Which version of `icc` did you use? Find this out by listing the current modules loaded.
- ▶ Compile and run the `hello.c` program in the exercises directory using a non-default C compiler.

## Exercise 6: Modules and Compilers (ctd)

- ▶ Try again to compile the `hello.c` program using the default `icc` compiler, and run it. You should see “*node* says: Hello, World!”.  
*Hints:* `icc hello.c -o hello` then run: `./hello`.
  - ▶ Which version of `icc` did you use? Find this out by listing the current modules loaded.
- 
- ▶ Compile and run the `hello.c` program in the exercises directory using a non-default C compiler.

## Exercise 6: Modules and Compilers (ctd)

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*Hints:* `icc hello.c -o hello` then run: `./hello`.
- ▶ Which version of `icc` did you use? Find this out by listing the current modules loaded.  
*Hints:* `module list` — the Intel compiler modules are named `intel/cce/version`. You can also work out the version directly by finding the location of the binary, e.g. doing `which icc` which should return the path:  
`/usr/local/Cluster-Apps/intel/cce/version/bin/icc`.
- ▶ Compile and run the `hello.c` program in the exercises directory using a non-default C compiler.

## Exercise 6: Modules and Compilers (ctd)

- ▶ Try again to compile the `hello.c` program using the default `icc` compiler, and run it. You should see “*node* says: Hello, World!”.

*Hints:* `icc hello.c -o hello` then run: `./hello`.

- ▶ Which version of `icc` did you use? Find this out by listing the current modules loaded.

*Hints:* `module list` — the Intel compiler modules are named `intel/cce/version`. You can also work out the version directly by finding the location of the binary, e.g. doing `which icc` which should return the path:

`/usr/local/Cluster-Apps/intel/cce/version/bin/icc`.

- ▶ Compile and run the `hello.c` program in the exercises directory using a non-default C compiler.

*Hints:* E.g. load the latest PGI C compiler (`pgcc`) with `module load pgi`. `module av` will show all possible choices (not all of which are compilers).

# Using HPC: Job Submission



# Using HPC: Job Submission

- ▶ Compute resources are managed by a scheduler:  
[SLURM](#)/PBS/SGE/LSF/...
- ▶ Jobs are submitted to the scheduler
  - analogous to submitting jobs to a print queue
  - a file (*submission script*) is copied and queued for processing.

# Using HPC: Job Submission

- ▶ Jobs are submitted from the **login nodes**
  - not themselves managed by the scheduler.
- ▶ Jobs may be either **non-interactive** (**batch**) or **interactive**.
- ▶ **Batch** jobs run a shell script on the first of a list of allocated nodes.
- ▶ **Interactive** jobs provide a command line on the first of a list of allocated nodes.

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# Using HPC: Job Submission

- ▶ The HPCS has used SLURM since February 2014.
- ▶ Jobs may use [part](#) or [all](#) of one or more nodes
  - the owner can specify --exclusive to force exclusive node access (automatic on Wilkes [but may change post-upgrade](#)).
- ▶ Template submission scripts are available under </usr/local/Cluster-Docs/SLURM>.

# Job Submission: Using SLURM

- ▶ Prepare a shell script and submit it to SLURM:

```
[abc123@login]$ sbatch slurm_submission_script  
Submitted batch job 790299
```

# Job Submission: Show Queue

- Submitted job scripts are copied and stored in a queue:

```
[abc123@login]$ squeue -u abc123
      JOBID PARTITION     NAME     USER   ST          TIME  NODES NODELIST(REASON)
    790299 sandybrid     Test3  abc123  PD          0:00      8  (Priority)
    790290 sandybrid     Test2  abc123   R  27:56:10      8 sand-6-[38-40],sand-7-[27-31]
```

# Job Submission: Show Queue

- Submitted job scripts are copied and stored in a queue:

```
[abc123@login]$ squeue -u abc123
      JOBID PARTITION     NAME     USER   ST          TIME  NODES NODELIST(REASON)
    790299 sandybrid     Test3  abc123  PD          0:00      8  (Resources)
    790290 sandybrid     Test2  abc123   R  27:56:10      8 sand-6-[38-40],sand-7-[27-31]
```

# Job Submission: Show Queue

- Submitted job scripts are copied and stored in a queue:

```
[abc123@login]$ squeue -u abc123
      JOBID PARTITION     NAME     USER   ST          TIME  NODES NODELIST(REASON)
    790299 sandybrid     Test3  abc123  PD          0:00      8  (AssocGrpCPUMinsLimit)
    790290 sandybrid     Test2  abc123   R  27:56:10      8 sand-6-[38-40],sand-7-[27-31]
```

# Job Submission: Monitor Job

- ▶ Examine a particular job:

```
[abc123@login]$ scontrol show job=790299
```

## Job Submission: Cancel Job

- ▶ Cancel a particular job:

```
[abc123@login]$ scancel 790299
```

# Job Submission: Scripts

## ► SLURM

See [slurm\\_submit.darwin](#), [slurm\\_submit.wilkes](#).

```
#!/bin/bash
#! Name of the job:
#SBATCH -J darwinjob
#! Which project should be charged:
#SBATCH -A CHANGEME
#! How many whole nodes should be allocated?
#SBATCH --nodes=1
#! How many tasks will there be in total? (<= nodes*16)
#SBATCH --ntasks=1
#! How much wallclock time will be required?
#SBATCH --time=02:00:00
#! Select partition:
#SBATCH -p sandybridge
...
```

## ► **#SBATCH** lines are *structured comments*

— correspond to sbatch command line options.

► The above job will be given 1 cpu on 1 node for 2 hours (by default there is 1 task per node, and 1 cpu per task).

# Job Submission: Scripts

## ► SLURM

See [slurm\\_submit.darwin](#), [slurm\\_submit.wilkes](#).

```
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# Job Submission: Scripts

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See [slurm\\_submit.darwin](#), [slurm\\_submit.wilkes](#).

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#! Select partition:
#SBATCH -p sandybridge
...
```

- **#SBATCH** lines are *structured comments*
  - correspond to sbatch command line options.
- The above job will be given 1 **cpu** on 1 node for 2 hours (by default there is 1 task per node, and 1 cpu per task).

# Job Submission: Scripts

## ► SLURM

See `slurm_submit.darwin`, `slurm_submit.wilkes`.

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#!/bin/bash
#! Name of the job:
#SBATCH -J darwinjob
#! Which project should be charged:
#SBATCH -A CHANGEME
#! How many whole nodes should be allocated?
#SBATCH --nodes=1
#! How many tasks will there be in total? (<= nodes*16)
#SBATCH --ntasks=16
#! How much wallclock time will be required?
#SBATCH --time=02:00:00
#! Select partition:
#SBATCH -p sandybridge
...
```

## ► **#SBATCH** lines are *structured comments*

— correspond to sbatch command line options.

► The above job will be given 16 cpus on 1 node for 2 hours (by default there is 1 task per node, and 1 cpu per task).

# Job Submission: Accounting Commands [HPCS]

- ▶ How many core hours available do I have?

```
mybalance
```

User	Usage	Account	Usage	Account Limit	Available (CPU hrs)
abc123	18	STARS	171	100,000	99,829
abc123	18	STARS-SL2	35	101,000	100,965
abc123	925	BLACKH	10,634	166,667	156,033

- ▶ How many core hours does some other project or user have?

```
gbalance -p HALOS
```

User	Usage	Account	Usage	Account Limit	Available (CPU hrs)
pq345	0	HALOS	317,656	600,000	282,344
xyz10	11,880	HALOS	317,656	600,000	282,344

(Use -u for user.)

- ▶ List all jobs charged to a project/user between certain times:

JobID	User	Account	JobName	Partition	End	NCPUS	CPUTimeRAW	ExitCode	State
14505	xyz10	halos	help	sandybrid+	2014-01-07T12:59:40	16	32	0:9	COMPLETED
14506	xyz10	halos	help	sandybrid+	2014-01-07T13:00:11	16	48	2:0	FAILED
...									

# Job Submission: Single Node Jobs

- ▶ Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=
#SBATCH --mem=3994
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS= # For OpenMP across cores
$application $options
...
```

# Job Submission: Single Node Jobs

- ▶ Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=
#SBATCH --mem=3994
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS= # For OpenMP across cores
$application $options
...
```

# Job Submission: Single Node Jobs

- ▶ Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=1
# Default is 1 cpu (core) per task
#SBATCH --mem=3994
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS= # For OpenMP across cores
$application $options
...
```

# Job Submission: Single Node Jobs

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```
#!/bin/bash
...
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# Default is 1 cpu (core) per task
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# Memory per node in MB - default is pro rata by cpu number
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...
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...
```

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#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=1
# Default is 1 cpu (core) per task
#SBATCH --mem=3994
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS= # For OpenMP across cores
$application $options
...
```

# Job Submission: Single Node Jobs

- ▶ Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=16 # Whole node

#SBATCH --mem=3994
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS=16 # For OpenMP across 16 cores
$application $options
...
```

# Job Submission: Single Node Jobs

- ▶ Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=8 # Half node

#SBATCH --mem=3994
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS=8 # For OpenMP across 8 cores
$application $options
...
```

# Job Submission: Single Node Jobs

- ▶ Serial jobs requiring large memory, or OpenMP codes.

```
#!/bin/bash
...
#SBATCH --nodes=1
#SBATCH --ntasks=1
# Default is 1 task per node
#SBATCH --cpus-per-task=16 # Whole node

#SBATCH --mem=3994
# Memory per node in MB - default is pro rata by cpu number
# Increasing --mem or --cpus-per-task implicitly increases the other
...
export OMP_NUM_THREADS=8    # For OpenMP across 8 cores (using all memory)
$application $options
...
```

# Job Submission: MPI Jobs

- ▶ Parallel job across multiple nodes.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=64      # i.e. 16x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
mpirun -np 64 $application $options
...
```

- ▶ SLURM-aware MPI launches remote tasks via SLURM.
- ▶ The template script uses \$SLURM\_TASKS\_PER\_NODE to set PPN.

# Job Submission: MPI Jobs

- ▶ Parallel job across multiple nodes.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=32      # i.e. 8x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
mpirun -ppn 8 -np 32 $application $options
...
```

- ▶ SLURM-aware MPI launches remote tasks via SLURM.
- ▶ The template script uses \$SLURM\_TASKS\_PER\_NODE to set PPN.

# Job Submission: MPI Jobs

- ▶ Parallel job across multiple nodes.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=32      # i.e. 8x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
mpirun -ppn 8 -np 32 $application $options
...
```

- ▶ SLURM-aware MPI launches remote tasks via SLURM.
- ▶ The template script uses \$SLURM\_TASKS\_PER\_NODE to set PPN.

# Job Submission: Hybrid Jobs

- ▶ Parallel jobs using both MPI and OpenMP.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=32      # i.e. 8x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
export OMP_NUM_THREADS=2  # i.e. 2 threads per MPI task.
mpirun -ppn 8 -np 32 $application $options
...
```

- ▶ This job uses 64 CPUs (each MPI task splits into 2 OpenMP threads).

# Job Submission: Hybrid Jobs

- ▶ Parallel jobs using both MPI and OpenMP.

```
#!/bin/bash
...
#SBATCH --nodes=4
#SBATCH --ntasks=32      # i.e. 8x4 MPI tasks in total.
#SBATCH --cpus-per-task=2
...
export OMP_NUM_THREADS=2  # i.e. 2 threads per MPI task.
mpirun -ppn 8 -np 32 $application $options
...
```

- ▶ This job uses 64 CPUs (each MPI task splits into 2 OpenMP threads).

# Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srun` to launch tasks (**job steps**) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=4
...
cd directory_for_job1
srun --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srun --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srun --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

# Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srun` to launch tasks (**job steps**) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=4
...
cd directory_for_job1
srun --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srun --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srun --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

# Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srun` to launch tasks (**job steps**) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=4
...
cd directory_for_job1
srun --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srun --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srun --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

# Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srun` to launch tasks (**job steps**) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=4
...
cd directory_for_job1
srun --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srun --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srun --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

# Job Submission: High Throughput Jobs

- ▶ Multiple serial jobs across multiple nodes.
- ▶ Use `srun` to launch tasks (**job steps**) within a job.

```
#!/bin/bash
...
#SBATCH --nodes=4
...
cd directory_for_job1
srun --exclusive -N 1 -n 1 $application $options_for_job1 > output 2> err &
cd directory_for_job2
srun --exclusive -N 1 -n 1 $application $options_for_job2 > output 2> err &
...
cd directory_for_job64
srun --exclusive -N 1 -n 1 $application $options_for_job64 > output 2> err &
wait
```

## Exercise 7: Submitting Jobs

- ▶ Submit a job which will run a copy of your hello program on all cores of 4 of the 24-core nodes which are available for training.

## Exercise 7: Submitting Jobs

- ▶ Submit a job which will run a copy of your hello program on all cores of 4 of the 24-core nodes which are available for training.

*Hints:* 1. Edit the script `job_script` in your exercises directory.

Set:

```
#SBATCH --nodes=4  
#SBATCH --ntasks=96  
application="../hello"
```

In the module section, make sure that the module you used to compile `hello` is also loaded (last).

2. Submit the job with `sbatch job_script`. The jobid is then printed.
3. Watch the job in the queue with `squeue`.
4. After it has disappeared, open the output file `slurm-jobid.out` in your editor. There should be 24 “Hello, World!” messages from 4 different nodes.

## Exercise 7: Submitting Jobs (ctd)

- ▶ Experiment with changing the number of nodes and tasks by changing and submitting job\_script (you are limited to 8 nodes in total).

# Job Submission: Interactive [HPCS]

- ▶ Compute nodes are accessible via SSH while you have a job running on them.
- ▶ Alternatively, submit an interactive job:  
`sintr -A MYPROJECT -N2 -n16 -t 2:0:0`
- ▶ Within the window (screen session):
  - \* Launches a shell on the first node (when the job starts).
  - \* Graphical applications should display correctly.
  - \* Create new shells with `ctrl-a c`, navigate with `ctrl-a n` and `ctrl-a p`.
  - \* `ssh` or `srun` can be used to start processes on any nodes in the job.
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# Job Submission: Array Jobs

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- ▶ [http://slurm.schedmd.com/job\\_array.html](http://slurm.schedmd.com/job_array.html)
- ▶ Used for submitting and managing large sets of similar jobs.
- ▶ Each job in the array has the same **initial** options.
- ▶ SLURM

```
[abc123@login]$ sbatch --array=1-7 -A STARS-SL2 submission_script
Submitted batch job 791609
```

```
[abc123@login-sand2]$ squeue -u abc123
      JOBID PARTITION     NAME     USER ST      TIME  NODES NODELIST(REASON)
    791609_1 sandybrid     hpl  abc123 R   0:06      1 sand-6-32
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- ▶ Updates can be applied to specific array elements using  
 `${SLURM_ARRAY_JOB_ID} -${SLURM_ARRAY_TASK_ID}`
- ▶ Alternatively operate on the entire array via  
 `${SLURM_ARRAY_JOB_ID}`.
- ▶ Some commands still require the SLURM\_JOB\_ID (sacct, sreport, sshare, sstat and a few others).

# Scheduling

- ▶ SLURM scheduling is multifactor:
  - ▶ QoS — payer or non-payer?
  - ▶ Age — how long has the job waited?
    - Don't cancel jobs that seem to wait too long.
  - ▶ Fair Share — how much recent usage?
    - Payers with little recent usage receive boost.
  - ▶ `sprio -j jobid`
- ▶ Backfilling
  - ▶ Promote lower priority jobs into gaps left by higher priority jobs.
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- ▶ The cluster is currently very busy and we do not preempt.
- ▶ 36 or 12 hour job walltimes are permitted.
- ▶ This sets the timescale at busy times (*without* backfilling).
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# Checkpointing

- ▶ Insurance against failures during long jobs.
- ▶ Restart from checkpoints to work around finite job length.
- ▶ Application native methods are best. Failing that ...
- ▶ Darwin & Wilkes nodes currently provide Berkeley Lab Checkpoint/Restart (BLCR)  
`cr_run, cr_checkpoint, cr_restart`
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# Job Submission: Scheduling Top Dos & Don'ts

## ► Do ...

- ▶ Give reasonably accurate wall times (allows [backfilling](#)).
- ▶ Check your balance occasionally ([mybalance](#)).
- ▶ Test on a small scale first.
- ▶ Implement [checkpointing](#) if possible (reduces resource wastage).

## ► Don't ...

- ▶ Request more than you need
  - you will wait longer and use more credits.
- ▶ Cancel jobs unnecessarily
  - priority increases over time.