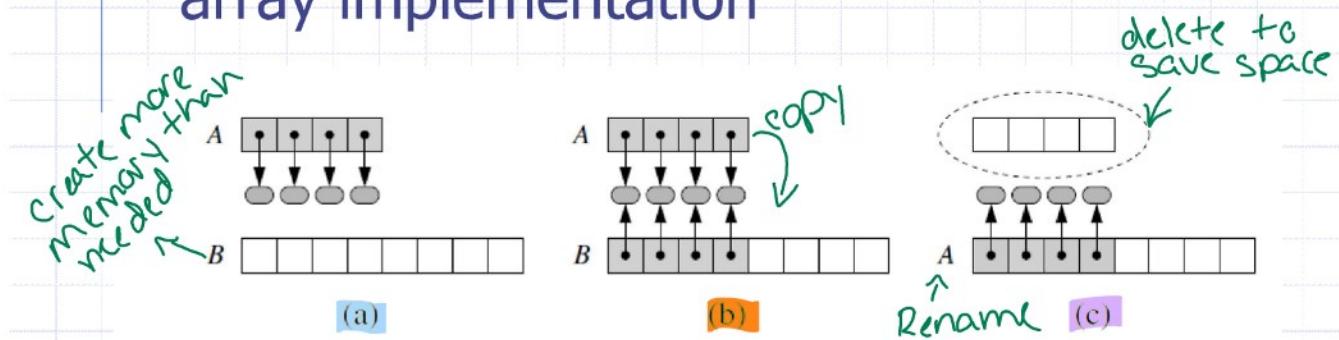


Dynamic Array

- A Python list can grow.
- Python list class uses an efficient dynamic array implementation



(a) create new array B (b) store elements of A in B (c)
reassign reference A to the new array.

$B = [] * 2n \# \text{doubling space}$

only need 1 extra cell, why double?
why not open up as needed (*append*)
to eliminate repetitive steps/operations

for $i=0 : n-1 \# \text{whole array from } i=0 \text{ up to } n-1$

$B[i] = A[i] \# \text{Dump elements of } A \rightarrow B$

$A = B \# \text{Rename}$

if full \rightarrow then append

$\textcircled{1} \quad [] * n \quad \# \text{create empty list of size } n$

OR

$\textcircled{2} \quad \text{extend list all at once}$

$\hookrightarrow \text{list.extend(another_list)} \quad \# \text{copy \& extend}$

$\boxed{1 \ 2 \ 14 \ 12}$

$\text{data.index}(4)$

#loop

\hookrightarrow start at index 0

\hookrightarrow compare each value to 4

$O(k+1) \quad \nearrow \# \text{operations}$

$\hookrightarrow k = \text{first occurrence}$

\hookrightarrow index location

Complexity is at most n

$\hookrightarrow O(n)$

$k = \text{index}$

data 1

$\begin{array}{|c|c|c|c|} \hline 0 & 1 & 2 & 3 \\ \hline 1 & 2 & 3 & 4 \\ \hline \end{array}$

Want to see if they are equivalent

data 2

$\begin{array}{|c|c|c|c|} \hline 0 & 1 & 2 & 3 \\ \hline 1 & 2 & 4 & 3 \\ \hline \end{array}$

\hookrightarrow compare each element

$$3 = k+1$$
$$O(k+1)$$

$\left\{ \begin{array}{l} k=0 \hookrightarrow 1 == 1 \checkmark \\ k=1 \hookrightarrow 2 == 2 \checkmark \\ k=2 \hookrightarrow 3 == 4 \times \text{stops} \\ \quad \quad \quad +1=3 \end{array} \right.$

Worst case:

\hookrightarrow go through whole list

$\hookrightarrow O(n)$

Slicing → exclusive?

data[j:k] → non-inclusive

$s[2:4]$ → $s[0:4]$ → index

$s[2] \rightarrow 2$ $O(1)$
 $s[3] \rightarrow 3$ $O(1)$

2 operations

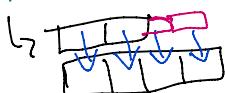
$k-j+1$ → storing value/slice
in another array
 $O(n)$

data1 + data2

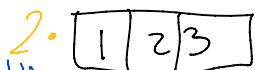
↳ add all indices up

↳ same size → n operations → $O(n)$ → best case

↳ not same size → append zeros → $O(2n)$ → worst case



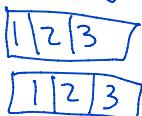
$c * data$



↳ not multiplying data by constant

↳ duplicating the lists

↳ $O(cn)$ → build list c times



↳ addition of sets

↳ merges lists together

`data[i] = val`
↳ setting value

`[1, 2, 3, 4]` now points to 100,
 ↑
`data[2] = 100 → [1, 2, 100, 4]`

`data.append(value)`

↳ does list have enough space?

↳ best case $O(1) \rightarrow 1$ operation \rightarrow over time $\xrightarrow{\text{experimentally}}$

↳ worst case $\rightarrow O(n) \rightarrow$ when it needs to grow
↳ mathematically

↳ based on dynamic array

↳ doesn't add for every growth spurt

↳ doesn't add every time \rightarrow not n^2 (arithmetic)

Something in between
↳ dynamic array

↳ either start w/:
↳ size = 0 \rightarrow build

OR

↳ size = n / empty
↳ make $2n$ if needed

`data.insert(k, value)`

$O(n-k+1)$

`[1, 2, 100, 4]` moving 3 cells right

↳ worst case \rightarrow insert at index 0
↳ n shifts $\rightarrow O(n)$

`data.insert(1, 2000) → [1, 2000, 2, 100, 4]`

index 1 \rightarrow 4

`data.pop() → insert at end`

↳ easiest b/c all other indices don't change \rightarrow only last

`data.pop(k)`

↳ pop data @ index k
↳ method of class list

`del data[k]`

↳ delete data @ index k
↳ works not just on lists

`[1, 2000, 2, 100]`

↳ `pop(2)`

`[1, 2000, X, 100]`
 ↑
 more

↳ `del data[2]`

`[1, 2000, X, 100]`
 ↑
 move

$O(n-k)$

↓
4
 ↓
2
 ↓
 2 operations
 ① Delete
 ② moving

`data.remove(value)`

↳ searches for value → found → delete
↳ based on values not indices

`data1.extend(data2)`

$O(n_2)$
 ↳ size of data2
 adding list2 to list1
 ↳ arithmetic version:
 ↳ `data1 += data2`

`data.reverse()`

`[1, 2000, 2, 100] → [100, 2, 2000, 1]`

↳ can recreate w/ recursion pattern

↳ best case sorting algorithm: $O(n \log n)$
 ↳ memorize
 ↳ Tim sort