# *Platelet*

## Team Reference Material

### (unlimited version)

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# 目录

# Chapter 1

# Graph Theory

## 1.1  2-SAT (ct)

```cpp
struct Edge {
    Edge *next;
    int to;
} *last[maxn << 1], e[maxn << 2], *ecnt = e;
inline void link(int a, int b)
{
    *++ecnt = (Edge) {last[a], b}; last[a] = ecnt;
}
int dfn[maxn], low[maxn], timer, st[maxn], top, id[maxn], colcnt, n;
bool fail, used[maxn];
void tarjan(int x, int fa)
{
    dfn[x] = low[x] = ++timer; st[++top] = x;
    for (R Edge *iter = last[x]; iter; iter = iter -> next)
        if (iter -> to != fa)
        {
            if (!dfn[iter -> to])
            {
                tarjan(iter -> to, x);
                cmin(low[x], low[iter -> to]);
            }
            else if (!id[iter -> to]) cmin(low[x], dfn[iter -> to]);
        }
    if (dfn[x] == low[x])
    {
        ++colcnt; bool flag = 1;
        for (; ;)
        {
            int now = st[top--];
            id[now] = colcnt;
            if (now <= 2 * n)
            {
                flag &= !used[id[now <= n ? now + n : now - n]];
                now <= n ? fail |= (id[now + n] == id[now]) : fail |= (id[now - n] == id[now]);
            }
            if (now == x) break;
        }
        used[colcnt] = flag;
    }
}
int ans[maxn], tot;
int main()
```

```cpp
{
    /*
        build your graph here.
    */
    for (R int i = 1; !fail && i <= n; ++i) if (!dfn[i]) tarjan(i, 0);
    if (fail)
    {
        puts("Impossible");
        return 0;
    }
    for (R int i = 1; i <= n; ++i) if (used[id[i]]) ans[++tot] = i;
    printf("%d\n", tot);
    std::sort(ans + 1, ans + tot + 1);
    for (R int i = 1; i <= tot; ++i) printf("%d ", ans[i]);
    return 0;
}
```

## 1.2   割点与桥 (ct)

### 割点

```cpp
int dfn[maxn], low[maxn], timer, ans, num;
void tarjan(int x, int fa)
{
    dfn[x] = low[x] = ++timer;
    for (Edge *iter = last[x]; iter; iter = iter -> next)
        if (iter -> to != fa)
        {
            if (!dfn[iter -> to])
            {
                tarjan(iter -> to, x);
                cmin(low[x], low[iter -> to]);
                if (dfn[x] <= low[iter -> to])
                {
                    cut[x] = 1;
                    if (!fa && dfn[x] < low[iter -> to]) num = 233;
                    else if (!fa) ++num;
                }
            }
            else cmin(low[x], dfn[iter -> to]);
        }
}
int main()
{
    for (int i = 1; i <= n; ++i)
        if (!dfn[i])
        {
            num = 0;
            tarjan(i, 0);
            if (num == 1) cut[i] = 0;
        }
}
```

### 桥

```cpp
int dfn[maxn], low[maxn], timer;
void tarjan(int x, int fa)
{
```

```
 4        dfn[x] = low[x] = ++timer;
 5        for (R Edge *iter = last[x]; iter; iter = iter -> next)
 6            if (iter -> to != fa)
 7            {
 8                if (!dfn[iter -> to])
 9                {
10                    dfs(iter -> to, x);
11                    cmin(low[x], low[iter -> to]);
12                    if (dfn[x] < low[iter -> to]) ans[x][iter -> to] = ans[iter -> to][x] = 1;
13                }
14                else cmin(low[x], dfn[iter -> to]);
15            }
16 }
```

## 1.3   Steiner tree (lhy)

```
 1 void Steiner_Tree()
 2 {
 3     memset(f, 0x3f, sizeof(f));
 4     for(int i = 1; i <= n; i++)
 5         f[0][i] = 0;
 6     for(int i = 1; i <= p; i++)
 7         f[1 << (i - 1)][idx[i]] = 0;
 8     int S = 1 << p;
 9     for(int s = 1; s < S; s++)
10     {
11         for(int i = 1; i <= n; i++)
12         {
13             for(int k = (s - 1) & s; k; k = (k - 1) & s)
14                 f[s][i] = min(f[s][i], f[k][i] + f[s ^ k][i]);
15         }
16         SPFA(f[s]);
17     }
18     int ans = inf;
19     for(int i = 1; i <= n; i++)
20         ans = min(ans, f[S - 1][i]);
21 }
```

## 1.4   K 短路 (lhy)

```
 1 const int MAXNODE = MAXN + MAXM * 2;

 2 bool used[MAXN];
 3 int n, m, cnt, S, T, Kth, N, TT;
 4 int rt[MAXN], seq[MAXN], adj[MAXN], from[MAXN], dep[MAXN];
 5 LL dist[MAXN], w[MAXM], ans[MAXK];

 6 struct GivenEdge{
 7     int u, v, w;
 8     GivenEdge() {};
 9     GivenEdge(int _u, int _v, int _w) : u(_u), v(_v), w(_w){};
10 }edge[MAXM];

11 struct Edge{
12     int v, nxt, w;
13     Edge() {};
14     Edge(int _v, int _nxt, int _w) : v(_v), nxt(_nxt), w(_w) {};
15 }e[MAXM];
```

```
16  inline void addedge(int u, int v, int w)
17  {
18      e[++cnt] = Edge(v, adj[u], w); adj[u] = cnt;
19  }

20  void dij(int S)
21  {
22      for(int i = 1; i <= N; i++)
23      {
24          dist[i] = INF;
25          dep[i] = 0x3f3f3f3f;
26          used[i] = false;
27          from[i] = 0;
28      }
29      static priority_queue<pair<LL, int>, vector<pair<LL, int> >, greater<pair<LL, int> > > hp;
30      while(!hp.empty())hp.pop();
31      hp.push(make_pair(dist[S] = 0, S));
32      dep[S] = 1;
33      while(!hp.empty())
34      {
35          pair<LL, int> now = hp.top();
36          hp.pop();
37          int u = now.second;
38          if(used[u])continue;
39          else used[u] = true;
40          for(int p = adj[u]; p; p = e[p].nxt)
41          {
42              int v = e[p].v;
43              if(dist[u] + e[p].w < dist[v])
44              {
45                  dist[v] = dist[u] + e[p].w;
46                  dep[v] = dep[u] + 1;
47                  from[v] = p;
48                  hp.push(make_pair(dist[v], v));
49              }
50          }
51      }
52      for(int i = 1; i <= m; i++)    w[i] = 0;
53      for(int i = 1; i <= N; i++)
54          if(from[i])w[from[i]] = -1;
55      for(int i = 1; i <= m; i++)
56      {
57          if(~w[i] && dist[edge[i].u] < INF && dist[edge[i].v] < INF)
58          {
59              w[i] = -dist[edge[i].u] + (dist[edge[i].v] + edge[i].w);
60          }
61          else
62          {
63              w[i] = -1;
64          }
65      }
66  }

67  inline bool cmp_dep(int p, int q)
68  {
69      return dep[p] < dep[q];
70  }

71  struct Heap{
72      LL key;
```

```
 73        int id, lc, rc, dist;
 74        Heap() {};
 75        Heap(LL k, int i, int l, int r, int d) : key(k), id(i), lc(l), rc(r), dist(d) {};
 76        inline void clear()
 77        {
 78            key = 0;
 79            id = lc = rc = dist = 0;
 80        }
 81   }hp[MAXNODE];

 82   inline int merge_simple(int u, int v)
 83   {
 84        if(!u)return v;
 85        if(!v)return u;
 86        if(hp[u].key > hp[v].key)
 87        {
 88            swap(u, v);
 89        }
 90        hp[u].rc = merge_simple(hp[u].rc, v);
 91        if(hp[hp[u].lc].dist < hp[hp[u].rc].dist)
 92        {
 93            swap(hp[u].lc, hp[u].rc);
 94        }
 95        hp[u].dist = hp[hp[u].rc].dist + 1;
 96        return u;
 97   }

 98   inline int merge_full(int u, int v)
 99   {
100        if(!u)return v;
101        if(!v)return u;
102        if(hp[u].key > hp[v].key)
103        {
104            swap(u, v);
105        }
106        int nownode = ++cnt;
107        hp[nownode] = hp[u];
108        hp[nownode].rc = merge_full(hp[nownode].rc, v);
109        if(hp[hp[nownode].lc].dist < hp[hp[nownode].rc].dist)
110        {
111            swap(hp[nownode].lc, hp[nownode].rc);
112        }
113        hp[nownode].dist = hp[hp[nownode].rc].dist + 1;
114        return nownode;
115   }

116   priority_queue<pair<LL, int>, vector<pair<LL, int> >, greater<pair<LL, int> > > Q;

117   int main()
118   {
119        while(scanf("%d%d", &n, &m) != EOF)
120        {
121            scanf("%d%d%d%d", &S, &T, &Kth, &TT);
122            for(int i = 1; i <= m; i++)
123            {
124                int u, v, w;
125                scanf("%d%d%d", &u, &v, &w);
126                edge[i] = {u, v, w};
127            }
128            N = n;
129            memset(adj, 0, sizeof(*adj) * (N + 1));
```

```
130         cnt = 0;
131         for(int i = 1; i <= m; i++)
132             addedge(edge[i].v, edge[i].u, edge[i].w);
133         dij(T);
134         if(dist[S] > TT)
135         {
136             puts("Whitesnake!");
137             continue;
138         }
139         for(int i = 1; i <= N; i++)
140             seq[i] = i;
141         sort(seq + 1, seq + N + 1, cmp_dep);

142         cnt = 0;
143         memset(adj, 0, sizeof(*adj) * (N + 1));
144         memset(rt, 0, sizeof(*rt) * (N + 1));
145         for(int i = 1; i <= m; i++)
146             addedge(edge[i].u, edge[i].v, edge[i].w);
147         rt[T] = cnt = 0;
148         hp[0].dist = -1;
149         for(int i = 1; i <= N; i++)
150         {
151             int u = seq[i], v = edge[from[u]].v;
152             rt[u] = 0;
153             for(int p = adj[u]; p; p = e[p].nxt)
154             {
155                 if(~w[p])
156                 {
157                     hp[++cnt] = Heap(w[p], p, 0, 0, 0);
158                     rt[u] = merge_simple(rt[u], cnt);
159                 }
160             }
161             if(i == 1)continue;
162             rt[u] = merge_full(rt[u], rt[v]);
163         }
164         while(!Q.empty())Q.pop();
165         Q.push(make_pair(dist[S], 0));
166         edge[0].v = S;
167         for(int kth = 1; kth <= Kth; kth++)
168         {
169             if(Q.empty())
170             {
171                 ans[kth] = -1;
172                 continue;
173             }
174             pair<LL, int> now = Q.top(); Q.pop();
175             ans[kth] = now.first;
176             int p = now.second;
177             if(hp[p].lc)
178             {
179                 Q.push(make_pair(+hp[hp[p].lc].key + now.first - hp[p].key, hp[p].lc));
180             }
181             if(hp[p].rc)
182             {
183                 Q.push(make_pair(+hp[hp[p].rc].key + now.first - hp[p].key, hp[p].rc));
184             }
185             if(rt[edge[hp[p].id].v])
186             {
187                 Q.push(make_pair(hp[rt[edge[hp[p].id].v]].key + now.first, rt[edge[hp[p].id].v]));
188             }
189     }
```

```
190          if(ans[Kth] == -1 || ans[Kth] > TT)
191          {
192              puts("Whitesnake!");
193          }
194          else
195          {
196              puts("yareyaredawa");
197          }
198      }
199 }
```

# 1.5  最大团

# 1.6  一般图最大匹配

# 1.7  KM 算法 (Nightfall)

$O(n^3)$，1-based，最大权匹配
不存在的边权值开到 $-n \times (|MAXV|)$，$\infty$ 为 $3n \times (|MAXV|)$
匹配为 $(lk_i, i)$

```
1  long long KM(int n, long long w[N][N])
2  {
3      long long ans = 0;
4      int x, py, p;
5      long long d;
6      for(int i = 1; i <= n; i++)
7          lx[i] = ly[i] = 0, lk[i] = -1;
8      for(int i = 1; i <= n; i++)
9          for(int j = 1; j <= n; j++)
10             lx[i] = max(lx[i], w[i][j]);
11     for(int i = 1; i <= n; i++)
12     {
13         for(int j = 1; j <= n; j++)
14             slk[j] = inf, vy[j] = 0;
15         for(lk[py = 0] = i; lk[py]; py = p)
16         {
17             vy[py] = 1; d = inf; x = lk[py];
18             for(int y = 1; y <= n; y++)
19                 if(!vy[y])
20                 {
21                     if(lx[x] + ly[y] - w[x][y] < slk[y])
22                         slk[y] = lx[x] + ly[y] - w[x][y], pre[y] = py;
23                     if(slk[y] < d)d = slk[y], p = y;
24                 }
25             for(int y = 0; y <= n; y++)
26                 if(vy[y])lx[lk[y]] -= d, ly[y] += d;
27                 else slk[y] -= d;
28         }
29         for(; py; py = pre[py])lk[py] = lk[pre[py]];
30     }
31     for(int i = 1; i <= n; i++)
32         ans += lx[i] + ly[i];
33     return ans;
34 }
```

## 1.8   支配树 (Nightfall,ct)

### DAG (ct)

```
1  struct Edge {
2      Edge *next;
3      int to;
4  } ;
5  Edge *last[maxn], e[maxm], *ecnt = e; // original graph
6  Edge *rlast[maxn], re[maxm], *recnt = re; // reversed-edge graph
7  Edge *tlast[maxn], te[maxn << 1], *tecnt = te; // dominate tree graph
8  int deg[maxn], q[maxn], fa[maxn][20], all_fa[maxn], fa_cnt, size[maxn], dep[maxn];
9  inline void link(int a, int b)
10 {
11     *++ecnt = (Edge) {last[a], b}; last[a] = ecnt; ++deg[b];
12 }
13 inline void link_rev(R int a, R int b)
14 {
15     *++recnt = (Edge) {rlast[a], b}; rlast[a] = recnt;
16 }
17 inline void link_tree(R int a, R int b)
18 {
19     *++tecnt = (Edge) {tlast[a], b}; tlast[a] = tecnt;
20 }
21 inline int getlca(R int a, R int b)
22 {
23     if (dep[a] < dep[b]) std::swap(a, b);
24     R int temp = dep[a] - dep[b];
25     for (R int i; temp; temp -= 1 << i)
26         a = fa[a][i = __builtin_ctz(temp)];
27     for (R int i = 16; ~i; --i)
28         if (fa[a][i] != fa[b][i])
29             a = fa[a][i], b = fa[b][i];
30     if (a == b) return a;
31     return fa[a][0];
32 }
33 void dfs(R int x)
34 {
35     size[x] = 1;
36     for (R Edge *iter = tlast[x]; iter; iter = iter -> next)
37         dfs(iter -> to), size[x] += size[iter -> to];
38 }
39 int main()
40 {
41     q[1] = 0;
42     R int head = 0, tail = 1;
43     while (head < tail)
44     {
45         R int now = q[++head];
46         fa_cnt = 0;
47         for (R Edge *iter = rlast[now]; iter; iter = iter -> next)
48             all_fa[++fa_cnt] = iter -> to;
49         for (; fa_cnt > 1; --fa_cnt)
50             all_fa[fa_cnt - 1] = getlca(all_fa[fa_cnt], all_fa[fa_cnt - 1]);
51         fa[now][0] = all_fa[fa_cnt];
52         dep[now] = dep[all_fa[fa_cnt]] + 1;
53         if (now) link_tree(fa[now][0], now);
54
55         for (R int i = 1; i <= 16; ++i)
56             fa[now][i] = fa[fa[now][i - 1]][i - 1];
```

```
56          for (R Edge *iter = last[now]; iter; iter = iter -> next)
57              if (--deg[iter -> to] == 0) q[++tail] = iter -> to;
58      }
59      dfs(0);
60      for (R int i = 1; i <= n; ++i) printf("%d\n", size[i] - 1 );
61      return 0;
62  }
```

## 一般图 (Nightfall)

```
1  struct Dominator_Tree{
2      int n, s, cnt;
3      int dfn[N], id[N], pa[N], semi[N], idom[N], p[N], mn[N];
4      vector<int> e[N], dom[N], be[N];
5
6      void ins(int x, int y){e[x].push_back(y);}
7
8      void dfs(int x)
9      {
10         dfn[x] = ++cnt; id[cnt] = x;
11         for(auto i:e[x])
12         {
13             if(!dfn[i])dfs(i), pa[dfn[i]] = dfn[x];
14             be[dfn[i]].push_back(dfn[x]);
15         }
16     }
17
18     int get(int x)
19     {
20         if(p[x] != p[p[x]])
21         {
22             if(semi[mn[x]] > semi[get(p[x])])mn[x] = get(p[x]);
23             p[x] = p[p[x]];
24         }
25         return mn[x];
26     }
27
28     void LT()
29     {
30         for(int i = cnt; i > 1; i--)
31         {
32             for(auto j:be[i])semi[i] = min(semi[i], semi[get(j)]);
33             dom[semi[i]].push_back(i);
34             int x = p[i] = pa[i];
35             for(auto j:dom[x])
36                 idom[j] = (semi[get(j)] < x ? get(j) : x);
37             dom[x].clear();
38         }
39         for(int i = 2; i <= cnt; i++)
40         {
41             if(idom[i] != semi[i])idom[i] = idom[idom[i]];
42             dom[id[idom[i]]].push_back(id[i]);
43         }
44     }
45     void build()
46     {
47         for(int i = 1; i <= n; i++)
48             dfn[i] = 0, dom[i].clear(), be[i].clear(), p[i] = mn[i] = semi[i] = i;
49         cnt = 0, dfs(s), LT();
50     }
```

```
47  };
```

## 1.9   虚树 (ct)

```
1   struct Edge {
2       Edge *next;
3       int to;
4   } *last[maxn], e[maxn << 1], *ecnt = e;
5   inline void link(int a, int b)
6   {
7       *++ecnt = (Edge) {last[a], b}; last[a] = ecnt;
8       *++ecnt = (Edge) {last[b], a}; last[b] = ecnt;
9   }
10  int a[maxn], n, dfn[maxn], pos[maxn], timer, inv[maxn], st[maxn];
11  int fa[maxn], size[maxn], dep[maxn], son[maxn], top[maxn];
12  bool vis[maxn];
13  void dfs1(int x)
14  {
15      vis[x] = 1; size[x] = 1; dep[x] = dep[fa[x]] + 1;
16      for (R Edge *iter = last[x]; iter; iter = iter -> next)
17          if (!vis[iter -> to])
18          {
19              fa[iter -> to] = x;
20              dfs1(iter -> to);
21              size[x] += size[iter -> to];
22              size[son[x]] < size[iter -> to] ? son[x] = iter -> to : 0;
23          }
24  }
25  void dfs2(int x)
26  {
27      vis[x] = 0; top[x] = x == son[fa[x]] ? top[fa[x]] : x;
28      dfn[x] = ++timer; pos[timer] = x;
29      if (son[x]) dfs2(son[x]);
30      for (R Edge *iter = last[x]; iter; iter = iter -> next)
31          if (vis[iter -> to]) dfs2(iter -> to);
32      inv[x] = timer;
33  }
34  inline int getlca(int a, int b)
35  {
36      while (top[a] != top[b])
37          dep[top[a]] < dep[top[b]] ? b = fa[top[b]] : a = fa[top[a]];
38      return dep[a] < dep[b] ? a : b;
39  }
40  inline bool cmp(int a, int b)
41  {
42      return dfn[a] < dfn[b];
43  }
44  inline bool isson(int a, int b)
45  {
46      return dfn[a] <= dfn[b] && dfn[b] <= inv[a];
47  }
48  typedef long long ll;
49  bool imp[maxn];
50  struct sEdge {
51      sEdge *next;
52      int to, w;
53  } *slast[maxn], se[maxn << 1], *secnt = se;
54  inline void slink(int a, int b, int w)
55  {
```

```
56      *++secnt = (sEdge) {slast[a], b, w}; slast[a] = secnt;
57 }
58 int main()
59 {
60      scanf("%d", &n);
61      for (int i = 1; i < n; ++i)
62      {
63          int a, b; scanf("%d%d", &a, &b);
64          link(a, b);
65      }
66      int m; scanf("%d", &m);
67      dfs1(1); dfs2(1);
68      memset(size, 0, (n + 1) << 2);
69      for (; m; --m)
70      {
71          int top = 0; scanf("%d", &k);
72          for (int i = 1; i <= k; ++i) scanf("%d", &a[i]), vis[a[i]] = imp[a[i]] = 1;
73          std::sort(a + 1, a + k + 1, cmp);
74          int p = k;
75          for (int i = 1; i < k; ++i)
76          {
77              int lca = getlca(a[i], a[i + 1]);
78              if (!vis[lca]) vis[a[++p]] = lca] = 1;
79          }
80          std::sort(a + 1, a + p + 1, cmp);
81          st[++top] = a[1];
82          for (int i = 2; i <= p; ++i)
83          {
84              while (!isson(st[top], a[i])) --top;
85              slink(st[top], a[i], dep[a[i]] - dep[st[top]]);
86              st[++top] = a[i];
87          }
88          /*
89              write your code here.
90          */
91          for (int i = 1; i <= p; ++i) vis[a[i]] = imp[a[i]] = 0, slast[a[i]] = 0;
92          secnt = se;
93      }
94      return 0;
95 }
```

## 1.10   点分治 (ct)

```
1  int root, son[maxn], size[maxn], sum;
2  bool vis[maxn];
3  void dfs_root(int x, int fa)
4  {
5      size[x] = 1; son[x] = 0;
6      for (R Edge *iter = last[x]; iter; iter = iter -> next)
7      {
8          if (iter -> to == fa || vis[iter -> to]) continue;
9          dfs_root(iter -> to, x);
10         size[x] += size[iter -> to];
11         cmax(son[x], size[iter -> to]);
12     }
13     cmax(son[x], sum - size[x]);
14     if (!root || son[x] < son[root]) root = x;
15 }
16 void dfs_chain(int x, int fa, int st1, int st2)
```

```
17  {
18      /*
19          write your code here.
20      */
21      for (Edge *iter = last[x]; iter; iter = iter -> next)
22      {
23          if (vis[iter -> to] || iter -> to == fa) continue;
24          dfs_chain(iter -> to, x);
25      }
26  }
27  void calc(int x)
28  {
29      for (Edge *iter = last[x]; iter; iter = iter -> next)
30      {
31          if (vis[iter -> to]) continue;
32          dfs_chain(iter -> to, x);
33          /*
34              write your code here.
35          */
36      }
37  }
38  void work(int x)
39  {
40      vis[x] = 1;
41      calc(x);
42      for (R Edge *iter = last[x]; iter; iter = iter -> next)
43      {
44          if (vis[iter -> to]) continue;
45          root = 0;
46          sum = size[iter -> to];
47          dfs_root(iter -> to, 0);
48          work(root);
49      }
50  }
51  int main()
52  {
53      root = 0; sum = n;
54      dfs_root(1, 0);
55      work(root);
56      return 0;
57  }
```

## 1.11　树上倍增 (ct)

```
1   int fa[maxn][17], mn[maxn][17], dep[maxn];
2   bool vis[maxn];
3   void dfs(int x)
4   {
5       vis[x] = 1;
6       for (int i = 1; i <= 16; ++i)
7       {
8           if (dep[x] < (1 << i)) break;
9           fa[x][i] = fa[fa[x][i - 1]][i - 1];
10          mn[x][i] = dmin(mn[x][i - 1], mn[fa[x][i - 1]][i - 1]);
11      }
12      for (Edge *iter = last[x]; iter; iter = iter -> next)
13          if (!vis[iter -> to])
14          {
15              fa[iter -> to][0] = x;
```

12

```
16        mn[iter -> to][0] = iter -> w;
17        dep[iter -> to] = dep[x] + 1;
18        dfs(iter -> to);
19      }
20 }
21 inline int getlca(int x, int y)
22 {
23     if (dep[x] < dep[y]) std::swap(x, y);
24     int t = dep[x] - dep[y];
25     for (int i = 0; i <= 16 && t; ++i)
26         if ((1 << i) & t)
27             x = fa[x][i], t ^= 1 << i;
28     for (int i = 16; i >= 0; --i)
29         if (fa[x][i] != fa[y][i])
30         {
31             x = fa[x][i];
32             y = fa[y][i];
33         }
34     if (x == y) return x;
35     return fa[x][0];
36 }
37 inline int getans(int x, int f)
38 {
39     int ans = inf, t = dep[x] - dep[f];
40     for (int i = 0; i <= 16 && t; ++i)
41         if (t & (1 << i))
42         {
43             cmin(ans, mn[x][i]);
44             x = fa[x][i];
45             t ^= 1 << i;
46         }
47     return ans;
48 }
```

## 1.12   Prufer 编码

## 1.13   Link-Cut Tree (ct)

```
1 struct Node *null;
2 struct Node {
3     Node *ch[2], *fa, *pos;
4     int val, mn, l, len; bool rev;
5     // min_val in chain
6     inline bool type()
7     {
8         return fa -> ch[1] == this;
9     }
10     inline bool check()
11     {
12         return fa -> ch[type()] == this;
13     }
14     inline void pushup()
15     {
16         pos = this; mn = val;
17         ch[0] -> mn < mn ? mn = ch[0] -> mn, pos = ch[0] -> pos : 0;
18         ch[1] -> mn < mn ? mn = ch[1] -> mn, pos = ch[1] -> pos : 0;
19         len = ch[0] -> len + ch[1] -> len + l;
20     }
```

```
21    inline void pushdown()
22    {
23        if (rev)
24        {
25            ch[0] -> rev ^= 1;
26            ch[1] -> rev ^= 1;
27            std::swap(ch[0], ch[1]);
28            rev ^= 1;
29        }
30    }
31    inline void pushdownall()
32    {
33        if (check()) fa -> pushdownall();
34        pushdown();
35    }
36    inline void rotate()
37    {
38        bool d = type(); Node *f = fa, *gf = f -> fa;
39        (fa = gf, f -> check()) ? fa -> ch[f -> type()] = this : 0;
40        (f -> ch[d] = ch[!d]) != null ? ch[!d] -> fa = f : 0;
41        (ch[!d] = f) -> fa = this;
42        f -> pushup();
43    }
44    inline void splay(bool need = 1)
45    {
46        if (need) pushdownall();
47        for (; check(); rotate())
48            if (fa -> check())
49                (type() == fa -> type() ? fa : this) -> rotate();
50        pushup();
51    }
52    inline Node *access()
53    {
54        Node *i = this, *j = null;
55        for (; i != null; i = (j = i) -> fa)
56        {
57            i -> splay();
58            i -> ch[1] = j;
59            i -> pushup();
60        }
61        return j;
62    }
63    inline void make_root()
64    {
65        access();
66        splay();
67        rev ^= 1;
68    }
69    inline void link(Node *that)
70    {
71        make_root();
72        fa = that;
73        splay(0);
74    }
75    inline void cut(Node *that)
76    {
77        make_root();
78        that -> access();
79        that -> splay(0);
80        that -> ch[0] = fa = null;
81        that -> pushup();
```

```
82          }
83  } mem[maxn];
84  inline Node *query(Node *a, Node *b)
85  {
86      a -> make_root(); b -> access(); b -> splay(0);
87      return b -> pos;
88  }
89  inline int dist(Node *a, Node *b)
90  {
91      a -> make_root(); b -> access(); b -> splay(0);
92      return b -> len;
93  }
```

## 1.14  圆方树 (ct)

```
1   int dfn[maxn], low[maxn], timer, st[maxn], top, id[maxn], scc;
2   void dfs(int x)
3   {
4       dfn[x] = low[x] = ++timer; st[++top] = x;
5       for (Edge *iter = last[x]; iter; iter = iter -> next)
6           if (!dfn[iter -> to])
7           {
8               dfs(iter -> to);
9               cmin(low[x], low[iter -> to]);
10              if (dfn[x] == low[iter->to])
11              {
12                  int now, elder = top, minn = c[x];
13                  ++scc;
14                  do
15                  {
16                      now = st[top--];
17                      cmin(minn, c[now]);
18                  }
19                  while (iter -> to != now);
20                  for (int i = top + 1; i <= elder; ++i)
21                      add(scc, st[i], minn);
22                  add(scc, x, minn);
23              }
24          }
25          else if (!id[iter -> to]) cmin(low[x], dfn[iter -> to]);
26  }
```

## 1.15  最小割

## 1.16  最大流 (ct)

### Dinic (ct)

```
1   struct Edge {
2       Edge *next, *rev;
3       int to, cap;
4   } *last[maxn], *cur[maxn], e[maxm], *ecnt = e;
5   inline void link(R int a, R int b, R int w)
6   {
7       *++ecnt = (Edge) {last[a], ecnt + 1, b, w}; last[a] = ecnt;
8       *++ecnt = (Edge) {last[b], ecnt - 1, a, 0}; last[b] = ecnt;
9   }
```

```
10  int ans, s, t, q[maxn], dep[maxn];
11  inline bool bfs()
12  {
13      memset(dep, -1, (t + 1) << 2);
14      dep[q[1] = t] = 0; int head = 0, tail = 1;
15      while (head < tail)
16      {
17          int now = q[++head];
18          for (Edge *iter = last[now]; iter; iter = iter -> next)
19              if (dep[iter -> to] == -1 && iter -> rev -> cap)
20                  dep[q[++tail] = iter -> to] = dep[now] + 1;
21      }
22      return dep[s] != -1;
23  }
24  int dfs(int x, int f)
25  {
26      if (x == t) return f;
27      int used = 0;
28      for (Edge* &iter = cur[x]; iter; iter = iter -> next)
29          if (iter -> cap && dep[iter -> to] + 1 == dep[x])
30          {
31              int v = dfs(iter -> to, dmin(f - used, iter -> cap));
32              iter -> cap -= v;
33              iter -> rev -> cap += v;
34              used += v;
35              if (used == f) return f;
36          }
37      return used;
38  }
39  inline void dinic()
40  {
41      while (bfs())
42      {
43          memcpy(cur, last, sizeof cur);
44          ans += dfs(s, inf);
45      }
46  }
```

## SAP (lhy)

```
1   void SAP(int n, int st, int ed)
2   {
3       for(int i = 1; i <= n; i++)
4           now[i] = son[i];
5       sumd[0] = n;
6       int flow = inf, x = st;
7       while(dis[st] < n)
8       {
9           back[x] = flow;
10          int flag = 0;
11          for(int i = now[x]; i != -1; i = edge[i].next)
12          {
13              int y = edge[i].y;
14              if(edge[i].f && dis[y] + 1 == dis[x])
15              {
16                  flag = 1;
17                  now[x] = i;
18                  pre[y] = i;
19                  flow = min(flow, edge[i].f);
20                  x = y;
```

```
21              if(x == ed)
22              {
23                  ans += flow;
24                  while(x != st)
25                  {
26                      edge[pre[x]].f -= flow;
27                      edge[pre[x] ^ 1].f += flow;
28                      x = edge[pre[x]].x;
29                  }
30                  flow = inf;
31              }
32              break;
33          }
34      }
35      if(flag)continue;
36      int minn = n - 1, tmp;
37      for(int i = son[x]; i != -1; i = edge[i].next)
38      {
39          int y = edge[i].y;
40          if(edge[i].f && dis[y] < minn)
41          {
42              minn = dis[y];
43              tmp = i;
44          }
45      }
46      now[x] = tmp;
47      if(!(--sumd[dis[x]]))return;
48      sumd[dis[x] = minn + 1]++;
49      if(x != st)flow = back[x = edge[pre[x]].x];
50  }
51 }
```

## 1.17  费用流 (ct)

### Dinic(ct)

```
1  struct Edge {
2      Edge *next, *rev;
3      int from, to, cap, cost;
4  } *last[maxn], *prev[maxn], e[maxm], *ecnt = e;
5  inline void link(int a, int b, int w, int c)
6  {
7      *++ecnt = (Edge) {last[a], ecnt + 1, a, b, w, c}; last[a] = ecnt;
8      *++ecnt = (Edge) {last[b], ecnt - 1, b, a, 0, -c}; last[b] = ecnt;
9  }
10 int s, t, q[maxn << 2], dis[maxn];
11 ll ans;
12 bool inq[maxn];
13 #define inf 0x7fffffff
14 inline bool spfa()
15 {
16     for (int i = 1; i <= t; ++i) dis[i] = inf;
17     int head = 0, tail = 1; dis[q[1] = s] = 0;
18     while (head < tail)
19     {
20         int now = q[++head]; inq[now] = 0;
21         for (Edge *iter = last[now]; iter; iter = iter -> next)
22             if (iter -> cap && dis[iter -> to] > dis[now] + iter -> cost)
23             {
```

```
24              dis[iter -> to] = dis[now] + iter -> cost;
25              prev[iter -> to] = iter;
26              !inq[iter -> to] ? inq[q[++tail] = iter -> to] = 1 : 0;
27          }
28      }
29      return dis[t] != inf;
30 }
31 inline void mcmf()
32 {
33      int x = inf;
34      for (Edge *iter = prev[t]; iter; iter = prev[iter -> from]) cmin(x, iter -> cap);
35      for (Edge *iter = prev[t]; iter; iter = prev[iter -> from])
36      {
37          iter -> cap -= x;
38          iter -> rev -> cap += x;
39          ans += 1ll * x * iter -> cost;
40      }
41 }
```

## zkw(lhy)

```
1  int aug(int no, int res)
2  {
3      if(no == ED)return mincost += 1ll * pil * res, res;
4      v[no] = 1;
5      int flow = 0;
6      for(int i = son[no]; i != -1; i = edge[i].next)
7          if(edge[i].f && !v[edge[i].y] && !edge[i].c)
8          {
9              int d = aug(edge[i].y, min(res, edge[i].f));
10             edge[i].f -= d, edge[i ^ 1].f += d, flow += d, res -= d;
11             if(!res)return flow;
12         }
13     return flow;
14 }

15 bool modlabel()
16 {
17     long long d = 0x3f3f3f3f3f3f3f3fll;
18     for(int i = 1; i <= cnt; i++)
19         if(v[i])
20         {
21             for(int j = son[i]; j != -1; j = edge[j].next)
22                 if(edge[j].f && !v[edge[j].y] && edge[j].c < d)d = edge[j].c;
23         }
24     if(d == 0x3f3f3f3f3f3f3f3fll)return 0;
25     for(int i = 1; i <= cnt; i++)
26         if(v[i])
27         {
28             for(int j = son[i]; j != -1; j = edge[j].next)
29                 edge[j].c -= d, edge[j ^ 1].c += d;
30         }
31     pil += d;
32     return 1;
33 }

34 void minimum_cost_flow_zkw()
35 {
36     pil = 0;
37     int nowans = 0;
```

```
38    nowf = 0;
39    do{
40        do{
41            for(int i = 1; i <= cnt; i++)
42                v[i] = 0;
43            nowans = aug(ST, inf);
44            nowf += nowans;
45        }while(nowans);
46    }while(modlabel());
47 }
```

## 1.18   有上下界的网络流 (Durandal)

$B(u,v)$ 表示边 $(u,v)$ 流量的下界，$C(u,v)$ 表示边 $(u,v)$ 流量的上界，设 $F(u,v)$ 表示边 $(u,v)$ 的实际流量
设 $G(u,v) = F(u,v) - B(u,v)$，则 $0 \le G(u,v) \le C(u,v) - B(u,v)$

- 无源汇的上下界可行流
  建立超级源点 $S^*$ 和超级汇点 $T^*$，对于原图每一条边 $(u,v)$ 在新网络中连如下三条边：$S^* \to v$，容量为 $B(u,v)$；$u \to T^*$，容量为 $B(u,v)$；$u \to v$，容量为 $C(u,v) - B(u,v)$。最后求新网络的最大流，判断从超级源点 $S^*$ 出发的边是否都满流即可，边 $(u,v)$ 的最终解中的实际流量为 $G(u,v) + B(u,v)$。

- 有源汇的上下界可行流
  从汇点 $T$ 到源点 $S$ 连一条上界为 $\infty$，下界为 $0$ 的边。按照无源汇的上下界可行流一样做即可，流量即为 $T \to S$ 边上的流量。

- 有源汇的上下界最大流
  - 在有源汇的上下界可行流中，从汇点 $T$ 到源点 $S$ 的边改为连一条上界为 $\infty$，下界为 $x$ 的边。$x$ 满足二分性质，找到最大的 $x$ 使得新网络存在有源汇的上下界可行流即为原图的最大流。
  - 从汇点 $T$ 到源点 $S$ 连一条上界为 $\infty$，下界为 $0$ 的边，变成无源汇的网络。按照无源汇的上下界可行流的方法，建立超级源点 $S^*$ 与超级汇点 $T^*$，求一遍 $S^* \to T^*$ 的最大流，再将从汇点 $T$ 到源点 $S$ 的这条边拆掉，求一次 $S \to T$ 的最大流即可。

- 有源汇的上下界最小流
  - 在有源汇的上下界可行流中，从汇点 $T$ 到源点 $S$ 的边改为连一条上界为 $x$，下界为 $0$ 的边。$x$ 满足二分性质，找到最小的 $x$ 使得新网络存在有源汇的上下界可行流即为原图的最大流。
  - 按照无源汇的上下界可行流的方法，建立超级源点 $S^*$ 与超级汇点 $T^*$，求一遍 $S^* \to T^*$ 的最大流，但是注意不加上汇点 $T$ 到源点 $S$ 的这条边，即不使之改为无源汇的网络去求解。求完后，再加上那条汇点 $T$ 到源点 $S$ 的边，上界为 $\infty$ 的边。因为这条边的下界为 $0$，所以 $S^*, T^*$ 无影响，再求一次 $S^* \to T^*$ 的最大流。若超级源点 $S^*$ 出发的边全部满流，则 $T \to S$ 边上的流量即为原图的最小流，否则无解。

## 1.19   差分约束

## 1.20   图论知识 (gy,lhy)

### 弦图

弦图：任意点数 $\ge 4$ 的环皆有弦的无向图
单纯点：与其相邻的点的诱导子图为完全图的点
完美消除序列：每次选择一个单纯点删去的序列
弦图必有完美消除序列
$O(m + n)$ 求弦图的完美消除序列：每次选择未选择的标号最大的点，并将与其相连的点标号 $+1$，得到完美消除序列的反序
最大团数 = 最小染色数：按完美消除序列从后往前贪心地染色
最小团覆盖 = 最大点独立集：按完美消除序列从前往后贪心地选点加入点独立集

## 计数问题

- **有根树计数**
  $a_1 = 1$

  $$a_{n+1} = \frac{\sum\limits_{j=1}^{n} j \cdot a_j \cdot S_{n,j}}{n}$$

  $$S_{n,j} = \sum_{i=1}^{n/j} a_{n+1-ij} = S_{n-j,j} + a_{n+1-j}$$

- **无根树计数**

  $$\begin{cases} a_n - \sum\limits_{i=1}^{n/2} a_i a_{n-i} & n \text{ is odd} \\ a_n - \sum\limits_{i=1}^{n/2} a_i a_{n-i} + \frac{1}{2} a_{\frac{n}{2}} \left( a_{\frac{n}{2}} + 1 \right) & n \text{ is even} \end{cases}$$

- **完全图生成树计数**
  $n^{n-2}$

- **矩阵-树定理**
  设 $\mathbf{A}[G]$ 为图 $G$ 的邻接矩阵、$\mathbf{D}[G]$ 为图 $G$ 的度数矩阵，则图 $G$ 的不同生成树的个数为 $\mathbf{C}[G] = \mathbf{D}[G] - \mathbf{A}[G]$ 的任意一个 $n-1$ 阶主子式的行列式值。

- **偶数点完全图完备匹配计数**
  $(n-1)!!$

- **无根二叉树计数**
  $(2n-5)!!$

- **有根二叉树计数**
  $(2n-3)!!$

## 最大权闭合子图

给定一个带点权的有向图，求其最大权闭合子图。
从源点 $S$ 向每一条正权点连一条容量为权值的边，每个负权点向汇点 $T$ 连一条容量为权值绝对值的边，有向图原来的边容量为 $\infty$。求它的最小割，与源点 $S$ 连通的点构成最大权闭合子图，权值为正权值和 $-$ 最小割。

## 最大密度子图

给定一个无向图，求其一个子图，使得子图的边数 $|E|$ 和点数 $|V|$ 满足 $\frac{|E|}{|V|}$ 最大。
二分答案 $k$，使得 $|E| - k|V| \geq 0$ 有解，将原图边和点都看作点，边 $(u,v)$ 分别向 $u$ 和 $v$ 连边求最大权闭合子图。

# Chapter 2

# Math

## 2.1 int64 相乘取模 (Durandal)

```
int64_t mul(int64_t x, int64_t y, int64_t p) {
    int64_t t = (x * y - (int64_t) ((long double) x / p * y + 1e-3) * p) % p;
    return t < 0 ? t + p : t;
}
```

## 2.2 ex-Euclid (gy)

```
// return gcd(a, b)
// ax+by=gcd(a,b)
int extend_gcd(int a, int b, int &x, int &y) {
    if (b == 0) {
        x = 1, y = 0;
        return a;
    }
    int res = extend_gcd(b, a % b, x, y);
    int t = y;
    y = x - a / b * y;
    x = t;
    return res;
}

// return minimal positive integer x so that ax+by=c
// or -1 if such x does not exist
int solve_equ(int a, int b, int c) {
    int x, y, d;
    d = extend_gcd(a, b, x, y);
    if (c % d)
        return -1;
    int t = c / d;
    x *= t;
    y *= t;
    int k = b / d;
    x = (x % k + k) % k;
    return x;
}

// return minimal positive integer x so that ax==b(mod p)
// or -1 if such x does not exist
int solve(int a, int b, int p) {
    a = (a % p + p) % p;
    b = (b % p + p) % p;
```

```
33      return solve_equ(a, p, b);
34  }
```

## 2.3   中国剩余定理 (Durandal)

返回是否可行，余数和模数结果为 $r_1, m_1$

```
1   bool CRT(int &r1, int &m1, int r2, int m2) {
2       int x, y, g = extend_gcd(m1, m2, x, y);
3       if ((r2 - r1) % g != 0) return false;
4       x = 1ll * (r2 - r1) * x % m2;
5       if (x < 0) x += m2;
6       x /= g;
7       r1 += m1 * x;
8       m1 *= m2 / g;
9       return true;
10  }
```

## 2.4   线性同余不等式 (Durandal)

必须满足 $0 \le d < m, 0 \le l \le r < m$，返回 $\min\{x \ge 0 \mid l \le x \cdot d \bmod m \le r\}$，无解返回 $-1$

```
1   int64_t calc(int64_t d, int64_t m, int64_t l, int64_t r) {
2       if (l == 0) return 0;
3       if (d == 0) return -1;
4       if (d * 2 > m) return calc(m - d, m, m - r, m - l);
5       if ((l - 1) / d < r / d) return (l - 1) / d + 1;
6       int64_t k = calc((-m % d + d) % d, d, l % d, r % d);
7       if (k == -1) return -1;
8       return (k * m + l - 1) / d + 1;
9   }
```

## 2.5   组合数

## 2.6   高斯消元 (ct)

增广矩阵大小为 $m \times (n + 1)$

```
1   db a[maxn][maxn], x[maxn];
2   int main()
3   {
4       int rank = 0;
5       for (int i = 1, now = 1; i <= m && now <= n; ++now)
6       {
7           if (fabs(a[i][now]) < eps)
8           {
9               for (int j = i + 1; j <= m; ++j)
10                  if (fabs(a[j][now]) > fabs(a[i][now]))
11                  {
12                      for (int k = now; k <= n + 1; ++k)
13                          std::swap(a[i][k], a[j][k]);
14                  }
15          }
16          if (fabs(a[i][now]) < eps) continue;
17
18          for (int j = i + 1; j <= m; ++j)
```

```
18          {
19              db temp = a[j][now] / a[i][now];
20              for (int k = now; k <= n + 1; ++k)
21                  a[j][k] -= temp * a[i][k];
22          }
23          ++i; ++rank;
24      }
25
26      if (rank == n)
27      {
28          x[n] = a[n][n + 1] / a[n][n];
29          for (int i = n - 1; i; --i)
30          {
31              for (int j = i + 1; j <= n; ++j)
32                  a[i][n + 1] -= x[j] * a[i][j];
33              x[i] = a[i][n + 1] / a[i][i];
34          }
35      }
36      else puts("Infinite Solution!");
37      return 0;
38  }
```

## 2.7   Miller Rabin & Pollard Rho (gy)

In Java, use BigInteger.isProbablePrime(int certainty) to replace miller_rabin(BigInteger number)

| Test Set | First Wrong Answer |
|---|---|
| $2$ | $2047$ |
| $2, 3$ | $1,373,653$ |
| $31, 73$ | $9,080,191$ |
| $2, 3, 5$ | $25,326,001$ |
| $2, 3, 5, 7$ | (INT32_MAX)$3,215,031,751$ |
| $2, 7, 61$ | $4,759,123,141$ |
| $2, 13, 23, 1662803$ | $1,122,004,669,633$ |
| $2, 3, 5, 7, 11$ | $2,152,302,898,747$ |
| $2, 3, 5, 7, 11, 13$ | $3,474,749,660,383$ |
| $2, 3, 5, 7, 11, 13, 17$ | $341,550,071,728,321$ |
| $2, 3, 5, 7, 11, 13, 17, 19, 23$ | $3,825,123,056,546,413,051$ |
| $2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37$ | (INT64_MAX)$318,665,857,834,031,151,167,461$ |
| $2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41$ | $3,317,044,064,679,887,385,961,981$ |

```
1  const int test_case_size = 12;
2  const int test_cases[test_case_size] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
3
4  int64_t multiply_mod(int64_t x, int64_t y, int64_t p) {
5      int64_t t = (x * y - (int64_t) ((long double) x / p * y + 1e-3) * p) % p;
6      return t < 0 ? t + p : t;
7  }
8
9  int64_t add_mod(int64_t x, int64_t y, int64_t p) {
10     return (0ull + x + y) % p;
11 }
12
13 int64_t power_mod(int64_t x, int64_t exp, int64_t p) {
14     int64_t ans = 1;
15     while (exp) {
16         if (exp & 1)
17             ans = multiply_mod(ans, x, p);
18         x = multiply_mod(x, x, p);
```

```
16          exp >>= 1;
17      }
18      return ans;
19  }

20  bool miller_rabin_check(int64_t prime, int64_t base) {
21      int64_t number = prime - 1;
22      for (; ~number & 1; number >>= 1)
23          continue;
24      int64_t result = power_mod(base, number, prime);
25      for (; number != prime - 1 && result != 1 && result != prime - 1; number <<= 1)
26          result = multiply_mod(result, result, prime);
27      return result == prime - 1 || (number & 1) == 1;
28  }

29  bool miller_rabin(int64_t number) {
30      if (number < 2)
31          return false;
32      if (number < 4)
33          return true;
34      if (~number & 1)
35          return false;
36      for (int i = 0; i < test_case_size && test_cases[i] < number; i++)
37          if (!miller_rabin_check(number, test_cases[i]))
38              return false;
39      return true;
40  }

41  int64_t gcd(int64_t x, int64_t y) {
42      return y == 0 ? x : gcd(y, x % y);
43  }

44  int64_t pollard_rho_test(int64_t number, int64_t seed) {
45      int64_t x = rand() % (number - 1) + 1, y = x;
46      int head = 1, tail = 2;
47      while (true) {
48          x = multiply_mod(x, x, number);
49          x = add_mod(x, seed, number);
50          if (x == y)
51              return number;
52          int64_t answer = gcd(std::abs(x - y), number);
53          if (answer > 1 && answer < number)
54              return answer;
55          if (++head == tail) {
56              y = x;
57              tail <<= 1;
58          }
59      }
60  }

61  void factorize(int64_t number, std::vector<int64_t> &divisor) {
62      if (number > 1) {
63          if (miller_rabin(number)) {
64              divisor.push_back(number);
65          } else {
66              int64_t factor = number;
67              while (factor >= number)
68                  factor = pollard_rho_test(number, rand() % (number - 1) + 1);
69              factorize(number / factor, divisor);
70              factorize(factor, divisor);
71          }
```

```
72        }
73 }
```

## 2.8    $O(m^2 \log n)$ **线性递推** (lhy)

```
1  typedef vector<int> poly;
2  //{1, 3} {2, 1} an = 2an-1 + an-2, calc(3) = 7
3  struct LinearRec{
4      int n, LOG;
5      poly first, trans;
6      vector<poly> bin;
7      poly add(poly &a, poly &b)
8      {
9          poly res(n * 2 + 1, 0);
10         for(int i = 0; i <= n; i++)
11             for(int j = 0; j <= n; j++)
12                 (res[i + j] += 1ll * a[i] * b[j] % mo) %= mo;
13         for(int i = 2 * n; i > n; i--)
14         {
15             for(int j = 0; j < n; j++)
16                 (res[i - 1 - j] += 1ll * res[i] * trans[j] % mo) %= mo;
17             res[i] = 0;
18         }
19         res.erase(res.begin() + n + 1, res.end());
20         return res;
21     }
22     LinearRec(poly &first, poly &trans, int LOG): LOG(LOG), first(first), trans(trans)
23     {
24         n = first.size();
25         poly a(n + 1, 0);
26         a[1] = 1;
27         bin.push_back(a);
28         for(int i = 1; i < LOG; i++)
29             bin.push_back(add(bin[i - 1], bin[i - 1]));
30     }
31     int calc(long long k)
32     {
33         poly a(n + 1, 0);
34         a[0] = 1;
35         for(int i = 0; i < LOG; i++)
36             if((k >> i) & 1)a = add(a, bin[i]);
37         int ret = 0;
38         for(int i = 0; i < n; i++)
39             if((ret += 1ll * a[i + 1] * first[i] % mo) >= mo)ret -= mo;
40         return ret;
41     }
42 };
```

## 2.9    **线性基** (ct)

```
1  int main()
2  {
3      for (int i = 1; i <= n; ++i)
4      {
5          ull x = F();
6          cmax(m, 63 - __builtin_clzll(x));
7          for ( ; x; )
8          {
```

```
 9              tmp = __builtin_ctzll(x);
10              if (!b[tmp])
11              {
12                  b[tmp] = x;
13                  break;
14              }
15              x ^= b[tmp];
16          }
17      }
18  }
```

## 2.10   FFT NTT FWT (lhy,ct,gy)

### FFT (ct)

0-based

```
 1  #include <cstdio>
 2  #include <cmath>
 3  #include <algorithm>

 4  #define R register
 5  #define maxn 262144
 6  typedef double db;
 7  const db pi = acos(-1);

 8  char S[1 << 20], *T = S;
 9  inline int F()
10  {
11      R char ch; R int cnt = 0;
12      while (ch = *T++, ch < '0' || ch > '9') ;
13      cnt = ch - '0';
14      while (ch = *T++, ch >= '0' && ch <= '9') cnt = cnt * 10 + ch - '0';
15      return cnt;
16  }
17  struct Complex {
18      db x, y;
19      inline Complex operator * (const Complex &that) const {return (Complex) {x * that.x - y *
           that.y, x * that.y + y * that.x};}
20      //inline Complex operator + (const Complex &that) const {return (Complex) {x + that.x, y +
           that.y};}
21      inline Complex operator += (const Complex &that){x+=that.x;y+=that.y;}
22      inline Complex operator - (const Complex &that) const {return (Complex) {x - that.x, y -
           that.y};}
23  } buf_a[maxn], buf_b[maxn], buf_c[maxn], w[maxn], c[maxn], a[maxn], b[maxn];

24  int n;
25  void bit_reverse(R Complex *x, R Complex *y)
26  {
27      for (R int i = 0; i < n; ++i) y[i] = x[i];
28      Complex tmp;
29      for (R int i = 0, j = 0; i < n; ++i)
30      {
31          (i>j)?tmp=y[i],y[i]=y[j],y[j]=tmp,0:1;
32          for (R int l = n >> 1; (j ^= l) < l; l >>= 1);
33      }
34  }
35  void init()
36  {
37      R int h=n>>1;
```

```
38          for (R int i = 0; i < h; ++i) w[i+h] = (Complex) {cos(2 * pi * i / n), sin(2 * pi * i / n)};
39          for (R int i = h; i--; )w[i]=w[i<<1];
40  }
41  void dft(R Complex *a)
42  {
43          R Complex tmp;
44          for(R int p = 2, m = 1; m != n; p = (m = p) << 1)
45              for(R int i = 0; i != n; i += p) for(R int j = 0; j != m; ++j)
46              {
47                  tmp = a[i + j + m] * w[j + m];
48                  a[i + j + m] = a[i + j] - tmp;
49                  a[i + j] += tmp;
50              }
51  }

52  int main()
53  {
54          fread(S, 1, 1 << 20, stdin);
55          R int na = F(), nb = F(), x;
56          for (R int i = 0; i <= na; ++i) a[i].x=F();
57          for (R int i = 0; i <= nb; ++i) b[i].x=F();
58          for (n = 1; n < na + nb + 1; n <<= 1) ;
59          bit_reverse(a, buf_a);
60          bit_reverse(b, buf_b);
61          init();
62          dft(buf_a);
63          dft(buf_b);
64          for (R int i = 0; i < n; ++i) c[i] = buf_a[i] * buf_b[i];
65          std::reverse(c + 1, c + n);
66          bit_reverse(c, buf_c);
67          dft(buf_c);
68          for (R int i = 0; i <= na + nb; ++i) printf("%d%c", int(buf_c[i].x / n + 0.5), "
             ↪ \n"[i==na+nb]);
69          return 0;
70  }
```

## NTT (gy)

0-based

```
1   const int N = 1e6 + 10;
2   const int64_t MOD = 998244353, G = 3;
3   int rev[N];

4   int64_t powMod(int64_t a, int64_t exp) {
5       int64_t ans = 1;
6       while (exp) {
7           if (exp & 1)
8               (ans *= a) %= MOD;
9           (a *= a) %= MOD;
10          exp >>= 1;
11      }
12      return ans;
13  }

14  void number_theoretic_transform(int64_t *p, int n, int idft) {
15      for (int i = 0; i < n; i++)
16          if (i < rev[i])
17              std::swap(p[i], p[rev[i]]);
18      for (int j = 1; j < n; j <<= 1) {
19          static int64_t wn1, w, t0, t1;
```

```
20          wn1 = powMod(G, (MOD - 1) / (j << 1));
21          if (idft == -1)
22              wn1 = powMod(wn1, MOD - 2);
23          for (int i = 0; i < n; i += j << 1) {
24              w = 1;
25              for (int k = 0; k < j; k++) {
26                  t0 = p[i + k];
27                  t1 = w * p[i + j + k] % MOD;
28                  p[i + k] = (t0 + t1) % MOD;
29                  p[i + j + k] = (t0 - t1 + MOD) % MOD;
30                  (w *= wn1) %= MOD;
31              }
32          }
33      }
34      if (idft == -1) {
35          int nInv = powMod(n, MOD - 2);
36          for (int i = 0; i < n; i++)
37              (p[i] *= nInv) %= MOD;
38      }
39  }

40  int64_t *ntt_main(int64_t *a, int64_t *b, int n, int m) {
41      static int64_t aa[N], bb[N];
42      static int nn, len;
43      len = 0;
44      for (nn = 1; nn < m + n; nn <<= 1)
45          len++;
46      for (int i = 0; i < nn; i++) {
47          aa[i] = a[i];
48          bb[i] = b[i];
49      }
50      rev[0] = 0;
51      for (int i = 1; i < nn; i++)
52          rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (len - 1));
53      number_theoretic_transform(aa, nn, 1);
54      number_theoretic_transform(bb, nn, 1);
55      for (int i = 0; i < nn; i++)
56          (aa[i] *= bb[i]) %= MOD;
57      number_theoretic_transform(aa, nn, -1);
58      return aa;
59  }
```

## FWT (lhy)

0-based

```
1   void fwt(int n, int *x, bool inv = false)
2   {
3       for(int i = 0; i < n; i++)
4           for(int j = 0; j < (1 << n); j++)
5               if((j >> i) & 1)
6               {
7                   int p = x[j ^ (1 << i)], q = x[j];
8                   if(!inv)
9                   {
10                      //xor
11                      x[j ^ (1 << i)] = p - q;
12                      x[j] = p + q;
13                      //or
14                      x[j ^ (1 << i)] = p;
15                      x[j] = p + q;
```

```
16                      //and
17                      x[j ^ (1 << i)] = p + q;
18                      x[j] = q;
19                  }
20                  else
21                  {
22                      //xor
23                      x[j ^ (1 << i)] = (p + q) >> 1;
24                      x[j] = (q - p) >> 1;
25                      //or
26                      x[j ^ (1 << i)] = p;
27                      x[j] = q - p;
28                      //and
29                      x[j ^ (1 << i)] = p - q;
30                      x[j] = q;
31                  }
32              }
33  }
34  void solve(int n, int *a, int *b, int *c)
35  {
36      fwt(n, a);
37      fwt(n, b);
38      for(int i = 0; i < (1 << n); i++)
39          c[i] = a[i] * b[i];
40      fwt(n, c, 1);
41  }
```

## 2.11  Lagrange 插值 (ct)

求解 $\sum\limits_{i=1}^{n} i^k \bmod (10^9 + 7)$

```
1   const int mod = 1e9 + 7;
2   int f[maxn], pre[maxn], suf[maxn], inp[maxn], p[maxn];
3   inline int qpow(int base, int power)
4   {
5       int ret = 1;
6       for (; power; power >>= 1, base = 1ll * base * base % mod)
7           power & 1 ? ret = 1ll * ret * base % mod : 0;
8       return ret;
9   }
10  bool vis[maxn];
11  int pr[maxn], prcnt, fpow[maxn];
12  int main()
13  {
14      int n = F(), k = F();
15      // ************
16      fpow[1] = 1;
17      for (int i = 2; i <= k + 2; ++i)
18      {
19          if (!vis[i]) pr[++prcnt] = i, fpow[i] = qpow(i, k);
20          for (int j = 1; j <= prcnt && i * pr[j] <= k + 2; ++j)
21          {
22              vis[i * pr[j]] = 1;
23              fpow[i * pr[j]] = 1ll * fpow[i] * fpow[pr[j]] % mod;
24              if (i % pr[j] == 0) break;
25          }
26      }
27      // ************* pre-processing
```

```
28      for (int i = 1; i <= k + 2; ++i) f[i] = (f[i - 1] + fpow[i]) % mod;
29      if (n <= k + 2) return !printf("%d\n", f[n] );
30      pre[0] = 1;
31      for (int i = 1; i <= k + 3; ++i) pre[i] = 1ll * pre[i - 1] * (n - i) % mod;
32      suf[k + 3] = 1;
33      for (int i = k + 2; i >= 0; --i) suf[i] = 1ll * suf[i + 1] * (n - i) % mod;

34      p[0] = 1;
35      for (int i = 1; i <= k + 2; ++i) p[i] = (1ll * p[i - 1] * i) % mod;

36      inp[k + 2] = qpow(p[k + 2], mod - 2);

37      for (int i = k + 1; i >= 0; --i) inp[i] = (1ll * inp[i + 1] * (i + 1)) % mod;

38      int ans = 0;
39      for (int i = 1; i <= k + 2; ++i)
40      {
41          int temp = inp[k + 2 - i]; if ((k + 2 - i) & 1) temp = mod - temp;
42          int tmp = 1ll * pre[i - 1] * suf[i + 1] % mod * temp % mod * inp[i - 1] % mod * f[i] % mod;
43          ans = (ans + tmp) % mod;
44      }
45      printf("%d\n", ans );
46      return 0;
47  }
```

## 2.12  杜教筛 (ct)

```
1   int phi[maxn], pr[maxn / 10], prcnt;
2   ll sph[maxn];
3   bool vis[maxn];
4   const int moha = 3333331;
5   struct Hash {
6       Hash *next;
7       int ps; ll ans;
8   } *last1[moha], mem[moha], *tot = mem;
9   inline ll S1(int n)
10  {
11      if (n < maxn) return sph[n];
12      for (R Hash *iter = last1[n % moha]; iter; iter = iter -> next)
13          if (iter -> ps == n) return iter -> ans;
14      ll ret = 1ll * n * (n + 1ll) / 2;
15      for (ll i = 2, j; i <= n; i = j + 1)
16      {
17          j = n / (n / i);
18          ret -= S1(n / i) * (j - i + 1);
19      }
20      *++tot = (Hash) {last1[n % moha], n, ret}; last1[n % moha] = tot;
21      return ret;
22  }
23  int main()
24  {
25      int T; scanf("%d", &T);
26      phi[1] = sph[1] = 1;
27      for (int i = 2; i < maxn; ++i)
28      {
29          if (!vis[i]) pr[++prcnt] = i, phi[i] = i - 1;
30          sph[i] = sph[i - 1] + phi[i];
31          for (int j = 1; j <= prcnt && 1ll * i * pr[j] < maxn; ++j)
```

```
32          {
33              vis[i * pr[j]] = 1;
34              if (i % pr[j])
35                  phi[i * pr[j]] = phi[i] * (pr[j] - 1);
36              else
37              {
38                  phi[i * pr[j]] = phi[i] * pr[j];
39                  break;
40              }
41          }
42      }
43      for (; T; --T)
44      {
45          int N; scanf("%d", &N);
46          printf("%lld\n", S1(N));
47      }
48      return 0;
49  }
```

## 2.13  BSGS (ct,Durandal)

### 2.13.1  BSGS (ct)

$p$ 是素数, 返回 $\min\{x \geq 0 \mid y^x \equiv z \pmod{p}\}$

```
1   const int mod = 19260817;
2   struct Hash
3   {
4       Hash *next;
5       int key, val;
6   } *last[mod], mem[100000], *tot = mem;
7   inline void insert(R int x, R int v)
8   {
9       *++tot = (Hash) {last[x % mod], x, v}; last[x % mod] = tot;
10  }
11  inline int query(R int x)
12  {
13      for (R Hash *iter = last[x % mod]; iter; iter = iter -> next)
14          if (iter -> key == x) return iter -> val;
15      return -1;
16  }
17  inline void del(R int x)
18  {
19      last[x % mod] = 0;
20  }
21  int main()
22  {
23      for (; T; --T)
24      {
25          R int y, z, p; scanf("%d%d%d", &y, &z, &p);
26          R int m = (int) sqrt(p * 1.0);
27          y %= p; z %= p;
28          if (!y && !z) {puts("0"); continue;}
29          if (!y) {puts("Orz, I cannot find x!"); continue;}
30          R int pw = 1;
31          for (R int i = 0; i < m; ++i, pw = 1ll * pw * y % p) insert(1ll * z * pw % p, i);
32          R int ans = -1;
33          for (R int i = 1, t, pw2 = pw; i <= p / m + 1; ++i, pw2 = 1ll * pw2 * pw % p)
34              if ((t = query(pw2)) != -1) {ans = i * m - t; break;}
35          if (ans == -1) puts("Orz, I cannot find x!");
```

```
36      else printf("%d\n", ans );
37      tot = mem; pw = 1;
38      for (R int i = 0; i < m; ++i, pw = 1ll * pw * y % p) del(1ll * z * pw % p);
39    }
40    return 0;
41  }
```

### 2.13.2  ex-BSGS (Durandal)

必须满足 $0 \le a < p, 0 \le b < p$, 返回 $\min\{x \ge 0 \mid a^x \equiv b \pmod{p}\}$

```
1   int64_t ex_bsgs(int64_t a, int64_t b, int64_t p) {
2       if (b == 1)
3           return 0;
4       int64_t t, d = 1, k = 0;
5       while ((t = std::__gcd(a, p)) != 1) {
6           if (b % t) return -1;
7           k++, b /= t, p /= t, d = d * (a / t) % p;
8           if (b == d) return k;
9       }
10      map.clear();
11      int64_t m = std::ceil(std::sqrt((long double) p));
12      int64_t a_m = pow_mod(a, m, p);
13      int64_t mul = b;
14      for (int j = 1; j <= m; j++) {
15          (mul *= a) %= p;
16          map[mul] = j;
17      }
18      for (int i = 1; i <= m; i++) {
19          (d *= a_m) %= p;
20          if (map.count(d))
21              return i * m - map[d] + k;
22      }
23      return -1;
24  }

25  int main() {
26      int64_t a, b, p;
27      while (scanf("%lld%lld%lld", &a, &b, &p) != EOF)
28          printf("%lld\n", ex_bsgs(a, b, p));
29      return 0;
30  }
```

## 2.14  直线下整点个数 (gy)

必须满足 $a \ge 0, b \ge 0, m > 0$, 返回 $\sum\limits_{i=0}^{n-1} \frac{a+bi}{m}$

```
1   int64_t count(int64_t n, int64_t a, int64_t b, int64_t m) {
2       if (b == 0)
3           return n * (a / m);
4       if (a >= m)
5           return n * (a / m) + count(n, a % m, b, m);
6       if (b >= m)
7           return (n - 1) * n / 2 * (b / m) + count(n, a, b % m, m);
8       return count((a + b * n) / m, (a + b * n) % m, m, b);
9   }
```

## 2.15  单纯形

## 2.16  辛普森积分

## 2.17  数学知识 (gy)

### 求和公式

- $\sum\limits_{k=1}^{n}(2k-1)^2 = \frac{1}{3}n(4n^2-1)$

- $\sum\limits_{k=1}^{n}k^3 = \frac{1}{4}n^2(n+1)^2$

- $\sum\limits_{k=1}^{n}(2k-1)^3 = n^2(2n^2-1)$

- $\sum\limits_{k=1}^{n}k^4 = \frac{1}{30}n(n+1)(2n+1)(3n^2+3m-1)$

- $\sum\limits_{k=1}^{n}k^5 = \frac{1}{12}n^2(n+1)^2(2n^2+2n-1)$

- $\sum\limits_{k=1}^{n}k(k+1) = \frac{1}{3}n(n+1)(n+2)$

- $\sum\limits_{k=1}^{n}k(k+1)(k+2) = \frac{1}{4}n(n+1)(n+2)(n+3)$

- $\sum\limits_{k=1}^{n}k(k+1)(k+2)(k+3) = \frac{1}{5}n(n+1)(n+2)(n+3)(n+4)$

### 错排公式

$D_n$ 表示 $n$ 个元素错位排列的方案数
$D_1 = 0, D_2 = 1$
$D_n = (n-1)(D_{n-2} + D_{n-1}), n \geq 3$
$D_n = n! \cdot (1 - \frac{1}{1!} + \frac{1}{2!} - \cdots + (-1)^n \frac{1}{n!})$

### Fibonacci sequence

$F_0 = 0, F_1 = 1$
$F_n = F_{n-1} + F_{n-2}$
$F_{n+1} \cdot F_{n-1} - F_n^2 = (-1)^n$
$F_{-n} = (-1)^n F_n$
$F_{n+k} = F_k \cdot F_{n+1} + F_{k-1} \cdot F_n$
$\gcd(F_m, F_n) = F_{\gcd(m,n)}$
$F_m \mid F_n^2 \Leftrightarrow nF_n \mid m$
$F_n = \frac{\varphi^n - \Psi^n}{\sqrt{5}}, \varphi = \frac{1+\sqrt{5}}{2}, \Psi = \frac{1-\sqrt{5}}{2}$
$F_n = \lfloor \frac{\varphi^n}{\sqrt{5}} + \frac{1}{2} \rfloor, n \geq 0$
$n(F) = \lfloor \log_\varphi(F \cdot \sqrt{5} + \frac{1}{2}) \rfloor$

### Stirling number (1st kind)

用 $\begin{bmatrix} n \\ k \end{bmatrix}$ 表示 Stirling number (1st kind)，为将 $n$ 个元素分成 $k$ 个环的方案数
$\begin{bmatrix} n+1 \\ k \end{bmatrix} = n\begin{bmatrix} n \\ k \end{bmatrix} + \begin{bmatrix} n \\ k-1 \end{bmatrix}, k > 0$
$\begin{bmatrix} 0 \\ 0 \end{bmatrix} = 1, \begin{bmatrix} n \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ n \end{bmatrix} = 0, n > 0$

$\left[\begin{smallmatrix}n\\k\end{smallmatrix}\right]$ 为将 $n$ 个元素分成 $k$ 个环的方案数

$$\left[\begin{matrix}x\\x-n\end{matrix}\right] = \sum_{k=0}^{n} \left\langle\!\!\left\langle\begin{matrix}n\\k\end{matrix}\right\rangle\!\!\right\rangle \binom{x+k}{2n}$$

## Stirling number (2nd kind)

用 $\left\{\begin{smallmatrix}n\\k\end{smallmatrix}\right\}$ 表示 Stirling number (2nd kind)，为将 $n$ 个元素划分成 $k$ 个非空集合的方案数

$\left\{\begin{smallmatrix}n+1\\k\end{smallmatrix}\right\} = k\left\{\begin{smallmatrix}n\\k\end{smallmatrix}\right\} + \left\{\begin{smallmatrix}n\\k-1\end{smallmatrix}\right\}, k > 0$

$\left\{\begin{smallmatrix}0\\0\end{smallmatrix}\right\} = 1, \left\{\begin{smallmatrix}n\\0\end{smallmatrix}\right\} = \left\{\begin{smallmatrix}0\\n\end{smallmatrix}\right\} = 0, n > 0$

$\left\{\begin{smallmatrix}n\\k\end{smallmatrix}\right\} = \frac{1}{k!}\sum_{j=0}^{k}(-1)^{k-j}\binom{k}{j}j^n$

$\left\{\begin{smallmatrix}n\\k\end{smallmatrix}\right\}$

$$\left\{\begin{matrix}x\\x-n\end{matrix}\right\} = \sum_{k=0}^{n} \left\langle\!\!\left\langle\begin{matrix}n\\k\end{matrix}\right\rangle\!\!\right\rangle \binom{x+n-k-1}{2n}$$

## Catalan number

$c_n$ 表示长度为 $2n$ 的合法括号序的数量

$c_1 = 1, c_{n+1} = \sum_{i=1}^{n} c_i \times c_{n+1-i}$

$c_n = \frac{\binom{2n}{n}}{n+1}$

## Bell number

$B_n$ 表示基数为 $n$ 的集合的划分方案数

$$B_i = \begin{cases} 1 & i = 0 \\ \sum_{k=0}^{n}\binom{n}{k}B_k & i > 0 \end{cases}$$

$B_n = \sum_{k=0}^{n}\left\{\begin{smallmatrix}n\\k\end{smallmatrix}\right\}$

$B_{p^m+n} \equiv mB_n + B_{n+1} \pmod{p}$

## 五边形数定理

$p(n)$ 表示将 $n$ 划分为若干个正整数之和的方案数

$p(n) = \sum_{k \in \mathbb{N}^*}(-1)^{k-1}p(n - \frac{k(3k-1)}{2})$

## Bernoulli number

$\sum_{j=0}^{m}\binom{m+1}{j}B_j = 0, m > 0$

$$B_i = \begin{cases} 1 & i = 0 \\ -\frac{\sum_{j=0}^{i-1}\binom{i+1}{j}B_j}{i+1} & i > 0 \end{cases}$$

$\sum_{k=1}^{n}k^m = \frac{1}{m+1}\sum_{k=0}^{m}\binom{m+1}{k}B_k n^{m+1-k}$

## Stirling permutation

$1, 1, 2, 2 \dots, n, n$ 的排列中，对于每个 $i$，都有两个 $i$ 之间的数大于 $i$

排列方案数为 $(2n-1)!!$

## Eulerian number

$\left\langle {n \atop k} \right\rangle$ 表示 1 到 $n$ 的排列中，恰有 $k$ 个数比前一个大的方案数

$\left\langle {n \atop 0} \right\rangle = \left\langle {n \atop n-1} \right\rangle = 1$

$\left\langle {0 \atop m} \right\rangle = [m = 0]$

$\left\langle {n \atop m} \right\rangle = \left\langle {n \atop n-1-m} \right\rangle$

$\left\langle {n \atop m} \right\rangle = (m+1)\left\langle {n-1 \atop m} \right\rangle + (n-m)\left\langle {n-1 \atop m-1} \right\rangle$

$\left\langle {n \atop m} \right\rangle = \sum\limits_{k=0}^{m} (-1)^k \binom{n+1}{k}(m+1-k)^n$

## Eulerian number (2nd kind)

$\left\langle\!\left\langle {n \atop k} \right\rangle\!\right\rangle$ 表示 Stirling permutation 中，恰有 $k$ 个数比前一个大的方案数

$\left\langle\!\left\langle {n \atop m} \right\rangle\!\right\rangle = (2n-m-1)\left\langle\!\left\langle {n-1 \atop m-1} \right\rangle\!\right\rangle + (m+1)\left\langle\!\left\langle {n-1 \atop m} \right\rangle\!\right\rangle$

$\left\langle\!\left\langle {n \atop 0} \right\rangle\!\right\rangle = 1$

$\left\langle\!\left\langle {0 \atop m} \right\rangle\!\right\rangle = [m = 0]$

## Burnside lemma

Let $G$ be a finite group that acts on a set $X$. For each $g$ in $G$ let $X^g$ denote the set of elements in $X$ that are fixed by $g$ (also said to be left invariant by $g$), i.e. $X^g = \{x \in X \mid g.x = x\}$. Burnside's lemma asserts the following formula for the number of orbits, denoted $|X/G|$:

$|X/G| = \frac{1}{|G|} \sum\limits_{g \in G} |X^g|$

Example application: The number of rotationally distinct colorings of the faces of a cube using $n$ colors
Let $X$ be the set of $n^6$ possible face colour combinations that can be applied to a cube in one particular orientation, and let the rotation group $G$ of the cube act on $X$ in the natural manner. Then two elements of $X$ belong to the same orbit precisely when one is simply a rotation of the other. The number of rotationally distinct colourings is thus the same as the number of orbits and can be found by counting the sizes of the fixed sets for the 24 elements of $G$.

- one identity element which leaves all $n^6$ elements of $X$ unchanged

- six 90-degree face rotations, each of which leaves $n^3$ of the elements of $X$ unchanged

- three 180-degree face rotations, each of which leaves $n^4$ of the elements of $X$ unchanged

- eight 120-degree vertex rotations, each of which leaves $n^2$ of the elements of $X$ unchanged

- six 180-degree edge rotations, each of which leaves $n^3$ of the elements of $X$ unchanged

The average fix size is thus $\frac{1}{24}(3^6 + 6 \cdot 3^3 + 3 \cdot 3^4 + 8 \cdot 3^2 + 6 \cdot 3^3) = 57$
Hence there are 57 rotationally distinct colorings of the faces of a cube in 3 colours.

## Möbius function

$$\mu(n) = \begin{cases} 1 & n \text{ is a square-free positive integer with an even number of prime factors} \\ -1 & n \text{ is a square-free positive integer with an odd number of prime factors} \\ 0 & n \text{ has a squared prime factor} \end{cases}$$

$$\sum_{d|n} \mu(d) = \begin{cases} 1 & n = 1 \\ 0 & n > 1 \end{cases}$$

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d) g(\tfrac{n}{d})$$

## Lagrange polynomial

给定次数为 $n$ 的多项式函数 $L(x)$ 上的 $n+1$ 个点 $(x_0, y_0), (x_1, y_1), \dots, (x_n, y_n)$

则 $L(x) = \sum\limits_{j=0}^{n} y_j \prod\limits_{0 \le m \le n, m \ne j} \frac{x - x_m}{x_j - x_m}$

# Chapter 3

# Geometry

## 3.1 点、直线、圆 (gy)

```
1  using number = long double;
2  const number eps = 1e-8;

3  number _sqrt(number x) {
4      return std::sqrt(std::max(x, (number) 0));
5  }
6  number _asin(number x) {
7      x = std::min(x, (number) 1), x = std::max(x, (number) -1);
8      return std::asin(x);
9  }
10 number _acos(number x) {
11     x = std::min(x, (number) 1), x = std::max(x, (number) -1);
12     return std::acos(x);
13 }

14 int sgn(number x) {
15     return (x > eps) - (x < -eps);
16 }
17 int cmp(number x, number y) {
18     return sgn(x - y);
19 }

20 struct point {
21     number x, y;
22     point() {}
23     point(number x, number y) : x(x), y(y) {}

24     number len2() const {
25         return x * x + y * y;
26     }
27     number len() const {
28         return _sqrt(len2());
29     }
30     point unit() const {
31         return point(x / len(), y / len());
32     }
33     point rotate90() const {
34         return point(-y, x);
35     }

36     friend point operator+(const point &a, const point &b) {
37         return point(a.x + b.x, a.y + b.y);
```

```
38          }
39          friend point operator-(const point &a, const point &b) {
40              return point(a.x - b.x, a.y - b.y);
41          }
42          friend point operator*(const point &a, number b) {
43              return point(a.x * b, a.y * b);
44          }
45          friend point operator/(const point &a, number b) {
46              return point(a.x / b, a.y / b);
47          }
48          friend number dot(const point &a, const point &b) {
49              return a.x * b.x + a.y * b.y;
50          }
51          friend number det(const point &a, const point &b) {
52              return a.x * b.y - a.y * b.x;
53          }
54          friend number operator==(const point &a, const point &b) {
55              return cmp(a.x, b.x) == 0 && cmp(a.y, b.y) == 0;
56          }
57      };
58      number dis2(const point &a, const point &b) {
59          return (a - b).len2();
60      }
61      number dis(const point &a, const point &b) {
62          return (a - b).len();
63      }

64      struct line {
65          point a, b;
66          line() {}
67          line(point a, point b) : a(a), b(b) {}
68          point value() const {
69              return b - a;
70          }
71      };

72      bool point_on_line(const point &p, const line &l) {
73          return sgn(det(p - l.a, p - l.b)) == 0;
74      }
75      // including endpoint
76      bool point_on_ray(const point &p, const line &l) {
77          return sgn(det(p - l.a, p - l.b)) == 0 &&
78              sgn(dot(p - l.a, l.b - l.a)) >= 0;
79      }
80      // including endpoints
81      bool point_on_seg(const point &p, const line &l) {
82          return sgn(det(p - l.a, p - l.b)) == 0 &&
83              sgn(dot(p - l.a, l.b - l.a)) >= 0 &&
84              sgn(dot(p - l.b, l.a - l.b)) >= 0;
85      }
86      bool seg_has_intersection(const line &a, const line &b) {
87          if (point_on_seg(a.a, b) || point_on_seg(a.b, b) ||
88                  point_on_seg(b.a, a) || point_on_seg(b.b, a))
89              return /* including endpoints */ true;
90          return sgn(det(a.a - b.a, b.b - b.a)) * sgn(det(a.b - b.a, b.b - b.a)) < 0
91              && sgn(det(b.a - a.a, a.b - a.a)) * sgn(det(b.b - a.a, a.b - a.a)) < 0;
92      }
93      point intersect(const line &a, const line &b) {
94          number s1 = det(a.b - a.a, b.a - a.a);
95          number s2 = det(a.b - a.a, b.b - a.a);
```

```
 96      return (b.a * s2 - b.b * s1) / (s2 - s1);
 97  }
 98  point projection(const point &p, const line &l) {
 99      return l.a + (l.b - l.a) * dot(p - l.a, l.b - l.a) / (l.b - l.a).len2();
100  }
101  number dis(const point &p, const line &l) {
102      return std::abs(det(p - l.a, l.b - l.a)) / (l.b - l.a).len();
103  }
104  point symmetry_point(const point &a, const point &o) {
105      return o + o - a;
106  }
107  point reflection(const point &p, const line &l) {
108      return symmetry_point(p, projection(p, l));
109  }
110  struct circle {
111      point o;
112      number r;
113      circle() {}
114      circle(point o, number r) : o(o), r(r) {}
115  };
116  bool intersect(const line &l, const circle &a, point &p1, point &p2) {
117      number x = dot(l.a - a.o, l.b - l.a);
118      number y = (l.b - l.a).len2();
119      number d = x * x - y * ((l.a - a.o).len2() - a.r * a.r);
120      if (sgn(d) < 0) return false;
121      point p = l.a - (l.b - l.a) * (x / y), delta = (l.b - l.a) * (_sqrt(d) / y);
122      p1 = p + delta, p2 = p - delta;
123      return true;
124  }
125  bool intersect(const circle &a, const circle &b, point &p1, point &p2) {
126      if (a.o == b.o && cmp(a.r, b.r) == 0)
127          return /* value for coincident circles */ false;
128      number s1 = (b.o - a.o).len();
129      if (cmp(s1, a.r + b.r) > 0 || cmp(s1, std::abs(a.r - b.r)) < 0)
130          return false;
131      number s2 = (a.r * a.r - b.r * b.r) / s1;
132      number aa = (s1 + s2) / 2, bb = (s1 - s2) / 2;
133      point p = (b.o - a.o) * (aa / (aa + bb)) + a.o;
134      point delta = (b.o - a.o).unit().rotate90() * _sqrt(a.r * a.r - aa * aa);
135      p1 = p + delta, p2 = p - delta;
136      return true;
137  }
138  bool tangent(const point &p0, const circle &c, point &p1, point &p2) {
139      number x = (p0 - c.o).len2();
140      number d = x - c.r * c.r;
141      if (sgn(d) < 0) return false;
142      if (sgn(d) == 0)
143          return /* value for point_on_line */ false;
144      point p = (p0 - c.o) * (c.r * c.r / x);
145      point delta = ((p0 - c.o) * (-c.r * _sqrt(d) / x)).rotate90();
146      p1 = c.o + p + delta;
147      p2 = c.o + p - delta;
148      return true;
149  }
150  bool ex_tangent(const circle &a, const circle &b, line &l1, line &l2) {
151      if (cmp(std::abs(a.r - b.r), (b.o - a.o).len()) == 0) {
152          point p1, p2;
153          intersect(a, b, p1, p2);
154          l1 = l2 = line(p1, p1 + (a.o - p1).rotate90());
```

```
155        return true;
156    } else if (cmp(a.r, b.r) == 0) {
157        point dir = b.o - a.o;
158        dir = (dir * (a.r / dir.len())).rotate90();
159        l1 = line(a.o + dir, b.o + dir);
160        l2 = line(a.o - dir, b.o - dir);
161        return true;
162    } else {
163        point p = (b.o * a.r - a.o * b.r) / (a.r - b.r);
164        point p1, p2, q1, q2;
165        if (tangent(p, a, p1, p2) && tangent(p, b, q1, q2)) {
166            l1 = line(p1, q1);
167            l2 = line(p2, q2);
168            return true;
169        } else {
170            return false;
171        }
172    }
173 }
174 bool in_tangent(const circle &a, const circle &b, line &l1, line &l2) {
175    if (cmp(a.r + b.r, (b.o - a.o).len()) == 0) {
176        point p1, p2;
177        intersect(a, b, p1, p2);
178        l1 = l2 = line(p1, p1 + (a.o - p1).rotate90());
179        return true;
180    } else {
181        point p = (b.o * a.r + a.o * b.r) / (a.r + b.r);
182        point p1, p2, q1, q2;
183        if (tangent(p, a, p1, p2) && tangent(p, b, q1, q2)) {
184            l1 = line(p1, q1);
185            l2 = line(p2, q2);
186            return true;
187        } else {
188            return false;
189        }
190    }
191 }
```

## 3.2   平面最近点对 (Grimoire)

```
1  bool byY(P a,P b){return a.y<b.y;}
2  LL solve(P *p,int l,int r){
3      LL d=1LL<<62;
4      if(l==r)
5          return d;
6      if(l+1==r)
7          return dis2(p[l],p[r]);
8      int mid=(l+r)>>1;
9      d=min(solve(l,mid),d);
10     d=min(solve(mid+1,r),d);
11     vector<P>tmp;
12     for(int i=l;i<=r;i++)
13         if(sqr(p[mid].x-p[i].x)<=d)
14             tmp.push_back(p[i]);
15     sort(tmp.begin(),tmp.end(),byY);
16     for(int i=0;i<tmp.size();i++)
17         for(int j=i+1;j<tmp.size()&&j-i<10;j++)
18             d=min(d,dis2(tmp[i],tmp[j]));
19     return d;
```

```
20 | }
```

## 3.3 凸包游戏 (Grimoire)

给定凸包，$O(n \log n)$ 完成询问:

- 点在凸包内

- 凸包外的点到凸包的两个切点

- 向量关于凸包的切点

- 直线与凸包的交点

传入凸包要求 1 号点为 $Pair(x, y)$ 最小的

```cpp
const int INF = 1000000000;
struct Convex
{
    int n;
    vector<Point> a, upper, lower;
    Convex(vector<Point> _a) : a(_a) {
        n = a.size();
        int ptr = 0;
        for(int i = 1; i < n; ++ i) if (a[ptr] < a[i]) ptr = i;
        for(int i = 0; i <= ptr; ++ i) lower.push_back(a[i]);
        for(int i = ptr; i < n; ++ i) upper.push_back(a[i]);
        upper.push_back(a[0]);
    }
    int sign(long long x) { return x < 0 ? -1 : x > 0; }
    pair<long long, int> get_tangent(vector<Point> &convex, Point vec) {
        int l = 0, r = (int)convex.size() - 2;
        for( ; l + 1 < r; ) {
            int mid = (l + r) / 2;
            if (sign((convex[mid + 1] - convex[mid]).det(vec)) > 0) r = mid;
            else l = mid;
        }
        return max(make_pair(vec.det(convex[r]), r)
            , make_pair(vec.det(convex[0]), 0));
    }
    void update_tangent(const Point &p, int id, int &i0, int &i1) {
        if ((a[i0] - p).det(a[id] - p) > 0) i0 = id;
        if ((a[i1] - p).det(a[id] - p) < 0) i1 = id;
    }
    void binary_search(int l, int r, Point p, int &i0, int &i1) {
        if (l == r) return;
        update_tangent(p, l % n, i0, i1);
        int sl = sign((a[l % n] - p).det(a[(l + 1) % n] - p));
        for( ; l + 1 < r; ) {
            int mid = (l + r) / 2;
            int smid = sign((a[mid % n] - p).det(a[(mid + 1) % n] - p));
            if (smid == sl) l = mid;
            else r = mid;
        }
        update_tangent(p, r % n, i0, i1);
    }
    int binary_search(Point u, Point v, int l, int r) {
        int sl = sign((v - u).det(a[l % n] - u));
        for( ; l + 1 < r; ) {
            int mid = (l + r) / 2;
            int smid = sign((v - u).det(a[mid % n] - u));
```

41

```
46              if (smid == sl) l = mid;
47              else r = mid;
48          }
49          return l % n;
50      }
51      // 判定点是否在凸包内，在边界返回 true
52      bool contain(Point p) {
53          if (p.x < lower[0].x || p.x > lower.back().x) return false;
54          int id = lower_bound(lower.begin(), lower.end()
55              , Point(p.x, -INF)) - lower.begin();
56          if (lower[id].x == p.x) {
57              if (lower[id].y > p.y) return false;
58          } else if ((lower[id - 1] - p).det(lower[id] - p) < 0) return false;
59          id = lower_bound(upper.begin(), upper.end(), Point(p.x, INF)
60              , greater<Point>()) - upper.begin();
61          if (upper[id].x == p.x) {
62              if (upper[id].y < p.y) return false;
63          } else if ((upper[id - 1] - p).det(upper[id] - p) < 0) return false;
64          return true;
65      }
66      // 求点 p 关于凸包的两个切点，如果在凸包外则有序返回编号
67      // 共线的多个切点返回任意一个，否则返回 false
68      bool get_tangent(Point p, int &i0, int &i1) {
69          if (contain(p)) return false;
70          i0 = i1 = 0;
71          int id = lower_bound(lower.begin(), lower.end(), p) - lower.begin();
72          binary_search(0, id, p, i0, i1);
73          binary_search(id, (int)lower.size(), p, i0, i1);
74          id = lower_bound(upper.begin(), upper.end(), p
75              , greater<Point>()) - upper.begin();
76          binary_search((int)lower.size() - 1, (int)lower.size() - 1 + id, p, i0, i1);
77          binary_search((int)lower.size() - 1 + id
78              , (int)lower.size() - 1 + (int)upper.size(), p, i0, i1);
79          return true;
80      }
81      // 求凸包上和向量 vec 叉积最大的点，返回编号，共线的多个切点返回任意一个
82      int get_tangent(Point vec) {
83          pair<long long, int> ret = get_tangent(upper, vec);
84          ret.second = (ret.second + (int)lower.size() - 1) % n;
85          ret = max(ret, get_tangent(lower, vec));
86          return ret.second;
87      }
88      // 求凸包和直线 u,v 的交点，如果无严格相交返回 false.
89      //如果有则是和 (i,next(i)) 的交点，两个点无序，交在点上不确定返回前后两条线段其中之一
90      bool get_intersection(Point u, Point v, int &i0, int &i1) {
91          int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
92          if (sign((v - u).det(a[p0] - u)) * sign((v - u).det(a[p1] - u)) < 0) {
93              if (p0 > p1) swap(p0, p1);
94              i0 = binary_search(u, v, p0, p1);
95              i1 = binary_search(u, v, p1, p0 + n);
96              return true;
97          } else {
98              return false;
99          }
100     }
101 };
```

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## 3.4  半平面交 (Grimoire)

```cpp
struct P{
    int quad() const { return sgn(y) == 1 || (sgn(y) == 0 && sgn(x) >= 0);}
};
struct L{
    bool onLeft(const P &p) const { return sgn((b - a)*( p - a)) > 0; }
    L push() const{ // push out eps
        const double eps = 1e-10;
        P delta = (b - a).turn90().norm() * eps;
        return L(a - delta, b - delta);
    }
};
bool sameDir(const L &l0, const L &l1) {
    return parallel(l0, l1) && sgn((l0.b - l0.a)^(l1.b - l1.a)) == 1;
}
bool operator < (const P &a, const P &b) {
    if (a.quad() != b.quad())
        return a.quad() < b.quad();
    else
        return sgn((a*b)) > 0;
}
bool operator < (const L &l0, const L &l1) {
    if (sameDir(l0, l1))
        return l1.onLeft(l0.a);
    else
        return (l0.b - l0.a) < (l1.b - l1.a);
}
bool check(const L &u, const L &v, const L &w) {
    return w.onLeft(intersect(u, v));
}
vector<P> intersection(vector<L> &l) {
    sort(l.begin(), l.end());
    deque<L> q;
    for (int i = 0; i < (int)l.size(); ++i) {
        if (i && sameDir(l[i], l[i - 1])) {
            continue;
        }
        while (q.size() > 1
            && !check(q[q.size() - 2], q[q.size() - 1], l[i]))
                q.pop_back();
        while (q.size() > 1
            && !check(q[1], q[0], l[i]))
                q.pop_front();
        q.push_back(l[i]);
    }
    while (q.size() > 2
        && !check(q[q.size() - 2], q[q.size() - 1], q[0]))
            q.pop_back();
    while (q.size() > 2
        && !check(q[1], q[0], q[q.size() - 1]))
            q.pop_front();
    vector<P> ret;
    for (int i = 0; i < (int)q.size(); ++i)
    ret.push_back(intersect(q[i], q[(i + 1) % q.size()]));
    return ret;
}
```

## 3.5　**点在多边形内** (Grimoire)

```
bool inPoly(P p,vector<P>poly){
    int cnt=0;
    for(int i=0;i<poly.size();i++){
        P a=poly[i],b=poly[(i+1)%poly.size()];
        if(onSeg(p,L(a,b)))
            return false;
        int x=sgn(det(a,p,b));
        int y=sgn(a.y-p.y);
        int z=sgn(b.y-p.y);
        cnt+=(x>0&&y<=0&&z>0);
        cnt-=(x<0&&z<=0&&y>0);
    }
    return cnt;
}
```

## 3.6　**最小圆覆盖** (Grimoire)

```
struct line{
    point p,v;
};
point Rev(point v){return point(-v.y,v.x);}
point operator*(line A,line B){
    point u=B.p-A.p;
    double t=(B.v*u)/(B.v*A.v);
    return A.p+A.v*t;
}
point get(point a,point b){
    return (a+b)/2;
}
point get(point a,point b,point c){
    if(a==b)return get(a,c);
    if(a==c)return get(a,b);
    if(b==c)return get(a,b);
    line ABO=(line){(a+b)/2,Rev(a-b)};
    line BCO=(line){(c+b)/2,Rev(b-c)};
    return ABO*BCO;
}
int main(){
    scanf("%d",&n);
    for(int i=1;i<=n;i++)scanf("%lf%lf",&p[i].x,&p[i].y);
    random_shuffle(p+1,p+1+n);
    O=p[1];r=0;
    for(int i=2;i<=n;i++){
        if(dis(p[i],O)<r+1e-6)continue;
        O=get(p[1],p[i]);r=dis(O,p[i]);
        for(int j=1;j<i;j++){
            if(dis(p[j],O)<r+1e-6)continue;
            O=get(p[i],p[j]);r=dis(O,p[i]);
            for(int k=1;k<j;k++){
                if(dis(p[k],O)<r+1e-6)continue;
                O=get(p[i],p[j],p[k]);r=dis(O,p[i]);
            }
        }
    }printf("%.2lf %.2lf %.2lf\n",O.x,O.y,r);
    return 0;
}
```

## 3.7   最小球覆盖 (Grimoire)

```cpp
bool equal(const double & x, const double & y) {
    return x + eps > y and y + eps > x;
}
double operator % (const Point & a, const Point & b) {
    return a.x * b.x + a.y * b.y + a.z * b.z;
}
Point operator * (const Point & a, const Point & b) {
    return Point(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z, a.x * b.y - a.y * b.x);
}
struct Circle {
    double r; Point o;
};
struct Plane {
    Point nor;
    double m;
    Plane(const Point & nor, const Point & a) : nor(nor){
        m = nor % a;
    }
};
Point intersect(const Plane & a, const Plane & b, const Plane & c) {
    Point c1(a.nor.x, b.nor.x, c.nor.x), c2(a.nor.y, b.nor.y, c.nor.y), c3(a.nor.z, b.nor.z,
        ↪ c.nor.z), c4(a.m, b.m, c.m);
    return 1 / ((c1 * c2) % c3) * Point((c4 * c2) % c3, (c1 * c4) % c3, (c1 * c2) % c4);
}
bool in(const Point & a, const Circle & b) {
    return sign((a - b.o).len() - b.r) <= 0;
}
bool operator < (const Point & a, const Point & b) {
    if(!equal(a.x, b.x)) {
        return a.x < b.x;
    }
    if(!equal(a.y, b.y)) {
        return a.y < b.y;
    }
    if(!equal(a.z, b.z)) {
        return a.z < b.z;
    }
    return false;
}
bool operator == (const Point & a, const Point & b) {
    return equal(a.x, b.x) and equal(a.y, b.y) and equal(a.z, b.z);
}
vector<Point> vec;
Circle calc() {
    if(vec.empty()) {
        return Circle(Point(0, 0, 0), 0);
    }else if(1 == (int)vec.size()) {
        return Circle(vec[0], 0);
    }else if(2 == (int)vec.size()) {
        return Circle(0.5 * (vec[0] + vec[1]), 0.5 * (vec[0] - vec[1]).len());
    }else if(3 == (int)vec.size()) {
        double r((vec[0] - vec[1]).len() * (vec[1] - vec[2]).len() * (vec[2] - vec[0]).len() / 2 /
            ↪ fabs(((vec[0] - vec[2]) * (vec[1] - vec[2])).len()));
        return Circle(intersect(Plane(vec[1] - vec[0], 0.5 * (vec[1] + vec[0])),
                        Plane(vec[2] - vec[1], 0.5 * (vec[2] + vec[1])),
                    Plane((vec[1] - vec[0]) * (vec[2] - vec[0]), vec[0])), r);
    }else {
        Point o(intersect(Plane(vec[1] - vec[0], 0.5 * (vec[1] + vec[0])),
```

```
57                       Plane(vec[2] - vec[0], 0.5 * (vec[2] + vec[0])),
58                       Plane(vec[3] - vec[0], 0.5 * (vec[3] + vec[0])))));
59           return Circle(o, (o - vec[0]).len());
60       }
61 }
62 Circle miniBall(int n) {
63     Circle res(calc());
64     for(int i(0); i < n; i++) {
65         if(!in(a[i], res)) {
66             vec.push_back(a[i]);
67             res = miniBall(i);
68             vec.pop_back();
69             if(i) {
70                 Point tmp(a[i]);
71                 memmove(a + 1, a, sizeof(Point) * i);
72                 a[0] = tmp;
73             }
74         }
75     }
76     return res;
77 }
78 int main() {
79     int n;
80     sort(a, a + n);
81     n = unique(a, a + n) - a;
82     vec.clear();
83     printf("%.10f\n", miniBall(n).r);
84 }
```

## 3.8  圆并 (Grimoire)

```
1  double ans[2001];
2  struct Point {
3      double x, y;
4      Point(){}
5      Point(const double & x, const double & y) : x(x), y(y) {}
6      void scan() {scanf("%lf%lf", &x, &y);}
7      double sqrlen() {return sqr(x) + sqr(y);}
8      double len() {return sqrt(sqrlen());}
9      Point rev() {return Point(y, -x);}
10     void print() {printf("%f %f\n", x, y);}
11     Point zoom(const double & d) {double lambda = d / len(); return Point(lambda * x, lambda * y);}
12 } dvd, a[2001];
13 Point centre[2001];
14 double atan2(const Point & x) {
15     return atan2(x.y, x.x);
16 }
17 Point operator - (const Point & a, const Point & b) {
18     return Point(a.x - b.x, a.y - b.y);
19 }
20 Point operator + (const Point & a, const Point & b) {
21     return Point(a.x + b.x, a.y + b.y);
22 }
23 double operator * (const Point & a, const Point & b) {
24     return a.x * b.y - a.y * b.x;
25 }
26 Point operator * (const double & a, const Point & b) {
27     return Point(a * b.x, a * b.y);
28 }
```

```
29  double operator % (const Point & a, const Point & b) {
30      return a.x * b.x + a.y * b.y;
31  }
32  struct circle {
33      double r; Point o;
34      circle() {}
35      void scan() {
36          o.scan();
37          scanf("%lf", &r);
38      }
39  } cir[2001];
40  struct arc {
41      double theta;
42      int delta;
43      Point p;
44      arc() {};
45      arc(const double & theta, const Point & p, int d) : theta(theta), p(p), delta(d) {}
46  } vec[4444];
47  int nV;
48  inline bool operator < (const arc & a, const arc & b) {
49      return a.theta + eps < b.theta;
50  }
51  int cnt;
52  inline void psh(const double t1, const Point p1, const double t2, const Point p2) {
53      if(t2 + eps < t1)
54          cnt++;
55      vec[nV++] = arc(t1, p1, 1);
56      vec[nV++] = arc(t2, p2, -1);
57  }
58  inline double cub(const double & x) {
59      return x * x * x;
60  }
61  inline void combine(int d, const double & area, const Point & o) {
62      if(sign(area) == 0) return;
63      centre[d] = 1 / (ans[d] + area) * (ans[d] * centre[d] + area * o);
64      ans[d] += area;
65  }
66  bool equal(const double & x, const double & y) {
67      return x + eps> y and y + eps > x;
68  }
69  bool equal(const Point & a, const Point & b) {
70      return equal(a.x, b.x) and equal(a.y, b.y);
71  }
72  bool equal(const circle & a, const circle & b) {
73      return equal(a.o, b.o) and equal(a.r, b.r);
74  }
75  bool f[2001];
76  int main() {
77      int n, m, index;
78      while(EOF != scanf("%d%d%d", &m, &n, &index)) {
79          index--;
80          for(int i(0); i < m; i++) {
81              a[i].scan();
82          }
83          for(int i(0); i < n; i++) {
84              cir[i].scan();//n 个圆
85          }
86          for(int i(0); i < n; i++) {//这一段在去重圆 能加速 删掉不会错
87              f[i] = true;
88              for(int j(0); j < n; j++) if(i != j) {
```

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```
89              if(equal(cir[i], cir[j]) and i < j or !equal(cir[i], cir[j]) and cir[i].r <
                 ↪ cir[j].r + eps and (cir[i].o - cir[j].o).sqrlen() < sqr(cir[i].r - cir[j].r) +
                 ↪ eps) {
90                  f[i] = false;
91                  break;
92              }
93          }
94      }
95      int n1(0);
96      for(int i(0); i < n; i++)
97          if(f[i])
98              cir[n1++] = cir[i];
99      n = n1;//去重圆结束
100     fill(ans, ans + n + 1, 0);//ans[i] 表示被圆覆盖至少 i 次的面积
101     fill(centre, centre + n + 1, Point(0, 0));//centre[i] 表示上面 ans[i] 部分的重心
102     for(int i(0); i < m; i++)
103         combine(0, a[i] * a[(i + 1) % m] * 0.5, 1. / 3 * (a[i] + a[(i + 1) % m]));
104     for(int i(0); i < n; i++) {
105         dvd = cir[i].o - Point(cir[i].r, 0);
106         nV = 0;
107         vec[nV++] = arc(-pi, dvd, 1);
108         cnt = 0;
109         for(int j(0); j < n; j++) if(j != i) {
110             double d = (cir[j].o - cir[i].o).sqrlen();
111             if(d < sqr(cir[j].r - cir[i].r) + eps) {
112                 if(cir[i].r + i * eps < cir[j].r + j * eps)
113                     psh(-pi, dvd, pi, dvd);
114             }else if(d + eps < sqr(cir[j].r + cir[i].r)) {
115                 double lambda = 0.5 * (1 + (sqr(cir[i].r) - sqr(cir[j].r)) / d);
116                 Point cp(cir[i].o + lambda * (cir[j].o - cir[i].o));
117                 Point nor((cir[j].o - cir[i].o).rev().zoom(sqrt(sqr(cir[i].r) - (cp -
                     ↪ cir[i].o).sqrlen()))));
118                 Point frm(cp + nor);
119                 Point to(cp - nor);
120                 psh(atan2(frm - cir[i].o), frm, atan2(to - cir[i].o), to);
121             }
122         }
123         sort(vec + 1, vec + nV);
124         vec[nV++] = arc(pi, dvd, -1);
125         for(int j = 0; j + 1 < nV; j++) {
126             cnt += vec[j].delta;
127             //if(cnt == 1) {//如果只算 ans[1] 和 centre[1], 可以加这个 if 加速.
128                 double theta(vec[j + 1].theta - vec[j].theta);
129                 double area(sqr(cir[i].r) * theta * 0.5);
130                 combine(cnt, area, cir[i].o + 1. / area / 3 * cub(cir[i].r) * Point(sin(vec[j +
                     ↪ 1].theta) - sin(vec[j].theta), cos(vec[j].theta) - cos(vec[j + 1].theta)));
131                 combine(cnt, -sqr(cir[i].r) * sin(theta) * 0.5, 1. / 3 * (cir[i].o + vec[j].p +
                     ↪ vec[j + 1].p));
132                 combine(cnt, vec[j].p * vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j +
                     ↪ 1].p));
133             //}
134         }
135     }
136     combine(0, -ans[1], centre[1]);
137     for(int i = 0; i < m; i++) {
138         if(i != index)
139             (a[index] - Point((a[i] - a[index]) * (centre[0] - a[index]), (a[i] - a[index]) %
                 ↪ (centre[0] - a[index])).zoom((a[i] - a[index]).len())).print();
140         else
141             a[i].print();
142     }
```

48

```
143        }
144        return 0;
145 }
```

## 3.9   圆与多边形并 (Grimoire)

```
1  double form(double x){
2      while(x>=2*pi)x-=2*pi;
3      while(x<0)x+=2*pi;
4      return x;
5  }
6  double calcCir(C cir){
7      vector<double>ang;
8      ang.push_back(0);
9      ang.push_back(pi);
10     double ans=0;
11     for(int i=1;i<=n;i++){
12         if(cir==c[i])continue;
13         P p1,p2;
14         if(intersect(cir,c[i],p1,p2)){
15             ang.push_back(form(cir.ang(p1)));
16             ang.push_back(form(cir.ang(p2)));
17         }
18     }
19     for(int i=1;i<=m;i++){
20         vector<P>tmp;
21         tmp=intersect(poly[i],cir);
22         for(int j=0;j<tmp.size();j++){
23             ang.push_back(form(cir.ang(tmp[j])));
24         }
25     }
26     sort(ang.begin(),ang.end());
27     for(int i=0;i<ang.size();i++){
28         double t1=ang[i],t2=(i+1==ang.size()?ang[0]+2*pi:ang[i+1]);
29         P p=cir.at((t1+t2)/2);
30         int ok=1;
31         for(int j=1;j<=n;j++){
32             if(cir==c[j])continue;
33             if(inC(p,c[j],true)){
34                 ok=0;
35                 break;
36             }
37         }
38         for(int j=1;j<=m&&ok;j++){
39             if(inPoly(p,poly[j],true)){
40                 ok=0;
41                 break;
42             }
43         }
44         if(ok){
45             double r=cir.r,x0=cir.o.x,y0=cir.o.y;
46             ans+=(r*r*(t2-t1)+r*x0*(sin(t2)-sin(t1))-r*y0*(cos(t2)-cos(t1)))/2;
47         }
48     }
49     return ans;
50 }
51 P st;
```

49

```
52 bool bySt(P a,P b){
53     return dis(a,st)<dis(b,st);
54 }
55 double calcSeg(L l){
56     double ans=0;
57     vector<P>pt;
58     pt.push_back(l.a);
59     pt.push_back(l.b);
60     for(int i=1;i<=n;i++){
61         P p1,p2;
62         if(intersect(c[i],l,p1,p2)){
63             if(onSeg(p1,l))
64                 pt.push_back(p1);
65             if(onSeg(p2,l))
66                 pt.push_back(p2);
67         }
68     }
69     st=l.a;
70     sort(pt.begin(),pt.end(),bySt);
71     for(int i=0;i+1<pt.size();i++){
72         P p1=pt[i],p2=pt[i+1];
73         P p=(p1+p2)/2;
74         int ok=1;
75         for(int j=1;j<=n;j++){
76             if(sgn(dis(p,c[j].o),c[j].r)<0){
77                 ok=0;
78                 break;
79             }
80         }
81         if(ok){
82             double x1=p1.x,y1=p1.y,x2=p2.x,y2=p2.y;
83             double res=(x1*y2-x2*y1)/2;
84             ans+=res;
85         }
86     }
87     return ans;
88 }
```

## 3.10   三角剖分 (Grimoire)

Triangulation::find 返回包含某点的三角形
Triangulation::add_point 将某点加入三角剖分
某个 *Triangle* 在三角剖分中当且仅当它的 *has_children* 为 0
如果要找到三角形 u 的邻域，则枚举它的所有 u.edge[i].tri，该条边的两个点为 u.p[(i + 1) % 3], u.p[(i + 2) % 3]
通过三角剖分构造 V 图：连接相邻三角形外接圆圆心
注意初始化内存池和 *Triangulation :: LOTS*
复杂度 $O(n \log n)$

```
1 const int N = 100000 + 5, MAX_TRIS = N * 6;
2 const double eps = 1e-6, PI = acos(-1.0);
3 struct P {
4     double x,y; P():x(0),y(0){}
5     P(double x, double y):x(x),y(y){}
6     bool operator ==(P const& that)const {return x==that.x&&y==that.y;}
7 };
8 inline double sqr(double x) { return x*x; }
9 double dist_sqr(P const& a, P const& b){return sqr(a.x-b.x)+sqr(a.y-b.y);}
10 bool in_circumcircle(P const& p1, P const& p2, P const& p3, P const& p4) {//p4 in C(p1,p2,p3)
```

```
11          double u11 = p1.x - p4.x, u21 = p2.x - p4.x, u31 = p3.x - p4.x;
12          double u12 = p1.y - p4.y, u22 = p2.y - p4.y, u32 = p3.y - p4.y;
13          double u13 = sqr(p1.x) - sqr(p4.x) + sqr(p1.y) - sqr(p4.y);
14          double u23 = sqr(p2.x) - sqr(p4.x) + sqr(p2.y) - sqr(p4.y);
15          double u33 = sqr(p3.x) - sqr(p4.x) + sqr(p3.y) - sqr(p4.y);
16          double det = -u13*u22*u31 + u12*u23*u31 + u13*u21*u32 - u11*u23*u32 - u12*u21*u33 +
                u11*u22*u33;
17          return det > eps;
18      }
19      double side(P const& a, P const& b, P const& p) { return (b.x-a.x)*(p.y-a.y) -
            (b.y-a.y)*(p.x-a.x);}
20      typedef int SideRef; struct Triangle; typedef Triangle* TriangleRef;
21      struct Edge {
22          TriangleRef tri; SideRef side; Edge() : tri(0), side(0) {}
23          Edge(TriangleRef tri, SideRef side) : tri(tri), side(side) {}
24      };
25      struct Triangle {
26          P p[3]; Edge edge[3]; TriangleRef children[3]; Triangle() {}
27          Triangle(P const& p0, P const& p1, P const& p2) {
28              p[0] = p0; p[1] = p1; p[2] = p2;
29              children[0] = children[1] = children[2] = 0;
30          }
31          bool has_children() const { return children[0] != 0; }
32          int num_children() const {
33              return children[0] == 0 ? 0
34                  : children[1] == 0 ? 1
35                  : children[2] == 0 ? 2 : 3;
36          }
37          bool contains(P const& q) const {
38              double a=side(p[0],p[1],q), b=side(p[1],p[2],q), c=side(p[2],p[0],q);
39              return a >= -eps && b >= -eps && c >= -eps;
40          }
41      } triange_pool[MAX_TRIS], *tot_triangles;
42      void set_edge(Edge a, Edge b) {
43          if (a.tri) a.tri->edge[a.side] = b;
44          if (b.tri) b.tri->edge[b.side] = a;
45      }
46      class Triangulation {
47          public:
48              Triangulation() {
49                  const double LOTS = 1e6;//初始为极大三角形
50                  the_root = new(tot_triangles++) Triangle(P(-LOTS,-LOTS),P(+LOTS,-LOTS),P(0,+LOTS));
51              }
52              TriangleRef find(P p) const { return find(the_root,p); }
53              void add_point(P const& p) { add_point(find(the_root,p),p); }
54          private:
55              TriangleRef the_root;
56              static TriangleRef find(TriangleRef root, P const& p) {
57                  for( ; ; ) {
58                      if (!root->has_children()) return root;
59                      else for (int i = 0; i < 3 && root->children[i] ; ++i)
60                          if (root->children[i]->contains(p))
61                              {root = root->children[i]; break;}
62                  }
63              }
64              void add_point(TriangleRef root, P const& p) {
65                  TriangleRef tab,tbc,tca;
66                  tab = new(tot_triangles++) Triangle(root->p[0], root->p[1], p);
67                  tbc = new(tot_triangles++) Triangle(root->p[1], root->p[2], p);
68                  tca = new(tot_triangles++) Triangle(root->p[2], root->p[0], p);
69                  set_edge(Edge(tab,0),Edge(tbc,1)); set_edge(Edge(tbc,0),Edge(tca,1));
```

```
70          set_edge(Edge(tca,0),Edge(tab,1)); set_edge(Edge(tab,2),root->edge[2]);
71          set_edge(Edge(tbc,2),root->edge[0]); set_edge(Edge(tca,2),root->edge[1]);
72          root->children[0]=tab; root->children[1]=tbc; root->children[2]=tca;
73          flip(tab,2); flip(tbc,2); flip(tca,2);
74      }
75      void flip(TriangleRef tri, SideRef pi) {
76          TriangleRef trj = tri->edge[pi].tri; int pj = tri->edge[pi].side;
77          if(!trj || !in_circumcircle(tri->p[0],tri->p[1],tri->p[2],trj->p[pj])) return;
78          TriangleRef trk = new(tot_triangles++) Triangle(tri->p[(pi+1)%3], trj->p[pj],
                ↪ tri->p[pi]);
79          TriangleRef trl = new(tot_triangles++) Triangle(trj->p[(pj+1)%3], tri->p[pi],
                ↪ trj->p[pj]);
80          set_edge(Edge(trk,0), Edge(trl,0));
81          set_edge(Edge(trk,1), tri->edge[(pi+2)%3]); set_edge(Edge(trk,2), trj->edge[(pj+1)%3]);
82          set_edge(Edge(trl,1), trj->edge[(pj+2)%3]); set_edge(Edge(trl,2), tri->edge[(pi+1)%3]);
83          tri->children[0]=trk; tri->children[1]=trl; tri->children[2]=0;
84          trj->children[0]=trk; trj->children[1]=trl; trj->children[2]=0;
85          flip(trk,1); flip(trk,2); flip(trl,1); flip(trl,2);
86      }
87 };
88 int n; P ps[N];
89 void build(){
90     tot_triangles = triange_pool; cin >> n;
91     for(int i = 0; i < n; ++ i) scanf("%lf%lf",&ps[i].x,&ps[i].y);
92     random_shuffle(ps, ps + n); Triangulation tri;
93     for(int i = 0; i < n; ++ i) tri.add_point(ps[i]);
94 }
```

## 3.11   三维几何基础 (Grimoire)

```
1  struct P {
2      double x, y, z;
3      P(){}
4      P(double _x,double _y,double _z):x(_x),y(_y),z(_z){}
5      double len2(){
6          return (x*x+y*y+z*z);
7      }
8      double len(){
9          return sqrt(x*x+y*y+z*z);
10     }
11 };
12 bool operator==(P a,P b){
13     return sgn(a.x-b.x)==0 && sgn(a.y-b.y)==0 && sgn(a.z-b.z)==0 ;
14 }
15 bool operator<(P a,P b){
16     return sgn(a.x-b.x) ? a.x<b.x :(sgn(a.y-b.y)?a.y<b.y :a.z<b.z);
17 }
18 P operator+(P a,P b){
19     return P(a.x+b.x,a.y+b.y,a.z+b.z);
20 }
21 P operator-(P a,P b){
22     return P(a.x-b.x,a.y-b.y,a.z-b.z);
23 }
24 P operator*(P a,double b){
25     return P(a.x*b,a.y*b,a.z*b);
26 }
27 P operator/(P a,double b){
28     return P(a.x/b,a.y/b,a.z/b);
29 }
```

```
30  P operator*(const P &a, const P &b) {
31      return P(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z, a.x * b.y - a.y * b.x);
32  }
33  double operator^(const P &a, const P &b) {
34      return a.x*b.x+a.y*b.y+a.z*b.z;
35  }

36  double dis(P a,P b){return (b-a).len();}
37  double dis2(P a,P b){return (b-a).len2();}

38  // 3D line intersect
39  P intersect(const P &a0, const P &b0, const P &a1, const P &b1) {
40      double t = ((a0.x - a1.x) * (a1.y - b1.y) - (a0.y - a1.y) * (a1.x - b1.x)) / ((a0.x - b0.x) *
        ↪ (a1.y - b1.y) - (a0.y - b0.y) * (a1.x - b1.x));
41      return a0 + (b0 - a0) * t;
42  }
43  // area-line intersect
44  P intersect(const P &a, const P &b, const P &c, const P &l0, const P &l1) {

45      P p = (b-a)*(c-a); // 平面法向量
46      double t = (p^(a-l0)) / (p^(l1-l0));
47      return l0 + (l1 - l0) * t;
48  }
```

## 3.12  三维凸包 (Grimoire)

```
1   int mark[1005][1005],n, cnt;;
2   double mix(const P &a, const P &b, const P &c) {
3       return a^(b*c);
4   }
5   double area(int a, int b, int c) {
6       return ((info[b] - info[a])*(info[c] - info[a])).len();
7   }
8   double volume(int a, int b, int c, int d) {
9       return mix(info[b] - info[a], info[c] - info[a], info[d] - info[a]);
10  }
11  struct Face {
12      int a, b, c; Face() {}
13      Face(int a, int b, int c): a(a), b(b), c(c) {}
14      int &operator [](int k) {
15          if (k == 0) return a; if (k == 1) return b; return c;
16      }
17  };
18  vector <Face> face;
19  inline void insert(int a, int b, int c) {
20      face.push_back(Face(a, b, c));
21  }
22  void add(int v) {
23      vector <Face> tmp; int a, b, c; cnt++;
24      for (int i = 0; i < SIZE(face); i++) {
25          a = face[i][0]; b = face[i][1]; c = face[i][2];
26          if (sgn(volume(v, a, b, c)) < 0)
27          mark[a][b] = mark[b][a] = mark[b][c] = mark[c][b] = mark[c][a] = mark[a][c] = cnt;
28          else tmp.push_back(face[i]);
29      } face = tmp;
30      for (int i = 0; i < SIZE(tmp); i++) {
31          a = face[i][0]; b = face[i][1]; c = face[i][2];
32          if (mark[a][b] == cnt) insert(b, a, v);
33          if (mark[b][c] == cnt) insert(c, b, v);
```

```
34          if (mark[c][a] == cnt) insert(a, c, v);
35      }
36  }
37  int Find() {
38      for (int i = 2; i < n; i++) {
39          P ndir = (info[0] - info[i])*(info[1] - info[i]);
40          if (ndir == P()) continue; swap(info[i], info[2]);
41          for (int j = i + 1; j < n; j++) if (sgn(volume(0, 1, 2, j)) != 0) {
42              swap(info[j], info[3]); insert(0, 1, 2); insert(0, 2, 1); return 1;
43          }
44      }
45      return 0;
46  }
47  //find the weight center
48  double calcDist(const P &p, int a, int b, int c) {
49      return fabs(mix(info[a] - p, info[b] - p, info[c] - p) / area(a, b, c));
50  }
51  //compute the minimal distance of center of any faces
52  P findCenter() { //compute center of mass
53      double totalWeight = 0;
54      P center(.0, .0, .0);
55      P first = info[face[0][0]];
56      for (int i = 0; i < SIZE(face); ++i) {
57          P p = (info[face[i][0]]+info[face[i][1]]+info[face[i][2]]+first)*.25;
58          double weight = mix(info[face[i][0]] - first, info[face[i][1]] - first, info[face[i][2]] -
                ↪ first);
59          totalWeight += weight; center = center + p * weight;
60      }
61      center = center / totalWeight;
62      return center;
63  }
64  double minDis(P p) {
65      double res = 1e100; //compute distance
66      for (int i = 0; i < SIZE(face); ++i)
67          res = min(res, calcDist(p, face[i][0], face[i][1], face[i][2]));
68      return res;
69  }
70  void findConvex(P *info,int n) {
71      sort(info, info + n); n = unique(info, info + n) - info;
72      face.clear(); random_shuffle(info, info + n);
73      if (!Find())return abort();
74      memset(mark, 0, sizeof(mark)); cnt = 0;
75      for (int i = 3; i < n; i++) add(i);
76  }
```

## 3.13    三维绕轴旋转 (gy)

右手大拇指指向 $axis$ 方向，四指弯曲方向旋转 $w$ 弧度

```
1  P rotate(const P& s, const P& axis, double w) {
2      double x = axis.x, y = axis.y, z = axis.z;
3      double s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
4          cosw = cos(w), sinw = sin(w);
5      double a[4][4];
6      memset(a, 0, sizeof a);
7      a[3][3] = 1;
8      a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
9      a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
10     a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
```

```
11      a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
12      a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
13      a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
14      a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
15      a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
16      a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;
17      double ans[4] = {0, 0, 0, 0}, c[4] = {s.x, s.y, s.z, 1};
18      for (int i = 0; i < 4; ++i)
19          for (int j = 0; j < 4; ++j)
20              ans[i] += a[j][i] * c[j];
21      return P(ans[0], ans[1], ans[2]);
22  }
```

## 3.14  几何知识 (gy)

### Pick theorem

顶点为整点的简单多边形，其面积 $A$，内部格点数 $i$，边上格点数 $b$ 满足：
$A = i + \frac{b}{2} - 1$

### 欧拉示性数

- 三维凸包的顶点个数 $V$，边数 $E$，面数 $F$ 满足：
  $V - E + F = 2$

- 平面图的顶点个数 $V$，边数 $E$，平面被划分的区域数 $F$，组成图形的连通部分的数目 $C$ 满足：
  $V - E + F = C + 1$

### 几何公式

- **三角形**
  半周长 $p = \frac{a+b+c}{2}$
  面积 $S = \frac{1}{2}aH_a = \frac{1}{2}ab \cdot \sin C = \sqrt{p(p-a)(p-b)(p-c)} = pr = \frac{abc}{4R}$
  中线长 $M_a = \frac{1}{2}\sqrt{2(b^2+c^2)-a^2} = \frac{1}{2}\sqrt{b^2+c^2+2bc \cdot \cos A}$
  角平分线长 $T_a = \frac{\sqrt{bc((b+c)^2-a^2)}}{b+c} = \frac{2bc}{b+c}\cos\frac{A}{2}$
  高 $H_a = b\sin C = \sqrt{b^2 - (\frac{a^2+b^2-c^2}{2a})^2}$
  内切圆半径 $r = \frac{S}{p} = 4R\sin\frac{A}{2}\sin\frac{B}{2}\sin\frac{C}{2} = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}} = p\tan\frac{A}{2}\tan\frac{B}{2}\tan\frac{C}{2}$
  外接圆半径 $R = \frac{abc}{4S} = \frac{a}{2\sin A}$
  旁切圆半径 $r_A = \frac{2S}{-a+b+c}$
  重心 $(\frac{x_1+x_2+x_3}{3}, \frac{y_1+y_2+y_3}{3})$
  外心 $\left(\dfrac{\begin{vmatrix} x_1^2+y_1^2 & y_1 & 1 \\ x_2^2+y_2^2 & y_2 & 1 \\ x_3^2+y_3^2 & y_3 & 1 \end{vmatrix}}{2\begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix}}, \dfrac{\begin{vmatrix} x_1 & x_1^2+y_1^2 & 1 \\ x_2 & x_2^2+y_2^2 & 1 \\ x_3 & x_3^2+y_3^2 & 1 \end{vmatrix}}{2\begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix}}\right)$
  内心 $(\frac{ax_1+bx_2+cx_3}{a+b+c}, \frac{ay_1+by_2+cy_3}{a+b+c})$
  垂心 $\left(\dfrac{\begin{vmatrix} x_2x_3+y_2y_3 & 1 & y_1 \\ x_3x_1+y_3y_1 & 1 & y_2 \\ x_1x_2+y_1y_2 & 1 & y_3 \end{vmatrix}}{2\begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix}}, \dfrac{\begin{vmatrix} x_2x_3+y_2y_3 & x_1 & 1 \\ x_3x_1+y_3y_1 & x_2 & 1 \\ x_1x_2+y_1y_2 & x_3 & 1 \end{vmatrix}}{2\begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix}}\right)$
  旁心 $(\frac{-ax_1+bx_2+cx_3}{-a+b+c}, \frac{-ay_1+by_2+cy_3}{-a+b+c})$

- **圆**
  弧长 $l = rA$
  弦长 $a = 2\sqrt{2hr - h^2} = 2r \cdot \sin\frac{A}{2}$
  弓形高 $h = r - \sqrt{r^2 - \frac{a^2}{4}} = r(1 - \cos\frac{A}{2})$
  扇形面积 $S_1 = \frac{1}{2}lr = \frac{1}{2}Ar^2$
  弓形面积 $S_2 = \frac{1}{2}r^2(A - \sin A)$

- **Circles of Apollonius**
  已知三个两两相切的圆，半径为 $r_1, r_2, r_3$
  与它们外切的圆半径为 $\left| \frac{r_1 r_2 r_3}{r_1 r_2 + r_2 r_3 + r_3 r_1 - 2\sqrt{r_1 r_2 r_3 (r_1 + r_2 + r_3)}} \right|$
  与它们内切的圆半径为 $\frac{r_1 r_2 r_3}{r_1 r_2 + r_2 r_3 + r_3 r_1 + 2\sqrt{r_1 r_2 r_3 (r_1 + r_2 + r_3)}}$

- **棱台**
  体积 $V = \frac{1}{3}h(A_1 + A_2 + \sqrt{A_1 A_2})$
  正棱台侧面积 $S = \frac{1}{2}(p_1 + p_2)l$，$l$ 为侧高

- **球**
  体积 $V = \frac{4}{3}\pi r^3$
  表面积 $S = 4\pi r^2$

- **球台**
  侧面积 $S = 2\pi rh$
  体积 $V = \frac{1}{6}\pi h(3(r_1^2 + r_2^2) + h_h)$

- **球扇形**
  球面面积 $S = 2\pi rh$
  体积 $V = \frac{2}{3}\pi r^2 h = \frac{2}{3}\pi r^3 h(1 - \cos\varphi)$

- **球面三角形**
  考虑单位球上的球面三角形，$a, b, c$ 表示三边长（弧所对球心角），$A, B, C$ 表示三角大小（切线夹角）
  余弦定理 $\cos a = \cos b \cdot \cos c + \sin a \cdot \sin b \cdot \cos A$
  正弦定理 $\frac{\sin A}{\sin a} = \frac{\sin B}{\sin b} = \frac{\sin C}{\sin c}$
  球面面积 $S = A + B + C - \pi$

- **四面体**
  体积 $V = \frac{1}{6}\left| \overrightarrow{AB} \cdot (\overrightarrow{AC} \times \overrightarrow{AD}) \right|$

# Chapter 4

# String

## 4.1 KMP (ct)

### KMP

```
int main()
{
    for (int i = 2, j = 0; i <= n; ++i)
    {
        for ( ; j && s[j + 1] != s[i]; j = fail[j]) ;
        s[i] == s[j + 1] ? ++j : 0;
        fail[i] = j;
    }
    return 0;
}
```

### exKMP

$extend_i$ 表示 $T$ 与 $S_{i,n}$ 的最长公共前缀

```
int next[maxn], extend[maxn], fail[maxn];
void getnext(R char *s, R int len)
{
    fail[1] = 0;
    R int p = 0;
    memset(next, 0, (len + 2) << 2);
    for (R int i = 2; i <= len; ++i)
    {
        while (p && s[p + 1] != s[i]) p = fail[p];
        s[p + 1] == s[i] ? ++p : 0;
        fail[i] = p;
        p ? cmax(next[i - p + 1], p) : 0;
    }
}
void getextend(R char *s, R int lens, R char *t, R int lent)
{
    getnext(t, lent);
    R int a = 1, p = 0;

    for (R int i = 1; i <= lens; ++i)
    {
        if (i + next[i - a + 1] - 1 >= p)
        {
            cmax(p, i - 1);
```

```
24          while (p < lens && p - i + 1 < lent && s[p + 1] == t[p - i + 2]) ++p;
25          a = i;
26          extend[i] = p - i + 1;
27      }
28      else extend[i] = next[i - a + 1];
29  }
30 }
```

## 4.2   AC 自动机

## 4.3   后缀数组 (ct)

```
1  char s[maxn];
2  int sa[maxn], rank[maxn], wa[maxn], wb[maxn], cnt[maxn], height[maxn];
3  inline void build(int n, int m)
4  {
5      int *x = wa, *y = wb, *t;
6      for (int i = 1; i <= n; ++i) cnt[x[i] = s[i] - 'a' + 1]++;
7      for (int i = 1; i <= m; ++i) cnt[i] += cnt[i - 1];
8      for (int i = n; i; --i) sa[cnt[x[i]]--] = i;
9      for (int j = 1; j < n || (j == 1 && m < n); j <<= 1, t = x, x = y, y = t)
10     {
11         memset(cnt + 1, 0, m << 2);
12         int p = 0;
13         for (int i = n - j + 1; i <= n; ++i) y[++p] = i;
14         for (int i = 1; i <= n; ++i)
15         {
16             ++cnt[x[i]];
17             sa[i] > j ? y[++p] = sa[i] - j : 0;
18         }
19         for (int i = 1; i <= m; ++i) cnt[i] += cnt[i - 1];
20         for (int i = n; i; --i) sa[cnt[x[y[i]]]--] = y[i];
21             m = 0;
22         for (int i = 1; i <= n; ++i)
23             y[sa[i]] = (i == 1 || x[sa[i]] != x[sa[i - 1]] || x[sa[i - 1] + j] != x[sa[i] + j]) ?
                   ↪ ++m : m;
24     }
25     for (int i = 1; i <= n; ++i) rank[sa[i]] = i;
26     for (int i = 1, j, k = 0; i <= n; height[rank[i++]] = k)
27         for (k ? --k : 0, j = sa[rank[i] - 1]; s[i + k] == s[j + k]; ++k);
28 }
```

## 4.4   后缀自动机 (ct,lhy)

### 后缀自动机 (lhy)

```
1  struct SAM {
2      SAM *next[26], *fa;
3      int val;
4  } mem[maxn], *last = mem, *tot = mem;
5  void extend(int c)
6  {
7      R SAM *p = last, *np;
8      last = np = ++tot; np -> val = p -> val + 1;
9      for (; p && !p -> next[c]; p = p -> fa) p -> next[c] = np;
10     if (!p) np -> fa = rt[id];
```

```
11      else
12      {
13          SAM *q = p -> next[c];
14          if (q -> val == p -> val + 1) np -> fa = q;
15          else
16          {
17              SAM *nq = ++tot;
18              memcpy(nq -> next, q -> next, sizeof nq -> next);
19              nq -> val = p -> val + 1;
20              nq -> fa = q -> fa;
21              q -> fa = np -> fa = nq;
22              for (; p && p -> next[c] == q; p = p -> fa) p -> next[c] = nq;
23          }
24      }
25  }
```

## 后缀自动机 (ct)

```
1  struct SAM {
2      SAM *next[26], *fa;
3      int val;
4  } mem[maxn], *last = mem, *tot = mem;
5  void extend(int c)
6  {
7      R SAM *p = last, *np;
8      last = np = ++tot; np -> val = p -> val + 1;
9      for (; p && !p -> next[c]; p = p -> fa) p -> next[c] = np;
10     if (!p) np -> fa = rt[id];
11     else
12     {
13         SAM *q = p -> next[c];
14         if (q -> val == p -> val + 1) np -> fa = q;
15         else
16         {
17             SAM *nq = ++tot;
18             memcpy(nq -> next, q -> next, sizeof nq -> next);
19             nq -> val = p -> val + 1;
20             nq -> fa = q -> fa;
21             q -> fa = np -> fa = nq;
22             for (; p && p -> next[c] == q; p = p -> fa) p -> next[c] = nq;
23         }
24     }
25 }
```

## 广义后缀自动机 (ct)

```
1  struct sam {
2      sam *next[26], *fa;
3      int val;
4  } mem[maxn << 1], *tot = mem;
5  inline sam *extend(R sam *p, R int c)
6  {
7      if (p -> next[c])
8      {
9          R sam *q = p -> next[c];
10         if (q -> val == p -> val + 1)
11             return q;
12         else
13         {
```

```
14          R sam *nq = ++tot;
15          memcpy(nq -> next, q -> next, sizeof nq -> next);
16          nq -> val = p -> val + 1;
17          nq -> fa = q -> fa;
18          q -> fa = nq;
19          for ( ; p && p -> next[c] == q; p = p -> fa)
20              p -> next[c] = nq;
21          return nq;
22      }
23  }
24  R sam *np = ++tot;
25  np -> val = p -> val + 1;
26  for ( ; p && !p -> next[c]; p = p -> fa) p -> next[c] = np;
27  if (!p)
28      np -> fa = mem;
29  else
30  {
31      R sam *q = p -> next[c];
32      if (q -> val == p -> val + 1)
33          np -> fa = q;
34      else
35      {
36          R sam *nq = ++tot;
37          memcpy(nq -> next, q -> next, sizeof nq -> next);
38          nq -> val = p -> val + 1;
39          nq -> fa = q -> fa;
40          q -> fa = np -> fa = nq;
41          for ( ; p && p -> next[c] == q; p = p -> fa)
42              p -> next[c] = nq;
43      }
44  }
45  return np;
46  }
```

## 4.5   Manacher (ct)

```
1  char str[maxn];
2  int p1[maxn], p2[maxn], n;
3  void manacher1()
4  {
5      int mx = 0, id;
6      for(int i = 1; i <= n; ++i)
7      {
8          if (mx >= i) p1[i] = dmin(mx - i, p1[(id << 1) - i]);
9          else p1[i] = 1;
10         for (; str[i + p1[i]] == str[i - p1[i]]; ++p1[i]) ;
11         if (p1[i] + i - 1 > mx) id = i, mx = p1[i] + i - 1;
12     }
13 }
14 void manacher2()
15 {
16     int mx = 0, id;
17     for(int i = 1; i <= n; i++)
18     {
19         if (mx >= i) p2[i] = dmin(mx - i, p2[(id << 1) - i]) ;
20         else p2[i] = 0;
21         for (; str[i + p2[i] + 1] == str[i - p2[i]]; ++p2[i]);
22         if (p2[i] + i > mx) id = i, mx = p2[i] + i;
23     }
```

```
24  }
25  int main()
26  {
27      scanf("%s", str + 1);
28      n = strlen(str + 1);
29      str[0] = '#';
30      str[n + 1] = '$';
31      manacher1();
32      manacher2();
33      return 0;
34  }
```

## 4.6   回文树 (ct)

```
1   char str[maxn];
2   int next[maxn][26], fail[maxn], len[maxn], cnt[maxn], last, tot, n;
3   inline int new_node(int l)
4   {
5       len[++tot] = l;
6       return tot;
7   }
8   inline void init()
9   {
10      tot = -1;
11      new_node(0);
12      new_node(-1);
13      str[0] = -1;
14      fail[0] = 1;
15  }
16  inline int get_fail(int x)
17  {
18      while (str[n - len[x] - 1] != str[n]) x = fail[x];
19      return x;
20  }
21  inline void extend(int c)
22  {
23      ++n;
24      int cur = get_fail(last);
25      if (!next[cur][c])
26      {
27          int now = new_node(len[cur] + 2);
28          fail[now] = next[get_fail(fail[cur])][c];
29          next[cur][c] = now;
30      }
31      last = next[cur][c];
32      ++cnt[last];
33  }
34  long long ans;
35  inline void count()
36  {
37      for (int i = tot; i; --i)
38      {
39          cnt[fail[i]] += cnt[i];
40          cmax(ans, 1ll * len[i] * cnt[i]);
41      }
42  }
43  int main()
44  {
45      scanf("%s", str + 1);
```

```
46     init();
47     for (int i = 1; str[i]; ++i)
48         extend(str[i] - 'a');
49     count();
50     printf("%lld\n", ans );
51     return 0;
52 }
```

## 4.7　最小表示法 (ct)

```
 1 int main()
 2 {
 3     int i = 0, j = 1, k = 0;
 4     while (i < n && j < n && k < n)
 5     {
 6         int tmp = a[(i + k) % n] - a[(j + k) % n];
 7         if (!tmp) k++;
 8         else
 9         {
10             if (tmp > 0) i += k + 1;
11             else j += k + 1;
12             if (i == j) ++j;
13             k = 0;
14         }
15     }
16     j = dmin(i, j);
17     for (int i = j; i < n; ++i) printf("%d ", a[i]);
18     for (int i = 0; i < j - 1; ++i) printf("%d ", a[i]);
19     if (j > 0) printf("%d\n", a[j - 1]);
20     return 0;
21 }
```

# Chapter 5

# Data Structure

## 5.1 莫队 (ct)

```
1  int size;
2  struct Query {
3      int l, r, id;
4      inline bool operator < (const Queuy &that) const {return l / size != that.l / size ? l < that.l
          ↪: ((l / size) & 1 ? r < that.r : r > that.r);}
5  } q[maxn];
6  int main()
7  {
8      size = (int) sqrt(n * 1.0);
9      std::sort(q + 1, q + m + 1);
10     int l = 1, r = 0;
11     for (int i = 1; i <= m; ++i)
12     {
13         for (; r < q[i].r; ) add(++r);
14         for (; r > q[i].r; ) del(r--);
15         for (; l < q[i].l; ) del(l++);
16         for (; l > q[i].l; ) add(--l);
17         /*
18             write your code here.
19         */
20     }
21     return 0;
22 }
```

## 5.2 ST 表 (ct)

```
1  int a[maxn], f[20][maxn], n;
2  int Log[maxn];

3  void build()
4  {
5      for (int i = 1; i <= n; ++i) f[0][i] = a[i];

6      int lim = Log[n];
7      for (int j = 1; j <= lim; ++j)
8      {
9          int *fj = f[j], *fj1 = f[j - 1];
10         for (int i = 1; i <= n - (1 << j) + 1; ++i)
11             fj[i] = dmax(fj1[i], fj1[i + (1 << (j - 1))]);
12     }
13 }
```

```
14  int Query(int l, int r)
15  {
16      int k = Log[r - l + 1];
17      return dmax(f[k][l], f[k][r - (1 << k) + 1]);
18  }
19  int main()
20  {
21      scanf("%d", &n);
22      Log[0] = -1;
23      for (int i = 1; i <= n; ++i)
24      {
25          scanf("%d", &a[i]);
26          Log[i] = Log[i >> 1] + 1;
27      }
28      build();
29      int q;
30      scanf("%d", &q);
31      for (; q; --q)
32      {
33          int l, r; scanf("%d%d", &l, &r);
34          printf("%d\n", Query(l, r) );
35      }
36  }
```

## 5.3   带权并查集 (ct)

```
1   struct edge
2   {
3       int a, b, w;
4       inline bool operator < (const edge &that) const {return w > that.w;}
5   } e[maxm];
6   int fa[maxn], f1[maxn], f2[maxn], f1cnt, f2cnt, val[maxn], size[maxn];
7   int main()
8   {
9       int n, m; scanf("%d%d", &n, &m);
10      for (int i = 1; i <= m; ++i)
11          scanf("%d%d%d", &e[i].a, &e[i].b, &e[i].w);
12      for (int i = 1; i <= n; ++i) size[i] = 1;
13      std::sort(e + 1, e + m + 1);
14      for (int i = 1; i <= m; ++i)
15      {
16          int x = e[i].a, y = e[i].b;
17          for ( ; fa[x]; x = fa[x]) ;
18          for ( ; fa[y]; y = fa[y]) ;
19          if (x != y)
20          {
21              if (size[x] < size[y]) std::swap(x, y);
22              size[x] += size[y];
23              val[y] = e[i].w;
24              fa[y] = x;
25          }
26      }
27
28      int q; scanf("%d", &q);
29      for (; q; --q)
30      {
31          int a, b; scanf("%d%d", &a, &b); f1cnt = f2cnt = 0;
32          for (; fa[a]; a = fa[a]) f1[++f1cnt] = a;
33          for (; fa[b]; b = fa[b]) f2[++f2cnt] = b;
```

```
33        if (a != b) {puts("-1"); continue;}
34        while (f1cnt && f2cnt && f1[f1cnt] == f2[f2cnt]) --f1cnt, --f2cnt;
35        int ret = 0x7fffffff;
36        for (; f1cnt; --f1cnt) cmin(ret, val[f1[f1cnt]]);
37        for (; f2cnt; --f2cnt) cmin(ret, val[f2[f2cnt]]);
38        printf("%d\n", ret);
39    }
40    return 0;
41 }
```

## 5.4   可并堆 (ct)

```
1  struct Node {
2      Node *ch[2];
3      ll val; int size;
4      inline void update()
5      {
6          size = ch[0] -> size + ch[1] -> size + 1;
7      }
8  } mem[maxn], *rt[maxn];
9  Node *merge(Node *a, Node *b)
10 {
11     if (a == mem) return b;
12     if (b == mem) return a;
13     if (a -> val < b -> val) std::swap(a, b);
14     // a -> pushdown();
15     std::swap(a -> ch[0], a -> ch[1]);
16     a -> ch[1] = merge(a -> ch[1], b);
17     a -> update();
18     return a;
19 }
```

## 5.5   zkw 线段树 (ct)

0-based

```
1  inline void build()
2  {
3      for (int i = M - 1; i; --i) tr[i] = dmax(tr[i << 1], tr[i << 1 | 1]);
4  }
5  inline void Change(int x, int v)
6  {
7      x += M; tr[x] = v; x >>= 1;
8      for (; x; x >>= 1) tr[x] = dmax(tr[x << 1], tr[x << 1 | 1]);
9  }
10 inline int Query(int s, int t)
11 {
12     int ret = -0x7fffffff;
13     for (s = s + M - 1, t = t + M + 1; s ^ t ^ 1; s >>= 1, t >>= 1)
14     {
15         if (~s & 1) cmax(ret, tr[s ^ 1]);
16         if (t & 1) cmax(ret, tr[t ^ 1]);
17     }
18     return ret;
19 }
20 int main()
21 {
22     int n; scanf("%d", &n);
```

65

```
23    for (M = 1; M < n; M <<= 1) ;
24    for (int i = 0; i < n; ++i)
25        scanf("%d", &tr[i + M]);
26    for (int i = n; i < M; ++i) tr[i + M] = -0x7fffffff;
27    build();
28    int q; scanf("%d", &q);
29    for (; q; --q)
30    {
31        int l, r; scanf("%d%d", &l, &r); --l, --r;
32        printf("%d\n", Query(l, r));
33    }
34    return 0;
35 }
```

# 5.6   李超线段树 (ct)

```
1  int size[maxn], dep[maxn], son[maxn], fa[maxn], top[maxn], dfn[maxn], pos[maxn], timer, rig[maxn];
2  ll dis[maxn];
3  bool vis[maxn];
4  // 树链剖分 begin
5  void dfs1(int x);
6  void dfs2(int x){cmax(rig[top[x]], dfn[x]);}
7  inline int getlca(int a, int b);
8  // 树链剖分 end
9  struct Seg {
10     Seg *ls, *rs;
11     ll min, k, b, vl, vr;
12     // min 表示区间最小值
13     // k 表示区间内 直线标记的斜率
14     // b 表示区间内 直线标记的截距
15     // vl, vr 表示区间内 x 的最小值和最大值
16     inline void update()
17     {
18         min = dmin(ls -> min, rs -> min);
19         k > 0 ? cmin(min, k * vl + b) : cmin(min, k * vr + b);
20     }
21 } ssegg[maxn << 2], *scnt = ssegg, *rt[maxn];
22 void build(int l, int r)
23 {
24     R Seg *o = scnt; o -> k = 0; o -> b = inf;
25     o -> vl = dis[pos[l]]; o -> vr = dis[pos[r]]; o -> min = inf;
26     if (l == r) return ;
27     int mid = l + r >> 1;
28     o -> ls = ++scnt; build(l, mid);
29     o -> rs = ++scnt; build(mid + 1, r);
30     o -> update();
31 }
32 int ql, qr, qk;
33 ll qb;
34 void modify(R Seg *o, int l, int r, int k, ll b)
35 {
36     int mid = l + r >> 1;

37     if (ql <= l && r <= qr)
38     {
39         if (l == r)
40         {
41             cmin(o -> min, k * o -> vl + b);
42             return ;
```

```
43            }
44            ll
45            val = o -> vl * k + b,
46            var = o -> vr * k + b,
47            vbl = o -> vl * o -> k + o -> b,
48            vbr = o -> vr * o -> k + o -> b;
49            if (val <= vbl && var <= vbr)
50            {
51                o -> k = k; o -> b = b;
52                o -> update();
53                return ;
54            }
55            if (val >= vbl && var >= vbr) return ;
56            ll dam = dis[pos[mid]], vam = dam * k + b, vbm = dam * o -> k + o -> b;
57            if (val >= vbl && vam <= vbm)
58            {
59                modify(o -> ls, l, mid, o -> k, o -> b);
60                o -> k = k; o -> b = b;
61            }
62            else if (val <= vbl && vam >= vbm)
63                modify(o -> ls, l, mid, k, b);
64            else
65            {
66                if (vam <= vbm && var >= vbr)
67                {
68                    modify(o -> rs, mid + 1, r, o -> k, o -> b);
69                    o -> k = k; o -> b = b;
70                }
71                else
72                    modify(o -> rs, mid + 1, r, k, b);
73            }
74            o -> update();
75            return ;
76        }
77        if (ql <= mid) modify(o -> ls, l, mid, k, b);
78        if (mid < qr) modify(o -> rs, mid + 1, r, k, b);
79        o -> update();
80 }
81 ll query(R Seg *o, int l, int r)
82 {
83        if (ql <= l && r <= qr) return o -> min;
84        int mid = l + r >> 1; ll ret = inf, tmp;
85        cmin(ret, dis[pos[dmax(ql, l)]] * o -> k + o -> b);
86        cmin(ret, dis[pos[dmin(qr, r)]] * o -> k + o -> b);
87        if (ql <= mid) tmp = query(o -> ls, l, mid), cmin(ret, tmp);
88        if (mid < qr) tmp = query(o -> rs, mid + 1, r), cmin(ret, tmp);
89        return ret;
90 }
91 inline void tr_modify(int x, int f)
92 {
93        while (top[x] != top[f])
94        {
95            ql = dfn[top[x]]; qr = dfn[x];
96            modify(rt[top[x]], ql, rig[top[x]], qk, qb);
97            x = fa[top[x]];
98        }
99        ql = dfn[f]; qr = dfn[x];
100       modify(rt[top[x]], dfn[top[x]], rig[top[x]], qk, qb);
101 }
102 inline ll tr_query(int s, int t)
103 {
```

```
104     ll ret = inf, tmp;
105     while (top[s] != top[t])
106     {
107         if (dep[top[s]] < dep[top[t]])
108         {
109             ql = dfn[top[t]]; qr = dfn[t];
110             tmp = query(rt[top[t]], ql, rig[top[t]]);
111             cmin(ret, tmp);
112             t = fa[top[t]];
113         }
114         else
115         {
116             ql = dfn[top[s]]; qr = dfn[s];
117             tmp = query(rt[top[s]], ql, rig[top[s]]);
118             cmin(ret, tmp);
119             s = fa[top[s]];
120         }
121     }
122     ql = dfn[s]; qr = dfn[t]; ql > qr ? std::swap(ql, qr), 1 : 0;
123     tmp = query(rt[top[s]], dfn[top[s]], rig[top[s]]);
124     cmin(ret, tmp);
125     return ret;
126 }
127 int main()
128 {
129     int n, m; scanf("%d%d", &n, &m);
130     for (int i = 1; i < n; ++i)
131     {
132         int a, b, w; scanf("%d%d%d", &a, &b, &w); link(a, b, w);
133     }
134     dfs1(1); dfs2(1);
135     for (int i = 1; i <= n; ++i)
136         if (top[i] == i)
137         {
138             rt[i] = ++scnt;
139             build(dfn[i], rig[i]);
140         }
141     for (; m; --m)
142     {
143         int opt, s, t, lca; scanf("%d%d%d", &opt, &s, &t);
144         lca = getlca(s, t);
145         if (opt == 1)
146         {
147             int a; ll b; scanf("%d%lld", &a, &b);
148             lca = getlca(s, t);
149             qk = -a; qb = a * dis[s] + b;
150             tr_modify(s, lca);
151             qk = a; qb = a * dis[s] - dis[lca] * 2 * a + b;
152             tr_modify(t, lca);
153         }
154         else
155         {
156             printf("%lld\n", tr_query(s, t));
157         }
158     }
159     return 0;
160 }
```

## 5.7　吉利线段树

## 5.8　二进制分组 (ct)

用线段树维护时间的操作序列，每次操作一个一个往线段树里面插，等到一个线段被插满的时候用归并来维护区间的信息。查询的时候如果一个线段没有被插满就递归下去。定位到一个区间的时候在区间里面归并出来的信息二分。

```
1  int x[maxn], tnum;
2  struct Seg {
3      int l, r, a, b;
4  } p[maxn * 200];
5  int lef[maxm << 2], rig[maxm << 2], pcnt, ta, tb, ql, qr, n, m, k, ans;
6  void update(R int o, R int l, R int r)
7  {
8      lef[o] = pcnt + 1;
9      for (R int i = lef[o << 1], j = lef[o << 1 | 1], head = 1; i <= rig[o << 1] || j <= rig[o << 1
         ↪ | 1]; )
10         if (p[i].r <= p[j].r)
11         {
12             p[++pcnt] = (Seg) {head, p[i].r, 1ll * p[i].a * p[j].a % m, (1ll * p[j].a * p[i].b +
                  ↪ p[j].b) % m};
13             head = p[i].r + 1;
14             p[i].r == p[j].r ? ++j : 0; ++i;
15         }
16         else
17         {
18             p[++pcnt] = (Seg) {head, p[j].r, 1ll * p[i].a * p[j].a % m, (1ll * p[j].a * p[i].b +
                  ↪ p[j].b) % m};
19             head = p[j].r + 1; ++j;
20         }
21     rig[o] = pcnt;
22 }
23 int find(R int o, R int t, R int &s)
24 {
25     R int l = lef[o], r = rig[o];
26     while (l < r)
27     {
28         R int mid = l + r >> 1;
29         if (t <= p[mid].r) r = mid;
30         else l = mid + 1;
31     }
32 //    printf("%d %d t %d s %d %d %d\n", p[l].l, p[l].r, t, s, p[l].a, p[l].b);
33     s = (1ll * s * p[l].a + p[l].b) % m;
34 }
35 void modify(R int o, R int l, R int r, R int t)
36 {
37     if (l == r)
38     {
39         lef[o] = pcnt + 1;
40         ql > 1 ? p[++pcnt] = (Seg) {1, ql - 1, 1, 0}, 1: 0;
41         p[++pcnt] = (Seg) {ql, qr, ta, tb};
42         qr < n ? p[++pcnt] = (Seg) {qr + 1, n, 1, 0}, 1: 0;
43         rig[o] = pcnt;
44         return ;
45     }
46     R int mid = l + r >> 1;
47     if (t <= mid) modify(o << 1, l, mid, t);
48     else modify(o << 1 | 1, mid + 1, r, t);
```

69

```
49        if (t == r) update(o, l, r);
50    }
51    void query(R int o, R int l, R int r)
52    {
53        if (ql <= l && r <= qr)
54        {
55            find(o, k, ans);
56            return ;
57        }
58        R int mid = l + r >> 1;
59        if (ql <= mid) query(o << 1, l, mid);
60        if (mid < qr) query(o << 1 | 1, mid + 1, r);
61    }
62    int main()
63    {
64        R int type; scanf("%d%d%d", &type, &n, &m);
65        for (R int i = 1; i <= n; ++i) scanf("%d", &x[i]);
66        R int Q; scanf("%d", &Q);
67        for (R int QQ = 1; QQ <= Q; ++QQ)
68        {
69            R int opt, l, r; scanf("%d%d%d", &opt, &l, &r);
70            type & 1 ? l ^= ans, r ^= ans : 0;
71            if (opt == 1)
72            {
73                scanf("%d%d", &ta, &tb); ++tnum; ql = l; qr = r;
74                modify(1, 1, Q, tnum);
75            }
76            else
77            {
78                scanf("%d", &k); type & 1 ? k ^= ans : 0; ql = l; qr = r;
79                ans = x[k];
80                query(1, 1, Q);
81                printf("%d\n", ans);
82            }
83        }
84        return 0;
85    }
```

## 5.9　Splay (ct)

### 指针版

```
1    struct Node *null;
2    struct Node {
3        Node *ch[2], *fa;
4        int val; bool rev;
5        inline bool type()
6        {
7            return fa -> ch[1] == this;
8        }
9        inline void pushup()
10       {
11       }
12       inline void pushdown()
13       {
14           if (rev)
15           {
16               ch[0] -> rev ^= 1;
17               ch[1] -> rev ^= 1;
```

```
18              std::swap(ch[0], ch[1]);
19              rev ^= 1;
20          }
21      }
22      inline void rotate()
23      {
24          bool d = type(); Node *f = fa, *gf = f -> fa;
25          (fa = gf, f -> fa != null) ? fa -> ch[f -> type()] = this : 0;
26          (f -> ch[d] = ch[!d]) != null ? ch[!d] -> fa = f : 0;
27          (ch[!d] = f) -> fa = this;
28          f -> pushup();
29      }
30      inline void splay()
31      {
32          for (; fa != null; rotate())
33              if (fa -> fa != null)
34                  (type() == fa -> type() ? fa : this) -> rotate();
35          pushup();
36      }
37  } mem[maxn];
```

## 维修序列

```
1  int fa[maxn], ch[maxn][2], a[maxn], size[maxn], cnt;
2  int sum[maxn], lmx[maxn], rmx[maxn], mx[maxn], v[maxn], id[maxn], root;
3  bool rev[maxn], tag[maxn];
4  inline void update(R int x)
5  {
6      R int ls = ch[x][0], rs = ch[x][1];
7      size[x] = size[ls] + size[rs] + 1;
8      sum[x] = sum[ls] + sum[rs] + v[x];
9      mx[x] = gmax(mx[ls], mx[rs]);
10     cmax(mx[x], lmx[rs] + rmx[ls] + v[x]);
11     lmx[x] = gmax(lmx[ls], sum[ls] + v[x] + lmx[rs]);
12     rmx[x] = gmax(rmx[rs], sum[rs] + v[x] + rmx[ls]);
13 }
14 inline void pushdown(R int x)
15 {
16     R int ls = ch[x][0], rs = ch[x][1];
17     if (tag[x])
18     {
19         rev[x] = tag[x] = 0;
20         if (ls) tag[ls] = 1, v[ls] = v[x], sum[ls] = size[ls] * v[x];
21         if (rs) tag[rs] = 1, v[rs] = v[x], sum[rs] = size[rs] * v[x];
22         if (v[x]>= 0)
23         {
24             if (ls) lmx[ls] = rmx[ls] = mx[ls] = sum[ls];
25             if (rs) lmx[rs] = rmx[rs] = mx[rs] = sum[rs];
26         }
27         else
28         {
29             if (ls) lmx[ls] = rmx[ls] = 0, mx[ls] = v[x];
30             if (rs) lmx[rs] = rmx[rs] = 0, mx[rs] = v[x];
31         }
32     }
33     if (rev[x])
34     {
35         rev[x] ^= 1; rev[ls] ^= 1; rev[rs] ^= 1;
36         swap(lmx[ls], rmx[ls]);swap(lmx[rs], rmx[rs]);
37         swap(ch[ls][0], ch[ls][1]); swap(ch[rs][0], ch[rs][1]);
```

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```
38            }
39      }
40      inline void rotate(R int x)
41      {
42          R int f = fa[x], gf = fa[f], d = ch[f][1] == x;
43          if (f == root) root = x;
44          (ch[f][d] = ch[x][d ^ 1]) > 0 ? fa[ch[f][d]] = f : 0;
45          (fa[x] = gf) > 0 ? ch[gf][ch[gf][1] == f] = x : 0;
46          fa[ch[x][d ^ 1] = f] = x;
47          update(f);
48      }
49      inline void splay(R int x, R int rt)
50      {
51          while (fa[x] != rt)
52          {
53              R int f = fa[x], gf = fa[f];
54              if (gf != rt) rotate((ch[gf][1] == f) ^ (ch[f][1] == x) ? x : f);
55              rotate(x);
56          }
57          update(x);
58      }
59      void build(R int l, R int r, R int rt)
60      {
61          if (l > r) return ;
62          R int mid = l + r >> 1, now = id[mid], last = id[rt];
63          if (l == r)
64          {
65              sum[now] = a[l];
66              size[now] = 1;
67              tag[now] = rev[now] = 0;
68              if (a[l] >= 0) lmx[now] = rmx[now] = mx[now] = a[l];
69              else lmx[now] = rmx[now] = 0, mx[now] = a[l];
70          }
71          else
72          {
73              build(l, mid - 1, mid);
74              build(mid + 1, r, mid);
75          }
76          v[now] = a[mid];
77          fa[now] = last;
78          update(now);
79          ch[last][mid >= rt] = now;
80      }
81      int find(R int x, R int rank)
82      {
83          if (tag[x] || rev[x]) pushdown(x);
84          R int ls = ch[x][0], rs = ch[x][1], lsize = size[ls];
85          if (lsize + 1 == rank) return x;
86          if (lsize >= rank)
87              return find(ls, rank);
88          else
89              return find(rs, rank - lsize - 1);
90      }
91      inline int prepare(R int l, R int tot)
92      {
93          R int x = find(root, l - 1), y = find(root, l + tot);
94          splay(x, 0);
95          splay(y, x);
96          return ch[y][0];
97      }
98      std::queue <int> q;
```

```
99   inline void Insert(R int left, R int tot)
100  {
101      for (R int i = 1; i <= tot; ++i ) a[i] = FastIn();
102      for (R int i = 1; i <= tot; ++i )
103          if (!q.empty()) id[i] = q.front(), q.pop();
104          else id[i] = ++cnt;
105      build(1, tot, 0);
106      R int z = id[(1 + tot) >> 1];
107      R int x = find(root, left), y = find(root, left + 1);
108      splay(x, 0);
109      splay(y, x);
110      fa[z] = y;
111      ch[y][0] = z;
112      update(y);
113      update(x);
114  }
115  void rec(R int x)
116  {
117      if (!x) return ;
118      R int ls = ch[x][0], rs = ch[x][1];
119      rec(ls); rec(rs); q.push(x);
120      fa[x] = ch[x][0] = ch[x][1] = 0;
121      tag[x] = rev[x] = 0;
122  }
123  inline void Delete(R int l, R int tot)
124  {
125      R int x = prepare(l, tot), f = fa[x];
126      rec(x); ch[f][0] = 0;
127      update(f); update(fa[f]);
128  }
129  inline void Makesame(R int l, R int tot, R int val)
130  {
131      R int x = prepare(l, tot), y = fa[x];
132      v[x] = val; tag[x] = 1; sum[x] = size[x] * val;
133      if (val >= 0) lmx[x] = rmx[x] = mx[x] = sum[x];
134      else lmx[x] = rmx[x] = 0, mx[x] = val;
135      update(y); update(fa[y]);
136  }
137  inline void Reverse(R int l, R int tot)
138  {
139      R int x = prepare(l, tot), y = fa[x];
140      if (!tag[x])
141      {
142          rev[x] ^= 1;
143          swap(ch[x][0], ch[x][1]);
144          swap(lmx[x], rmx[x]);
145          update(y); update(fa[y]);
146      }
147  }
148  inline void Query(R int l, R int tot)
149  {
150      R int x = prepare(l, tot);
151      printf("%d\n",sum[x] );
152  }
153  #define inf ((1 << 30))
154  int main()
155  {
156      R int n = FastIn(),  m = FastIn(), l, tot, val;
157      R char op, op2;
158      mx[0] = a[1] = a[n + 2] = -inf;
159      for (R int i = 2; i <= n + 1; i++ )
```

```
160        {
161            a[i] = FastIn();
162        }
163        for (R int i = 1; i <= n + 2; ++i) id[i] = i;
164        n += 2; cnt = n; root = (n + 1) >> 1;
165        build(1, n, 0);
166        for (R int i = 1; i <= m; i++ )
167        {
168            op = getc();
169            while (op < 'A' || op > 'Z') op = getc();
170            getc(); op2 = getc();getc();getc();getc();getc();
171            if (op == 'M' && op2 == 'X')
172            {
173                printf("%d\n",mx[root] );
174            }
175            else
176            {
177                l = FastIn() + 1; tot = FastIn();
178                if (op == 'I') Insert(l, tot);
179                if (op == 'D') Delete(l, tot);
180                if (op == 'M') val = FastIn(), Makesame(l, tot, val);
181                if (op == 'R')
182                    Reverse(l, tot);
183                if (op == 'G')
184                    Query(l, tot);
185            }
186        }
187        return 0;
188 }
```

## 5.10   Treap (ct)

```
1  struct Treap {
2      Treap *ls, *rs;
3      int size;
4      bool rev;
5      inline void update()
6      {
7          size = ls -> size + rs -> size + 1;
8      }
9      inline void set_rev()
10     {
11         rev ^= 1;
12         std::swap(ls, rs);
13     }
14     inline void pushdown()
15     {
16         if (rev)
17         {
18             ls -> set_rev();
19             rs -> set_rev();
20             rev = 0;
21         }
22     }
23 } mem[maxn], *root, *null = mem;
24 struct Pair {
25     Treap *fir, *sec;
26 };
27 Treap *build(R int l, R int r)
```

```
28  {
29      if (l > r) return null;
30      R int mid = l + r >> 1;
31      R Treap *now = mem + mid;
32      now -> rev = 0;
33      now -> ls = build(l, mid - 1);
34      now -> rs = build(mid + 1, r);
35      now -> update();

36      return now;
37  }
38  inline Treap *Find_kth(R Treap *now, R int k)
39  {
40      if (!k) return mem;
41      if (now -> ls -> size >= k) return Find_kth(now -> ls, k);
42      else if (now -> ls -> size + 1 == k) return now;
43      else return Find_kth(now -> rs, k - now -> ls -> size - 1);
44  }
45  Treap *merge(R Treap *a, R Treap *b)
46  {
47      if (a == null) return b;
48      if (b == null) return a;
49      if (rand() % (a -> size + b -> size) < a -> size)
50      {
51          a -> pushdown();
52          a -> rs = merge(a -> rs, b);
53          a -> update();
54          return a;
55      }
56      else
57      {
58          b -> pushdown();
59          b -> ls = merge(a, b -> ls);
60          b -> update();
61          return b;
62      }
63  }
64  Pair split(R Treap *now, R int k)
65  {
66      if (now == null) return (Pair) {null, null};
67      R Pair t = (Pair) {null, null};
68      now -> pushdown();
69      if (k <= now -> ls -> size)
70      {
71          t = split(now -> ls, k);
72          now -> ls = t.sec;
73          now -> update();
74          t.sec = now;
75      }
76      else
77      {
78          t = split(now -> rs, k - now -> ls -> size - 1);
79          now -> rs = t.fir;
80          now -> update();
81          t.fir = now;
82      }
83      return t;
84  }
85  inline void set_rev(int l, int r)
86  {
87      R Pair x = split(root, l - 1);
```

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```
88        R Pair y = split(x.sec, r - l + 1);
89        y.fir -> set_rev();
90        root = merge(x.fir, merge(y.fir, y.sec));
91    }
```

## 5.11  可持久化平衡树 (ct)

```
1   char str[maxn];
2   struct Treap
3   {
4       Treap *ls, *rs;
5       char data; int size;
6       inline void update()
7       {
8           size = ls -> size + rs -> size + 1;
9       }
10  } *root[maxn], mem[maxcnt], *tot = mem, *last = mem, *null = mem;
11  inline Treap* new_node(char ch)
12  {
13      *++tot = (Treap) {null, null, ch, 1};
14      return tot;
15  }
16  struct Pair
17  {
18      Treap *fir, *sec;
19  };
20  inline Treap *copy(Treap *x)
21  {
22      if (x == null) return null;
23      if(x > last) return x;
24      *++tot = *x;
25      return tot;
26  }
27  Pair Split(Treap *x, int k)
28  {
29      if (x == null) return (Pair) {null, null};
30      Pair y;
31      Treap *nw = copy(x);
32      if (nw -> ls -> size >= k)
33      {
34          y = Split(nw -> ls, k);
35          nw -> ls = y.sec;
36          nw -> update();
37          y.sec = nw;
38      }
39      else
40      {
41          y = Split(nw -> rs, k - nw -> ls -> size - 1);
42          nw -> rs = y.fir;
43          nw -> update();
44          y.fir = nw;
45      }
46      return y;
47  }
48  Treap *Merge(Treap *a, Treap *b)
49  {
50      if (a == null) return b;
51      if (b == null) return a;
52      Treap *nw;
```

```
53      if (rand() % (a -> size + b -> size) < a -> size)
54      {
55          nw = copy(a);
56          nw -> rs = Merge(nw -> rs, b);
57      }
58      else
59      {
60          nw = copy(b);
61          nw -> ls = Merge(a, nw -> ls);
62      }
63      nw -> update();
64      return nw;
65  }
66  Treap *Build(int l, int r)
67  {
68      if (l > r) return null;
69      R int mid = l + r >> 1;
70      Treap *nw = new_node(str[mid]);
71      nw -> ls = Build(l, mid - 1);
72      nw -> rs = Build(mid + 1, r);
73      nw -> update();
74      return nw;
75  }
76  int now;
77  inline void Insert(int k, char ch)
78  {
79      Pair x = Split(root[now], k);
80      Treap *nw = new_node(ch);
81      root[++now] = Merge(Merge(x.fir, nw), x.sec);
82  }
83  inline void Del(int l, int r)
84  {
85      Pair x = Split(root[now], l - 1);
86      Pair y = Split(x.sec, r - l + 1);
87      root[++now] = Merge(x.fir, y.sec);
88  }
89  inline void Copy(int l, int r, int ll)
90  {
91      Pair x = Split(root[now], l - 1);
92      Pair y = Split(x.sec, r - l + 1);
93      Pair z = Split(root[now], ll);
94      Treap *ans = y.fir;
95      root[++now] = Merge(Merge(z.fir, ans), z.sec);
96  }
97  void Print(Treap *x, int l, int r)
98  {
99      if (!x) return ;
100     if (l > r) return;
101     R int mid = x -> ls -> size + 1;
102     if (r < mid)
103     {
104         Print(x -> ls, l, r);
105         return ;
106     }
107     if (l > mid)
108     {
109         Print(x -> rs, l - mid, r - mid);
110         return ;
111     }
112     Print(x -> ls, l, mid - 1);
113     printf("%c", x -> data );
```

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```
114         Print(x -> rs, 1, r - mid);
115 }
116 void Printtree(Treap *x)
117 {
118     if (!x) return;
119     Printtree(x -> ls);
120     printf("%c", x -> data );
121     Printtree(x -> rs);
122 }
123 int main()
124 {
125     srand(time(0) + clock());
126     null -> ls = null -> rs = null; null -> size = 0; null -> data = 0;
127     int n = F();
128     gets(str + 1);
129     int len = strlen(str + 1);
130     root[0] = Build(1, len);
131     while (1)
132     {
133         last = tot;
134         R char opt = getc();
135         while (opt < 'A' || opt > 'Z')
136         {
137             if (opt == EOF) return 0;
138             opt = getc();
139         }
140         if (opt == 'I')
141         {
142             R int x = F();
143             R char ch = getc();
144             Insert(x, ch);
145         }
146         else if (opt == 'D')
147         {
148             R int l = F(), r = F();
149             Del(l, r);
150         }
151         else if (opt == 'C')
152         {
153             R int x = F(), y = F(), z = F();
154             Copy(x, y, z);
155         }
156         else if (opt == 'P')
157         {
158             R int x = F(), y = F(), z = F();
159             Print(root[now - x], y, z);
160             puts("");
161         }
162     }
163     return 0;
164 }
```

## 5.12   CDQ 分治 (ct)

```
1 struct event
2 {
3     int x, y, id, opt, ans;
4 } t[maxn], q[maxn];
5 void cdq(int left, int right)
```

```cpp
6  {
7      if (left == right) return ;
8      R int mid = left + right >> 1;
9      cdq(left, mid);
10     cdq(mid + 1, right);
11     //分成若干个子问题
12     ++now;
13     for (int i = left, j = mid + 1; j <= right; ++j)
14     {
15         for (; i <= mid && q[i].x <= q[j].x; ++i)
16             if (!q[i].opt)
17                 add(q[i].y, q[i].ans);
18         //考虑前面的修改操作对后面的询问的影响
19         if (q[j].opt)
20             q[j].ans += query(q[j].y);
21     }
22     R int i, j, k = 0;
23     //以下相当于归并排序
24     for (i = left, j = mid + 1; i <= mid && j <= right; )
25     {
26         if (q[i].x <= q[j].x)
27             t[k++] = q[i++];
28         else
29             t[k++] = q[j++];
30     }
31     for (; i <= mid; )
32         t[k++] = q[i++];
33     for (; j <= right; )
34         t[k++] = q[j++];
35     for (int i = 0; i < k; ++i)
36         q[left + i] = t[i];
37 }
```

## 5.13   Bitset (ct)

```cpp
1  namespace Game {
2  #define maxn 300010
3  #define maxs 30010
4  uint b1[32][maxs], b2[32][maxs];
5  int popcnt[256];
6  inline void set(R uint *s, R int pos)
7  {
8      s[pos >> 5] |= 1u << (pos & 31);
9  }
10 inline int popcount(R uint x)
11 {
12     return popcnt[x >> 24 & 255]
13         + popcnt[x >> 16 & 255]
14         + popcnt[x >> 8  & 255]
15         + popcnt[x       & 255];
16 }
17 void main() {
18     int n, q;
19     scanf("%d%d", &n, &q);
20     char *s1 = new char[n + 1];
21     char *s2 = new char[n + 1];
22     scanf("%s%s", s1, s2);
```

```
23      uint *anss = new uint[q];

24      for (R int i = 1; i < 256; ++i) popcnt[i] = popcnt[i >> 1] + (i & 1);

25      #define modify(x, _p)\
26      {\
27          for (R int j = 0; j < 32 && j <= _p; ++j)\
28              set(b##x[j], _p - j);\
29      }
30      for (R int i = 0; i < n; ++i)
31          if (s1[i] == '0') modify(1, 3 * i)
32          else if (s1[i] == '1') modify(1, 3 * i + 1)
33          else modify(1, 3 * i + 2)

34      for (R int i = 0; i < n; ++i)
35          if (s2[i] == '1') modify(2, 3 * i)
36          else if (s2[i] == '2') modify(2, 3 * i + 1)
37          else modify(2, 3 * i + 2)

38      for (int Q = 0; Q < q; ++Q) {
39          R int x, y, l;
40          scanf("%d%d%d", &x, &y, &l); x *= 3; y *= 3; l *= 3;
41          uint *f1 = b1[x & 31], *f2 = b2[y & 31], ans = 0;
42          R int i = x >> 5, j = y >> 5, p, lim;
43          for (p = 0, lim = l >> 5; p + 8 < lim; p += 8, i += 8, j += 8)
44          {
45              ans += popcount(f1[i + 0] & f2[j + 0]);
46              ans += popcount(f1[i + 1] & f2[j + 1]);
47              ans += popcount(f1[i + 2] & f2[j + 2]);
48              ans += popcount(f1[i + 3] & f2[j + 3]);
49              ans += popcount(f1[i + 4] & f2[j + 4]);
50              ans += popcount(f1[i + 5] & f2[j + 5]);
51              ans += popcount(f1[i + 6] & f2[j + 6]);
52              ans += popcount(f1[i + 7] & f2[j + 7]);
53          }
54          for (; p < lim; ++p, ++i, ++j) ans += popcount(f1[i] & f2[j]);
55          R uint S = (1u << (l & 31)) - 1;
56          ans += popcount(f1[i] & f2[j] & S);
57          anss[Q] = ans;
58      }

59      output_arr(anss, q * sizeof(uint));
60  }
61  }
```

## 5.14  斜率优化 (ct)

对于斜截式 $y = kx + b$，如果把 $k_i$ 看成斜率，那 dp 时需要最小化截距，把斜截式转化为 $b_i = -k_i x_j + y_j$，就可以把可以转移到这个状态的点看作是二维平面上的点 $(-x_j, y_j)$，问题转化为了在平面上找一个点使得斜率为 $k_i$ 的直线的截距最小。这样的点一定在凸包上，这样的点在凸包上和前一个点的斜率 $\leq k_i$，和后面一个点的斜率 $\geq k_i$。这样就可以在凸包上二分来加速转移。当点的横坐标 $x_i$ 和斜率 $k_i$ 都是单调的，还可以用单调队列来维护凸包。

### 单调队列

```
1   int a[maxn], n, l;
2   ll sum[maxn], f[maxn];
3   inline ll sqr(ll x) {return x * x;}
4   #define y(_i) (f[_i] + sqr(sum[_i] + l))
```

```cpp
#define x(_i) (2 * sum[_i])
inline double slope(int i, int j)
{
    return (y(i) - y(j)) / (1.0 * (x(i) - x(j)));
}
int q[maxn];
int main()
{
    n = F(), l = F() + 1;
    for (int i = 1; i <= n; ++i) a[i] = F(), sum[i] = sum[i - 1] + a[i];
    for (int i = 1; i <= n; ++i) sum[i] += i;
    f[0] = 0;
/*
    memset(f, 63, sizeof (f));
    for (int i = 1; i <= n; ++i)
    {
        int pos;
        for (int j = 0; j < i; ++j)
        {
            long long tmp = f[j] + sqr(sum[i] - sum[j] - l);
            f[i] > tmp ? f[i] = tmp, pos = j : 0;
        }
    }
*/
    int h = 1, t = 1;
    q[h] = 0;
    for (int i = 1; i <= n; ++i)
    {
        while (h < t && slope(q[h], q[h + 1]) <= sum[i]) ++h;
        f[i] = f[q[h]] + sqr(sum[i] - sum[q[h]] - l);
        while (h < t && slope(q[t - 1], i) < slope(q[t - 1], q[t])) --t;
        q[++t] = i;
    }
    printf("%lld\n", f[n] );
    return 0;
}
```

# 线段树

```cpp
// NOI 2014 购票
int dep[maxn], fa[maxn], son[maxn], dfn[maxn], timer, pos[maxn], size[maxn], n, top[maxn];
ll d[maxn], p[maxn], q[maxn], l[maxn], f[maxn];
int stcnt;
void dfs1(int x);
void dfs2(int x);
#define P pair<ll, ll>
#define mkp make_pair
#define x first
#define y second
#define inf ~0ULL >> 2
inline double slope(const P &a, const P &b)
{
    return (b.y - a.y) / (double) (b.x - a.x);
}
struct Seg
{
    vector<P> v;
    inline void add(const P &that)
    {
        int top = v.size();
```

```
22          P *v = this -> v.data() - 1;
23          while (top > 1 && slope(v[top - 1], v[top]) > slope(v[top], that)) --top;
24          this -> v.erase(this -> v.begin() + top, this -> v.end());
25          this -> v.push_back(that);
26      }
27      inline ll query(ll k)
28      {
29          if (v.empty()) return inf;
30          int l = 0, r = v.size() - 1;
31          while (l < r)
32          {
33              int mid = l + r >> 1;
34              if (slope(v[mid], v[mid + 1]) > k) r = mid;
35              else l = mid + 1;
36          }
37          cmin(l, v.size() - 1);
38          return v[l].y - v[l].x * k;
39      }
40  } tr[1 << 19];
41  void Change(int o, int l, int r, int x, P val)
42  {
43      tr[o].add(val);
44      if (l == r) return;
45      int mid = l + r >> 1;
46      if (x <= mid) Change(o << 1, l, mid, x, val);
47      else Change(o << 1 | 1, mid + 1, r, x, val);
48  }
49  int ql, qr, now, tmp;
50  ll len;
51  inline ll Query(int o, int l, int r)
52  {
53      if (ql <= l && r <= qr && d[tmp] - d[pos[r]] > len) return inf;
54      if (ql <= l && r <= qr && d[tmp] - d[pos[l]] <= len)
55          return tr[o].query(p[now]);
56      ll ret = inf, temp;
57      int mid = l + r >> 1;
58      if (ql <= mid) temp = Query(o << 1, l, mid), cmin(ret, temp);
59      if (mid < qr) temp = Query(o << 1 | 1, mid + 1, r), cmin(ret, temp);
60      return ret;
61  }
62  inline ll calc()
63  {
64      ll ret = inf;
65      ll lx = l[now];
66      tmp = now;
67      while (lx >= 0 && tmp)
68      {
69          len = lx;
70          ql = dfn[top[tmp]];
71          qr = dfn[tmp];
72          ll g = Query(1, 1, n);
73          cmin(ret, g);
74          lx -= d[tmp] - d[fa[top[tmp]]];
75          tmp = fa[top[tmp]];
76      }
77      return ret;
78  }
79  int main()
80  {
81      n = F(); int t = F();
82      for (int i = 2; i <= n; ++i)
```

```
83      {
84          fa[i] = F(); ll dis = F(); p[i] = F(), q[i] = F(), l[i] = F();
85          link(fa[i], i); d[i] = d[fa[i]] + dis;
86      }
87      dfs1(1);
88      dfs2(1);
89      Change(1, 1, n, 1, mkp(0, 0));
90      for (now = 2; now <= n; ++now)
91      {
92          f[now] = calc() + q[now] + d[now] * p[now];
93          Change(1, 1, n, dfn[now], mkp(d[now], f[now]));
94          printf("%lld\n", f[now] );
95      }
96      return 0;
97  }
```

## 5.15  树分块

树分块套分块：给定一棵有点权的树，每次询问链上不同点权个数

```
1   int col[maxn], hash[maxn], hcnt, n, m;
2   int near[maxn];
3   bool vis[maxn];
4   int mark[maxn], mcnt, tcnt[maxn], tans;
5   int pre[256][maxn];
6   struct Block {
7       int cnt[256];
8   } mem[maxn], *tot = mem;
9   inline Block *nw(Block *last, int v)
10  {
11      Block *ret = ++tot;
12      memcpy(ret -> cnt, last -> cnt, sizeof (ret -> cnt));
13      ++ret -> cnt[v & 255];
14      return ret;
15  }
16  struct Arr {
17      Block *b[256];
18      inline int v(int c) {return b[c >> 8] -> cnt[c & 255];}
19  } c[maxn];
20  inline Arr cp(Arr last, int v)
21  {
22      Arr ret;
23      memcpy(ret.b, last.b, sizeof (ret.b));
24      ret.b[v >> 8] = nw(last.b[v >> 8], v);
25      return ret;
26  }
27  void bfs()
28  {
29      int head = 0, tail = 1; q[1] = 1;
30      while (head < tail)
31      {
32          int now = q[++head]; size[now] = 1; vis[now] = 1; dep[now] = dep[fa[now]] + 1;
33          for (Edge *iter = last[now]; iter; iter = iter -> next)
34              if (!vis[iter -> to])
35                  fa[q[++tail] = iter -> to] = now;
36      }
37      for (int i = n; i; --i)
38      {
39          int now = q[i];
40          size[fa[now]] += size[now];
```

```
41              size[son[fa[now]]] < size[now] ? son[fa[now]] = now : 0;
42          }
43          for (int i = 0; i < 256; ++i) c[0].b[i] = mem;
44          for (int i = 1; i <= n; ++i)
45          {
46              int now = q[i];
47              c[now] = cp(c[fa[now]], col[now]);
48              top[now] = son[fa[now]] == now ? top[fa[now]] : now;
49          }
50  }
51  inline int getlca(int a, int b) ;
52  void dfs_init(int x)
53  {
54      vis[x] = 1; ++tcnt[col[x]] == 1 ? ++tans : 0;
55      pre[mcnt][x] = tans;
56      for (Edge *iter = last[x]; iter; iter = iter -> next)
57          if (!vis[iter -> to]) dfs_init(iter -> to);
58      --tcnt[col[x]] == 0 ? --tans : 0;
59  }
60  int jp[maxn];
61  int main()
62  {
63      scanf("%d%d", &n, &m);
64      for (int i = 1; i <= n; ++i) scanf("%d", &col[i]), hash[++hcnt] = col[i];
65      std::sort(hash + 1, hash + hcnt + 1);
66      hcnt = std::unique(hash + 1, hash + hcnt + 1) - hash - 1;
67      for (int i = 1; i <= n; ++i) col[i] = std::lower_bound(hash + 1, hash + hcnt + 1, col[i]) -
        ↪ hash;
68      for (int i = 1; i < n; ++i)
69      {
70          int a, b; scanf("%d%d", &a, &b); link(a, b);
71      }
72      bfs();
73      int D = sqrt(n);
74      for (int i = 1; i <= n; ++i)
75          if (dep[i] % D == 0 && size[i] >= D)
76          {
77              memset(vis, 0, n + 1);
78              mark[i] = ++mcnt;
79              dfs_init(i);
80          }
81      for (int i = 1; i <= n; ++i) near[q[i]] = mark[q[i]] ? q[i] : near[fa[q[i]]];
82      int ans = 0;
83      memset(vis, 0, n + 1);
84      for (; m; --m)
85      {
86          int x, y; scanf("%d%d", &x, &y);
87          x ^= ans; ans = 0;
88          int lca = getlca(x, y);
89          if (dep[near[x]] < dep[lca]) std::swap(x, y);
90          if (dep[near[x]] >= dep[lca])
91          {
92              Arr *_a = c + near[x];
93              Arr *_b = c + y;
94              Arr *_c = c + lca;
95              Arr *_d = c + fa[lca];
96              for (; !mark[x]; x = fa[x])
97                  if (_a -> v(col[x]) + _b -> v(col[x]) == _c -> v(col[x]) + _d -> v(col[x]) &&
                    ↪ !vis[col[x]])
98                      vis[jp[++ans] = col[x]] = 1;
99              for (int i = 1; i <= ans; ++i) vis[jp[i]] = 0;
```

```
100            ans += pre[mark[near[x]]][y];
101        }
102        else
103        {
104            for (; x != lca; x = fa[x]) !vis[col[x]] ? vis[jp[++ans] = col[x]] = 1 : 0;
105            for (; y != lca; y = fa[y]) !vis[col[y]] ? vis[jp[++ans] = col[y]] = 1 : 0;
106            !vis[col[lca]] ? vis[jp[++ans] = col[lca]] = 1 : 0;
107            for (int i = 1; i <= ans; ++i) vis[jp[i]] = 0;
108        }
109        printf("%d\n", ans);
110    }
111    return 0;
112 }
```

## 5.16   DLX

# Chapter 6

# Others

## 6.1  vimrc (gy)

```
se et ts=4 sw=4 sts=4 nu sc sm lbr is hls mouse=a
sy on
ino <tab> <c-n>
ino <s-tab> <tab>
au winnew * winc L

nm <f6> ggVG"+y
nm <f7> :w<cr>:!rm ##<cr>:make<cr>
nm <f8> :!@@<cr>
nm <f9> :!@@ < in<cr>
nm <s-f9> :!(time @@ < in &> out) &>> out<cr>:sp out<cr>

au filetype cpp cm @@ ./a.out | cm ## a.out | se cin fdm=syntax mp=g++\ %\ -std=c++11\ -Wall\
    ↪ -Wextra\ -Wconversion\ -O2

map <c-p> :ha<cr>
se pheader=%n\ %f

au filetype java cm @@ java %< | cm ## %<.class | se cin fdm=syntax mp=javac\ %
au filetype python cm @@ python % | se si fdm=indent
au bufenter *.kt setf kotlin
au filetype kotlin cm @@ kotlin _%<Kt | cm ## _%<Kt.class | se si mp=kotlinc\ %
```

## 6.2  STL 释放内存 (Durandal)

```
template <typename T>
__inline void clear(T &container) {
    container.clear();
    T(container).swap(container);
}
```

## 6.3  开栈 (Durandal)

```
register char *_sp __asm__("rsp");
int main() {
    const int size = 400 << 20; // 400 MB
    static char *sys, *mine(new char[size] + size - 4096);
    sys = _sp; _sp = mine;
    _main(); // main method
```

```
7       _sp = sys;
8       return 0;
9  }
```

## 6.4   O3 (gy)

```
1  __attribute__((optimize("-O3"))) int main() { return 0; }
```

## 6.5   Java Template (gy)

```java
1  import java.io.*;
2  import java.math.*;
3  import java.util.*;
4
5  public class Template {
6      // Input
7      private static BufferedReader reader;
8      private static StringTokenizer tokenizer;
9      private static String next() {
10         try {
11             while (tokenizer == null || !tokenizer.hasMoreTokens())
12                 tokenizer = new StringTokenizer(reader.readLine());
13         } catch (IOException e) {
14             // do nothing
15         }
16         return tokenizer.nextToken();
17     }
18     private static int nextInt() {
19         return Integer.parseInt(next());
20     }
21     private static double nextDouble() {
22         return Double.parseDouble(next());
23     }
24     private static BigInteger nextBigInteger() {
25         return new BigInteger(next());
26     }
27
28     public static void main(String[] args) {
29         reader = new BufferedReader(new InputStreamReader(System.in));
30         Scanner scanner = new Scanner(System.in);
31         while (scanner.hasNext())
32             scanner.next();
33     }
34
35     // BigInteger & BigDecimal
36     private static void bigDecimal() {
37         BigDecimal a = BigDecimal.valueOf(1.0);
38         BigDecimal b = a.setScale(50, RoundingMode.HALF_EVEN);
39         BigDecimal c = b.abs();
40         // if scale omitted, b.scale is used
41         BigDecimal d = c.divide(b, 50, RoundingMode.HALF_EVEN);
42         // since Java 9
43         BigDecimal e = d.sqrt(new MathContext(50, RoundingMode.HALF_EVEN));
44         BigDecimal x = new BigDecimal(BigInteger.ZERO);
45         BigInteger y = BigDecimal.ZERO.toBigInteger(); // RoundingMode.DOWN
46         y = BigDecimal.ZERO.setScale(0, RoundingMode.HALF_EVEN).unscaledValue();
47     }
48     // sqrt for Java 8
```

```
46      // can solve scale=100 for 10000 times in about 1 second
47      private static BigDecimal sqrt(BigDecimal a, int scale) {
48          if (a.compareTo(BigDecimal.ZERO) < 0)
49              return BigDecimal.ZERO.setScale(scale, RoundingMode.HALF_EVEN);
50          int length = a.precision() - a.scale();
51          BigDecimal ret = new BigDecimal(BigInteger.ONE, -length / 2);
52          for (int i = 1; i <= Integer.highestOneBit(scale) + 10; i++)
53              ret = ret.add(a.divide(ret, scale,
                 ↪ RoundingMode.HALF_EVEN)).divide(BigDecimal.valueOf(2), scale,
                 ↪ RoundingMode.HALF_EVEN);
54          return ret;
55      }
56      // can solve a=2^10000 for 100000 times in about 1 second
57      private static BigInteger sqrt(BigInteger a) {
58          int length = a.bitLength() - 1;
59          BigInteger l = BigInteger.ZERO.setBit(length / 2), r = BigInteger.ZERO.setBit(length / 2);
60          while (!l.equals(r)) {
61              BigInteger m = l.add(r).shiftRight(1);
62              if (m.multiply(m).compareTo(a) < 0)
63                  l = m.add(BigInteger.ONE);
64              else
65                  r = m;
66          }
67          return l;
68      }
69      
70      // Collections
71      private static void arrayList() {
72          List<Integer> list = new ArrayList<>();
73          // Generic array is banned
74          List[] lists = new List[100];
75          lists[0] = new ArrayList<Integer>();
76          // for List<Integer>, remove(Integer) stands for element, while remove(int) stands for
              ↪ index
77          list.remove(list.get(1));
78          list.remove(list.size() - 1);
79          list.clear();
80          Queue<Integer> queue = new LinkedList<>();
81          // return the value without popping
82          queue.peek();
83          // pop and return the value
84          queue.poll();
85          Queue<Integer> priorityQueue = new PriorityQueue<>();
86          Deque<Integer> deque = new ArrayDeque<>();
87          deque.peekFirst();
88          deque.peekLast();
89          deque.pollFirst();
90          TreeSet<Integer> set = new TreeSet<>();
91          TreeSet<Integer> anotherSet = new TreeSet<>(Comparator.reverseOrder());
92          set.ceiling(1);
93          set.floor(1);
94          set.lower(1);
95          set.higher(1);
96          set.contains(1);
97          HashSet<Integer> hashSet = new HashSet<>();
98          HashMap<String, Integer> map = new HashMap<>();
99          map.put("", 1);
100         map.get("");
101         map.forEach((string, integer) -> System.out.println(string + integer));
102         TreeMap<String, Integer> treeMap = new TreeMap<>();
            Arrays.sort(new int[10]);
```

```
103        Arrays.sort(new Integer[10], (a, b) -> {
104            if (a.equals(b)) return 0;
105            if (a > b) return -1;
106            return 1;
107        });
108        Arrays.sort(new Integer[10], Comparator.comparingInt((a) -> (int) a).reversed());
109        long a = 1_000_000_000_000_000_000L;
110        int b = Integer.MAX_VALUE;
111        int c = 'a';
112    }
113 }
```

## 6.6  Big Fraction (gy)

```
1  fun gcd(a: Long, b: Long): Long = if (b == 0L) a else gcd(b, a % b)
2
3  class Fraction(val a: BigInteger, val b: BigInteger) {
4      constructor(a: Long, b: Long) : this(BigInteger.valueOf(a / gcd(a, b)), BigInteger.valueOf(b /
        ↪ gcd(a, b)))
5
6      operator fun plus(o: Fraction): Fraction {
7          var gcd = b.gcd(o.b)
8          val tempProduct = (b / gcd) * (o.b / gcd)
9          var ansA = a * (o.b / gcd) + o.a * (b / gcd)
10         val gcd2 = ansA.gcd(gcd)
11         ansA /= gcd2
12         gcd /= gcd2
13         return Fraction(ansA, gcd * tempProduct)
14     }
15
16     operator fun minus(o: Fraction): Fraction {
17         var gcd = b.gcd(o.b)
18         val tempProduct = (b / gcd) * (o.b / gcd)
19         var ansA = a * (o.b / gcd) - o.a * (b / gcd)
20         val gcd2 = ansA.gcd(gcd)
21         ansA /= gcd2
22         gcd /= gcd2
23         return Fraction(ansA, gcd * tempProduct)
24     }
25
26     operator fun times(o: Fraction): Fraction {
27         val gcd1 = a.gcd(o.b)
28         val gcd2 = b.gcd(o.a)
29         return Fraction((a / gcd1) * (o.a / gcd2), (b / gcd2) * (o.b / gcd1))
30     }
31 }
```

## 6.7  模拟退火 (ct)

```
1  db ans_x, fans;
2  inline double rand01() {return rand() / 2147483647.0;}
3  inline double randp() {return (rand() & 1 ? 1 : -1) * rand01();}
4  inline double f(double x)
5  {
6      /*
7          write your function here.
8      */
9      if (maxx < fans) {fans = maxx; ans_x = x;}
```

89

```
10        return maxx;
11  }
12  int main()
13  {
14      srand(time(NULL) + clock());
15      db x = 0, fnow = f(x);
16      fans = 1e30;
17      for (db T = 1e4; T > 1e-4; T *= 0.997)
18      {
19          db nx = x + randp() * T, fnext = f(nx);
20          db delta = fnext - fnow;
21          if (delta < 1e-9 || exp(-delta / T) > rand01())
22          {
23              x = nx;
24              fnow = fnext;
25          }
26      }
27      return 0;
28  }
```

## 6.8    三分 (ct)

```
1   inline db cubic_search()
2   {
3       double l = -1e4, r = 1e4;
4       for (int i = 1; i <= 100; ++i)
5       {
6           double ll = (l + r) * 0.5;
7           double rr = (ll + r) * 0.5;
8           if (check(ll) < check(rr)) r = rr;
9           else l = ll;
10      }
11      return (l + r) * 0.5;
12  }
```

## 6.9    Zeller Congruence (gy)

```
1   int day_in_week(int year, int month, int day) {
2       if (month == 1 || month == 2)
3           month += 12, year--;
4       int c = year / 100, y = year % 100, m = month, d = day;
5       int ret = (y + y / 4 + c / 4 + 5 * c + 13 * (m + 1) / 5 + d + 6) % 7;
6       return ret >= 0 ? ret : ret + 7;
7   }
```

## 6.10    博弈论模型 (gy)

- Wythoff's game
  给定两堆石子，每次可以从任意一堆中取至少一个石子，或从两堆中取相同的至少一个石子，取走最后石子的胜
  先手胜当且仅当石子数满足：
  $\lfloor (b - a) \times \phi \rfloor = a, (a \le b, \phi = \frac{\sqrt{5}+1}{2})$
  先手胜对应的石子数构成两个序列：
  Lower Wythoff sequence: $a_n = \lfloor n \times \phi \rfloor$
  Upper Wythoff sequence: $b_n = \lfloor n \times \phi^2 \rfloor$

- Fibonacci nim
  给定一堆石子，第一次可以取至少一个、少于石子总数数量的石子，之后每次可以取至少一个、不超过上次取石子数量两倍的石子，取走最后石子的胜
  先手胜当且仅当石子数为斐波那契数

## 6.11    积分表 (integral-table.com)

$$\int x^n dx = \frac{1}{n+1} x^{n+1}, \ n \neq -1$$

$$\int \frac{1}{x} dx = \ln |x|$$

$$\int u\, dv = uv - \int v\, du$$

$$\int \frac{1}{ax+b} dx = \frac{1}{a} \ln |ax+b|$$

$$\int \frac{1}{(x+a)^2} dx = -\frac{1}{x+a}$$

$$\int (x+a)^n dx = \frac{(x+a)^{n+1}}{n+1}, \ n \neq -1$$

$$\int x(x+a)^n dx = \frac{(x+a)^{n+1}((n+1)x-a)}{(n+1)(n+2)}$$

$$\int \frac{1}{1+x^2} dx = \tan^{-1} x$$

$$\int \frac{1}{a^2+x^2} dx = \frac{1}{a} \tan^{-1} \frac{x}{a}$$

$$\int \frac{x}{a^2+x^2} dx = \frac{1}{2} \ln |a^2+x^2|$$

$$\int \frac{x^2}{a^2+x^2} dx = x - a \tan^{-1} \frac{x}{a}$$

$$\int \frac{x^3}{a^2+x^2} dx = \frac{1}{2} x^2 - \frac{1}{2} a^2 \ln |a^2+x^2|$$

$$\int \frac{1}{ax^2+bx+c} dx = \frac{2}{\sqrt{4ac-b^2}} \tan^{-1} \frac{2ax+b}{\sqrt{4ac-b^2}}$$

$$\int \frac{1}{(x+a)(x+b)} dx = \frac{1}{b-a} \ln \frac{a+x}{b+x}, \ a \neq b$$

$$\int \frac{x}{(x+a)^2} dx = \frac{a}{a+x} + \ln |a+x|$$

$$\int \frac{x}{ax^2+bx+c} dx = \frac{1}{2a} \ln |ax^2+bx+c| - \frac{b}{a\sqrt{4ac-b^2}} \tan^{-1} \frac{2ax+b}{\sqrt{4ac-b^2}}$$

$$\int \sqrt{x-a}\, dx = \frac{2}{3} (x-a)^{3/2}$$

$$\int \frac{1}{\sqrt{x \pm a}}\, dx = 2\sqrt{x \pm a}$$

$$\int \frac{1}{\sqrt{a-x}}\, dx = -2\sqrt{a-x}$$

$$\int x\sqrt{x-a}\, dx = \begin{cases} \frac{2a}{3}(x-a)^{3/2} + \frac{2}{5}(x-a)^{5/2}, \ \text{or} \\ \frac{2}{3}x(x-a)^{3/2} - \frac{4}{15}(x-a)^{5/2}, \ \text{or} \\ \frac{2}{15}(2a+3x)(x-a)^{3/2} \end{cases}$$

$$\int \sqrt{ax+b}\, dx = \left( \frac{2b}{3a} + \frac{2x}{3} \right) \sqrt{ax+b}$$

$$\int (ax+b)^{3/2} dx = \frac{2}{5a} (ax+b)^{5/2}$$

$$\int \frac{x}{\sqrt{x \pm a}}\, dx = \frac{2}{3} (x \mp 2a) \sqrt{x \pm a}$$

$$\int \sqrt{\frac{x}{a-x}}\, dx = -\sqrt{x(a-x)} - a \tan^{-1} \frac{\sqrt{x(a-x)}}{x-a}$$

$$\int \sqrt{\frac{x}{a+x}}\, dx = \sqrt{x(a+x)} - a \ln \left( \sqrt{x} + \sqrt{x+a} \right)$$

$$\int x\sqrt{ax+b}\, dx = \frac{2}{15a^2} (-2b^2 + abx + 3a^2 x^2) \sqrt{ax+b}$$

$$\int \sqrt{x(ax+b)}\, dx = $$
$$\frac{1}{4a^{3/2}} \left( (2ax+b)\sqrt{ax(ax+b)} - b^2 \ln \left| a\sqrt{x} + \sqrt{a(ax+b)} \right| \right)$$

$$\int \sqrt{x^3(ax+b)}\, dx = $$
$$\left( \frac{b}{12a} - \frac{b^2}{8a^2 x} + \frac{x}{3} \right) \sqrt{x^3(ax+b)} + \frac{b^3}{8a^{5/2}} \ln \left| a\sqrt{x} + \sqrt{a(ax+b)} \right|$$

$$\int \sqrt{x^2 \pm a^2}\, dx = \frac{1}{2} x\sqrt{x^2 \pm a^2} \pm \frac{1}{2} a^2 \ln \left| x + \sqrt{x^2 \pm a^2} \right|$$

$$\int \sqrt{a^2 - x^2}\, dx = \frac{1}{2} x\sqrt{a^2 - x^2} + \frac{1}{2} a^2 \tan^{-1} \frac{x}{\sqrt{a^2 - x^2}}$$

$$\int x\sqrt{x^2 \pm a^2}\, dx = \frac{1}{3} \left( x^2 \pm a^2 \right)^{3/2}$$

$$\int \frac{1}{\sqrt{x^2 \pm a^2}}\, dx = \ln \left| x + \sqrt{x^2 \pm a^2} \right|$$

$$\int \frac{1}{\sqrt{a^2 - x^2}}\, dx = \sin^{-1} \frac{x}{a}$$

$$\int \frac{x}{\sqrt{x^2 \pm a^2}}\, dx = \sqrt{x^2 \pm a^2}$$

$$\int \frac{x}{\sqrt{a^2 - x^2}}\, dx = -\sqrt{a^2 - x^2}$$

$$\int \frac{x^2}{\sqrt{x^2 \pm a^2}}\, dx = \frac{1}{2} x\sqrt{x^2 \pm a^2} \mp \frac{1}{2} a^2 \ln \left| x + \sqrt{x^2 \pm a^2} \right|$$

$$\int \sqrt{ax^2+bx+c}\, dx = $$
$$\frac{b+2ax}{4a} \sqrt{ax^2+bx+c} + \frac{4ac-b^2}{8a^{3/2}} \ln \left| 2ax+b+2\sqrt{a(ax^2+bx+c)} \right|$$

$$\int x\sqrt{ax^2+bx+c}\, dx = $$
$$\frac{1}{48a^{5/2}} \left( 2\sqrt{a}\sqrt{ax^2+bx+c} \left( -3b^2 + 2abx + 8a(c+ax^2) \right) \right.$$
$$\left. + 3(b^3 - 4abc) \ln \left| b+2ax+2\sqrt{a}\sqrt{ax^2+bx+c} \right| \right)$$

$$\int \frac{1}{\sqrt{ax^2+bx+c}}\, dx = \frac{1}{\sqrt{a}} \ln \left| 2ax+b+2\sqrt{a(ax^2+bx+c)} \right|$$

$$\int \frac{x}{\sqrt{ax^2+bx+c}}\, dx = $$
$$\frac{1}{a} \sqrt{ax^2+bx+c} - \frac{b}{2a^{3/2}} \ln \left| 2ax+b+2\sqrt{a(ax^2+bx+c)} \right|$$

$$\int \frac{dx}{(a^2+x^2)^{3/2}} = \frac{x}{a^2\sqrt{a^2+x^2}}$$

$$\int \sin ax\, dx = -\frac{1}{a} \cos ax$$

$$\int \sin^2 ax\, dx = \frac{x}{2} - \frac{\sin 2ax}{4a}$$

$$\int \sin^3 ax \; dx = -\frac{3\cos ax}{4a} + \frac{\cos 3ax}{12a}$$

$$\int \cos ax \; dx = \frac{1}{a} \sin ax$$

$$\int \cos^2 ax \; dx = \frac{x}{2} + \frac{\sin 2ax}{4a}$$

$$\int \cos^3 ax\,dx = \frac{3\sin ax}{4a} + \frac{\sin 3ax}{12a}$$

$$\int \cos x \sin x \; dx = \frac{1}{2} \sin^2 x + c_1 = -\frac{1}{2} \cos^2 x + c_2 = -\frac{1}{4} \cos 2x + c_3$$

$$\int \cos ax \sin bx \; dx = \frac{\cos[(a-b)x]}{2(a-b)} - \frac{\cos[(a+b)x]}{2(a+b)}, a \neq b$$

$$\int \sin^2 ax \cos bx \; dx = -\frac{\sin[(2a-b)x]}{4(2a-b)} + \frac{\sin bx}{2b} - \frac{\sin[(2a+b)x]}{4(2a+b)}$$

$$\int \sin^2 x \cos x \; dx = \frac{1}{3} \sin^3 x$$

$$\int \cos^2 ax \sin bx \; dx = \frac{\cos[(2a-b)x]}{4(2a-b)} - \frac{\cos bx}{2b} - \frac{\cos[(2a+b)x]}{4(2a+b)}$$

$$\int \cos^2 ax \sin ax \; dx = -\frac{1}{3a} \cos^3 ax$$

$$\int \sin^2 ax \cos^2 bx\,dx =$$
$$\frac{x}{4} - \frac{\sin 2ax}{8a} - \frac{\sin[2(a-b)x]}{16(a-b)} + \frac{\sin 2bx}{8b} - \frac{\sin[2(a+b)x]}{16(a+b)}$$

$$\int \sin^2 ax \cos^2 ax \; dx = \frac{x}{8} - \frac{\sin 4ax}{32a}$$

$$\int \tan ax \; dx = -\frac{1}{a} \ln \cos ax$$

$$\int \tan^2 ax \; dx = -x + \frac{1}{a} \tan ax$$

$$\int \tan^3 ax\,dx = \frac{1}{a} \ln \cos ax + \frac{1}{2a} \sec^2 ax$$

$$\int \sec x \; dx = \ln |\sec x + \tan x| = 2 \tanh^{-1} \left( \tan \frac{x}{2} \right)$$

$$\int \sec^2 ax \; dx = \frac{1}{a} \tan ax$$

$$\int \sec^3 x \; dx = \frac{1}{2} \sec x \tan x + \frac{1}{2} \ln |\sec x + \tan x|$$

$$\int \sec x \tan x \; dx = \sec x$$

$$\int \sec^2 x \tan x \; dx = \frac{1}{2} \sec^2 x$$

$$\int \sec^n x \tan x \; dx = \frac{1}{n} \sec^n x, n \neq 0$$

$$\int \csc x \; dx = \ln \left| \tan \frac{x}{2} \right| = \ln |\csc x - \cot x| + C$$

$$\int \csc^2 ax \; dx = -\frac{1}{a} \cot ax$$

$$\int \csc^3 x \; dx = -\frac{1}{2} \cot x \csc x + \frac{1}{2} \ln |\csc x - \cot x|$$

$$\int \csc^n x \cot x \; dx = -\frac{1}{n} \csc^n x, n \neq 0$$

$$\int \sec x \csc x \; dx = \ln |\tan x|$$

$$\int x \cos x \; dx = \cos x + x \sin x$$

$$\int x \cos ax \; dx = \frac{1}{a^2} \cos ax + \frac{x}{a} \sin ax$$

$$\int x^2 \cos x \; dx = 2x \cos x + \left( x^2 - 2 \right) \sin x$$

$$\int x^2 \cos ax \; dx = \frac{2x \cos ax}{a^2} + \frac{a^2 x^2 - 2}{a^3} \sin ax$$

$$\int x \sin x \; dx = -x \cos x + \sin x$$

$$\int x \sin ax \; dx = -\frac{x \cos ax}{a} + \frac{\sin ax}{a^2}$$

$$\int x^2 \sin x \; dx = \left( 2 - x^2 \right) \cos x + 2x \sin x$$

$$\int x^2 \sin ax \; dx = \frac{2 - a^2 x^2}{a^3} \cos ax + \frac{2x \sin ax}{a^2}$$

$$\int x \cos^2 x \; dx = \frac{x^2}{4} + \frac{1}{8} \cos 2x + \frac{1}{4} x \sin 2x$$

$$\int x \sin^2 x \; dx = \frac{x^2}{4} - \frac{1}{8} \cos 2x - \frac{1}{4} x \sin 2x$$