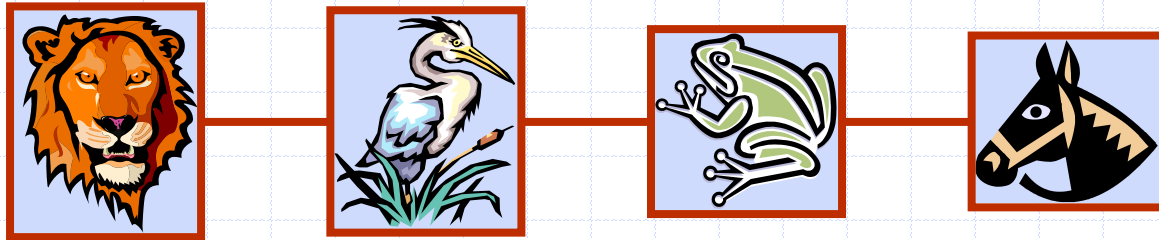


# Linked Lists



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# Array Sequences

- ◆ Python's list class is highly optimized, and often a great choice for storage.

- ◆ However, there are some notable disadvantages:

- 1. The length of a dynamic array might be longer than the number of elements that it stores.*

*↳ using more memory than needed  
↳ "expensive" to access index element*

- 2. Amortized bounds for operations may be unacceptable in real-time systems. (.append, .extend)*

- 3. Insertions and deletions at interior positions of an array are expensive.*

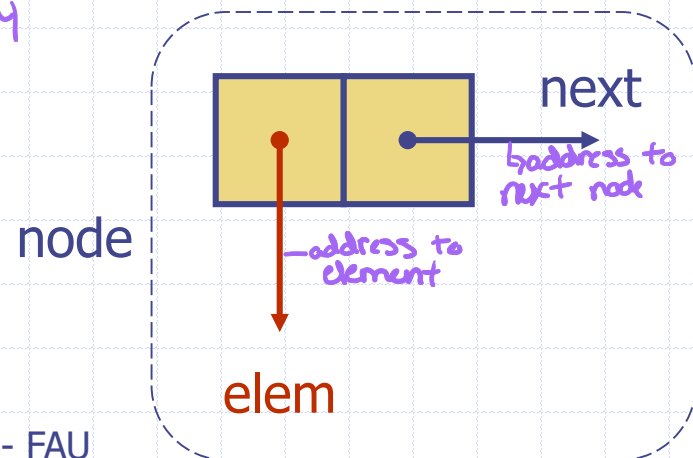
# Linked-List (I)

- ◆ Linked list makes it easy to add, remove or delete elements from anywhere in the sequence.
- ◆ It does however need more memory space!
- ◆ Linked list consists of nodes.
- ◆ Nodes reference the item stored and the item that is next in line.
- ◆ Nodes are not stored in memory sequentially.

↳ optimize memory

hexadecimals

0-15  
F  $\times 16^2$  16<sup>1</sup> 16<sup>0</sup>  
A E 0 C



# Linked List (II)

- ◆ Elements of a linked list do not use the indexing system.
- ◆ We cannot tell if an element is 1<sup>st</sup>, 2<sup>nd</sup> or 5<sup>th</sup> .
- ◆ Linked lists are considered as a class of **non-contiguous** data structures.
  - They occupy various locations in memory.
  - So, we need to browse the entire list to find a specific item in a linked list.

# 2 Types of Advanced Data Structures

## Contiguous

↳ elements are stored continuously, in order

↳ Stack

↳ Queue

↳ Built-ins

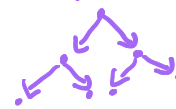
↳ list

↳ tuple

## Non-contiguous

↳ Social media

↳ Tree



↳ graph



↳ internet



↳ what actions can a link list support?

↳ create a node @ head or tail

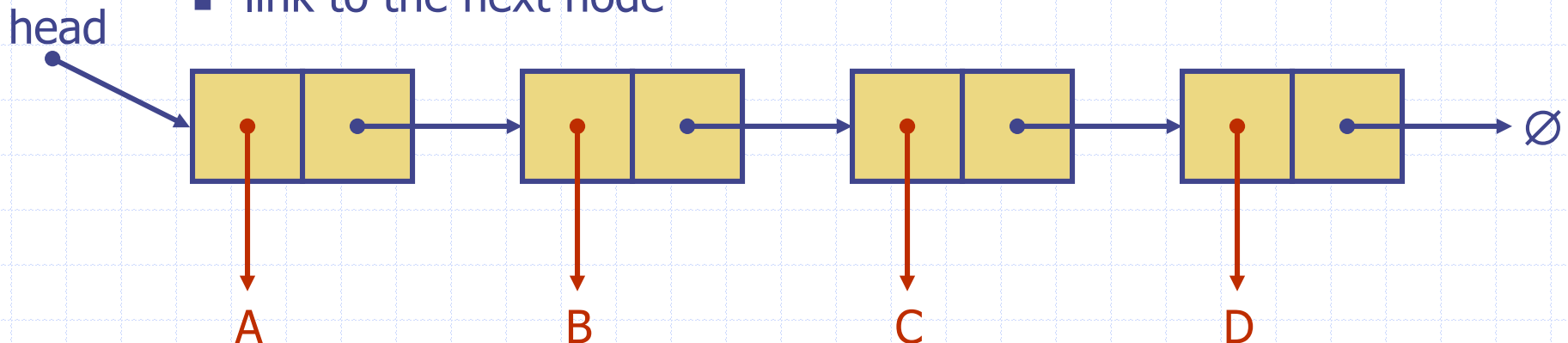
↳ insert data

↳ removal of node

↳ traversing link list

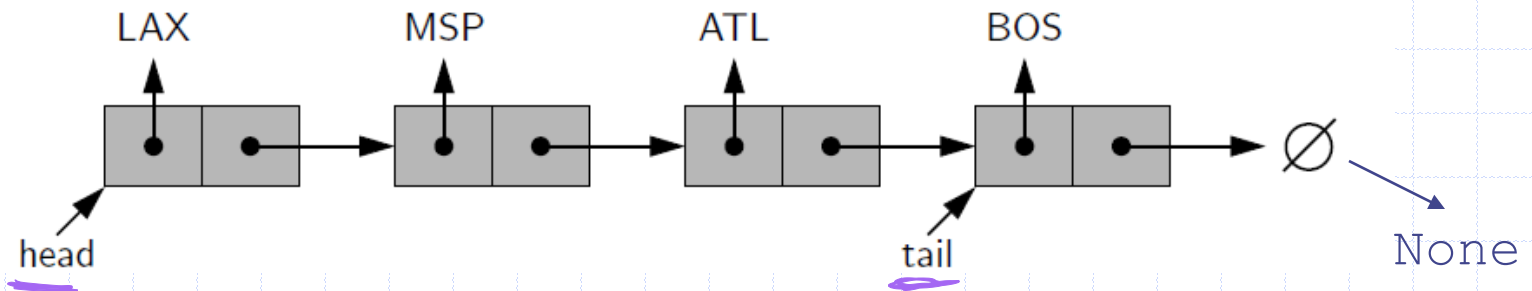
# Singly Linked List

- ◆ A singly linked list is a concrete data structure consisting of a sequence of nodes, starting from a head pointer
- ◆ Each node stores
  - element
  - link to the next node



# Linked list Operations

*Use more memory than array*



- ◆ Creation of a node
- ◆ Insertion of a node to the begin, middle or end.
- ◆ Deletion of a node
- ◆ Traversing the linked list: Begin at head and use each reference to reach the end (Indicated by None)



expand  
list of  
nodes

① Create Node

② Add Data to Node  
(MIA-airport code)

③  $mia.next = L.head$   
list

④  $L.head = L.mia$

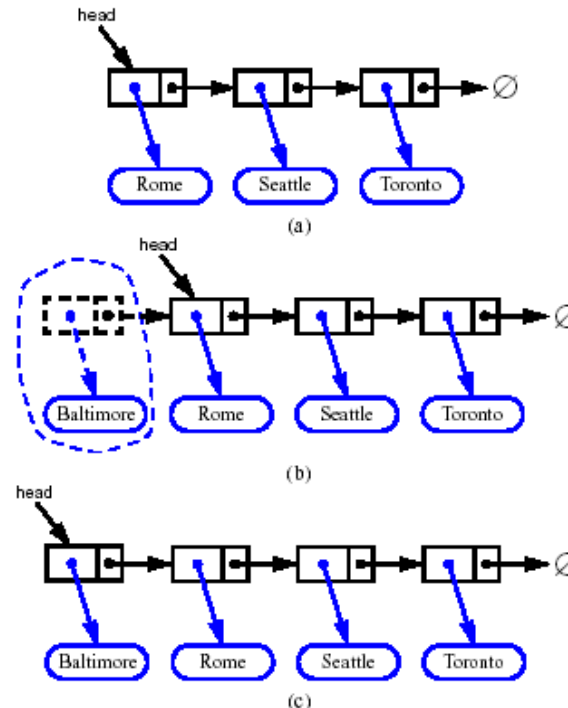
⑤  $L.size = L.size + 1$



# Inserting at the Head

Stack:  $\rightarrow$  LIFO  
 $\hookrightarrow$  treat stack like link list  
 $\hookrightarrow$  add on stack  $\rightarrow$  add to head  
 $\hookrightarrow$  use Node for stack implementation  
 $\hookrightarrow$  not array

1. Allocate a new node
2. Insert new element
3. Have new node point to old head
4. Update head to point to new node



**Algorithm** add\_first(L, e):  $e = \text{new element (MIA)}$

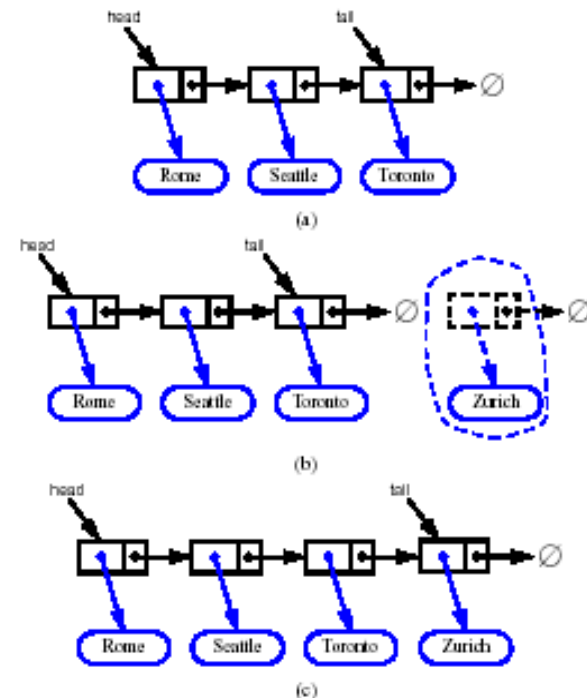
```

newest = Node(e) {create new node instance storing reference to element e}
newest.next = L.head {set new node's next to reference the old head node}
L.head = newest {set variable head to reference the new node}
L.size = L.size + 1 {increment the node count}
    
```

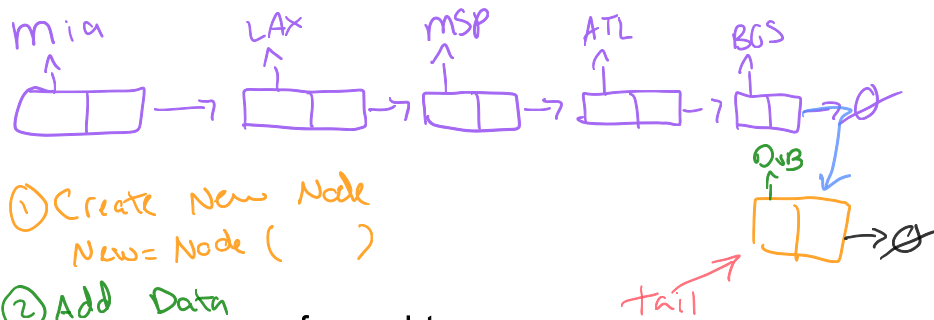
$\hookrightarrow$  keep track of size  
 $\hookrightarrow$  or have to traverse whole list to count  
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# Inserting at the Tail

1. Allocate a new node
2. Insert new element
3. Have new node point to null
4. Have old last node point to new node
5. Update tail to point to new node



## Insert @ tail



① Create New Node  
 $\text{New} = \text{Node} ( \quad )$

② Add Data  
 $\text{New} = \text{DUB}$  reference data

③  $\text{tail}.\text{next} = \text{New}$   
Connect current tail to New

④  $\text{New}.\text{next} = \text{None}$

⑤ move tail (optional)  
 $\text{L}.\text{tail} = \text{new}$

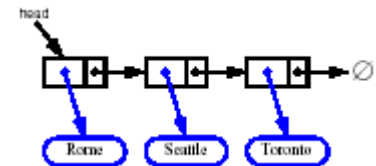
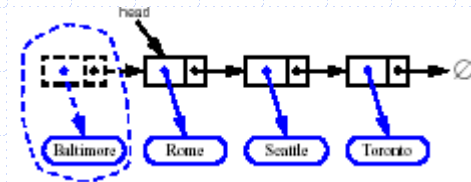
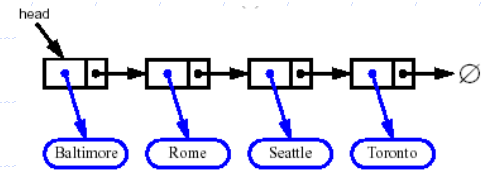
⑥  $\text{L}.\text{size} = \text{L}.\text{size} + 1$   
add to size

# Removing at the Head

1. Update head to point to next node in the list
2. Allow garbage collector to reclaim the former first node

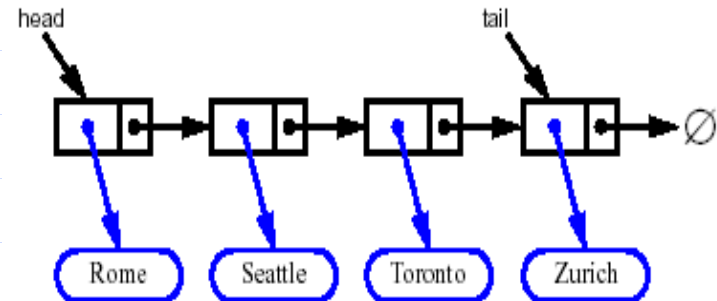
$L.head = L.head.next$

$O(1)$   
complexity



# Removing at the Tail

- ◆ Removing at the tail of a singly linked list is not efficient!
- ◆ There is no constant-time way to update the tail to point to the previous node



*Not efficient for single linked list  
↳ have to start at beginning and go through all until end*

# The Node Class

```
class Node:
    '''Lightweight class for storing a singly linked node.'''

    def __init__(self, element, next):          # initialize node's field
        self.element = element                # reference to user's element
        self.next = next                      # reference to next node

    def getElement(self):                      # Accessor methods
        return self.element

    def getNext(self):                         # Accessor methods
        return self.next

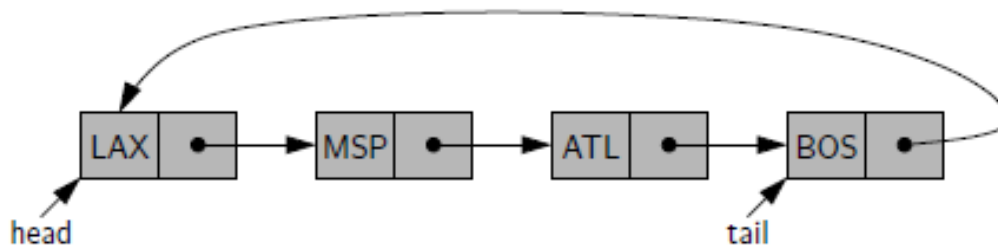
    def setElement(self, new):                # Modifier methods
        self.element = new

    def setNext(self, newNext):
        self.next = newNext
```

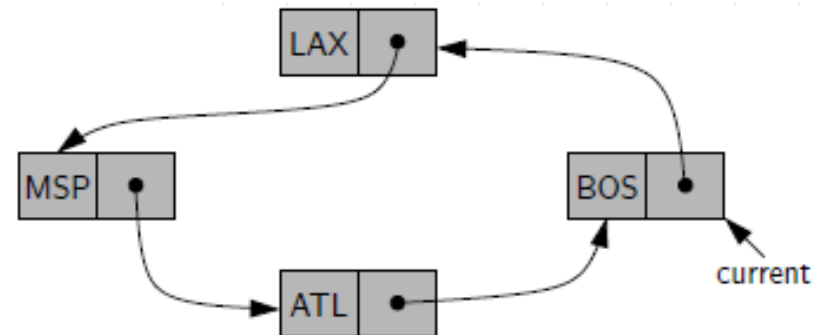
Refer to the python code

# Circularly Linked Lists

- ◆ Used for data sets that are cyclic : players taking turn in a game
  - ◆ No begin or end node
  - ◆ The 'current' identifier indicates the reference to a designated node.
  - ◆ `current = current.next` helps advance through nodes
- ↳ head / tail : depends on who's turn it is*

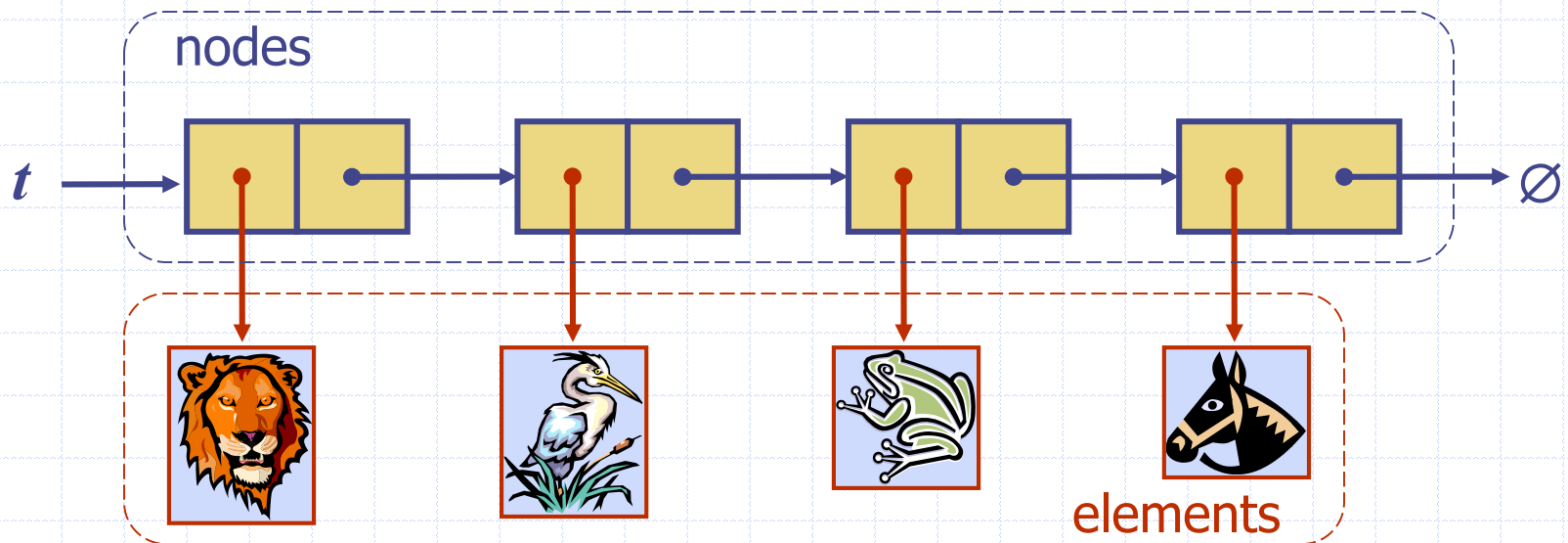


*↳ how to tell if circular?*  
*↳ circular : no none at end*



# Stack as a Linked List

- ◆ We can implement a stack with a singly linked list
- ◆ The top element is stored at the first node of the list
- ◆ The space used is  $O(n)$  and each operation of the Stack ADT takes  $O(1)$  time





Queue FIFO easier to remove head than tail

# Linked-List Stack in Python

```
1 class LinkedStack:
2     """LIFO Stack implementation using a singly linked list for storage."""
3
4     #----- nested _Node class -----
5     class _Node:
6         """Lightweight, nonpublic class for storing a singly linked node."""
7         __slots__ = '_element', '_next' # streamline memory usage
8
9         def __init__(self, element, next): # initialize node's fields
10             self._element = element # reference to user's element
11             self._next = next # reference to next node
12
13     #----- stack methods -----
14     def __init__(self):
15         """Create an empty stack."""
16         self._head = None # reference to the head node
17         self._size = 0 # number of stack elements
18
19     def __len__(self):
20         """Return the number of elements in the stack."""
21         return self._size
22
```

```
23     def is_empty(self):
24         """Return True if the stack is empty."""
25         return self._size == 0
26
27     def push(self, e):
28         """Add element e to the top of the stack."""
29         self._head = self._Node(e, self._head) # create and link a new node
30         self._size += 1
31
32     def top(self):
33         """Return (but do not remove) the element at the top of the stack.
34
35         Raise Empty exception if the stack is empty.
36         """
37         if self.is_empty():
38             raise Empty('Stack is empty')
39         return self._head._element # top of stack is at head of list

```

```
40     def pop(self):
41         """Remove and return the element from the top of the stack (i.e., LIFO).
42
43         Raise Empty exception if the stack is empty.
44         """
45         if self.is_empty():
46             raise Empty('Stack is empty')
47         answer = self._head._element
48         self._head = self._head._next # bypass the former top node
49         self._size -= 1
50         return answer

```

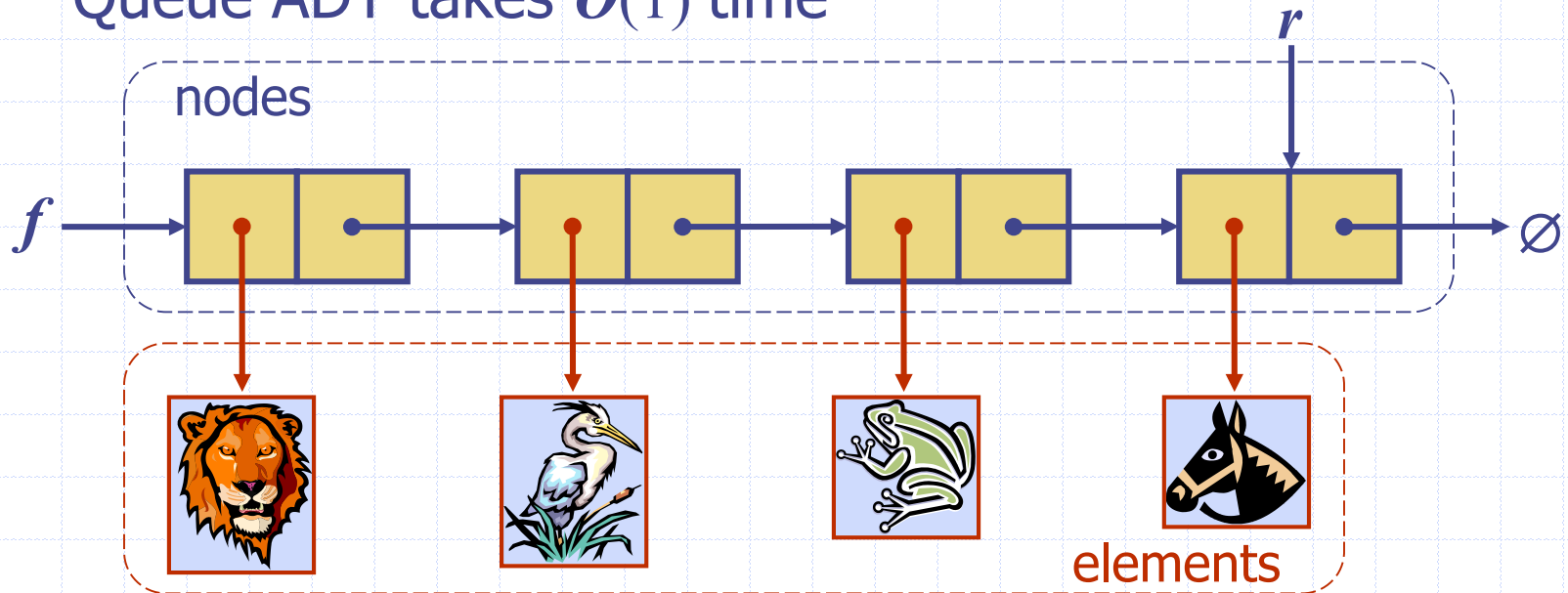
# Linked Stack Implementation

- ◆ All methods complete in worst case  $O(1)$  time.
- ◆ Huge improvement over array-based stack ADT.

Operation	Running Time
S.push(e)	$O(1)$
S.pop()	$O(1)$
S.top()	$O(1)$
len(S)	$O(1)$
S.is_empty()	$O(1)$

# Queue as a Linked List

- ◆ We can implement a queue with a singly linked list
  - The front element is stored at the first node
  - The rear element is stored at the last node
- ◆ The space used is  $O(n)$  and each operation of the Queue ADT takes  $O(1)$  time



# Linked-List Queue in Python

```
1 class LinkedQueue:
2     """FIFO queue implementation using a singly linked list for storage."""
3
4     class _Node:
5         """Lightweight, nonpublic class for storing a singly linked node."""
6         (omitted here; identical to that of LinkedStack._Node)
7
8     def __init__(self):
9         """Create an empty queue."""
10        self._head = None
11        self._tail = None
12        self._size = 0                # number of queue elements
13
14    def __len__(self):
15        """Return the number of elements in the queue."""
16        return self._size
17
18    def is_empty(self):
19        """Return True if the queue is empty."""
20        return self._size == 0
21
22    def first(self):
23        """Return (but do not remove) the element at the front of the queue."""
24        if self.is_empty():
25            raise Empty('Queue is empty')
26        return self._head._element    # front aligned with head of list
```

```
27    def dequeue(self):
28        """Remove and return the first element of the queue (i.e., FIFO).
29
30        Raise Empty exception if the queue is empty.
31        """
32        if self.is_empty():
33            raise Empty('Queue is empty')
34        answer = self._head._element
35        self._head = self._head._next
36        self._size -= 1
37        if self.is_empty():           # special case as queue is empty
38            self._tail = None         # removed head had been the tail
39        return answer
40
41    def enqueue(self, e):
42        """Add an element to the back of queue."""
43        newest = self._Node(e, None)   # node will be new tail node
44        if self.is_empty():
45            self._head = newest        # special case: previously empty
46        else:
47            self._tail._next = newest
48            self._tail = newest        # update reference to tail node
49            self._size += 1
```