Reo A Channel-Based Model for Component Composition

Seyed Mohammad Mehdi Ahmadpanah

smahmadpanah@aut.ac.ir

Apr. 10, 2018

Outline

- Motivation
- Introduction
- Syntax
- Semantics
- Tool
- Conclusion



Motivation

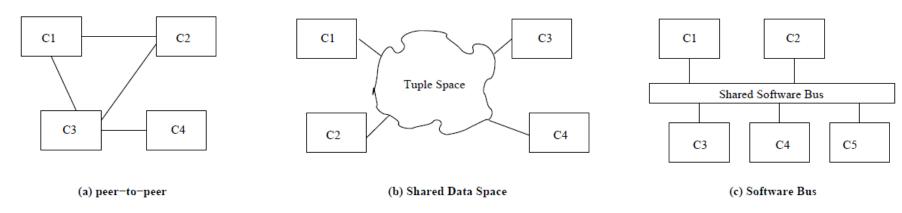


- Modular Design Popularity
- Need for a Glue-Language
 - interfacing gaps must be filled with additional code
- Bulk of Specialized Glue-Code in Complex Systems
 - interfacing code for compositional construction
- Maintenance Difficulties
 - hard to evolve, inflexible
- Alternative: construct the glue code <u>compositionally</u>, out of primitive <u>connectors</u>
- Reo introduced by Dr. Farhad Arbab in 2004
 - composition of channels to make complex connectors

Motivation (cont.)



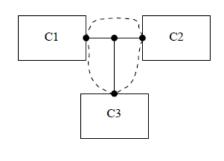
Channel-based communication models are complete

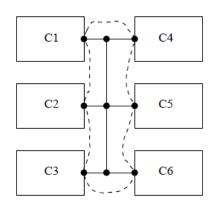


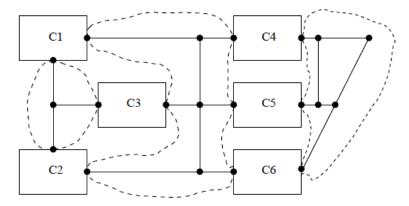
- Advantages over Other Communication Models
 - Efficiency
 - Security
 - Architectural Expressiveness
 - Anonymity

Introduction

- System consists of Components and Connectors
 - Emphasis on Connectors and their composition only!
 - Coordination from outside
 - Based on calculus of channels
 - complex connectors are constructed through composition of simpler ones

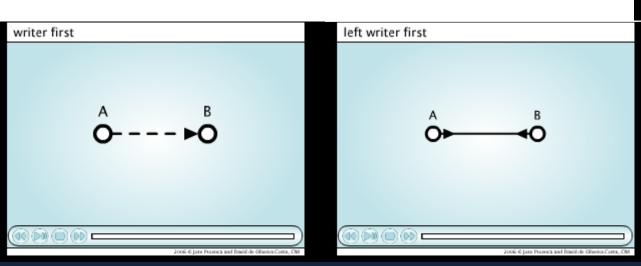


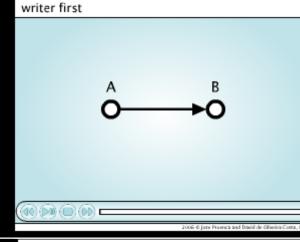


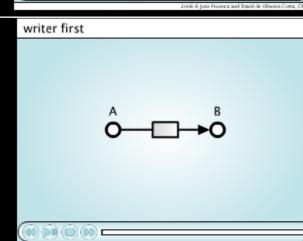


Channel Types

- Sync (write/take)
- Lossy
- SyncDrain (write/write)
- FIFO and FIFOn

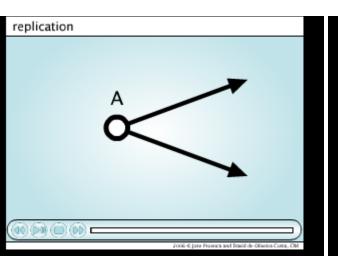


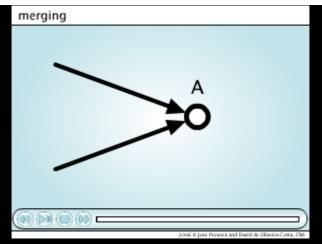


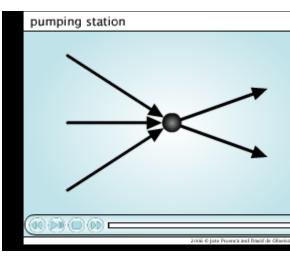


Node Types

- Source: accepts data into its channel
- Sink: dispenses data out of its channel
- Mixed: Source + Sink (no buffer)





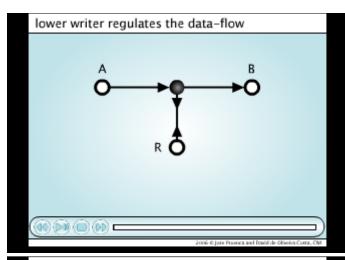


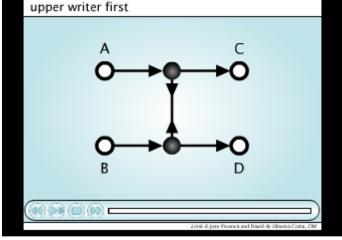
Syntax

- The set of primitive operations on channels and channel ends
 - create(ChanType)
 - forget(e)
 - move(e, loc)
 - connect([t,] e)
 - disconnect(e)
 - wait([t,] conds)
 - read([t,] inp[, v[, pat]])
 - take([t,] inp[, v[, pat]])
 - write([t,], outp, v)
 - join(e1, e2)
 - split(e[, quoin])
 - hide(e)

Channel Composition

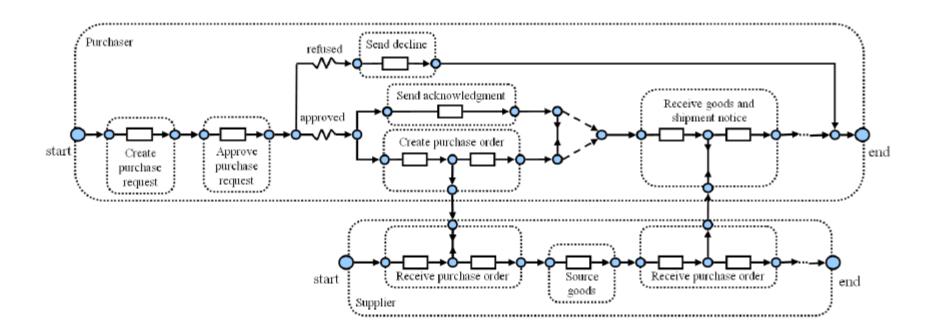
```
WCRegulator(n)
   \langle a, x1 \rangle = create(Sync)
   \langle x2, b \rangle = create(Sync)
   \langle x, y \rangle = create(SyncDrain)
  connect(x1)
  connect(x2)
  join(x, x1)
  join(x1, x2)
  hide(x)
  c = \langle \rangle
  for i = 1 to n do
      \langle u, w \rangle = create(Sync)
     c = c \circ \langle u \rangle
      connect(w)
     join(y, w)
  done
  hide(y)
  return \langle a, b, c \rangle
```



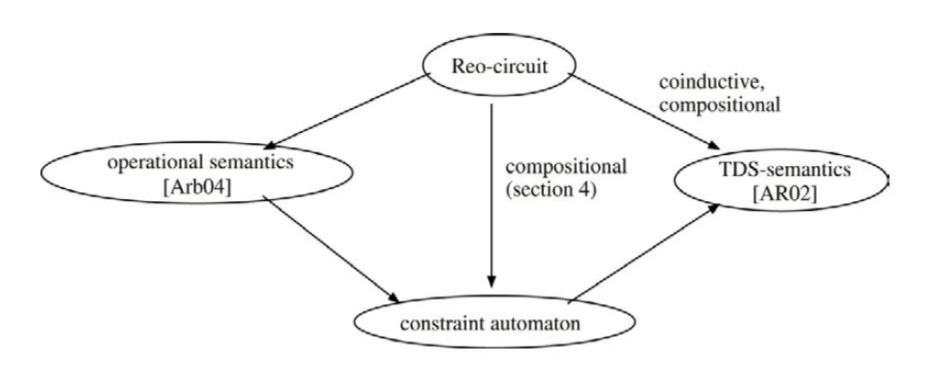


Real World Example

Reo Circuit for the Purchase Order Scenario



Semantics



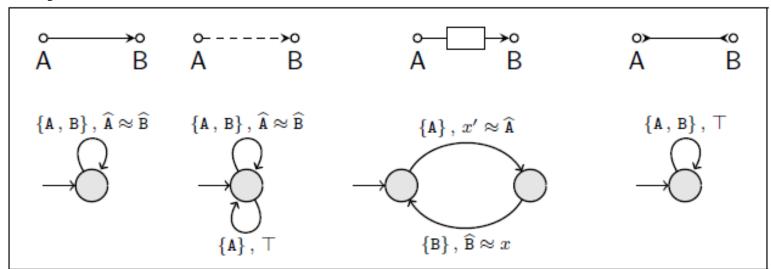
$$q \xrightarrow{N,g} p$$

$$g \in DC(N, Data)$$

$$data(A) = d_A : A \in N$$

Semantics (cont.)

- Constraint Automata Examples
 - Synchronous
 - Lossy Synchronous
 - Asynchronous FIFO
 - Synchronous Drain



Semantics (cont.)

Product Automata (Join)

$$egin{aligned} \mathcal{A}_1 &= \; (Q_1,\mathcal{N}_1, \longrightarrow_1, Q_{0,1}) \ \mathcal{A}_2 &= (Q_2,\mathcal{N}_2, \longrightarrow_2, Q_{0,2}) \ \\ \mathcal{A}_1 &\bowtie \mathcal{A}_2 &= \; (Q_1 \times Q_2, \mathcal{N}_1 \cup \mathcal{N}_2, \longrightarrow, Q_{0,1} \times Q_{0,2}) \ \\ & \underbrace{q_1 \stackrel{N_1,g_1}{\longrightarrow}_1 \; p_1, \; \; q_2 \stackrel{N_2,g_2}{\longrightarrow}_2 \; p_2, \; \; N_1 \cap \mathcal{N}_2 = N_2 \cap \mathcal{N}_1}_{\langle q_1,q_2 \rangle} \ \\ & \underbrace{\langle q_1,q_2 \rangle \stackrel{N_1 \cup N_2,g_1 \wedge g_2}{\longrightarrow} \langle p_1,p_2 \rangle}_{\langle q_1,q_2 \rangle} \ \\ & \underbrace{q_1 \stackrel{N,g}{\longrightarrow}_1 \; p_1, \; \; N \cap \mathcal{N}_2 = \emptyset}_{\langle q_1,q_2 \rangle} \ \\ & \underbrace{\langle q_1,q_2 \rangle \stackrel{N,g}{\longrightarrow}_1 \langle p_1,q_2 \rangle}_{\langle p_1,q_2 \rangle} \ \end{aligned}$$

Semantics (cont.)

Relation on Timed Data Streams (TDS)

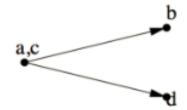
$$TDS = \left\{ \langle \alpha, a \rangle \in Data^{\omega} \times \mathbb{R}^{\omega}_{+} \mid \forall n \geq 0 : a_{n} < a_{n+1} \text{ and } \lim_{n \to \infty} a_{n} = \infty \right\}$$

- Examples
 - Channels

$$\langle \alpha, a \rangle$$
 Sync $\langle \beta, b \rangle \equiv \alpha = \beta \wedge a = b$
 $\langle \alpha, a \rangle$ SyncDrain $\langle \beta, b \rangle \equiv a = b$
 $\langle \alpha, a \rangle$ FIFO1 $\langle \beta, b \rangle \equiv \alpha = \beta \wedge a < b < a'$

Replicator

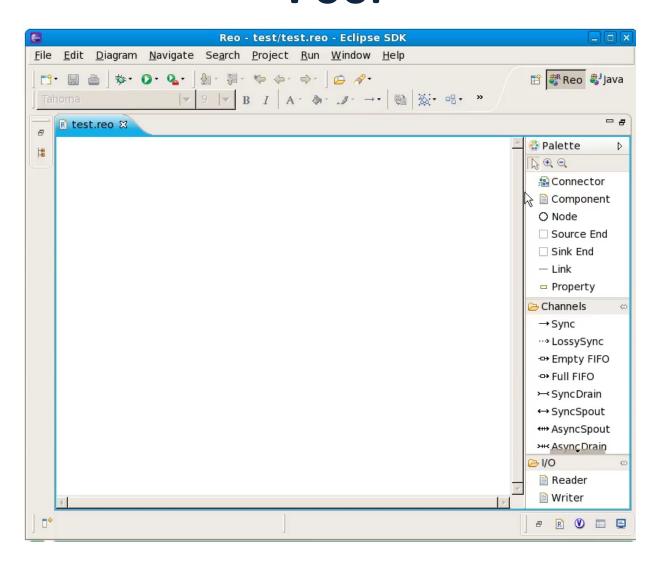
$$R(\langle \alpha, a \rangle; \langle \beta, b \rangle, \langle \gamma, c \rangle) \equiv \alpha = \beta = \gamma \land a = b = c$$



Tool

- Extensible Coordination Tools (ECT)
 - set of Eclipse Plugins
 - Graphical Editing of Reo connectors and constraint automata
 - Animation of Reo connectors
 - Code Generation from Reo connectors or constraint automata
 - Model Checking connectors using Vereofy and mCRL2

Tool



Conclusion

- Reo is a language for compositional construction of interaction protocols
- Interaction is the only first-class concept in Reo
- Protocols manifest as a connectors
- In its graphical syntax, connectors are graphs
 - Data items flow through channels represented as edges
 - Boundary nodes permit (components to perform) I/O operations
- Formal semantics given as Constraint Automata (and various other formalisms)
- Tool support: draw, animate, verify, compile

References

- [1] F. Arbab "Reo: A Channel-Based Coordination Model for Component Composition," Mathematical Structures in Computer Science, vol 14(03), pp. 329-366, 2004.
- [2] C. Baier, M. Sirjani, F. Arbab, and J. Rutten "Modeling Component Connectors in Reo by Constraint Automata," Science of Computer Programming, vol 61(2), pp. 75-113, 2006.
- [3] N. Kokash, and F. Arbab, "Formal Behavioal Modeling and Compliance Analysis for Service-Oriented Systems," Formal Methods for Components and Objects, pp. 21-41, 2008.
- [4] M. R. Mousavi, M. Sirjani, and F. Arbab, "Formal Semantics and Analysis of Component Connectors in Reo," Electronic Notes in Theoretical Computer Science, vol. 154, pp. 83–99, 2006.
- [5] Reo homepage, http://reo.project.cwi.nl/reo/
 Animations from http://reo.project.cwi.nl/webreo/

Any Questions?

