

CMPE-250 Assembly and Embedded Programming

Laboratory Exercise 3

Memory, Conditional Branching, and Debugging Tools

By submitting this report, I attest that its contents are wholly my individual writing about this exercise and that they reflect the submitted code. I further acknowledge that permitted collaboration for this exercise consists only of discussions of concepts with course staff and fellow students. Other than code provided by the instructor for this exercise, all code was developed by me.

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Lab Section 1

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Results

The activity involved calculation of the following equations:

$$F = 3P + 2Q - 75$$

$$G = 2P - 4Q + 63$$

$$\text{Result} = F + G$$

For the following two sets of inputs.

Input Set 1: P = 9, Q = 4

Figure 1 shows the values stored in the memory after the program written was successfully debugged. The first two memory addresses show the values of F and G, respectively. They are marked in Green because they were written to.

The next two memory spaces are marked in Red as they were read from. They held the values of P and Q, respectively. The last highlighted memory address is the Result variable, which too was updated at the end of the operation.

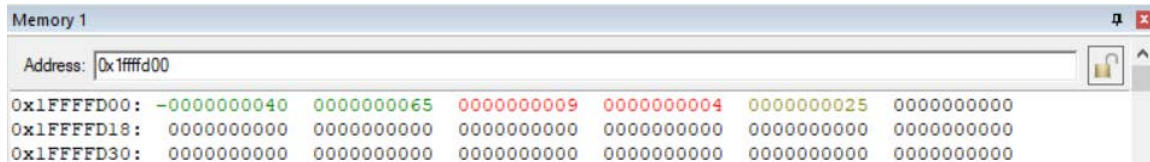


Figure 1. Memory variables result of Input Set 1

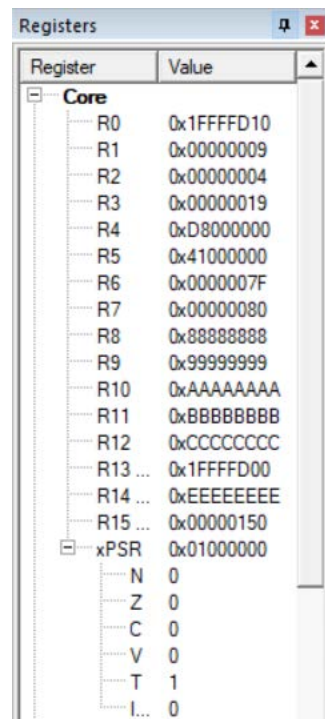


Figure 2. Register Content for Input Set 1

Manual Calculation for input set 1:

$$F = 3(9) + 2(4) - 75 = 27 + 8 - 75 = -40$$

$$\underline{F = -40}$$

$$G = 2(9) - 4(4) + 63 = 18 - 16 + 63 = 65$$

$$\underline{G = 65}$$

$$\text{Result} = -40 + 65 = 25$$

Input Set 2: P = 13, Q = -14

Figure 2 shows the values stored in the memory after the program written was successfully debugged. Like the first input set, the first two memory addresses show the value of F and G. The next two addresses are the inputs P and Q, and the last address is the Result.

Address	Value
0x1FFFD00:	-0000000064 0000000000 0000000013 -0000000014 -0000000064 0000000000
0x1FFFD18:	0000000000 0000000000 0000000000 0000000000 0000000000 0000000000
0x1FFFD30:	0000000000 0000000000 0000000000 0000000000 0000000000 0000000000

Figure 3. Memory variables result of Input Set 2

Register	Value
Core	
R0	0x1FFFD10
R1	0x000000D
R2	0xFFFFFFFF2
R3	0xFFFFF0C
R4	0xC0000000
R5	0x00000000
R6	0x0000007F
R7	0x00000080
R8	0x88888888
R9	0x99999999
R10	0xAAAAAAAA
R11	0xBBBBBBBB
R12	0CCCCCCC
R13	0x1FFFD00
R14	0xEEEEEEEE
R15	0x00000150
xPSR	0x81000000
N	1
Z	0
C	0
V	0
T	1
I...	0

Figure 4. Register Content for Input Set 2

Manual Calculation for input set 2:

$$F = 3(13) + 2(-14) - 75 = 39 - 28 - 75 = -64$$

$$\underline{F = -64}$$

$$G = 2(13) - 4(-14) + 63 = 26 + 56 + 63 = 145 \quad \leftarrow \text{Overflow}$$

$$\underline{G = 0}$$

$$\text{Result} = -64 + 0 = -64$$

Question

Could you reduce or eliminate overflow by changing the order of operations within the expressions (F, G, and/or Result)? Explain why or why not.

No, overflow cannot be eliminated by changing the order of operations within the expressions. No matter what we do, we are bound to get the same result which will cause an overflow.