

# It Wasn't Me, It Was the Prototype!

## Towards a Formal Model of JavaScript Prototype Pollution

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KTH



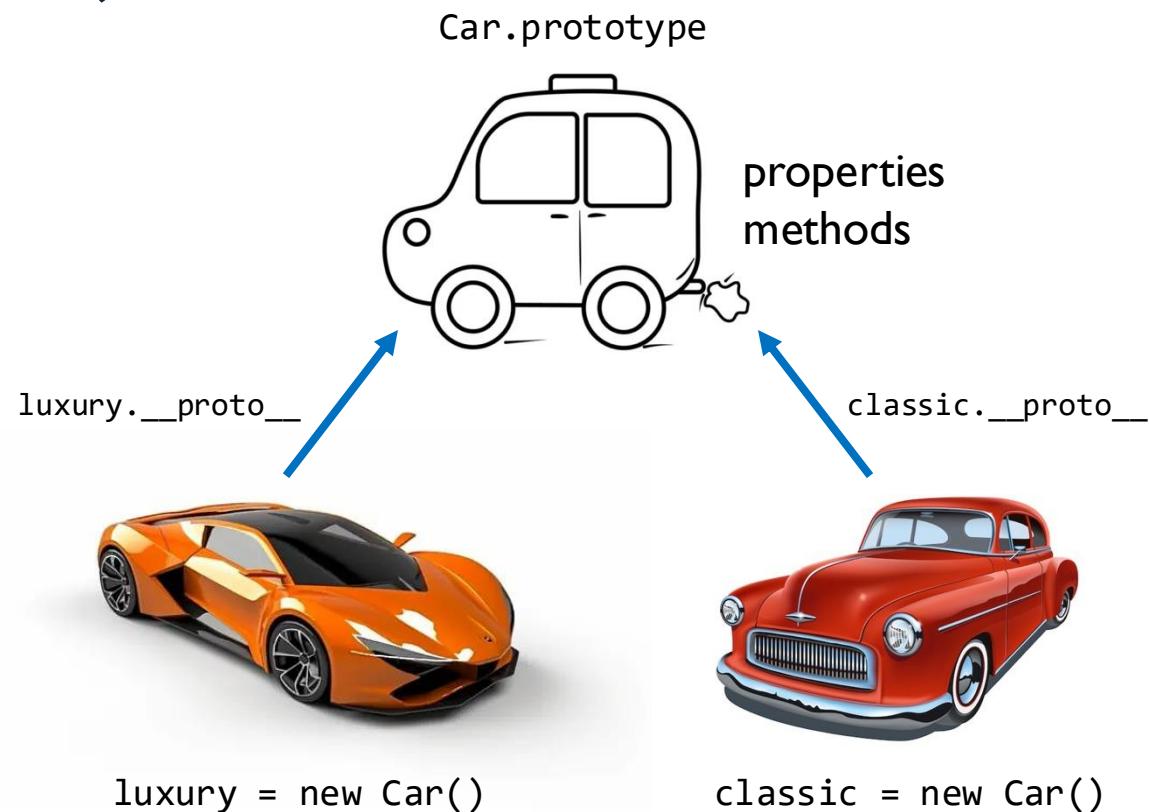
Chalmers

ShiftLeft Workshop, Gothenburg

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# JavaScript objects

- **Mutable** collection of properties
  - Values evaluated at **runtime**
- **Prototype: object blueprint**
  - Reusing existing objects
  - An object with a set of properties and functions **shared between all objects of the same type**
  - Exposed as regular programming construct  
`luxury["__proto__"], classic.__proto__, Car.prototype`



# Prototype-based Inheritance



snowy : Dog	
__proto__	●
age	5

Dog.prototype	
__proto__	●
numOfLegs	4

Dog	
prototype	●

Animal.prototype	
__proto__	●
eat()	function {...}

Animal	
prototype	●

snowy.age  
snowy.numOfLegs  
snowy.eat()  
snowy.toString()  
snowy.prop

Object.prototype	
__proto__	null
toString()	function {...}

Object	
prototype	●



# Prototype pollution (PP)

brutalicious : Dog	
__proto__	●
age	100

`brutalicious.__proto__.numOfLegs = 3`

snowy : Dog	
__proto__	●
age	5

Poor  
snowy! 😞



Dog.prototype	
__proto__	●
numOfLegs	3

Animal.prototype	
__proto__	●
eat()	function {...}

Object.prototype	
__proto__	null
toString()	function {...}

Object	
prototype	●

Dog	
prototype	●

Animal	
prototype	●

# Prototype pollution (cont.)

- A vulnerability where an attacker can modify an **object's prototype at runtime**
  - May then be inherited by user-defined objects
  - May then result in **unintended behavior**

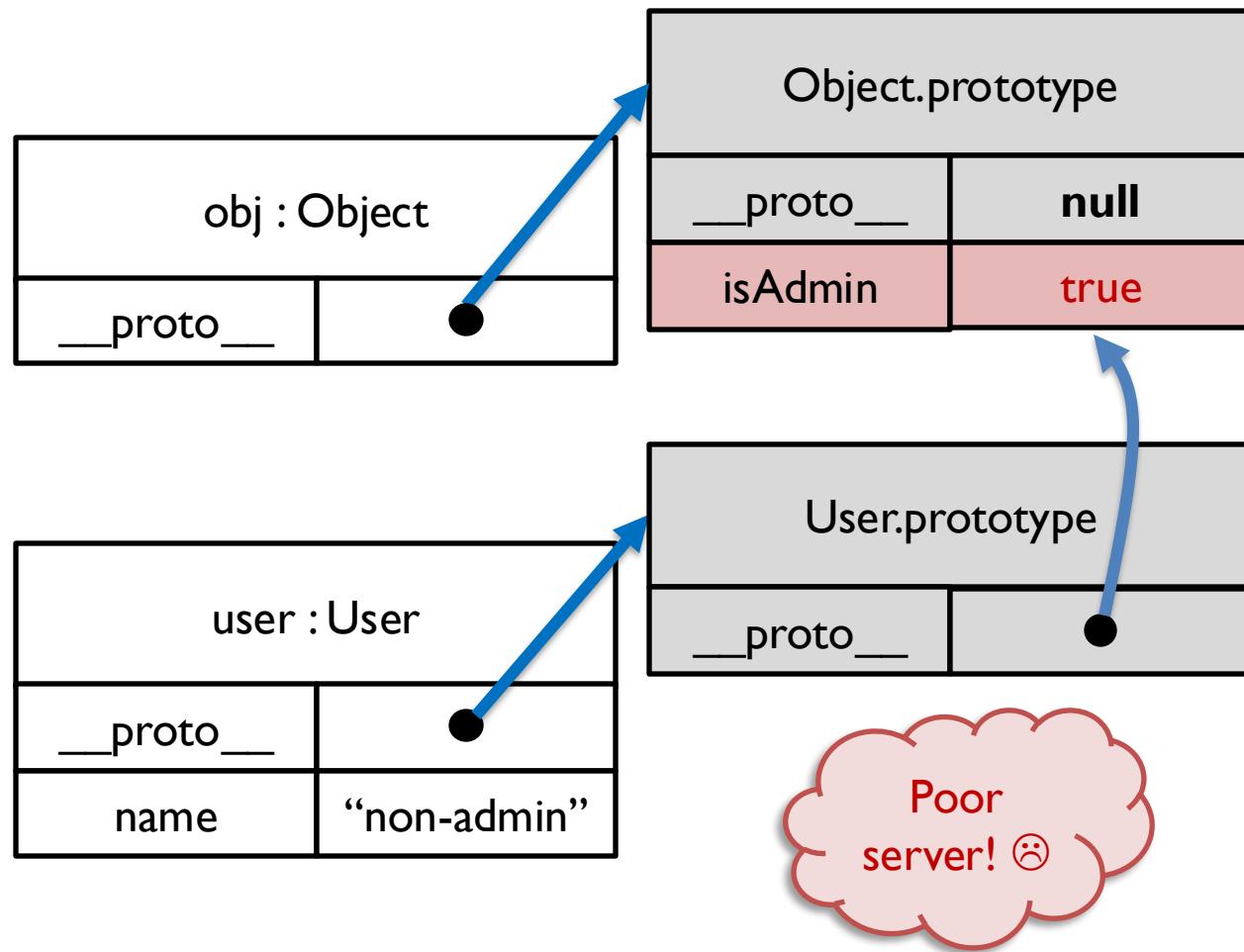


# Prototype pollution (cont.)

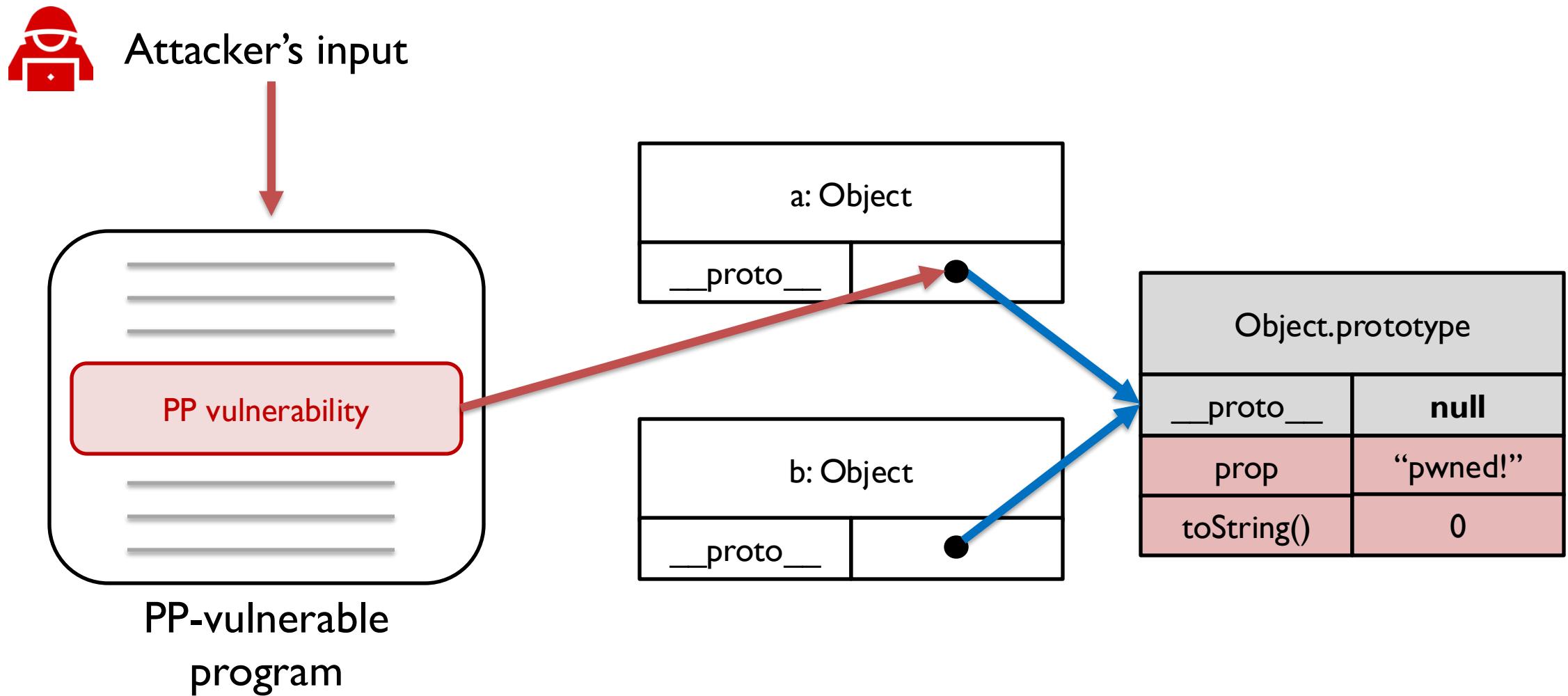
```
{ "__proto__": { "isAdmin": true } }
```



```
function checkAdmin(user) {
  if (user.isAdmin) {
    console.log("Welcome, admin user!!!");
  } else {
    console.log("Just a regular user...");
  }
}
...
let user = new User("non-admin");
...
let obj = JSON.parse(input);
checkAdmin(user);
```

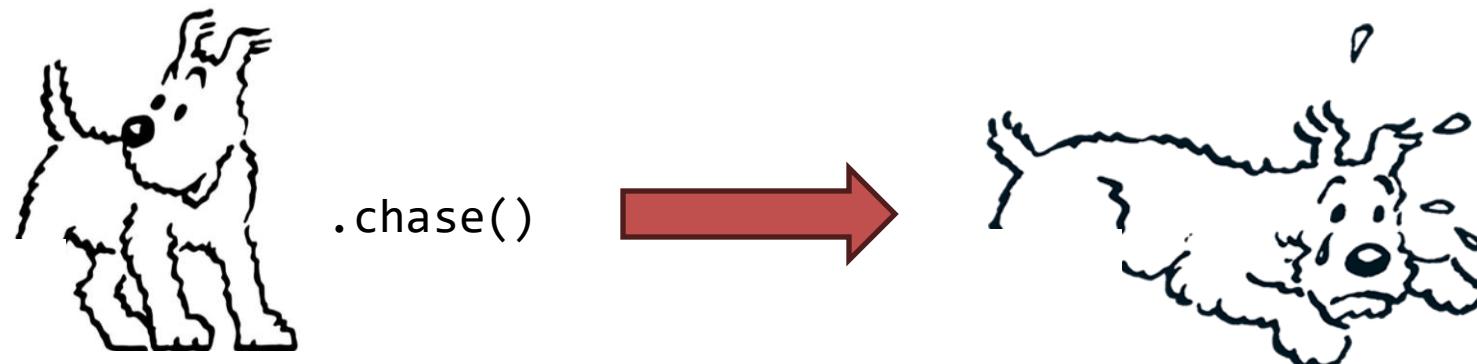


# PP-vulnerable program

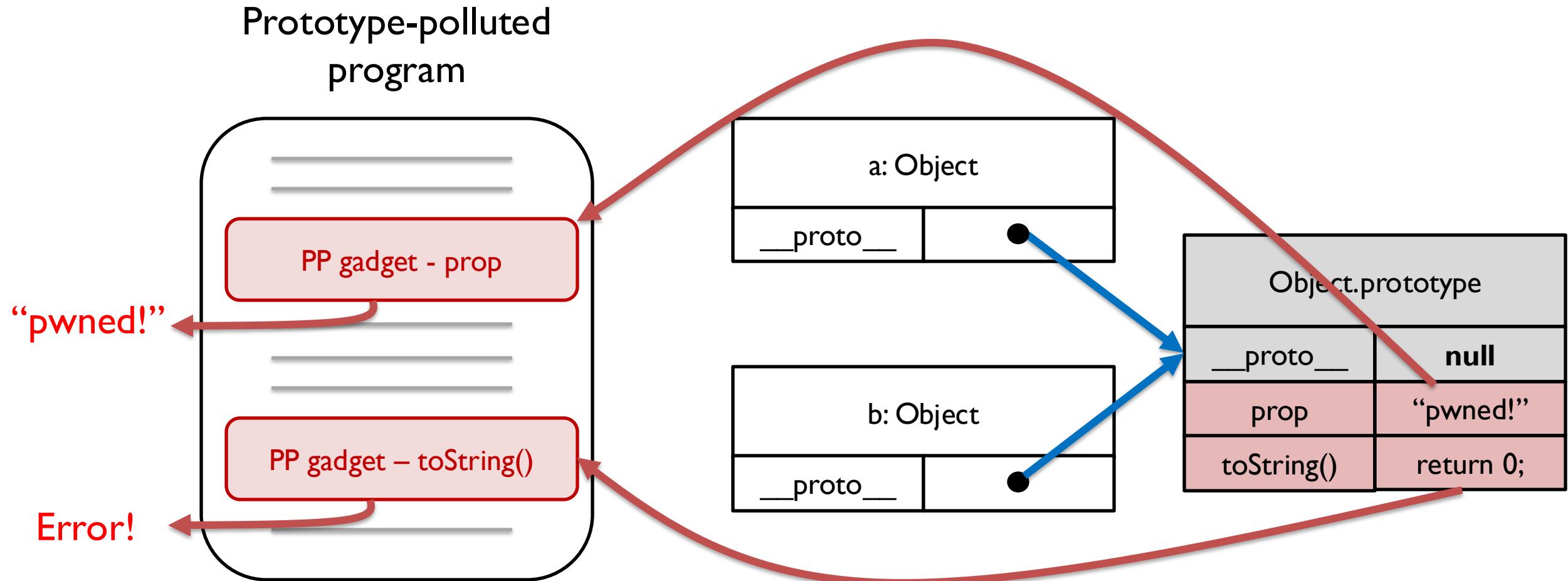


# Gadget

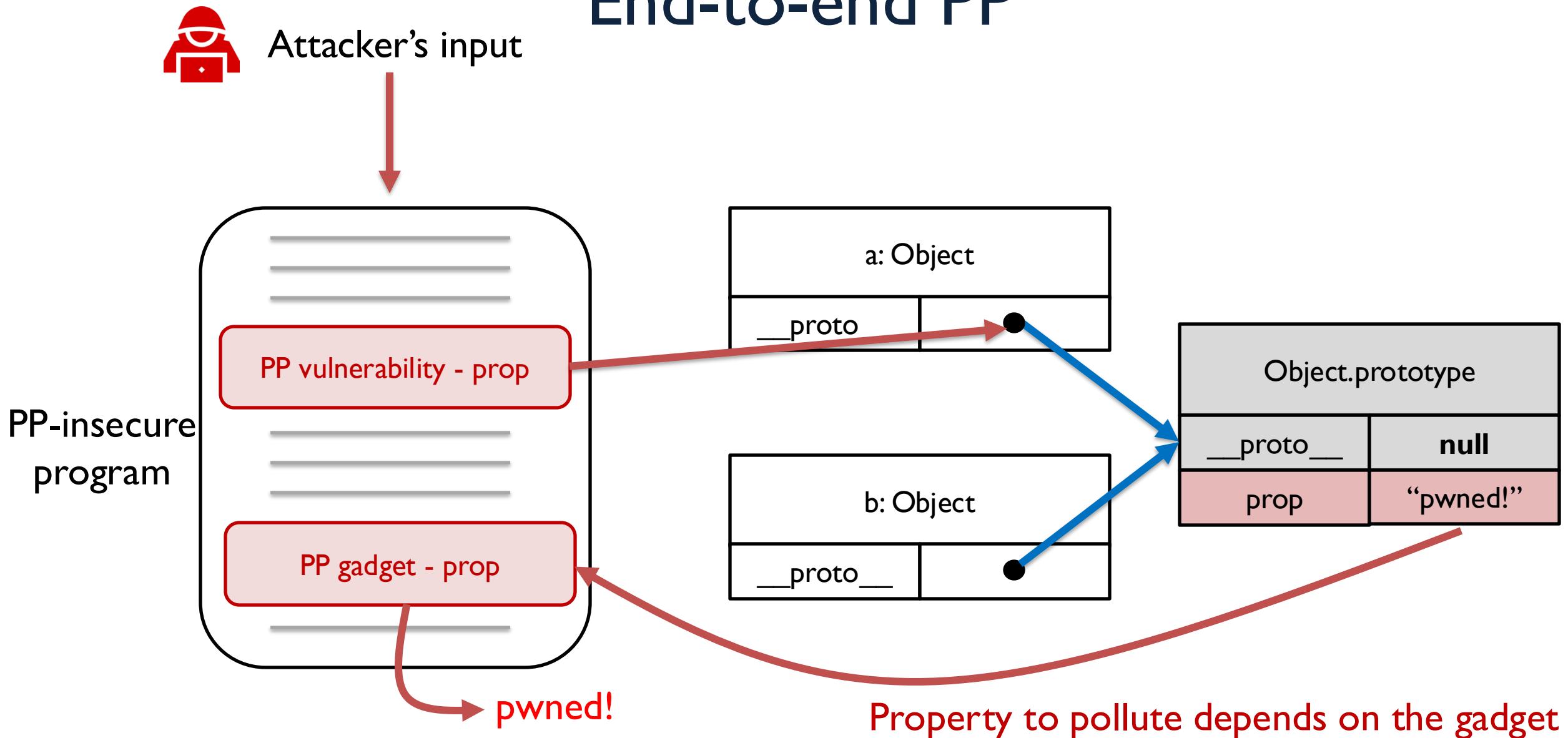
- PP's impact depends on the existence of **gadgets**
  - An **otherwise benign** piece of code which *inadvertently read from polluted properties* to **execute security-sensitive operations**
    - Triggering unintended behavior
- Examples: privilege elevation, reading secrets, RCE, DoS, log pollution



# Gadget (cont.)



# End-to-end PP



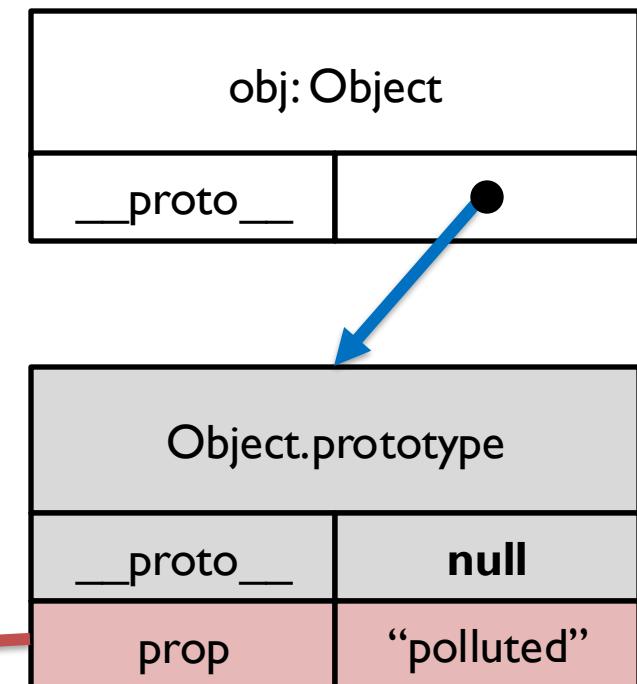
# Example 1

```
var inProto = "__proto__"  
var inProp = "prop"  
var inVal = "polluted"  
  
var obj = {}  
var p = obj[inProto]  
p[inProp] = inVal  
  
console.log({}.prop) // "polluted"
```

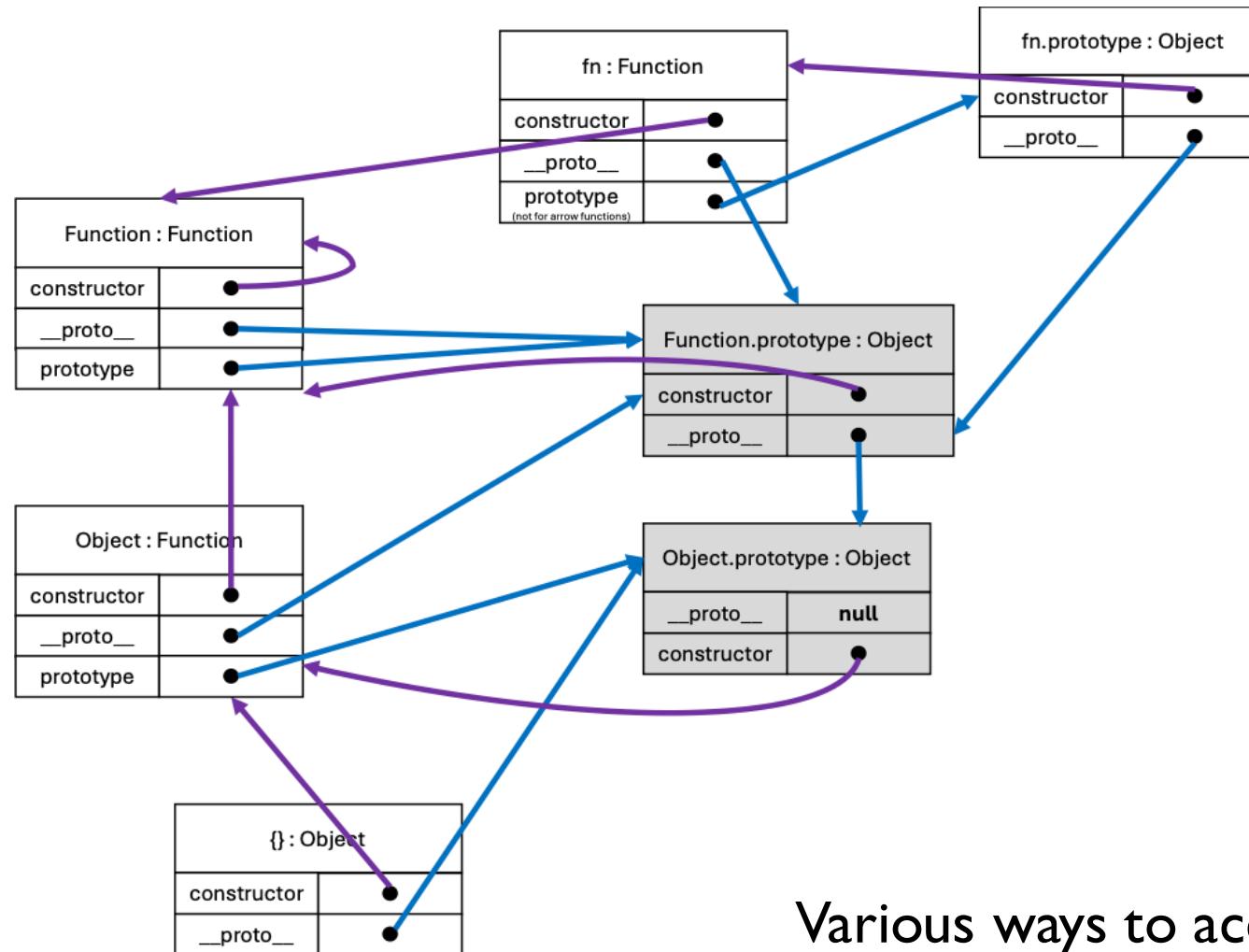
Malicious input

PP vulnerability

Gadget



# Who says JavaScript is hard to learn?! 😊

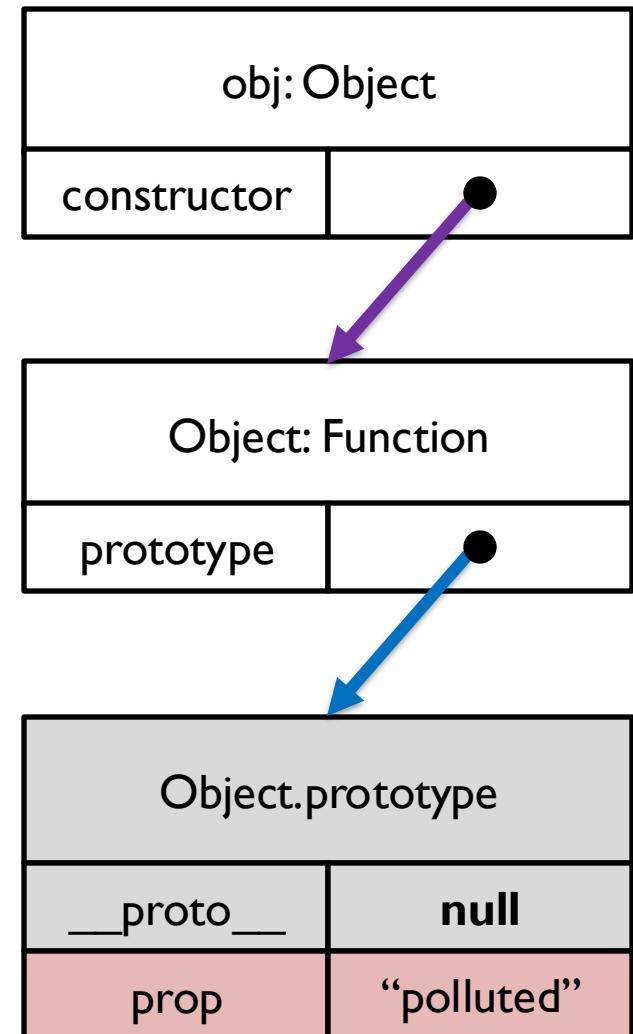
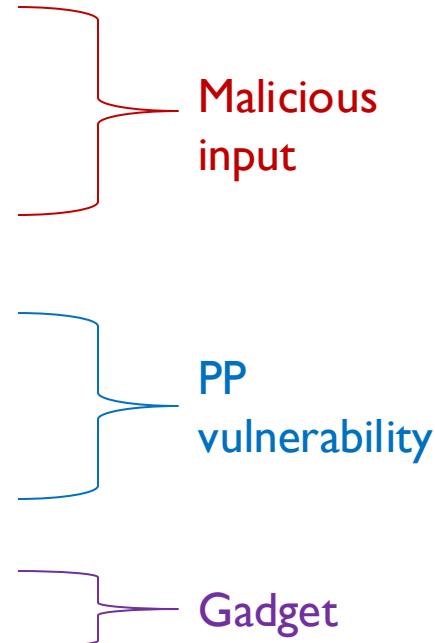


# Example 2

```
var inConstr = "constructor"
var inProto = "prototype"
var inProp = "prop"
var inVal = "polluted"

var obj = {}
var c = obj[inConstr]
var p = c[inProto]
p[inProp] = inVal

console.log({}.prop) // "polluted"
```



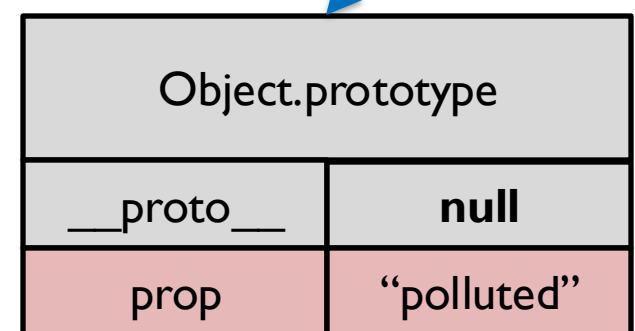
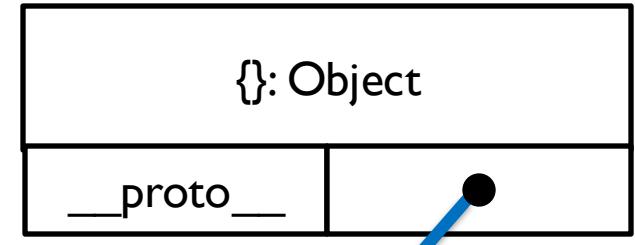
# Example 3

```
function merge(dst, src) {  
    for (let key in src) {  
        if (!src.hasOwnProperty(key)) continue;  
        if (typeof dst[key] === 'object') {  
            merge(dst[key], src[key]);  
        } else {  
            dst[key] = src[key];  
        }  
    }  
  
    var input = '{"__proto__": {"prop": "polluted"}}';  
    merge({}, JSON.parse(input));  
  
    console.log({}.prop); // "polluted"
```

PP  
vulnerability

Malicious  
input  
PP vulnerability

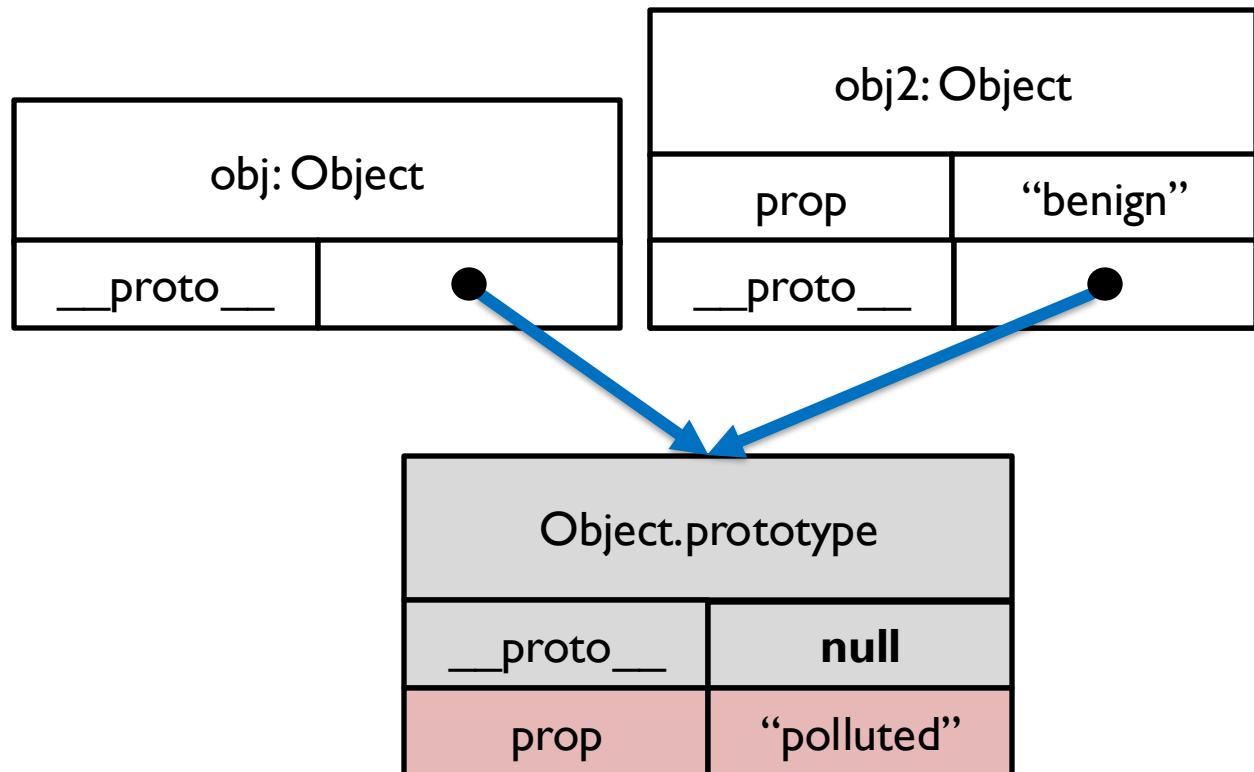
Gadget



# Example 4

```
var inProto = "__proto__"  
var inProp = "prop"  
var inVal = "polluted"  
  
var obj = {}  
var p = obj[inProto]  
p[inProp] = inVal  
  
var obj2 = {}  
obj2.prop = "benign"  
  
console.log(obj2.prop) // "benign"
```

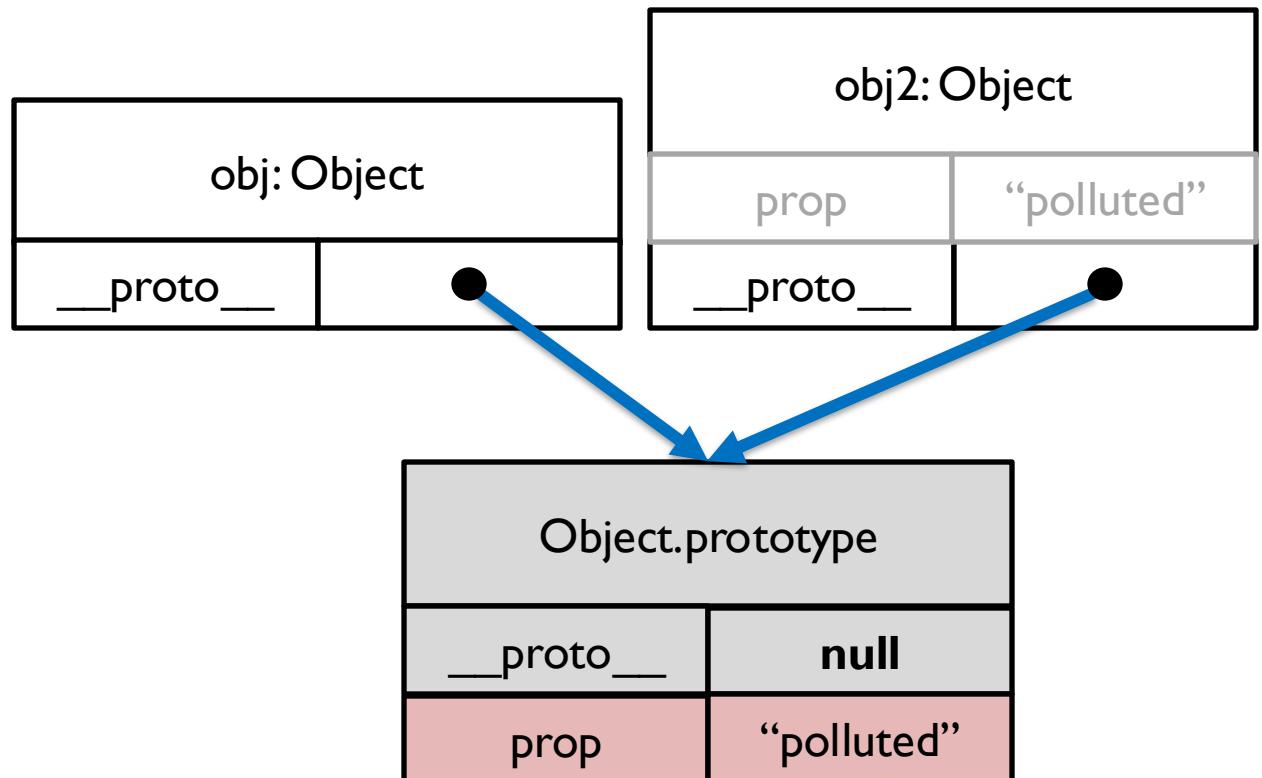
Observable behavior?



# Example 5

```
var inProto = "__proto__"  
var inProp = "prop"  
var inVar = true  
  
var obj, obj2 = {}  
if (inVar)  
  obj[inProto][inProp] = "polluted"  
else  
  obj2.prop = "polluted"  
  
console.log(obj2.prop) // "polluted"
```

Observable behavior?



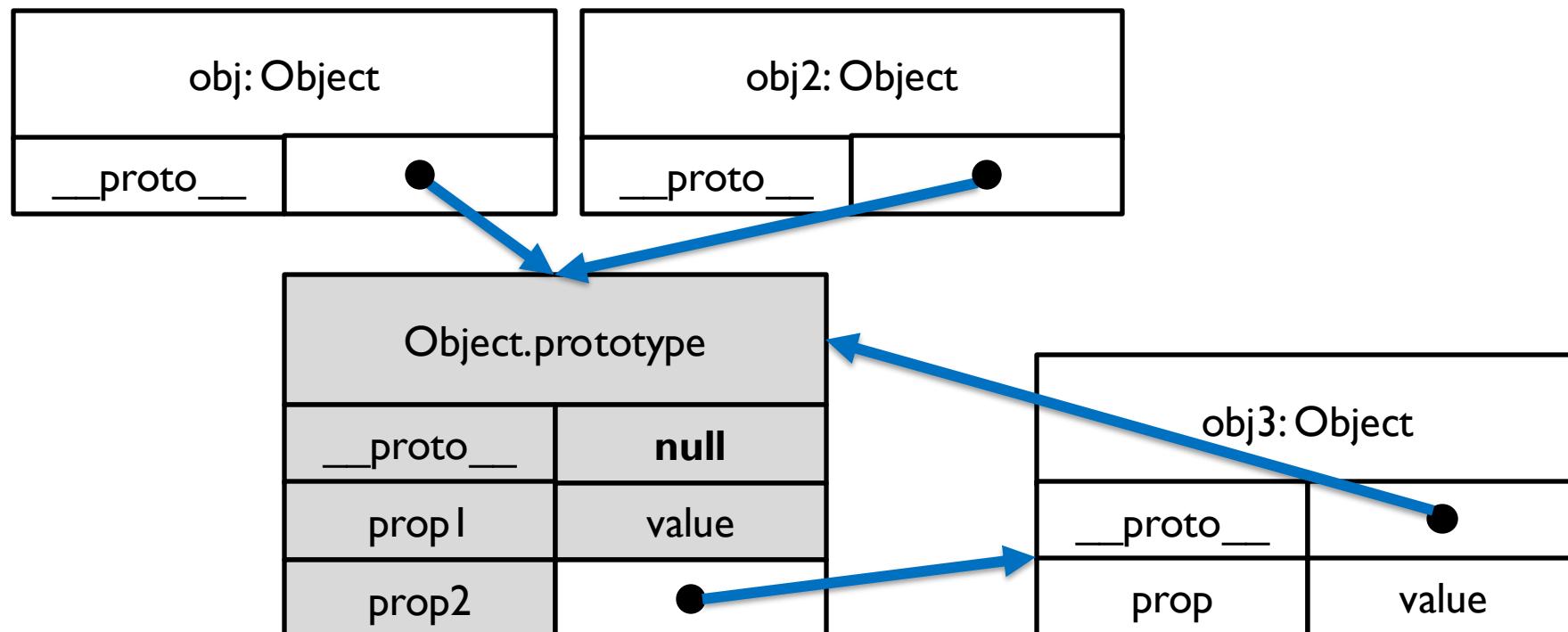
# From PP to unintended behavior

- The literature focuses on *detection* rather than *defining* PP and gadgets *formally*
  - Any untrusted access to `__proto__` is forbidden (*shallow definition*)
  - Any modification to an object property that is reachable from a prototype object is forbidden (*deep definition*)
- Gadgets
  - ACE, SSRF, privilege escalation, DoS, log pollution, cryptographic downgrade, ...

What are the formal definitions of **PP-secure** and **gadget-free** programs?

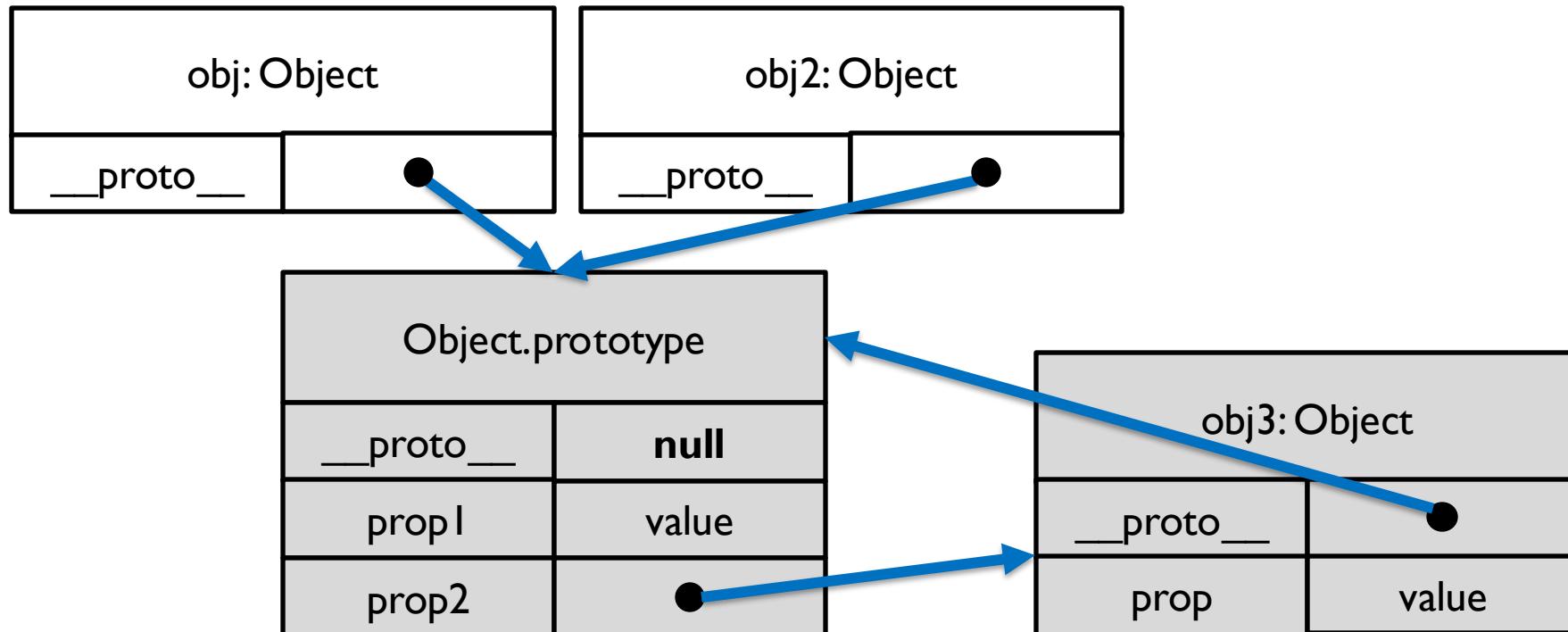
# Shallow definition

The gray cells are expected to remain same for any untrusted input



# Deep definition

The gray cells are expected to remain same for any untrusted input



# Model language

A subset of the non-strict semantics of ECMA-262 standard

$v ::= s \mid n \mid b \mid undefined \mid null$

$e ::= v \mid x \mid e \oplus e \mid e(e) \mid [e[e] \mid new\ e(e)] \mid function\ (x)\ c$

$i ::= x \mid [i[e]]$

$c ::= skip \mid [i = e] \mid if\ (e)\ c\ else\ c \mid while\ (e)\ c \mid c;\ c \mid return\ e \mid [out_L(e)]$

# PP security

**PP-secure( $c$ )**  $\stackrel{\text{def}}{=} \forall E_1, E_2, H_1, H_2.$

$$(E_1, H_1) \simeq_T (E_2, H_2) \wedge$$

$\Gamma \models \langle c, E_1, H_1, \emptyset \rangle \rightarrow^* \langle E_3, H_3, \tau_1 \rangle \wedge$

$\Gamma \models \langle c, E_2, H_2, \emptyset \rangle \rightarrow^* \langle E_4, H_4, \tau_2 \rangle$

$$\implies \tau_1 \simeq_{stut}^{proto} \tau_2.$$

Shallow vs.  
deep

# PP-Gadget freedom

$$\textbf{\textit{PP-Gadget-Free}}(c) \stackrel{\textit{def}}{=} \forall E_1, E_2, H_1, H_2, H_3^P, H_4^P.$$

$$\Gamma \models \langle c, E_1, H_1, \emptyset \rangle \rightarrow_v^* \langle E'_1, H'_1, \tau_1 \rangle \wedge$$

$$\Gamma \models \langle c, E_2, H_2, \emptyset \rangle \rightarrow_{v'}^* \langle E'_2, H'_2, \tau_2 \rangle \wedge$$

$$v \sim_\beta v'$$

$$\implies$$

$$\Gamma \models \langle c, E_1, [H_1^O \uplus H_3^P], \emptyset \rangle \rightarrow_{v''}^* \langle E''_1, H''_1, \tau_3 \rangle \wedge$$

$$\Gamma \models \langle c, E_2, [H_2^O \uplus H_4^P], \emptyset \rangle \rightarrow_{v'''}^* \langle E''_2, H''_2, \tau_4 \rangle \wedge$$

$$v'' \sim_\beta v'''.$$

# End-to-end PP

$E2E\text{-secure}(c) \stackrel{\text{def}}{=} \forall E_1, E_2, H_1, H_2.$

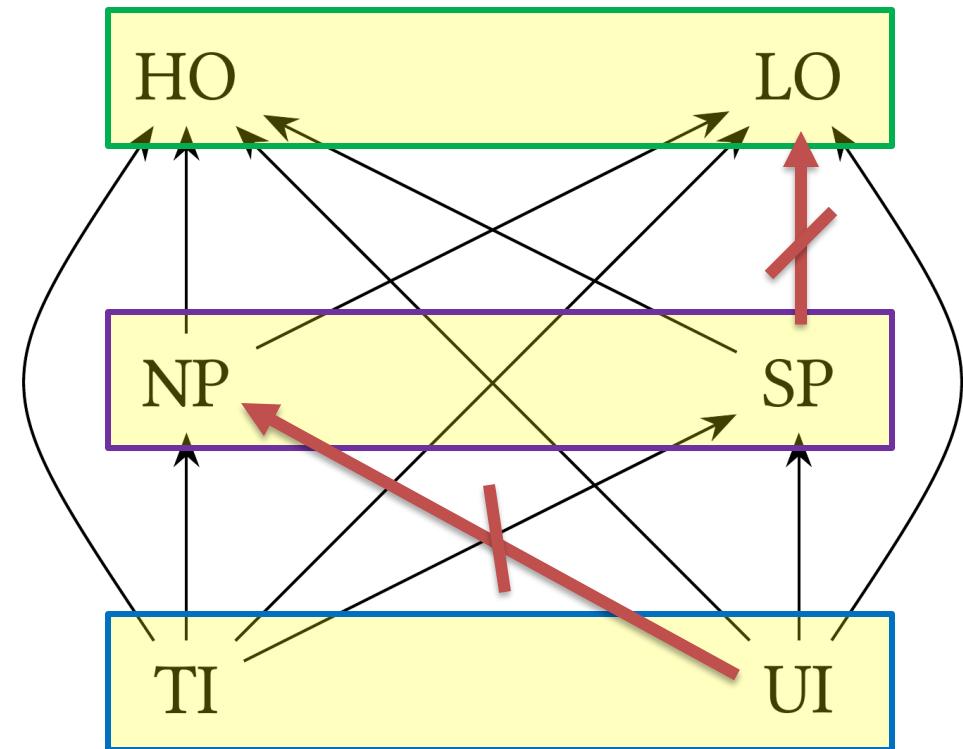
$$(E_1, H_1) \simeq_T (E_2, H_2) \wedge$$

$$H_1 \simeq^{\kappa, NP} H_2 \wedge$$

$\Gamma \models \langle c, E_1, H_1, \emptyset \rangle \rightarrow_v^* \langle E_3, H_3, \tau_1 \rangle \wedge$

$\Gamma \models \langle c, E_2, H_2, \emptyset \rangle \rightarrow_{v'}^* \langle E_4, H_4, \tau_2 \rangle$

$$\implies v \sim_{\beta}^{LO} v' \vee \tau_1 \simeq_{stut}^{\kappa, SP} \tau_2.$$



# Ongoing work

- Sound enforcement mechanisms
  - Runtime monitoring
- Non-root pollution
- Degree of exploitability and attacker's knowledge
- Modeling mitigations
  - `Object.freeze()`, `Object.seal()`
  - Use of `new Set()`, `new Map()`



# Takeaways

## Formal modeling of prototype-pollution security

