Platelet

 $\begin{array}{cc} Team \ Reference \ Material \\ {}_{(unlimited \ version)} \end{array}$



凌皓煜 陈 彤

顾 逸

目录

1	\mathbf{Gra}	ph Theory	1
	1.1	2-SAT (ct)	1
	1.2	割点与桥 (ct)	2
	1.3	Steiner tree (lhy)	3
	1.4	K 短路 (lhy)	3
	1.5	最大团	7
	1.6	一般图最大匹配	7
	1.7	带花树	7
	1.8	KM 算法	7
	1.9	支配树	7
	1.10		8
		Prufer 编码	
		Link-Cut Tree (ct)	
		圆方树 (ct)	
		最小割	
		最大流 (ct)	
		费用流 (ct)	
		有上下界的网络流 (Durandal)	
		差分约束 1	
		图论知识 (gy,lhy)	
	1.21	国记》(gy,my) · · · · · · · · · · · · · · · · · · ·	'
2	Mat	$^{\circ}$ h	9
	2.1	int64 相乘取模 (Durandal)	9
	2.2	扩展欧几里得 (gy)	
	2.3	中国剩余定理 (Durandal)	
	2.4	後性同余不等式 (Durandal)	
	2.5	组合数	
	2.6	高斯消元 (ct)	
	2.7	Miller Rabin & Pollard Rho (gy)	
	2.8	$O(m^2 \log n)$ 线性递推 $\dots \dots \dots$	
	2.9	Polynomial	
	2.3	2.9.1 FFT	
		2.9.2 NTT & 多项式求逆	
	2.10	Lagrange 插值 (ct)	
		Lagrange mm (ct)	
	2.12	BSGS (ct,Durandal)	
			,4
		2.12.1 BSGS(ct)	
	0.10	2.12.2 扩展 BSGS(Durandal)	5
		2.12.2 扩展 BSGS(Durandal) 2 直线下整点个数 (gy) 2	5
	2.14	2.12.2 扩展 BSGS(Durandal) 2 直线下整点个数 (gy) 2 单纯形 2	5 5 6
	$2.14 \\ 2.15$	2.12.2 扩展 BSGS(Durandal) 2 直线下整点个数 (gy) 2	5 5 6

3	Geo	netry	29
	3.1	点、直线、圆 (gy)	29
	3.2	平面最近点对 (Grimoire)	32
	3.3	コ包游戏 (Grimoire)	33
	3.4	半平面交 (Grimoire)	35
	3.5	点在多边形内 (Grimoire)	36
	3.6	最小圆覆盖 (Grimoire)	36
	3.7	最小球覆盖 (Grimoire)	37
	3.8	圆并 (Grimoire)	38
	3.9	圆与多边形并 (Grimoire)	41
		三角剖分 (Grimoire) 	42
	3.11	三维几何基础 (Grimoire)	44
		三维凸包 (Grimoire)	45
		三维绕轴旋转 (Grimoire)	46
		几何知识 (gy)	47
	0.11	101.1 VH W. (P2)	11
4	Stri	σ	49
_	4.1	0	49
	4.2	AC 自动机	50
		1,7,815	
	4.3	后缀数组 (ct)	50
	4.4	后缀自动机 (ct,lhy)	50
	4.5	Manacher (ct)	52
	4.6	回文树 (lhy)	53
	4.7	最小表示法 (lhy)	54
_	ъ.,		
5			55
5	5.1	莫队 (ct)	55
5	$5.1 \\ 5.2$	莫队 (ct)	55 55
5	5.1	莫队 (ct)	55 55 56
5	$5.1 \\ 5.2$	莫队 (ct)	55 55 56 57
5	5.1 5.2 5.3	莫队 (ct)	55 55 56
5	5.1 5.2 5.3 5.4	莫队 (ct)	55 55 56 57
5	5.1 5.2 5.3 5.4 5.5	莫队 (ct)	55 55 56 57 57
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7	莫队 (ct)	55 55 56 57 57 58 62
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8	莫队 (ct)	55 55 56 57 57 58 62 63
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9	莫队 (ct)	55 55 56 57 57 58 62 63 66
5	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9	莫队 (ct)	55 55 56 57 57 58 62 63
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10	莫队 (ct) ST 表 (ct) 滞权并查集 (ct) 可并堆 (ct) kw 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct)	55 55 56 57 57 58 62 63 66 67
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth	模队 (ct) ST 表 (ct) F权并查集 (ct) 可并堆 (ct) ckw 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct)	55 55 56 57 57 58 62 63 66 67 69
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1	模队 (ct) ST 表 (ct) 带权并查集 (ct) 可并堆 (ct) ckw 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct) rs vimrc (gy)	55 55 56 57 57 58 62 63 66 67 69
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1 6.2	莫队 (ct)	55 55 56 57 57 58 62 63 66 67 69 69
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1 6.2 6.3	模队 (ct) ST 表 (ct) 滞权并查集 (ct) 可并堆 (ct) ckw 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct) rs vimrc (gy) STL 释放内存 (Durandal) 开栈 (Durandal)	55 55 56 57 57 58 62 63 66 67 69 69
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1 6.2 6.3 6.4	莫队 (ct) ST 表 (ct) 滞权并查集 (ct) 可并堆 (ct) ckw 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct) rs vimrc (gy) STL 释放内存 (Durandal) 开栈 (Durandal) の3 (gy)	55 55 56 57 57 58 62 63 66 67 69 69 69 70
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1 6.2 6.3 6.4 6.5	模队 (ct) ST 表 (ct) 带权并查集 (ct) 可并堆 (ct) dkw 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct) rs vimrc (gy) STL 释放内存 (Durandal) 开栈 (Durandal) 开栈 (Durandal) O3 (gy) Java Template (gy)	55 55 56 57 57 58 62 63 66 67 69 69 69 70 70
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1 6.2 6.3 6.4 6.5 6.6	模队 (ct) ST 表 (ct) 带权并查集 (ct) 可并堆 (ct) 或w 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct) rs vimrc (gy) STL 释放内存 (Durandal) 开栈 (Durandal) 开栈 (Durandal) O3 (gy) Java Template (gy) Big Fraction (gy)	55 55 56 57 57 58 62 63 66 67 69 69 69 70 70 72
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1 6.2 6.3 6.4 6.5	模队 (ct) ST 表 (ct) 带权并查集 (ct) 可并堆 (ct) dkw 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct) rs vimrc (gy) STL 释放内存 (Durandal) 开栈 (Durandal) 开栈 (Durandal) O3 (gy) Java Template (gy)	55 55 56 57 57 58 62 63 66 67 69 69 69 70 70
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1 6.2 6.3 6.4 6.5 6.6	模队 (ct) ST 表 (ct) 带权并查集 (ct) 可并堆 (ct) 或w 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct) rs vimrc (gy) STL 释放内存 (Durandal) 开栈 (Durandal) 开栈 (Durandal) O3 (gy) Java Template (gy) Big Fraction (gy)	55 55 56 57 57 58 62 63 66 67 69 69 69 70 70 72
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1 6.2 6.3 6.4 6.5 6.6 6.7	模队 (ct) ST 表 (ct) 带权并查集 (ct) 可并堆 (ct) dkw 线段树 (ct) Splay (ct) Treap (ct) 可持久化平衡树 (ct) CDQ 分治 (ct) Bitset (ct) rs vimrc (gy) STL 释放内存 (Durandal) 开栈 (Durandal) O3 (gy) Java Template (gy) Big Fraction (gy) 模拟退火 (ct)	55 55 56 57 57 58 62 63 66 67 69 69 69 70 72 72
	5.1 5.2 5.3 5.4 5.5 5.6 5.7 5.8 5.9 5.10 Oth 6.1 6.2 6.3 6.4 6.5 6.6 6.7 6.8 6.9	模队 (ct)	55 55 56 57 57 58 62 63 66 67 69 69 69 70 72 72 73

Chapter 1

Graph Theory

1.1 2-SAT (ct)

```
struct Edge {
       Edge *next;
       int to;
  } *last[maxn << 1], e[maxn << 2], *ecnt = e;
5 inline void link(int a, int b)
       *++ecnt = (Edge) {last[a], b}; last[a] = ecnt;
  }
   int dfn[maxn], low[maxn], timer, st[maxn], top, id[maxn], colcnt, n;
   bool fail, used[maxn];
10
   void tarjan(int x, int fa)
11
12
       dfn[x] = low[x] = ++timer; st[++top] = x;
13
       for (R Edge *iter = last[x]; iter; iter = iter -> next)
14
           if (iter -> to != fa)
15
16
               if (!dfn[iter -> to])
17
18
                   tarjan(iter -> to, x);
19
                    cmin(low[x], low[iter -> to]);
20
21
               else if (!id[iter -> to]) cmin(low[x], dfn[iter -> to]);
22
23
       if (dfn[x] == low[x])
25
           ++colcnt; bool flag = 1;
           for (; ;)
27
28
               int now = st[top--];
29
               id[now] = colcnt;
30
               if (now \le 2 * n)
31
32
                   flag \&= !used[id[now <= n ? now + n : now - n]];
33
                   now <= n ? fail |= (id[now + n] == id[now]) : fail |= (id[now - n] == id[now]);</pre>
34
35
               if (now == x) break;
36
           }
37
           used[colcnt] = flag;
38
39
40 }
41 int ans[maxn], tot;
42 int main()
```

1.2. 割点与桥 (ct) 1. Graph Theory

```
43 {
44
           build your graph here.
45
       */
46
       for (R int i = 1; !fail && i <= n; ++i) if (!dfn[i]) tarjan(i, 0);</pre>
47
       if (fail)
48
       {
49
           puts("Impossible");
50
           return 0;
51
       }
52
       for (R int i = 1; i <= n; ++i) if (used[id[i]]) ans[++tot] = i;
53
       printf("%d\n", tot);
       std::sort(ans + 1, ans + tot + 1);
55
       for (R int i = 1; i <= tot; ++i) printf("%d ", ans[i]);</pre>
56
       return 0;
57
```

1.2 割点与桥 (ct)

割点

```
int dfn[maxn], low[maxn], timer, ans, num;
void tarjan(int x, int fa)
3 | {
       dfn[x] = low[x] = ++timer;
       for (Edge *iter = last[x]; iter; iter = iter -> next)
5
            if (iter -> to != fa)
6
7
                if (!dfn[iter -> to])
9
                     tarjan(iter -> to, x);
cmin(low[x], low[iter -> to]);
10
11
                     if (dfn[x] <= low[iter -> to])
12
13
                         cut[x] = 1;
14
                         if (!fa && dfn[x] < low[iter \rightarrow to]) num = 233;
15
                         else if (!fa) ++num;
16
17
18
                else cmin(low[x], dfn[iter -> to]);
19
20
21
  | }
22
  int main()
23
       for (int i = 1; i <= n; ++i)
24
            if (!dfn[i])
25
            ₹
26
                num = 0;
27
                tarjan(i, 0);
28
                if (num == 1) cut[i] = 0;
29
            }
30
```

桥

```
int dfn[maxn], low[maxn], timer;
void tarjan(int x, int fa)
{
```

1. Graph Theory 1.3. Steiner tree (lhy)

```
dfn[x] = low[x] = ++timer;
       for (R Edge *iter = last[x]; iter; iter = iter -> next)
5
           if (iter -> to != fa)
6
7
               if (!dfn[iter -> to])
8
9
                   dfs(iter -> to, x);
10
                    cmin(low[x], low[iter -> to]);
11
                    if (dfn[x] < low[iter -> to]) ans[x][iter -> to] = ans[iter -> to][x] = 1;
12
13
               else cmin(low[x], dfn[iter -> to]);
14
           }
15
16
```

1.3 Steiner tree (lhy)

```
void Steiner_Tree()
1
2
3
       memset(f, 0x3f, sizeof(f));
       for(int i = 1; i <= n; i++)
5
           f[0][i] = 0;
       for(int i = 1; i <= p; i++)</pre>
           f[1 << (i - 1)][idx[i]] = 0;
       int S = 1 << p;</pre>
       for(int s = 1; s < S; s++)</pre>
9
10
           for(int i = 1; i <= n; i++)
11
12
                for(int k = (s - 1) \& s; k; k = (k - 1) \& s)
13
                    f[s][i] = min(f[s][i], f[k][i] + f[s^k][i]);
15
           }
16
           SPFA(f[s]);
17
       }
18
       int ans = inf;
       for(int i = 1; i <= n; i++)
19
            ans = min(ans, f[S - 1][i]);
20
21
```

1.4 K 短路 (lhy)

```
const int MAXNODE = MAXN + MAXM * 2;
  bool used[MAXN];
int n, m, cnt, S, T, Kth, N, TT;
  int rt[MAXN], seq[MAXN], adj[MAXN], from[MAXN], dep[MAXN];
  LL dist[MAXN], w[MAXM], ans[MAXK];
  struct GivenEdge{
6
      int u, v, w;
      GivenEdge() {};
      GivenEdge(int _u, int _v, int _w) : u(_u), v(_v), w(_w){};
  }edge[MAXM];
  struct Edge{
11
      int v, nxt, w;
12
      Edge() {};
13
      Edge(int _v, int _nxt, int _w) : v(_v), nxt(_nxt), w(_w) \{\};
15 }e[MAXM];
```

1.4. K 短路 (lhy) 1. Graph Theory

```
inline void addedge(int u, int v, int w)
17 {
       e[++cnt] = Edge(v, adj[u], w); adj[u] = cnt;
18
19
   void dij(int S)
20
21
       for(int i = 1; i <= N; i++)</pre>
22
23
           dist[i] = INF;
24
           dep[i] = 0x3f3f3f3f;
25
           used[i] = false;
26
           from[i] = 0;
27
       }
28
       static priority_queue<pair<LL, int>, vector<pair<LL, int> >, greater<pair<LL, int> > hp;
29
       while(!hp.empty())hp.pop();
30
       hp.push(make_pair(dist[S] = 0, S));
31
       dep[S] = 1;
32
       while(!hp.empty())
33
           pair<LL, int> now = hp.top();
35
           hp.pop();
36
           int u = now.second;
37
           if(used[u])continue;
38
           else used[u] = true;
39
           for(int p = adj[u]; p; p = e[p].nxt)
40
41
               int v = e[p].v;
42
43
               if(dist[u] + e[p].w < dist[v])</pre>
44
                   dist[v] = dist[u] + e[p].w;
45
                   dep[v] = dep[u] + 1;
46
                   from[v] = p;
47
                   hp.push(make_pair(dist[v], v));
48
49
           }
50
51
       for(int i = 1; i <= m; i++)</pre>
                                       w[i] = 0;
52
       for(int i = 1; i <= N; i++)</pre>
53
           if(from[i])w[from[i]] = -1;
54
       for(int i = 1; i <= m; i++)
55
56
           57
58
               w[i] = -dist[edge[i].u] + (dist[edge[i].v] + edge[i].w);
59
           }
60
           else
61
           {
62
               w[i] = -1;
63
           }
64
       }
65
66
67 inline bool cmp_dep(int p, int q)
68
       return dep[p] < dep[q];</pre>
69
  }
70
71 struct Heap{
       LL key;
```

1. Graph Theory 1.4. K 短路 (lhy)

```
int id, lc, rc, dist;
73
        Heap() {};
74
        Heap(LL k, int i, int l, int r, int d) : key(k), id(i), lc(l), rc(r), dist(d) {};
75
        inline void clear()
76
77
            key = 0;
78
            id = lc = rc = dist = 0;
79
80
    }hp[MAXNODE];
81
    inline int merge_simple(int u, int v)
82
83
        if(!u)return v;
84
        if(!v)return u;
85
        if(hp[u].key > hp[v].key)
86
87
            swap(u, v);
88
        }
89
        hp[u].rc = merge_simple(hp[u].rc, v);
90
        if(hp[hp[u].lc].dist < hp[hp[u].rc].dist)</pre>
91
        {
92
            swap(hp[u].lc, hp[u].rc);
93
        }
94
        hp[u].dist = hp[hp[u].rc].dist + 1;
95
        return u;
96
97
    inline int merge_full(int u, int v)
98
99
        if(!u)return v;
100
101
        if(!v)return u;
102
        if(hp[u].key > hp[v].key)
103
        {
            swap(u, v);
104
        }
105
        int nownode = ++cnt;
106
        hp[nownode] = hp[u];
107
        hp[nownode].rc = merge_full(hp[nownode].rc, v);
108
        if(hp[hp[nownode].lc].dist < hp[hp[nownode].rc].dist)</pre>
109
110
            swap(hp[nownode].lc, hp[nownode].rc);
111
112
113
        hp[nownode].dist = hp[hp[nownode].rc].dist + 1;
114
        return nownode;
115
   priority_queue<pair<LL, int>, vector<pair<LL, int> >, greater<pair<LL, int> > Q;
116
    int main()
117
118
        while(scanf("%d%d", &n, &m) != EOF)
119
120
            scanf("%d%d%d%d", &S, &T, &Kth, &TT);
121
            for(int i = 1; i <= m; i++)</pre>
122
123
                 int u, v, w;
124
                 scanf("%d%d%d", &u, &v, &w);
125
                 edge[i] = \{u, v, w\};
126
127
            N = n;
128
            memset(adj, 0, sizeof(*adj) * (N + 1));
129
```

1.4. K 短路 (lhy) 1. Graph Theory

```
cnt = 0;
130
            for(int i = 1; i <= m; i++)
131
                 addedge(edge[i].v, edge[i].u, edge[i].w);
132
            dij(T);
133
            if(dist[S] > TT)
134
            {
135
                 puts("Whitesnake!");
136
                 continue;
137
            }
138
            for(int i = 1; i <= N; i++)
139
                 seq[i] = i;
140
            sort(seq + 1, seq + N + 1, cmp_dep);
141
            cnt = 0;
142
            memset(adj, 0, sizeof(*adj) * (N + 1));
143
            memset(rt, 0, sizeof(*rt) * (N + 1));
144
            for(int i = 1; i <= m; i++)</pre>
145
                 addedge(edge[i].u, edge[i].v, edge[i].w);
146
            rt[T] = cnt = 0;
147
            hp[0].dist = -1;
148
            for(int i = 1; i <= N; i++)</pre>
149
150
                 int u = seq[i], v = edge[from[u]].v;
151
                 rt[u] = 0;
152
                 for(int p = adj[u]; p; p = e[p].nxt)
153
154
                     if(~w[p])
155
                     {
156
                          hp[++cnt] = Heap(w[p], p, 0, 0, 0);
157
                          rt[u] = merge_simple(rt[u], cnt);
158
159
                 }
160
                 if(i == 1)continue;
161
                 rt[u] = merge_full(rt[u], rt[v]);
162
            }
163
            while(!Q.empty())Q.pop();
164
            Q.push(make_pair(dist[S], 0));
165
            edge[0].v = S;
166
            for(int kth = 1; kth <= Kth; kth++)</pre>
167
168
                 if(Q.empty())
169
                 {
170
171
                     ans[kth] = -1;
                     continue;
^{172}
173
                 pair<LL, int> now = Q.top(); Q.pop();
174
                 ans[kth] = now.first;
175
                 int p = now.second;
176
                 if(hp[p].lc)
177
                 {
178
                     Q.push(make_pair(+hp[hp[p].lc].key + now.first - hp[p].key, hp[p].lc));
179
                 }
180
                 if(hp[p].rc)
181
                 {
182
                     Q.push(make_pair(+hp[hp[p].rc].key + now.first - hp[p].key, hp[p].rc));
183
                 }
184
                 if(rt[edge[hp[p].id].v])
185
                 {
186
                     Q.push(make_pair(hp[rt[edge[hp[p].id].v]].key + now.first, rt[edge[hp[p].id].v]));
187
                 }
188
            }
189
```

1. Graph Theory 1.5. 最大团

```
if(ans[Kth] == -1 \mid \mid ans[Kth] > TT)
190
              {
191
                   puts("Whitesnake!");
192
              }
193
              else
194
              {
195
                   puts("yareyaredawa");
196
197
         }
198
199
```

- 1.5 最大团
- 1.6 一般图最大匹配
- 1.7 带花树
- 1.8 KM 算法
- 1.9 支配树

DAG (ct)

```
struct Edge {
       Edge *next;
       int to;
3
   } ;
   Edge *last[maxn], e[maxm], *ecnt = e; // original graph
  Edge *rlast[maxn], re[maxm], *recnt = re; // reversed-edge graph
   Edge *tlast[maxn], te[maxn << 1], *tecnt = te; // dominate tree graph</pre>
   int deg[maxn], q[maxn], fa[maxn][20], all_fa[maxn], fa_cnt, size[maxn], dep[maxn];
   inline void link(int a, int b)
9
10
       *++ecnt = (Edge) {last[a], b}; last[a] = ecnt; ++deg[b];
11
  }
12
   inline void link_rev(R int a, R int b)
13
   {
14
       *++recnt = (Edge) {rlast[a], b}; rlast[a] = recnt;
15
  }
16
   inline void link_tree(R int a, R int b)
17
       *++tecnt = (Edge) {tlast[a], b}; tlast[a] = tecnt;
19
  }
20
   inline int getlca(R int a, R int b)
21
22
       if (dep[a] < dep[b]) std::swap(a, b);</pre>
23
       R int temp = dep[a] - dep[b];
24
       for (R int i; temp; temp -= 1 << i)
25
       a = fa[a][i = __builtin_ctz(temp)];
for (R int i = 16; ~i; --i)
26
27
           if (fa[a][i] != fa[b][i])
28
                a = fa[a][i], b = fa[b][i];
29
       if (a == b) return a;
30
       return fa[a][0];
31
  }
32
   void dfs(R int x)
33
34 {
```

1.10. 虚树 (ct) 1. Graph Theory

```
size[x] = 1;
       for (R Edge *iter = tlast[x]; iter; iter = iter -> next)
36
           dfs(iter -> to), size[x] += size[iter -> to];
37
  }
38
  int main()
39
   {
40
       q[1] = 0;
41
       R int head = 0, tail = 1;
42
       while (head < tail)
43
44
           R int now = q[++head];
45
46
           fa_cnt = 0;
           for (R Edge *iter = rlast[now]; iter; iter = iter -> next)
47
               all_fa[++fa_cnt] = iter -> to;
48
           for (; fa_cnt > 1; --fa_cnt)
49
                all_fa[fa_cnt - 1] = getlca(all_fa[fa_cnt], all_fa[fa_cnt - 1]);
50
           fa[now][0] = all_fa[fa_cnt];
51
           dep[now] = dep[all_fa[fa_cnt]] + 1;
52
           if (now) link_tree(fa[now][0], now);
53
           for (R int i = 1; i \le 16; ++i)
               fa[now][i] = fa[fa[now][i - 1]][i - 1];
           for (R Edge *iter = last[now]; iter; iter = iter -> next)
56
               if (--deg[iter \rightarrow to] == 0) q[++tail] = iter \rightarrow to;
57
       }
58
       dfs(0);
59
       for (R int i = 1; i <= n; ++i) printf("%d\n", size[i] - 1 );
60
       return 0;
61
62
```

一般图

1.10 虚树 (ct)

```
struct Edge {
       Edge *next;
       int to;
  } *last[maxn], e[maxn << 1], *ecnt = e;</pre>
5 inline void link(int a, int b)
6
       *++ecnt = (Edge) {last[a], b}; last[a] = ecnt;
       *++ecnt = (Edge) {last[b], a}; last[b] = ecnt;
int a[maxn], n, dfn[maxn], pos[maxn], timer, inv[maxn], st[maxn];
int fa[maxn], size[maxn], dep[maxn], son[maxn], top[maxn];
  bool vis[maxn];
   void dfs1(int x)
13
   {
14
       vis[x] = 1; size[x] = 1; dep[x] = dep[fa[x]] + 1;
15
       for (R Edge *iter = last[x]; iter; iter = iter -> next)
16
           if (!vis[iter -> to])
17
18
               fa[iter -> to] = x;
19
               dfs1(iter -> to);
20
               size[x] += size[iter -> to];
21
               size[son[x]] < size[iter -> to] ? son[x] = iter -> to : 0;
22
23
24 }
void dfs2(int x)
```

1. Graph Theory 1.10. 虚树 (ct)

```
26 {
       vis[x] = 0; top[x] = x == son[fa[x]] ? top[fa[x]] : x;
27
       dfn[x] = ++timer; pos[timer] = x;
28
       if (son[x]) dfs2(son[x]);
29
       for (R Edge *iter = last[x]; iter; iter = iter -> next)
30
           if (vis[iter -> to]) dfs2(iter -> to);
31
       inv[x] = timer;
32
33
   inline int getlca(int a, int b)
34
35
       while (top[a] != top[b])
36
           dep[top[a]] < dep[top[b]] ? b = fa[top[b]] : a = fa[top[a]];
37
       return dep[a] < dep[b] ? a : b;</pre>
38
   }
39
   inline bool cmp(int a, int b)
40
41
       return dfn[a] < dfn[b];
42
  }
43
   inline bool isson(int a, int b)
44
45
       return dfn[a] <= dfn[b] && dfn[b] <= inv[a];</pre>
46
  }
47
   typedef long long 11;
48
   bool imp[maxn];
49
   struct sEdge {
50
       sEdge *next;
51
       int to, w;
52
   } *slast[maxn], se[maxn << 1], *secnt = se;</pre>
53
   inline void slink(int a, int b, int w)
54
55
       *++secnt = (sEdge) {slast[a], b, w}; slast[a] = secnt;
56
57
58
   int main()
59
       scanf("%d", &n);
60
       for (int i = 1; i < n; ++i)
61
       {
62
           int a, b; scanf("%d%d", &a, &b);
63
           link(a, b);
64
       }
65
       int m; scanf("%d", &m);
66
       dfs1(1); dfs2(1);
67
       memset(size, 0, (n + 1) << 2);
       for (; m; --m)
69
70
           int top = 0; scanf("%d", &k);
71
           for (int i = 1; i <= k; ++i) scanf("%d", \&a[i]), vis[a[i]] = imp[a[i]] = 1;
72
           std::sort(a + 1, a + k + 1, cmp);
73
           int p = k;
74
           for (int i = 1; i < k; ++i)
75
76
                int lca = getlca(a[i], a[i + 1]);
77
                if (!vis[lca]) vis[a[++p] = lca] = 1;
78
           }
79
           std::sort(a + 1, a + p + 1, cmp);
80
           st[++top] = a[1];
81
           for (int i = 2; i <= p; ++i)
82
83
                while (!isson(st[top], a[i])) --top;
84
                slink(st[top], a[i], dep[a[i]] - dep[st[top]]);
85
                st[++top] = a[i];
86
```

1.11. 树上点分治 (ct) 1. Graph Theory

1.11 树上点分治 (ct)

```
int root, son[maxn], size[maxn], sum;
bool vis[maxn];
void dfs_root(int x, int fa)
4 | {
       size[x] = 1; son[x] = 0;
       for (R Edge *iter = last[x]; iter; iter = iter -> next)
           if (iter -> to == fa || vis[iter -> to]) continue;
           dfs_root(iter -> to, x);
10
           size[x] += size[iter -> to];
11
           cmax(son[x], size[iter -> to]);
12
       cmax(son[x], sum - size[x]);
13
       if (!root || son[x] < son[root]) root = x;</pre>
14
15
   void dfs_chain(int x, int fa, int st1, int st2)
16
17
18
19
           write your code here.
20
       for (Edge *iter = last[x]; iter; iter = iter -> next)
21
22
           if (vis[iter -> to] || iter -> to == fa) continue;
23
           dfs_chain(iter -> to, x);
24
25
26
   void calc(int x)
27
28
       for (Edge *iter = last[x]; iter; iter = iter -> next)
30
           if (vis[iter -> to]) continue;
           dfs_chain(iter -> to, x);
32
33
               write your code here.
34
35
       }
36
37
   void work(int x)
38
39
       vis[x] = 1;
40
41
       calc(x);
       for (R Edge *iter = last[x]; iter; iter = iter -> next)
42
43
           if (vis[iter -> to]) continue;
44
           root = 0;
45
           sum = size[iter -> to];
46
47
           dfs_root(iter -> to, 0);
```

1. Graph Theory 1.12. 树上倍增 (ct)

```
work(root);
48
       }
49
  }
50
  int main()
51
52
       root = 0; sum = n;
53
       dfs_root(1, 0);
54
       work(root);
55
       return 0;
56
```

1.12 树上倍增 (ct)

```
int fa[maxn][17], mn[maxn][17], dep[maxn];
  bool vis[maxn];
  void dfs(int x)
3
  {
4
       vis[x] = 1;
       for (int i = 1; i <= 16; ++i)
6
           if (dep[x] < (1 << i)) break;
           fa[x][i] = fa[fa[x][i - 1]][i - 1];
           mn[x][i] = dmin(mn[x][i - 1], mn[fa[x][i - 1]][i - 1]);
10
11
       for (Edge *iter = last[x]; iter; iter = iter -> next)
12
           if (!vis[iter -> to])
13
           {
14
                fa[iter -> to][0] = x;
15
                mn[iter -> to][0] = iter -> w;
16
                dep[iter \rightarrow to] = dep[x] + 1;
17
               dfs(iter -> to);
18
           }
19
20
   inline int getlca(int x, int y)
21
22
       if (dep[x] < dep[y]) std::swap(x, y);
23
       int t = dep[x] - dep[y];
24
       for (int i = 0; i \le 16 \&\& t; ++i)
25
           if ((1 << i) & t)
26
               x = fa[x][i], t ^= 1 << i;
27
       for (int i = 16; i >= 0; --i)
           if (fa[x][i] != fa[y][i])
29
30
           {
               x = fa[x][i];
31
               y = fa[y][i];
32
33
       if (x == y) return x;
34
       return fa[x][0];
35
36
   inline int getans(int x, int f)
37
38
       int ans = inf, t = dep[x] - dep[f];
39
       for (int i = 0; i <= 16 && t; ++i)
40
           if (t & (1 << i))
41
42
           {
                cmin(ans, mn[x][i]);
43
               x = fa[x][i];
44
                t ^= 1 << i;
45
           }
46
```

1.13. Prufer 编码 1. Graph Theory

```
return ans;
48 }
```

1.13 Prufer 编码

1.14 Link-Cut Tree (ct)

```
struct Node *null;
   struct Node {
       Node *ch[2], *fa, *pos;
       int val, mn, l, len; bool rev;
       // min_val in chain
       inline bool type()
            return fa -> ch[1] == this;
       inline bool check()
10
11
            return fa -> ch[type()] == this;
13
       }
14
       inline void pushup()
15
            pos = this; mn = val;
16
            ch[0] \rightarrow mn < mn ? mn = ch[0] \rightarrow mn, pos = ch[0] \rightarrow pos : 0;
17
            ch[1] \rightarrow mn < mn ? mn = ch[1] \rightarrow mn, pos = ch[1] \rightarrow pos : 0;
18
            len = ch[0] -> len + ch[1] -> len + 1;
19
20
       inline void pushdown()
21
22
            if (rev)
23
24
                 ch[0] -> rev ^= 1;
25
                 ch[1] -> rev ^= 1;
26
                 std::swap(ch[0], ch[1]);
27
                 rev ^= 1;
28
29
30
       inline void pushdownall()
31
32
            if (check()) fa -> pushdownall();
34
            pushdown();
       }
35
       inline void rotate()
36
37
            bool d = type(); Node *f = fa, *gf = f -> fa;
38
            (fa = gf, f \rightarrow check()) ? fa \rightarrow ch[f \rightarrow type()] = this : 0;
39
            (f \rightarrow ch[d] = ch[!d]) != null ? ch[!d] \rightarrow fa = f : 0;
40
            (ch[!d] = f) \rightarrow fa = this;
41
            f -> pushup();
42
43
       inline void splay(bool need = 1)
44
45
            if (need) pushdownall();
46
            for (; check(); rotate())
47
                 if (fa -> check())
48
                      (type() == fa -> type() ? fa : this) -> rotate();
49
            pushup();
50
       }
```

1. Graph Theory 1.15. 圆方树 (ct)

```
inline Node *access()
52
53
           Node *i = this, *j = null;
54
           for (; i != null; i = (j = i) -> fa)
55
            {
56
                i -> splay();
57
                i \rightarrow ch[1] = j;
58
                i -> pushup();
59
            }
60
61
           return j;
62
       inline void make_root()
63
64
           access();
65
           splay();
66
           rev ^= 1;
67
68
       inline void link(Node *that)
69
70
71
           make_root();
72
           fa = that;
           splay(0);
73
       }
74
       inline void cut(Node *that)
75
76
           make_root();
77
           that -> access();
78
           that -> splay(0);
79
           that -> ch[0] = fa = null;
80
81
            that -> pushup();
       }
82
83
   } mem[maxn];
   inline Node *query(Node *a, Node *b)
84
85
       a -> make_root(); b -> access(); b -> splay(0);
86
       return b -> pos;
87
  }
88
  inline int dist(Node *a, Node *b)
89
90
       a -> make_root(); b -> access(); b -> splay(0);
91
       return b -> len;
92
```

1.15 圆方树 (ct)

```
int dfn[maxn], low[maxn], timer, st[maxn], top, id[maxn], scc;
  void dfs(int x)
2
  {
3
       dfn[x] = low[x] = ++timer; st[++top] = x;
       for (Edge *iter = last[x]; iter; iter = iter -> next)
5
           if (!dfn[iter -> to])
6
               dfs(iter -> to);
               cmin(low[x], low[iter -> to]);
9
               if (dfn[x] == low[iter->to])
10
11
                   int now, elder = top, minn = c[x];
12
                   ++scc;
13
                   do
14
```

1.16. 最小割 1. Graph Theory

```
{
                         now = st[top--];
16
                         cmin(minn, c[now]);
17
                    }
18
                    while (iter -> to != now);
19
                    for (int i = top + 1; i <= elder; ++i)</pre>
20
                         add(scc, st[i], minn);
21
                    add(scc, x, minn);
22
23
           }
24
            else if (!id[iter -> to]) cmin(low[x], dfn[iter -> to]);
25
26
```

1.16 最小割

1.17 最大流 (ct)

```
struct Edge {
       Edge *next, *rev;
       int to, cap;
  5 inline void link(R int a, R int b, R int w)
6
       *++ecnt = (Edge) {last[a], ecnt + 1, b, w}; last[a] = ecnt;
       *++ecnt = (Edge) {last[b], ecnt - 1, a, 0}; last[b] = ecnt;
9
  int ans, s, t, q[maxn], dep[maxn];
10
   inline bool bfs()
11
12
       memset(dep, -1, (t + 1) << 2);
13
       dep[q[1] = t] = 0; int head = 0, tail = 1;
14
       while (head < tail)
15
16
           int now = q[++head];
17
           for (Edge *iter = last[now]; iter; iter = iter -> next)
18
               if (dep[iter -> to] == -1 && iter -> rev -> cap)
19
                   dep[q[++tail] = iter \rightarrow to] = dep[now] + 1;
20
21
       return dep[s] != -1;
22
24 int dfs(int x, int f)
25
       if (x == t) return f;
26
       int used = 0;
27
       for (Edge* &iter = cur[x]; iter; iter = iter -> next)
28
           if (iter \rightarrow cap && dep[iter \rightarrow to] + 1 == dep[x])
29
30
               int v = dfs(iter -> to, dmin(f - used, iter -> cap));
31
               iter -> cap -= v;
32
               iter -> rev -> cap += v;
33
               used += v;
34
               if (used == f) return f;
35
          }
36
       return used;
37
  ١,
38
39 inline void dinic()
  {
40
41
       while (bfs())
```

1. Graph Theory 1.18. 费用流 (ct)

1.18 费用流 (ct)

Dinic(ct)

```
struct Edge {
       Edge *next, *rev;
2
       int from, to, cap, cost;
  } *last[maxn], *prev[maxn], e[maxm], *ecnt = e;
  inline void link(int a, int b, int w, int c)
5
6
       *++ecnt = (Edge) {last[a], ecnt + 1, a, b, w, c}; last[a] = ecnt;
       *++ecnt = (Edge) {last[b], ecnt - 1, b, a, 0, -c}; last[b] = ecnt;
   int s, t, q[maxn << 2], dis[maxn];</pre>
10
  ll ans;
  bool inq[maxn];
   #define inf Ox7fffffff
   inline bool spfa()
14
15
       for (int i = 1; i <= t; ++i) dis[i] = inf;</pre>
16
       int head = 0, tail = 1; dis[q[1] = s] = 0;
17
       while (head < tail)
18
19
20
            int now = q[++head]; inq[now] = 0;
21
            for (Edge *iter = last[now]; iter; iter = iter -> next)
                if (iter -> cap && dis[iter -> to] > dis[now] + iter -> cost)
22
23
                {
                    dis[iter -> to] = dis[now] + iter -> cost;
24
                    prev[iter -> to] = iter;
25
                    !inq[iter \rightarrow to] ? inq[q[++tail] = iter \rightarrow to] = 1 : 0;
26
27
28
       return dis[t] != inf;
29
30
   inline void mcmf()
31
32
33
       int x = inf;
34
       for (Edge *iter = prev[t]; iter; iter = prev[iter -> from]) cmin(x, iter -> cap);
       for (Edge *iter = prev[t]; iter; iter = prev[iter -> from])
35
36
            iter -> cap -= x;
37
            iter \rightarrow rev \rightarrow cap += x;
38
            ans += 111 * x * iter -> cost;
39
40
41
```

zkw(lhy)

```
int aug(int no, int res)
{
    if(no == ED)return mincost += 111 * pil * res, res;
    v[no] = 1;
```

```
int flow = 0;
       for(int i = son[no]; i != -1; i = edge[i].next)
6
7
           if(edge[i].f && !v[edge[i].y] && !edge[i].c)
                int d = aug(edge[i].y, min(res, edge[i].f));
9
                edge[i].f -= d, edge[i ^ 1].f += d, flow += d, res -= d;
10
                if(!res)return flow;
11
12
       return flow;
13
14
15 bool modlabel()
16
       long long d = 0x3f3f3f3f3f3f3f3f3f11;
17
       for(int i = 1; i <= cnt; i++)
18
           if(v[i])
19
           {
20
                for(int j = son[i]; j != -1; j = edge[j].next)
21
                    if(edge[j].f \&\& !v[edge[j].y] \&\& edge[j].c < d)d = edge[j].c;
22
           }
23
       if(d == 0x3f3f3f3f3f3f3f3f11)return 0;
       for(int i = 1; i <= cnt; i++)
25
           if(v[i])
26
           {
27
                for(int j = son[i]; j != -1; j = edge[j].next)
28
                    edge[j].c -= d, edge[j ^ 1].c += d;
29
30
       pil += d;
31
       return 1;
32
33
34
   void minimum_cost_flow_zkw()
35
       pil = 0;
36
       int nowans = 0;
37
       nowf = 0;
38
       do{
39
40
                for(int i = 1; i <= cnt; i++)
41
                    v[i] = 0;
42
                nowans = aug(ST, inf);
43
44
                nowf += nowans;
45
           }while(nowans);
       }while(modlabel());
46
47
```

1.19 有上下界的网络流 (Durandal)

B(u,v) 表示边 (u,v) 流量的下界,C(u,v) 表示边 (u,v) 流量的上界,设 F(u,v) 表示边 (u,v) 的实际流量设 G(u,v) = F(u,v) - B(u,v),则 $0 \le G(u,v) \le C(u,v) - B(u,v)$

- 无源汇的上下界可行流 建立超级源点 S^* 和超级汇点 T^* ,对于原图每一条边 (u,v) 在新网络中连如下三条边: $S^* \to v$,容量为 $B(u,v);\ u\to T^*$,容量为 $B(u,v);\ u\to v$,容量为 C(u,v)-B(u,v)。最后求新网络的最大流,判断从 超级源点 S^* 出发的边是否都满流即可,边 (u,v) 的最终解中的实际流量为 G(u,v)+B(u,v)。
- 有源汇的上下界可行流 从汇点 T 到源点 S 连一条上界为 ∞ ,下界为 0 的边。按照无源汇的上下界可行流一样做即可,流量即为 $T\to S$ 边上的流量。

1. Graph Theory 1.20. 差分约束

- 有源汇的上下界最大流
 - 在有源汇的上下界可行流中,从汇点 T 到源点 S 的边改为连一条上界为 ∞,下界为 x 的边。x 满足二分性质,找到最大的 x 使得新网络存在有源汇的上下界可行流即为原图的最大流。
 - 从汇点 T 到源点 S 连一条上界为 ∞ ,下界为 0 的边,变成无源汇的网络。按照无源汇的上下界可行流的方法,建立超级源点 S^* 与超级汇点 T^* ,求一遍 S^* → T^* 的最大流,再将从汇点 T 到源点 S 的这条边拆掉,求一次 S → T 的最大流即可。
- 有源汇的上下界最小流
 - 在有源汇的上下界可行流中,从汇点 T 到源点 S 的边改为连一条上界为 x,下界为 0 的边。x 满足二分性质,找到最小的 x 使得新网络存在有源汇的上下界可行流即为原图的最大流。
 - 按照无源汇的上下界可行流的方法,建立超级源点 S^* 与超级汇点 T^* ,求一遍 $S^* \to T^*$ 的最大流,但是注意不加上汇点 T 到源点 S 的这条边,即不使之改为无源汇的网络去求解。求完后,再加上那条汇点 T 到源点 S 的边,上界为 ∞ 的边。因为这条边的下界为 0,所以 S^* , T^* 无影响,再求一次 $S^* \to T^*$ 的最大流。若超级源点 S^* 出发的边全部满流,则 $T \to S$ 边上的流量即为原图的最小流,否则无解。

1.20 差分约束

1.21 图论知识 (gy,lhy)

弦图

弦图: 任意点数 ≥ 4 的环皆有弦的无向图

单纯点:与其相邻的点的诱导子图为完全图的点

完美消除序列:每次选择一个单纯点删去的序列

弦图必有完美消除序列

O(m+n) 求弦图的完美消除序列:每次选择未选择的标号最大的点,并将与其相连的点标号 +1,得到完美消除序列的反序

最大团数 = 最小染色数:按完美消除序列从后往前贪心地染色

最小团覆盖 = 最大点独立集:按完美消除序列从前往后贪心地选点加入点独立集

计数问题

• 有根树计数

$$\begin{aligned} a_1 &= 1 \\ a_{n+1} &= \frac{\sum\limits_{j=1}^n j \cdot a_j \cdot S_{n,j}}{n} \\ S_{n,j} &= \sum\limits_{i=1}^{n/j} a_{n+1-ij} = S_{n-j,j} + a_{n+1-j} \end{aligned}$$

• 无根树计数

$$\begin{cases} a_n - \sum_{i=1}^{n/2} a_i a_{n-i} & n \text{ is odd} \\ a_n - \sum_{i=1}^{n/2} a_i a_{n-i} + \frac{1}{2} a_{\frac{n}{2}} (a_{\frac{n}{2}} + 1) & n \text{ is even} \end{cases}$$

• 完全图生成树计数 n^{n-2}

• 矩阵-树定理

设 $\mathbf{A}[G]$ 为图 G 的邻接矩阵、 $\mathbf{D}[G]$ 为图 G 的度数矩阵,则图 G 的不同生成树的个数为 $\mathbf{C}[G] = \mathbf{D}[G] - \mathbf{A}[G]$ 的任意一个 n-1 阶主子式的行列式值。

• 偶数点完全图完备匹配计数 (n-1)!!

1.21. 图论知识 (gy,lhy) 1. Graph Theory

- 无根二叉树计数 (2n-5)!!
- 有根二叉树计数 (2n-3)!!

最大权闭合子图

给定一个带点权的有向图, 求其最大权闭合子图。

从源点 S 向每一条正权点连一条容量为权值的边,每个负权点向汇点 T 连一条容量为权值绝对值的边,有向图原来的边容量为 ∞ 。求它的最小割,与源点 S 连通的点构成最大权闭合子图,权值为正权值和 — 最小割。

最大密度子图

给定一个无向图,求其一个子图,使得子图的边数 |E| 和点数 |V| 满足 $\frac{|E|}{|V|}$ 最大。

二分答案 k,使得 $|E|-k|V| \ge 0$ 有解,将原图边和点都看作点,边 (u,v) 分别向 u 和 v 连边求最大权闭合 子图。

Chapter 2

Math

2.1 int64 相乘取模 (Durandal)

```
int64_t mul(int64_t x, int64_t y, int64_t p) {
   int64_t t = (x * y - (int64_t) ((long double) x / p * y + 1e-3) * p) % p;
   return t < 0 ? t + p : t;
}</pre>
```

2.2 扩展欧几里得 (gy)

```
// return gcd(a, b)
   // ax+by=gcd(a,b)
   int extend_gcd(int a, int b, int &x, int &y) {
       if (b == 0) \{
           x = 1, y = 0;
5
6
           return a;
       int res = extend_gcd(b, a % b, x, y);
       int t = y;
9
       y = x - a / b * y;
10
       x = t;
11
       return res;
12
13
   // return minimal positive integer x so that ax+by=c
   // or -1 if such x does not exist
   int solve_equ(int a, int b, int c) {
17
       int x, y, d;
       d = extend_gcd(a, b, x, y);
18
       if (c % d)
19
          return -1;
20
       int t = c / d;
21
       x *= t;
22
       y *= t;
23
       int k = b / d;
24
       x = (x \% k + k) \% k;
25
26
       return x;
27
   // return minimal positive integer x so that ax==b \pmod{p}
   // or -1 if such x does not exist
29
30 int solve(int a, int b, int p) {
      a = (a \% p + p) \% p;
31
       b = (b \% p + p) \% p;
```

```
return solve_equ(a, p, b);
34 }
```

2.3 中国剩余定理 (Durandal)

返回是否可行,余数和模数结果为 r_1, m_1

```
bool CRT(int &r1, int &m1, int r2, int m2) {
    int x, y, g = extend_gcd(m1, m2, x, y);
    if ((r2 - r1) % g != 0) return false;
    x = 111 * (r2 - r1) * x % m2;
    if (x < 0) x += m2;
    x /= g;
    r1 += m1 * x;
    m1 *= m2 / g;
    return true;
}</pre>
```

2.4 线性同余不等式 (Durandal)

必须满足 $0 \le d < m$, $0 \le l \le r < m$, 返回 $\min\{x \ge 0 \mid l \le x \cdot d \mod m \le r\}$, 无解返回 -1

```
int64_t calc(int64_t d, int64_t m, int64_t l, int64_t r) {
    if (1 == 0) return 0;
    if (d == 0) return -1;
    if (d * 2 > m) return calc(m - d, m, m - r, m - 1);
    if ((1 - 1) / d < r / d) return (1 - 1) / d + 1;
    int64_t k = calc((-m % d + d) % d, d, l % d, r % d);
    if (k == -1) return -1;
    return (k * m + 1 - 1) / d + 1;
}</pre>
```

2.5 组合数

2.6 高斯消元 (ct)

增广矩阵大小为 $m \times (n+1)$

```
db a[maxn][maxn], x[maxn];
  int main()
  |{
       int rank = 0;
       for (int i = 1, now = 1; i <= m && now <= n; ++now)
       {
           if (fabs(a[i][now]) < eps)</pre>
                for (int j = i + 1; j \le m; ++j)
9
                    if (fabs(a[j][now]) > fabs(a[i][now]))
10
11
                        for (int k = now; k \le n + 1; ++k)
12
                            std::swap(a[i][k], a[j][k]);
13
14
           }
15
           if (fabs(a[i][now]) < eps) continue;</pre>
16
17
           for (int j = i + 1; j \le m; ++j)
```

```
{
                db temp = a[j][now] / a[i][now];
19
                for (int k = now; k \le n + 1; ++k)
20
                    a[j][k] = temp * a[i][k];
21
22
           ++i; ++rank;
23
24
       if (rank == n)
25
26
           x[n] = a[n][n + 1] / a[n][n];
27
           for (int i = n - 1; i; --i)
28
29
                for (int j = i + 1; j \le n; ++j)
30
                    a[i][n + 1] -= x[j] * a[i][j];
31
                x[i] = a[i][n + 1] / a[i][i];
32
           }
33
       }
34
       else puts("Infinite Solution!");
35
       return 0;
36
37
```

2.7 Miller Rabin & Pollard Rho (gy)

In Java, use BigInteger.isProbablePrime(int certainty) to replace miller_rabin(BigInteger number)

Test Set	First Wrong Answer
2	2047
2,3	1,373,653
31,73	9,080,191
2, 3, 5	25, 326, 001
2, 3, 5, 7	(INT32_MAX)3,215,031,751
2, 7, 61	4,759,123,141
2, 13, 23, 1662803	1, 122, 004, 669, 633
2, 3, 5, 7, 11	2, 152, 302, 898, 747
2, 3, 5, 7, 11, 13	3, 474, 749, 660, 383
2, 3, 5, 7, 11, 13, 17	341, 550, 071, 728, 321
2, 3, 5, 7, 11, 13, 17, 19, 23	3,825,123,056,546,413,051
2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37	(INT64_MAX)318,665,857,834,031,151,167,461
2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41	3,317,044,064,679,887,385,961,981

```
const int test_case_size = 12;
   const int test_cases[test_case_size] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
   int64_t multiply_mod(int64_t x, int64_t y, int64_t p) {
       int64_t t = (x * y - (int64_t) ((long double) x / p * y + 1e-3) * p) % p;
       return t < 0 ? t + p : t;
   int64_t add_mod(int64_t x, int64_t y, int64_t p) {
       return (Oull + x + y) % p;
   int64_t power_mod(int64_t x, int64_t exp, int64_t p) {
10
       int64_t ans = 1;
11
       while (exp) {
12
           if (exp & 1)
13
               ans = multiply_mod(ans, x, p);
14
          x = multiply_mod(x, x, p);
15
```

```
exp >>= 1;
16
       }
17
18
       return ans;
  }
19
   bool miller_rabin_check(int64_t prime, int64_t base) {
20
       int64_t number = prime - 1;
21
       for (; ~number & 1; number >>= 1)
22
           continue;
23
       int64_t result = power_mod(base, number, prime);
24
       for (; number != prime - 1 && result != 1 && result != prime - 1; number <<= 1)
25
           result = multiply_mod(result, result, prime);
26
       return result == prime - 1 || (number & 1) == 1;
27
  }
28
   bool miller_rabin(int64_t number) {
29
       if (number < 2)
30
           return false;
31
       if (number < 4)
32
           return true;
33
       if (~number & 1)
35
           return false;
       for (int i = 0; i < test_case_size && test_cases[i] < number; i++)</pre>
36
           if (!miller_rabin_check(number, test_cases[i]))
37
               return false;
38
       return true;
39
   }
40
   int64_t gcd(int64_t x, int64_t y) {
41
42
       return y == 0 ? x : gcd(y, x % y);
43
   int64_t pollard_rho_test(int64_t number, int64_t seed) {
       int64_t x = rand() \% (number - 1) + 1, y = x;
45
       int head = 1, tail = 2;
46
       while (true) {
47
           x = multiply_mod(x, x, number);
48
           x = add_mod(x, seed, number);
49
           if (x == y)
50
               return number;
51
           int64_t answer = gcd(std::abs(x - y), number);
52
           if (answer > 1 && answer < number)
53
               return answer;
55
           if (++head == tail) {
56
               y = x;
               tail <<= 1;
57
           }
58
       }
59
   }
60
   void factorize(int64_t number, std::vector<int64_t> &divisor) {
61
       if (number > 1) {
62
           if (miller_rabin(number)) {
63
               divisor.push_back(number);
64
65
           } else {
               int64_t factor = number;
66
               while (factor >= number)
67
                    factor = pollard_rho_test(number, rand() % (number - 1) + 1);
68
               factorize(number / factor, divisor);
69
               factorize(factor, divisor);
70
           }
```

```
\begin{bmatrix} 72 \\ 73 \end{bmatrix}
```

2.8 $O(m^2 \log n)$ 线性递推

2.9 Polynomial

- 2.9.1 FFT
- 2.9.2 NTT & 多项式求逆
- 2.10 Lagrange 插值 (ct)

```
求解 \sum_{i=1}^{n} i^k \mod (10^9 + 7)
```

```
const int mod = 1e9 + 7;
1 int f[maxn], pre[maxn], suf[maxn], inp[maxn], p[maxn];
3 inline int qpow(int base, int power)
       int ret = 1;
       for (; power; power >>= 1, base = 111 * base * base \% mod)
           power & 1 ? ret = 111 * ret * base % mod : 0;
       return ret;
  }
9
  bool vis[maxn];
10
   int pr[maxn], prcnt, fpow[maxn];
11
   int main()
12
13
       int n = F(), k = F();
14
       // ******
15
       fpow[1] = 1;
16
       for (int i = 2; i \le k + 2; ++i)
17
18
           if (!vis[i]) pr[++prcnt] = i, fpow[i] = qpow(i, k);
19
           for (int j = 1; j <= prcnt && i * pr[j] <= k + 2; ++j)
20
21
               vis[i * pr[j]] = 1;
22
               fpow[i * pr[j]] = 111 * fpow[i] * fpow[pr[j]] % mod;
23
               if (i % pr[j] == 0) break;
           }
25
26
       }
       // ******** pre-processing
27
       for (int i = 1; i \le k + 2; ++i) f[i] = (f[i - 1] + fpow[i]) % mod;
28
       if (n \le k + 2) return !printf("%d\n", f[n]);
29
       pre[0] = 1;
30
       for (int i = 1; i \le k + 3; ++i) pre[i] = 111 * pre[i - 1] * (n - i) % mod;
31
       suf[k + 3] = 1;
32
       for (int i = k + 2; i >= 0; --i) suf[i] = 111 * suf[i + 1] * (n - i) % mod;
33
34
       for (int i = 1; i \le k + 2; ++i) p[i] = (111 * p[i - 1] * i) % mod;
35
       inp[k + 2] = qpow(p[k + 2], mod - 2);
36
       for (int i = k + 1; i \ge 0; --i) inp[i] = (111 * inp[i + 1] * (i + 1)) % mod;
37
       int ans = 0;
38
```

2.11. 社教筛 2. Math

```
for (int i = 1; i <= k + 2; ++i)
{
    int temp = inp[k + 2 - i]; if ((k + 2 - i) & 1) temp = mod - temp;
    int tmp = 1ll * pre[i - 1] * suf[i + 1] % mod * temp % mod * inp[i - 1] % mod * f[i] % mod;
    ans = (ans + tmp) % mod;
}

printf("%d\n", ans );
    return 0;
}</pre>
```

2.11 杜教筛

2.12 BSGS (ct,Durandal)

2.12.1 BSGS(ct)

p 是素数,返回 $\min\{x \ge 0 \mid y^x \equiv z \mod p\}$

```
1 const int mod = 19260817;
2 struct Hash
3 {
       Hash *next;
       int key, val;
  } *last[mod], mem[100000], *tot = mem;
6
  inline void insert(R int x, R int v)
   {
       *++tot = (Hash) {last[x \% mod], x, v}; last[x \% mod] = tot;
9
10
   inline int query(R int x)
11
12
       for (R Hash *iter = last[x % mod]; iter; iter = iter -> next)
13
           if (iter -> key == x) return iter -> val;
14
15
       return -1;
16
  inline void del(R int x)
17
18
       last[x \% mod] = 0;
19
20
  int main()
21
22
       for (; T; --T)
23
24
           R int y, z, p; scanf("%d%d%d", &y, &z, &p);
25
           R int m = (int) sqrt(p * 1.0);
26
           y %= p; z %= p;
27
           if (!y && !z) {puts("0"); continue;}
28
           if (!y) {puts("Orz, I cannot find x!"); continue;}
29
           R int pw = 1;
30
           for (R int i = 0; i < m; ++i, pw = 111 * pw * y % p) insert(111 * z * pw % p, i);
31
           R int ans = -1;
32
           for (R int i = 1, t, pw2 = pw; i \leq p / m + 1; ++i, pw2 = 111 * pw2 * pw % p)
33
               if ((t = query(pw2)) != -1) {ans = i * m - t; break;}
           if (ans == -1) puts("Orz, I cannot find x!");
35
           else printf("%d\n", ans );
36
           tot = mem; pw = 1;
37
           for (R int i = 0; i < m; ++i, pw = 111 * pw * y \% p) del(111 * z * pw \% p);
38
39
       return 0;
40
```

41 | }

2.12.2 扩展 BSGS(Durandal)

必须满足 $0 \le a < p$, $0 \le b < p$, 返回 $\min\{x \ge 0 \mid a^x \equiv b \mod p\}$

```
int64_t ex_bsgs(int64_t a, int64_t b, int64_t p) {
       if (b == 1)
           return 0;
       int64_t t, d = 1, k = 0;
       while ((t = std::__gcd(a, p)) != 1) {
           if (b \% t) return -1;
           k++, b /= t, p /= t, d = d * (a / t) % p;
           if (b == d) return k;
9
       map.clear();
10
       int64_t m = std::ceil(std::sqrt((long double) p));
11
12
       int64_t a_m = pow_mod(a, m, p);
13
       int64_t mul = b;
       for (int j = 1; j \le m; j++) {
14
           (mul *= a) %= p;
15
           map[mul] = j;
16
17
       for (int i = 1; i <= m; i++) {
18
           (d *= a_m) \% = p;
19
20
           if (map.count(d))
               return i * m - map[d] + k;
21
22
23
       return -1;
24
   int main() {
25
       int64_t a, b, p;
26
       while (scanf("%lld%lld", &a, &b, &p) != EOF)
27
           printf("%lld\n", ex_bsgs(a, b, p));
28
       return 0;
29
30
```

2.13 直线下整点个数 (gy)

必须满足 $a \ge 0, b \ge 0, m > 0$, 返回 $\sum_{i=0}^{n-1} \frac{a+bi}{m}$

```
int64_t count(int64_t n, int64_t a, int64_t b, int64_t m) {
   if (b == 0)
      return n * (a / m);
   if (a >= m)
      return n * (a / m) + count(n, a % m, b, m);
   if (b >= m)
      return (n - 1) * n / 2 * (b / m) + count(n, a, b % m, m);
   return count((a + b * n) / m, (a + b * n) % m, m, b);
}
```

2.14. 单纯形 2. Math

2.14 单纯形

2.15 辛普森积分

2.16 数学知识 (gy)

求和公式

•
$$\sum_{k=1}^{n} (2k-1)^2 = \frac{1}{3}n(4n^2-1)$$

•
$$\sum_{k=1}^{n} k^3 = \frac{1}{4}n^2(n+1)^2$$

•
$$\sum_{k=1}^{n} (2k-1)^3 = n^2(2n^2-1)$$

•
$$\sum_{k=1}^{n} k^4 = \frac{1}{30}n(n+1)(2n+1)(3n^2+3m-1)$$

•
$$\sum_{k=1}^{n} k^5 = \frac{1}{12}n^2(n+1)^2(2n^2+2n-1)$$

•
$$\sum_{k=1}^{n} k(k+1) = \frac{1}{3}n(n+1)(n+2)$$

•
$$\sum_{k=1}^{n} k(k+1)(k+2) = \frac{1}{4}n(n+1)(n+2)(n+3)$$

•
$$\sum_{k=1}^{n} k(k+1)(k+2)(k+3) = \frac{1}{5}n(n+1)(n+2)(n+3)(n+4)$$

错排公式

 D_n 表示 n 个元素错位排列的方案数

$$D_1 = 0, D_2 = 1$$

$$D_n = (n-1)(D_{n-2} + D_{n-1}), n \ge 3$$

$$D_n = n! \cdot (1 - \frac{1}{1!} + \frac{1}{2!} - \dots + (-1)^n \frac{1}{n!})$$

Fibonacci sequence

$$F_{0} = 0, F_{1} = 1$$

$$F_{n} = F_{n-1} + F_{n-2}$$

$$F_{n+1} \cdot F_{n-1} - F_{n}^{2} = (-1)^{n}$$

$$F_{-n} = (-1)^{n} F_{n}$$

$$F_{n+k} = F_{k} \cdot F_{n+1} + F_{k-1} \cdot F_{n}$$

$$\gcd(F_{m}, F_{n}) = F_{\gcd(m, n)}$$

$$F_{m} \mid F_{n}^{2} \Leftrightarrow nF_{n} \mid m$$

$$F_{n} = \frac{\varphi^{n} - \Psi^{n}}{\sqrt{5}}, \varphi = \frac{1 + \sqrt{5}}{2}, \Psi = \frac{1 - \sqrt{5}}{2}$$

$$F_{n} = \lfloor \frac{\varphi^{n}}{\sqrt{5}} + \frac{1}{2} \rfloor, n \ge 0$$

$$n(F) = \lfloor \log_{\wp}(F \cdot \sqrt{5} + \frac{1}{2}) \rfloor$$

Stirling number (1st kind)

用 $\begin{bmatrix}n\\k\end{bmatrix}$ 表示 Stirling number (1st kind),为将 n 个元素分成 k 个环的方案数 $\begin{bmatrix}n+1\\k\end{bmatrix}=n\begin{bmatrix}n\\k\end{bmatrix}+\begin{bmatrix}n\\k-1\end{bmatrix}, k>0$ $\begin{bmatrix}0\\0\end{bmatrix}=1,\begin{bmatrix}n\\0\end{bmatrix}=\begin{bmatrix}0\\n\end{bmatrix}=0, n>0$

2. Math 2.16. 数学知识 (gy)

$$\begin{bmatrix} n \\ k \end{bmatrix}$$
 为将 n 个元素分成 k 个环的方案数 $\begin{bmatrix} x \\ x-n \end{bmatrix} = \sum_{k=0}^{n} \left\langle {n \atop k} \right\rangle {x+k \choose 2n}$

Stirling number (2nd kind)

用 ${n \brace k}$ 表示 Stirling number (2nd kind),为将 n 个元素划分成 k 个非空集合的方案数 ${n+1 \brace k} = k \begin{Bmatrix} k \end{Bmatrix} + \begin{Bmatrix} n \cr k-1 \end{Bmatrix}, k > 0$ ${0 \brace 0} = 1, \begin{Bmatrix} n \cr 0 \end{Bmatrix} = \begin{Bmatrix} 0 \cr n \end{Bmatrix} = 0, n > 0$ ${n \brack k} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \brack j} j^n$ ${n \brack k}$ ${n \brack k} = \sum_{k=0}^{n} \left\langle {n \brack k} \right\rangle {x+n-k-1 \brack 2n}$

Catalan number

 c_n 表示长度为 2n 的合法括号序的数量 $c_1=1,\,c_{n+1}=\sum_{i=1}^n c_i\times c_{n+1-i}$ $c_n=\frac{\binom{2n}{n}}{n+1}$

Bell number

 B_n 表示基数为 n 的集合的划分方案数 $B_i = \begin{cases} 1 & i = 0\\ \sum\limits_{k=0}^{n} \binom{n}{k} B_k & i > 0 \end{cases}$ $B_n = \sum\limits_{k=0}^{n} \binom{n}{k}$

五边形数定理

p(n) 表示将 n 划分为若干个正整数之和的方案数 $p(n) = \sum_{k \in \mathbb{N}^*} (-1)^{k-1} p(n - \frac{k(3k-1)}{2})$

Bernoulli number

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0, m > 0$$

$$B_i = \begin{cases} 1 & i = 0 \\ -\sum_{j=0}^{i-1} {i+1 \choose j} B_j & i > 0 \end{cases}$$

$$\sum_{k=1}^{n} k^m = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_k n^{m+1-k}$$

Stirling permutation

1,1,2,2...,n,n 的排列中,对于每个 i,都有两个 i 之间的数大于 i 排列方案数为 (2n-1)!!

2.16. 数学知识 (gv) 2. Math

Eulerian number

Eulerian number (2nd kind)

Möbius function

$$\mu(n) = \begin{cases} 1 & n \text{ is a square-free positive integer with an even number of prime factors} \\ -1 & n \text{ is a square-free positive integer with an odd number of prime factors} \\ 0 & n \text{ has a squared prime factor} \end{cases}$$

$$\sum_{d|n} \mu(d) = \begin{cases} 1 & n=1 \\ 0 & n>1 \\ g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d) g(\frac{n}{d}) \end{cases}$$

Lagrange polynomial

给定次数为
$$n$$
 的多项式函数 $L(x)$ 上的 $n+1$ 个点 $(x_0,y_0),(x_1,y_1),\dots,(x_n,y_n)$ 则 $L(x)=\sum\limits_{j=0}^n y_j\prod\limits_{0\leq m\leq n,m\neq j}\frac{x-x_m}{x_j-x_m}$

Chapter 3

Geometry

3.1 点、直线、圆 (gy)

```
using number = long double;
  const number eps = 1e-8;
  number _sqrt(number x) {
      return std::sqrt(std::max(x, (number) 0));
  }
  number _asin(number x) {
      x = std::min(x, (number) 1), x = std::max(x, (number) -1);
      return std::asin(x);
9
  number _acos(number x) {
10
       x = std::min(x, (number) 1), x = std::max(x, (number) -1);
11
       return std::acos(x);
12
13
  int sgn(number x) {
       return (x > eps) - (x < -eps);
15
16
  int cmp(number x, number y) {
17
       return sgn(x - y);
18
19
   struct point {
20
       number x, y;
22
       point() {}
       point(number x, number y) : x(x), y(y) {}
       number len2() const {
24
           return x * x + y * y;
25
26
       number len() const {
27
           return _sqrt(len2());
28
29
       point unit() const {
30
           return point(x / len(), y / len());
31
32
       point rotate90() const {
33
           return point(-y, x);
34
35
       friend point operator+(const point &a, const point &b) {
36
           return point(a.x + b.x, a.y + b.y);
37
```

3.1. 点、直线、圆 (gy) 3. Geometry

```
friend point operator-(const point &a, const point &b) {
39
           return point(a.x - b.x, a.y - b.y);
40
41
       friend point operator*(const point &a, number b) {
42
           return point(a.x * b, a.y * b);
43
44
       friend point operator/(const point &a, number b) {
45
           return point(a.x / b, a.y / b);
46
47
       friend number dot(const point &a, const point &b) {
48
49
           return a.x * b.x + a.y * b.y;
50
       friend number det(const point \& a \,,\, const point \& b) {
51
           return a.x * b.y - a.y * b.x;
52
53
       friend number operator == (const point &a, const point &b) {
54
           return cmp(a.x, b.x) == 0 && cmp(a.y, b.y) == 0;
55
56
  };
57
  number dis2(const point &a, const point &b) {
       return (a - b).len2();
59
  }
60
  number dis(const point &a, const point &b) {
61
       return (a - b).len();
62
  }
63
   struct line {
64
       point a, b;
65
66
       line() {}
67
       line(point a, point b) : a(a), b(b) {}
       point value() const {
           return b - a;
69
70
71 };
72 bool point_on_line(const point &p, const line &l) {
       return sgn(det(p - 1.a, p - 1.b)) == 0;
73
74 }
   // including endpoint
  | bool point_on_ray(const point &p, const line &l) {
77
       return sgn(det(p - 1.a, p - 1.b)) == 0 &&
78
           sgn(dot(p - 1.a, 1.b - 1.a)) >= 0;
79
   // including endpoints
80
  bool point_on_seg(const point &p, const line &l) {
81
       return sgn(det(p - 1.a, p - 1.b)) == 0 &&
82
           sgn(dot(p - 1.a, 1.b - 1.a)) >= 0 &&
83
           sgn(dot(p - 1.b, 1.a - 1.b)) >= 0;
84
85
   bool seg_has_intersection(const line &a, const line &b) {
86
       if (point_on_seg(a.a, b) || point_on_seg(a.b, b) ||
87
               point_on_seg(b.a, a) || point_on_seg(b.b, a))
           return /* including endpoints */ true;
89
       return sgn(det(a.a - b.a, b.b - b.a)) * sgn(det(a.b - b.a, b.b - b.a)) < 0
90
           && sgn(det(b.a - a.a, a.b - a.a)) * sgn(det(b.b - a.a, a.b - a.a)) < 0;
91
92
  point intersect(const line &a, const line &b) {
93
       number s1 = det(a.b - a.a, b.a - a.a);
94
       number s2 = det(a.b - a.a, b.b - a.a);
```

3.1. 点、直线、圆 (gy)

```
return (b.a * s2 - b.b * s1) / (s2 - s1);
96
   }
97
   point projection(const point &p, const line &1) {
98
       return 1.a + (1.b - 1.a) * dot(p - 1.a, 1.b - 1.a) / (1.b - 1.a).len2();
99
100
   number dis(const point &p, const line &1) {
101
       return std::abs(det(p - 1.a, 1.b - 1.a)) / (1.b - 1.a).len();
102
103
   point symmetry_point(const point &a, const point &o) {
104
       return o + o - a;
105
106
   point reflection(const point &p, const line &l) {
107
108
       return symmetry_point(p, projection(p, 1));
   }
109
   struct circle {
110
       point o;
111
       number r;
112
       circle() {}
113
       circle(point o, number r) : o(o), r(r) {}
114
115 };
   bool intersect(const line &1, const circle &a, point &p1, point &p2) {
116
       number x = dot(1.a - a.o, 1.b - 1.a);
117
       number y = (1.b - 1.a).len2();
118
       number d = x * x - y * ((1.a - a.o).len2() - a.r * a.r);
119
       if (sgn(d) < 0) return false;</pre>
120
       point p = 1.a - (1.b - 1.a) * (x / y), delta = (1.b - 1.a) * (_sqrt(d) / y);
121
       p1 = p + delta, p2 = p - delta;
122
       return true;
123
124
125
   bool intersect(const circle &a, const circle &b, point &p1, point &p2) {
126
       if (a.o == b.o \&\& cmp(a.r, b.r) == 0)
           return /* value for coincident circles */ false;
127
       number s1 = (b.o - a.o).len();
128
       if (cmp(s1, a.r + b.r) > 0 \mid \mid cmp(s1, std::abs(a.r - b.r)) < 0)
129
            return false;
130
       number s2 = (a.r * a.r - b.r * b.r) / s1;
131
       number aa = (s1 + s2) / 2, bb = (s1 - s2) / 2;
132
       point p = (b.o - a.o) * (aa / (aa + bb)) + a.o;
133
       point delta = (b.o - a.o).unit().rotate90() * _sqrt(a.r * a.r - aa * aa);
134
       p1 = p + delta, p2 = p - delta;
135
       return true;
136
137
   }
138
   bool tangent(const point &p0, const circle &c, point &p1, point &p2) {
139
       number x = (p0 - c.o).len2();
       number d = x - c.r * c.r;
140
       if (sgn(d) < 0) return false;
141
       if (sgn(d) == 0)
142
            return /* value for point_on_line */ false;
143
       point p = (p0 - c.o) * (c.r * c.r / x);
144
       point delta = ((p0 - c.o) * (-c.r * \_sqrt(d) / x)).rotate90();
145
       p1 = c.o + p + delta;
146
       p2 = c.o + p - delta;
147
148
       return true;
149
   bool ex_tangent(const circle &a, const circle &b, line &l1, line &l2) {
150
       if (cmp(std::abs(a.r - b.r), (b.o - a.o).len()) == 0) {
151
            point p1, p2;
152
            intersect(a, b, p1, p2);
153
            11 = 12 = line(p1, p1 + (a.o - p1).rotate90());
154
```

```
155
            return true;
        } else if (cmp(a.r, b.r) == 0) {
156
            point dir = b.o - a.o;
157
            dir = (dir * (a.r / dir.len())).rotate90();
158
            11 = line(a.o + dir, b.o + dir);
159
            12 = line(a.o - dir, b.o - dir);
160
            return true;
161
        } else {
162
            point p = (b.o * a.r - a.o * b.r) / (a.r - b.r);
163
            point p1, p2, q1, q2;
164
            if (tangent(p, a, p1, p2) && tangent(p, b, q1, q2)) {
165
                11 = line(p1, q1);
166
                12 = line(p2, q2);
167
                return true;
168
            } else {
169
                return false;
170
            }
171
172
173
   bool in_tangent(const circle &a, const circle &b, line &11, line &12) {
174
        if (cmp(a.r + b.r, (b.o - a.o).len()) == 0) {
175
176
            point p1, p2;
            intersect(a, b, p1, p2);
177
            11 = 12 = line(p1, p1 + (a.o - p1).rotate90());
178
            return true;
179
        } else {
180
            point p = (b.o * a.r + a.o * b.r) / (a.r + b.r);
181
            point p1, p2, q1, q2;
182
            if (tangent(p, a, p1, p2) && tangent(p, b, q1, q2)) {
183
                11 = line(p1, q1);
184
                12 = line(p2, q2);
185
186
                return true;
            } else {
                return false;
188
189
        }
190
   }
191
```

3.2 平面最近点对 (Grimoire)

```
bool byY(P a,P b){return a.y<b.y;}</pre>
   LL solve(P *p,int l,int r){
       LL d=1LL<<62;
       if(l==r)
           return d;
       if(l+1==r)
           return dis2(p[1],p[r]);
       int mid=(1+r)>>1;
9
       d=min(solve(1,mid),d);
       d=min(solve(mid+1,r),d);
10
       vector<P>tmp;
11
       for(int i=1;i<=r;i++)</pre>
12
13
            if(sqr(p[mid].x-p[i].x) \le d)
                tmp.push_back(p[i]);
14
       sort(tmp.begin(),tmp.end(),byY);
15
       for(int i=0;i<tmp.size();i++)</pre>
16
            for(int j=i+1; j<tmp.size()&&j-i<10; j++)</pre>
17
                d=min(d,dis2(tmp[i],tmp[j]));
18
       return d;
19
```

20 }

3.3 凸包游戏 (Grimoire)

给定凸包, $O(n \log n)$ 完成询问:

- 点在凸包内
- 凸包外的点到凸包的两个切点
- 向量关于凸包的切点
- 直线与凸包的交点

传入凸包要求 1 号点为 Pair(x,y) 最小的

```
const int INF = 1000000000;
  struct Convex
  {
3
       int n;
       vector<Point> a, upper, lower;
5
       Convex(vector<Point> _a) : a(_a) {
           n = a.size();
           int ptr = 0;
           for(int i = 1; i < n; ++ i) if (a[ptr] < a[i]) ptr = i;</pre>
9
           for(int i = 0; i <= ptr; ++ i) lower.push_back(a[i]);</pre>
10
           for(int i = ptr; i < n; ++ i) upper.push_back(a[i]);</pre>
11
           upper.push_back(a[0]);
12
13
       int sign(long long x) { return x < 0 ? -1 : x > 0; }
14
       pair<long long, int> get_tangent(vector<Point> &convex, Point vec) {
15
           int 1 = 0, r = (int)convex.size() - 2;
16
           for(; 1 + 1 < r; ) {
17
               int mid = (1 + r) / 2;
18
               if (sign((convex[mid + 1] - convex[mid]).det(vec)) > 0) r = mid;
19
               else 1 = mid;
20
           }
21
           return max(make_pair(vec.det(convex[r]), r)
22
                , make_pair(vec.det(convex[0]), 0));
23
24
       void update_tangent(const Point &p, int id, int &i0, int &i1) {
25
           if ((a[i0] - p).det(a[id] - p) > 0) i0 = id;
26
           if ((a[i1] - p).det(a[id] - p) < 0) i1 = id;
27
28
       void binary_search(int 1, int r, Point p, int &i0, int &i1) {
29
           if (1 == r) return;
30
           update_tangent(p, 1 % n, i0, i1);
31
           int sl = sign((a[1 % n] - p).det(a[(1 + 1) % n] - p));
32
           for(; 1 + 1 < r; ) {
33
               int mid = (1 + r) / 2;
34
               int smid = sign((a[mid % n] - p).det(a[(mid + 1) % n] - p));
35
               if (smid == sl) l = mid;
36
               else r = mid;
37
           }
38
           update_tangent(p, r % n, i0, i1);
39
40
       int binary_search(Point u, Point v, int 1, int r) {
41
           int sl = sign((v - u).det(a[1 % n] - u));
42
           for(; 1 + 1 < r; ) {
43
               int mid = (1 + r) / 2;
44
               int smid = sign((v - u).det(a[mid % n] - u));
45
```

```
if (smid == sl) l = mid;
               else r = mid;
47
           }
48
           return 1 % n;
49
50
       // 判定点是否在凸包内, 在边界返回 true
51
       bool contain(Point p) {
52
           if (p.x < lower[0].x || p.x > lower.back().x) return false;
53
           int id = lower_bound(lower.begin(), lower.end()
54
               , Point(p.x, -INF)) - lower.begin();
55
           if (lower[id].x == p.x) {
56
               if (lower[id].y > p.y) return false;
57
           } else if ((lower[id - 1] - p).det(lower[id] - p) < 0) return false;</pre>
58
           id = lower_bound(upper.begin(), upper.end(), Point(p.x, INF)
59
               , greater<Point>()) - upper.begin();
60
           if (upper[id].x == p.x) {
61
               if (upper[id].y < p.y) return false;</pre>
62
           } else if ((upper[id - 1] - p).det(upper[id] - p) < 0) return false;</pre>
63
64
           return true;
65
       // 求点 p 关于凸包的两个切点, 如果在凸包外则有序返回编号
       // 共线的多个切点返回任意一个, 否则返回 false
67
       bool get_tangent(Point p, int &i0, int &i1) {
68
           if (contain(p)) return false;
69
           i0 = i1 = 0;
70
           int id = lower_bound(lower.begin(), lower.end(), p) - lower.begin();
71
           binary_search(0, id, p, i0, i1);
72
           binary_search(id, (int)lower.size(), p, i0, i1);
73
           id = lower_bound(upper.begin(), upper.end(), p
74
75
               , greater<Point>()) - upper.begin();
76
           binary_search((int)lower.size() - 1, (int)lower.size() - 1 + id, p, i0, i1);
77
           binary_search((int)lower.size() - 1 + id
               , (int)lower.size() - 1 + (int)upper.size(), p, i0, i1);
78
79
           return true;
       }
80
       // 求凸包上和向量 vec 叉积最大的点,返回编号,共线的多个切点返回任意一个
81
       int get_tangent(Point vec) {
82
           pair<long long, int> ret = get_tangent(upper, vec);
83
           ret.second = (ret.second + (int)lower.size() - 1) % n;
84
           ret = max(ret, get_tangent(lower, vec));
85
           return ret.second;
86
       }
87
       // 求凸包和直线 u,v 的交点, 如果无严格相交返回 false.
       //如果有则是和 (i,next(i)) 的交点,两个点无序,交在点上不确定返回前后两条线段其中之一
89
       bool get_intersection(Point u, Point v, int &i0, int &i1) {
90
           int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
91
           if (sign((v - u).det(a[p0] - u)) * sign((v - u).det(a[p1] - u)) < 0)  {
92
               if (p0 > p1) swap(p0, p1);
93
               i0 = binary_search(u, v, p0, p1);
94
               i1 = binary_search(u, v, p1, p0 + n);
95
               return true;
96
           } else {
97
               return false;
98
           }
99
       }
100
101
   ∣};
```

3.4 半平面交 (Grimoire)

```
int quad() const { return sgn(y) == 1 \mid \mid (sgn(y) == 0 \&\& sgn(x) >= 0);}
2
  };
3
   struct L{
4
       bool onLeft(const P &p) const { return sgn((b - a)*(p - a)) > 0; }
5
       L push() const{ // push out eps
6
           const double eps = 1e-10;
           P delta = (b - a).turn90().norm() * eps;
           return L(a - delta, b - delta);
9
10
  };
11
   bool sameDir(const L &10, const L &11) {
12
       return parallel(10, 11) && sgn((10.b - 10.a)^(11.b - 11.a)) == 1;
13
14
   bool operator < (const P &a, const P &b) {</pre>
15
       if (a.quad() != b.quad())
16
           return a.quad() < b.quad();</pre>
17
       else
18
           return sgn((a*b)) > 0;
19
20
21
   bool operator < (const L &10, const L &11) {</pre>
22
       if (sameDir(10, 11))
           return l1.onLeft(l0.a);
23
24
       else
           return (10.b - 10.a) < (11.b - 11.a);</pre>
25
26
   bool check(const L &u, const L &v, const L &w) {
27
       return w.onLeft(intersect(u, v));
28
29
   vector<P> intersection(vector<L> &1) {
30
       sort(1.begin(), 1.end());
31
       deque<L> q;
32
33
       for (int i = 0; i < (int)1.size(); ++i) {</pre>
           if (i && sameDir(l[i], l[i - 1])) {
34
35
                continue;
           }
36
           while (q.size() > 1
37
                && !check(q[q.size() - 2], q[q.size() - 1], l[i]))
38
                    q.pop_back();
39
           while (q.size() > 1
40
                && !check(q[1], q[0], l[i]))
41
                    q.pop_front();
42
           q.push_back(1[i]);
43
44
       while (q.size() > 2
45
           && !check(q[q.size() - 2], q[q.size() - 1], q[0]))
46
                q.pop_back();
47
       while (q.size() > 2
48
           && !check(q[1], q[0], q[q.size() - 1]))
49
                q.pop_front();
50
       vector<P> ret;
51
       for (int i = 0; i < (int)q.size(); ++i)</pre>
52
       ret.push_back(intersect(q[i], q[(i + 1) % q.size()]));
53
       return ret;
55
```

3.5 点在多边形内 (Grimoire)

```
bool inPoly(P p,vector<P>poly){
       int cnt=0;
       for(int i=0;i<poly.size();i++){</pre>
3
           P a=poly[i],b=poly[(i+1)%poly.size()];
           if(onSeg(p,L(a,b)))
               return false;
           int x=sgn(det(a,p,b));
           int y=sgn(a.y-p.y);
           int z=sgn(b.y-p.y);
10
           cnt+=(x>0&&y<=0&&z>0);
           cnt-=(x<0\&\&z<=0\&\&y>0);
11
       }
12
       return cnt;
13
14
```

3.6 最小圆覆盖 (Grimoire)

```
struct line{
       point p,v;
3
  point Rev(point v){return point(-v.y,v.x);}
  point operator*(line A,line B){
       point u=B.p-A.p;
       double t=(B.v*u)/(B.v*A.v);
       return A.p+A.v*t;
  }
   point get(point a,point b){
10
       return (a+b)/2;
11
12 }
point get(point a, point b, point c){
       if(a==b)return get(a,c);
14
       if(a==c)return get(a,b);
15
       if(b==c)return get(a,b);
16
17
       line ABO=(line)\{(a+b)/2, Rev(a-b)\};
18
       line BCO=(line)\{(c+b)/2,Rev(b-c)\};
       return ABO*BCO;
19
  }
20
   int main(){
21
       scanf("%d",&n);
22
       for(int i=1;i<=n;i++)scanf("%lf%lf",&p[i].x,&p[i].y);</pre>
23
       random_shuffle(p+1,p+1+n);
24
       0=p[1];r=0;
25
26
       for(int i=2;i<=n;i++){
            if(dis(p[i],0)<r+1e-6)continue;</pre>
27
            0=get(p[1],p[i]);r=dis(0,p[i]);
28
            for(int j=1;j<i;j++){</pre>
29
                if(dis(p[j],0)<r+1e-6)continue;</pre>
30
                0=get(p[i],p[j]);r=dis(0,p[i]);
31
                for(int k=1;k<j;k++){</pre>
32
                    if(dis(p[k],0)<r+1e-6)continue;</pre>
33
                    O=get(p[i],p[j],p[k]);r=dis(0,p[i]);
34
35
36
       }printf("%.21f %.21f %.21f\n",0.x,0.y,r);
37
       return 0;
38
39
```

3.7 最小球覆盖 (Grimoire)

```
bool equal(const double & x, const double & y) {
       return x + eps > y and y + eps > x;
2
   }
3
   double operator % (const Point & a, const Point & b) {
       return a.x * b.x + a.y * b.y + a.z * b.z;
5
6
   Point operator * (const Point & a, const Point & b) {
       return Point(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z, a.x * b.y - a.y * b.x);
10
   struct Circle {
       double r; Point o;
11
  };
12
   struct Plane {
13
       Point nor:
14
       double m;
15
       Plane(const Point & nor, const Point & a) : nor(nor){
16
           m = nor \% a;
17
18
  };
19
  Point intersect(const Plane & a, const Plane & b, const Plane & c) {
       Point c1(a.nor.x, b.nor.x, c.nor.x), c2(a.nor.y, b.nor.y, c.nor.y), c3(a.nor.z, b.nor.z,
21
         \rightarrow c.nor.z), c4(a.m, b.m, c.m);
       return 1 / ((c1 * c2) % c3) * Point((c4 * c2) % c3, (c1 * c4) % c3, (c1 * c2) % c4);
22
   }
23
   bool in(const Point & a, const Circle & b) {
24
       return sign((a - b.o).len() - b.r) <= 0;
25
26
   bool operator < (const Point & a, const Point & b) {</pre>
27
       if(!equal(a.x, b.x)) {
28
           return a.x < b.x;
29
30
       if(!equal(a.y, b.y)) {
31
32
           return a.y < b.y;
33
       if(!equal(a.z, b.z)) {
34
           return a.z < b.z;
35
36
       return false;
37
38
   bool operator == (const Point & a, const Point & b) {
39
       return equal(a.x, b.x) and equal(a.y, b.y) and equal(a.z, b.z);
40
41
   vector<Point> vec;
42
   Circle calc() {
43
       if(vec.empty()) {
44
           return Circle(Point(0, 0, 0), 0);
45
       }else if(1 == (int)vec.size()) {
46
           return Circle(vec[0], 0);
47
       }else if(2 == (int)vec.size()) {
48
           return Circle(0.5 * (vec[0] + vec[1]), 0.5 * (vec[0] - vec[1]).len());
49
       }else if(3 == (int)vec.size()) {
50
           double r((vec[0] - vec[1]).len() * (vec[1] - vec[2]).len() * (vec[2] - vec[0]).len() / 2 /
51
             \hookrightarrow fabs(((vec[0] - vec[2]) * (vec[1] - vec[2])).len()));
           return Circle(intersect(Plane(vec[1] - vec[0], 0.5 * (vec[1] + vec[0])),
52
                               Plane(vec[2] - vec[1], 0.5 * (vec[2] + vec[1])),
53
                        Plane((vec[1] - vec[0]) * (vec[2] - vec[0]), vec[0])), r);
54
       }else {
55
           Point o(intersect(Plane(vec[1] - vec[0], 0.5 * (vec[1] + vec[0])),
56
```

3.8. 圆并 (Grimoire) 3. Geometry

```
Plane(vec[2] - vec[0], 0.5 * (vec[2] + vec[0])),
57
                      Plane(vec[3] - vec[0], 0.5 * (vec[3] + vec[0])));
58
           return Circle(o, (o - vec[0]).len());
59
       }
60
   }
61
   Circle miniBall(int n) {
62
       Circle res(calc());
63
       for(int i(0); i < n; i++) {</pre>
64
            if(!in(a[i], res)) {
65
                vec.push_back(a[i]);
66
                res = miniBall(i);
67
                vec.pop_back();
                if(i) {
69
                    Point tmp(a[i]);
70
                    memmove(a + 1, a, sizeof(Point) * i);
71
                    a[0] = tmp;
72
                }
73
           }
74
75
       return res;
76
  }
77
   int main() {
79
       int n;
       sort(a, a + n);
80
       n = unique(a, a + n) - a;
81
       vec.clear():
82
       printf("%.10f\n", miniBall(n).r);
83
84
```

3.8 圆并 (Grimoire)

```
double ans[2001];
   struct Point {
2
       double x, y;
       Point(){}
       Point(const double & x, const double & y) : x(x), y(y) {}
       void scan() {scanf("%lf%lf", &x, &y);}
       double sqrlen() {return sqr(x) + sqr(y);}
       double len() {return sqrt(sqrlen());}
       Point rev() {return Point(y, -x);}
       void print() {printf("%f %f\n", x, y);}
       Point zoom(const double & d) {double lambda = d / len(); return Point(lambda * x, lambda * y);}
12 } dvd, a[2001];
Point centre[2001];
  double atan2(const Point & x) {
       return atan2(x.y, x.x);
15
16
  Point operator - (const Point \& a, const Point \& b) {
17
       return Point(a.x - b.x, a.y - b.y);
18
19
  Point operator + (const Point & a, const Point & b) {
20
       return Point(a.x + b.x, a.y + b.y);
22
  double operator * (const Point & a, const Point & b) {
23
24
       return a.x * b.y - a.y * b.x;
<sub>25</sub> | }
  Point operator * (const double & a, const Point & b) {
26
       return Point(a * b.x, a * b.y);
27
28 }
```

3. Geometry 3.8. 圆并 (Grimoire)

```
29 double operator % (const Point & a, const Point & b) {
       return a.x * b.x + a.y * b.y;
30
  }
31
   struct circle {
32
       double r; Point o;
33
       circle() {}
34
       void scan() {
35
           o.scan();
36
           scanf("%lf", &r);
37
       }
38
   } cir[2001];
39
40
   struct arc {
       double theta;
41
       int delta;
42
       Point p;
43
       arc() {};
44
       arc(const double & theta, const Point & p, int d) : theta(theta), p(p), delta(d) {}
45
  } vec[4444];
46
  int nV;
47
  inline bool operator < (const arc & a, const arc & b) {
       return a.theta + eps < b.theta;
49
50
51
  int cnt;
   inline void psh(const double t1, const Point p1, const double t2, const Point p2) {
52
       if(t2 + eps < t1)
53
           cnt++;
54
       vec[nV++] = arc(t1, p1, 1);
55
       vec[nV++] = arc(t2, p2, -1);
56
57
   inline double cub(const double & x) {
58
59
       return x * x * x;
60
   inline void combine(int d, const double & area, const Point & o) {
61
       if(sign(area) == 0) return;
62
       centre[d] = 1 / (ans[d] + area) * (ans[d] * centre[d] + area * o);
63
       ans[d] += area;
64
  ١}
65
   bool equal(const double & x, const double & y) {
66
       return x + eps> y and y + eps > x;
67
  }
68
   bool equal(const Point & a, const Point & b) {
69
       return equal(a.x, b.x) and equal(a.y, b.y);
70
71
  }
72
   bool equal(const circle & a, const circle & b) {
73
       return equal(a.o, b.o) and equal(a.r, b.r);
74
  }
   bool f[2001];
75
   int main() {
76
       int n, m, index;
77
       while(EOF != scanf("%d%d%d", &m, &n, &index)) {
78
79
           for(int i(0); i < m; i++) {
80
               a[i].scan();
81
           }
82
           for(int i(0); i < n; i++) {
83
               cir[i].scan();//n 个圆
84
85
           for(int i(0); i < n; i++) {//这一段在去重圆 能加速 删掉不会错
86
               f[i] = true;
87
               for(int j(0); j < n; j++) if(i != j) {</pre>
88
```

3.8. 圆并 (Grimoire) 3. Geometry

```
if(equal(cir[i], cir[j]) and i < j or !equal(cir[i], cir[j]) and cir[i].r <</pre>
                                                               \rightarrow cir[j].r + eps and (cir[i].o - cir[j].o).sqrlen() < sqr(cir[i].r - cir[j].r) +
                                                               → eps) {
                                                                     f[i] = false;
  90
                                                                     break;
 91
                                                         }
  92
                                             }
  93
                                  }
 94
                                  int n1(0);
  95
                                  for(int i(0); i < n; i++)
  96
                                             if(f[i])
  97
                                                         cir[n1++] = cir[i];
  98
                                 n = n1;//去重圆结束
  99
                                 fill(ans, ans + n + 1, 0);//ans[i] 表示被圆覆盖至少 i 次的面积
100
                                 fill(centre, centre + n + 1, Point(0, 0));//centre[i] 表示上面 ans[i] 部分的重心
101
                                 for(int i(0); i < m; i++)</pre>
102
                                             combine(0, a[i] * a[(i + 1) % m] * 0.5, 1. / 3 * (a[i] + a[(i + 1) % m]));
103
                                  for(int i(0); i < n; i++) {</pre>
104
                                             dvd = cir[i].o - Point(cir[i].r, 0);
105
                                             nV = 0;
106
                                             vec[nV++] = arc(-pi, dvd, 1);
107
                                             cnt = 0;
                                             for(int j(0); j < n; j++) if(j != i) {
109
                                                         double d = (cir[j].o - cir[i].o).sqrlen();
110
                                                         if(d < sqr(cir[j].r - cir[i].r) + eps) {
111
                                                                     if(cir[i].r + i * eps < cir[j].r + j * eps)
112
                                                                                psh(-pi, dvd, pi, dvd);
113
                                                         }else if(d + eps < sqr(cir[j].r + cir[i].r)) {</pre>
114
                                                                     double lambda = 0.5 * (1 + (sqr(cir[i].r) - sqr(cir[j].r)) / d);
115
                                                                     Point cp(cir[i].o + lambda * (cir[j].o - cir[i].o));
116
117
                                                                     Point nor((cir[j].o - cir[i].o).rev().zoom(sqrt(sqr(cir[i].r) - (cp -

    cir[i].o).sqrlen())));
                                                                     Point frm(cp + nor);
                                                                     Point to(cp - nor);
119
                                                                     psh(atan2(frm - cir[i].o), frm, atan2(to - cir[i].o), to);
120
                                                         }
121
                                             }
122
                                             sort(vec + 1, vec + nV);
123
                                             vec[nV++] = arc(pi, dvd, -1);
124
                                             for(int j = 0; j + 1 < nV; j++) {
125
                                                         cnt += vec[j].delta;
126
                                                          //if(cnt == 1) {//如果只算 ans[1] 和 centre[1], 可以加这个 if 加速.
127
                                                                     double theta(vec[j + 1].theta - vec[j].theta);
                                                                     double area(sqr(cir[i].r) * theta * 0.5);
129
                                                                     combine(cnt, area, cir[i].o + 1. / area / 3 * cub(cir[i].r) * Point(sin(vec[j +
130
                                                                           \rightarrow 1].theta) - sin(vec[j].theta), cos(vec[j].theta) - cos(vec[j + 1].theta)));
                                                                      combine(cnt, -sqr(cir[i].r) * sin(theta) * 0.5, 1. / 3 * (cir[i].o + vec[j].p + (cir[i].o + vec[i].p + (cir[i].o + vec[i].o + vec[i].p + (cir[i].o + vec[i].o + vec[i].p + (cir[i].o + vec[i].o + vec[i].o + 
131
                                                                           \hookrightarrow \text{vec}[j + 1].p));
                                                                     combine(cnt, vec[j].p * vec[j + 1].p * 0.5, 1. / 3 * (<math>vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1].p * 0.5, 1. / 3 * (vec[j].p + vec[j + 1]
132
                                                                           \hookrightarrow 1].p));
                                                         //}
133
                                             }
134
135
                                  combine(0, -ans[1], centre[1]);
                                  for(int i = 0; i < m; i++) {
137
                                             if(i != index)
138
                                                          (a[index] - Point((a[i] - a[index]) * (centre[0] - a[index]), (a[i] - a[index]) %
139
                                                               \hookrightarrow (centre[0] - a[index])).zoom((a[i] - a[index]).len())).print();
                                             else
140
                                                         a[i].print();
141
                                 }
142
```

```
143 }
144 return 0;
145 }
```

3.9 圆与多边形并 (Grimoire)

```
double form(double x){
       while(x \ge 2*pi)x = 2*pi;
3
       while(x<0)x+=2*pi;
       return x;
   }
   double calcCir(C cir){
       vector<double>ang;
       ang.push_back(0);
       ang.push_back(pi);
       double ans=0;
10
       for(int i=1;i<=n;i++){</pre>
11
            if(cir==c[i])continue;
12
            P p1,p2;
13
            if(intersect(cir,c[i],p1,p2)){
                 ang.push_back(form(cir.ang(p1)));
16
                 ang.push_back(form(cir.ang(p2)));
            }
17
       }
18
       for(int i=1;i<=m;i++){</pre>
19
            vector<P>tmp;
20
            tmp=intersect(poly[i],cir);
21
            for(int j=0;j<tmp.size();j++){</pre>
22
                 ang.push_back(form(cir.ang(tmp[j])));
23
            }
24
       }
25
       sort(ang.begin(),ang.end());
26
       for(int i=0;i<ang.size();i++){</pre>
27
            double t1=ang[i],t2=(i+1==ang.size()?ang[0]+2*pi:ang[i+1]);
28
            P p=cir.at((t1+t2)/2);
29
            int ok=1;
30
            for(int j=1; j<=n; j++) {</pre>
31
                 if(cir==c[j])continue;
32
                 if(inC(p,c[j],true)){
33
                     ok=0;
34
35
                     break;
                }
36
            }
37
             \texttt{for(int} \ j=1; j<=m\&\&ok; j++)\{
38
                if(inPoly(p,poly[j],true)){
39
                     ok=0;
40
                     break;
41
                }
42
            }
43
            if(ok){
44
45
                 double r=cir.r,x0=cir.o.x,y0=cir.o.y;
                 ans += (r*r*(t2-t1) + r*x0*(sin(t2) - sin(t1)) - r*y0*(cos(t2) - cos(t1)))/2;
46
            }
47
       }
48
       return ans;
49
   }
50
51 P st;
```

```
52 bool bySt(P a,P b){
       return dis(a,st) < dis(b,st);</pre>
53
54 }
   double calcSeg(L 1){
55
       double ans=0;
56
       vector<P>pt;
57
       pt.push_back(1.a);
58
       pt.push_back(1.b);
59
        for(int i=1;i<=n;i++){</pre>
60
            P p1,p2;
61
            if(intersect(c[i],1,p1,p2)){
62
63
                 if(onSeg(p1,1))
                     pt.push_back(p1);
64
                 if(onSeg(p2,1))
65
                     pt.push_back(p2);
66
            }
67
       }
68
       st=l.a;
69
       sort(pt.begin(),pt.end(),bySt);
70
       for(int i=0;i+1<pt.size();i++){</pre>
71
            P p1=pt[i],p2=pt[i+1];
72
            P p=(p1+p2)/2;
73
            int ok=1;
74
            for(int j=1; j<=n; j++){</pre>
75
                 if(sgn(dis(p,c[j].o),c[j].r)<0){
76
                     ok=0;
77
                     break;
78
                 }
79
            }
80
81
            if(ok){
82
                 double x1=p1.x,y1=p1.y,x2=p2.x,y2=p2.y;
83
                 double res=(x1*y2-x2*y1)/2;
84
                 ans+=res;
            }
85
       }
86
       return ans;
87
```

3.10 三角剖分 (Grimoire)

Triangulation:: find 返回包含某点的三角形 $Triangulation:: add_point$ 将某点加入三角剖分 某个 Triangle 在三角剖分中当且仅当它的 $has_children$ 为 0 如果要找到三角形 u 的邻域,则枚举它的所有 u.edge[i].tri,该条边的两个点为 u.p[(i+1)%3],u.p[(i+2)%3] 通过三角剖分构造 V 图:连接相邻三角形外接圆圆心 注意初始化内存池和 Triangulation:: LOTS 复杂度 $O(n\log n)$

```
const int N = 100000 + 5, MAX_TRIS = N * 6;
const double eps = 1e-6, PI = acos(-1.0);
struct P {
    double x,y; P():x(0),y(0){}
    P(double x, double y):x(x),y(y){}
    bool operator ==(P const& that)const {return x==that.x&&y==that.y;}
};
inline double sqr(double x) { return x*x; }
double dist_sqr(P const& a, P const& b){return sqr(a.x-b.x)+sqr(a.y-b.y);}
bool in_circumcircle(P const& p1, P const& p2, P const& p3, P const& p4) {//p4 in C(p1,p2,p3)}
double u11 = p1.x - p4.x, u21 = p2.x - p4.x, u31 = p3.x - p4.x;
```

```
double u12 = p1.y - p4.y, u22 = p2.y - p4.y, u32 = p3.y - p4.y;
12
      double u13 = sqr(p1.x) - sqr(p4.x) + sqr(p1.y) - sqr(p4.y);
13
      double u23 = sqr(p2.x) - sqr(p4.x) + sqr(p2.y) - sqr(p4.y);
14
      double u33 = sqr(p3.x) - sqr(p4.x) + sqr(p3.y) - sqr(p4.y);
15
      16
         \rightarrow u11*u22*u33;
      return det > eps;
17
18
   double side(P const& a, P const& b, P const& p) { return (b.x-a.x)*(p.y-a.y) -
19
    \hookrightarrow (b.y-a.y)*(p.x-a.x);}
   typedef int SideRef; struct Triangle; typedef Triangle* TriangleRef;
20
   struct Edge {
21
      TriangleRef tri; SideRef side; Edge() : tri(0), side(0) {}
22
      Edge(TriangleRef tri, SideRef side) : tri(tri), side(side) {}
23
  }:
24
   struct Triangle {
25
      P p[3]; Edge edge[3]; TriangleRef children[3]; Triangle() {}
26
      Triangle(P const& p0, P const& p1, P const& p2) {
27
           p[0] = p0; p[1] = p1; p[2] = p2;
28
           children[0] = children[1] = children[2] = 0;
29
30
      bool has_children() const { return children[0] != 0; }
31
      int num_children() const {
32
          return children[0] == 0 ? 0
33
               : children[1] == 0 ? 1
34
               : children[2] == 0 ? 2 : 3;
35
36
      bool contains(P const& q) const {
37
           double a=side(p[0],p[1],q), b=side(p[1],p[2],q), c=side(p[2],p[0],q);
38
           return a >= -eps && b >= -eps && c >= -eps;
39
40
   } triange_pool[MAX_TRIS], *tot_triangles;
41
42
   void set_edge(Edge a, Edge b) {
      if (a.tri) a.tri->edge[a.side] = b;
43
      if (b.tri) b.tri->edge[b.side] = a;
44
  }
45
   class Triangulation {
46
      public:
47
           Triangulation() {
48
               const double LOTS = 1e6;//初始为极大三角形
49
               the_root = new(tot_triangles++) Triangle(P(-LOTS,-LOTS),P(+LOTS,-LOTS));
50
           }
51
           TriangleRef find(P p) const { return find(the_root,p); }
52
           void add_point(P const& p) { add_point(find(the_root,p),p); }
53
      private:
54
55
           TriangleRef the_root;
           static TriangleRef find(TriangleRef root, P const& p) {
56
               for(;;) {
57
                   if (!root->has_children()) return root;
58
                   else for (int i = 0; i < 3 && root->children[i]; ++i)
59
                           if (root->children[i]->contains(p))
60
                               {root = root->children[i]; break;}
61
62
           }
63
           void add_point(TriangleRef root, P const& p) {
64
               TriangleRef tab,tbc,tca;
65
               tab = new(tot_triangles++) Triangle(root->p[0], root->p[1], p);
66
               tbc = new(tot_triangles++) Triangle(root->p[1], root->p[2], p);
67
               tca = new(tot_triangles++) Triangle(root->p[2], root->p[0], p);
68
               set_edge(Edge(tab,0),Edge(tbc,1)); set_edge(Edge(tbc,0),Edge(tca,1));
69
               set_edge(Edge(tca,0),Edge(tab,1)); set_edge(Edge(tab,2),root->edge[2]);
70
```

```
set_edge(Edge(tbc,2),root->edge[0]); set_edge(Edge(tca,2),root->edge[1]);
71
               root->children[0]=tab; root->children[1]=tbc; root->children[2]=tca;
72
               flip(tab,2); flip(tbc,2); flip(tca,2);
73
           }
74
           void flip(TriangleRef tri, SideRef pi) {
75
               TriangleRef trj = tri->edge[pi].tri; int pj = tri->edge[pi].side;
76
               if(!trj || !in_circumcircle(tri->p[0],tri->p[1],tri->p[2],trj->p[pj])) return;
77
               TriangleRef trk = new(tot_triangles++) Triangle(tri->p[(pi+1)%3], trj->p[pj],
78

    tri->p[pi]);
               TriangleRef trl = new(tot_triangles++) Triangle(trj->p[(pj+1)%3], tri->p[pi],
79
                 → trj->p[pj]);
               set_edge(Edge(trk,0), Edge(trl,0));
80
               set\_edge(Edge(trk,1), tri->edge[(pi+2)\%3]); set\_edge(Edge(trk,2), trj->edge[(pj+1)\%3]);
81
               set\_edge(Edge(trl,1), trj->edge([pj+2)\%3]); set\_edge(Edge(trl,2), tri->edge([pi+1)\%3]);
82
               tri->children[0]=trk; tri->children[1]=trl; tri->children[2]=0;
83
               trj->children[0]=trk; trj->children[1]=trl; trj->children[2]=0;
84
               flip(trk,1); flip(trk,2); flip(trl,1); flip(trl,2);
85
86
87
  };
  int n; P ps[N];
   void build(){
       tot_triangles = triange_pool; cin >> n;
       for(int i = 0; i < n; ++ i) scanf("%lf%lf",&ps[i].x,&ps[i].y);</pre>
91
       random_shuffle(ps, ps + n); Triangulation tri;
92
       for(int i = 0; i < n; ++ i) tri.add_point(ps[i]);</pre>
93
94
```

3.11 三维几何基础 (Grimoire)

```
struct P {
        double x, y, z;
2
        P(){}
        P(double _x, double _y, double _z):x(_x),y(_y),z(_z){}
        double len2(){
            return (x*x+y*y+z*z);
6
        double len(){
            return sqrt(x*x+y*y+z*z);
10
   };
11
   bool operator==(P a,P b){
        return sgn(a.x-b.x)==0 && sgn(a.y-b.y)==0 && sgn(a.z-b.z)==0;
  | }
  bool operator<(P a,P b){</pre>
15
        \texttt{return sgn}(\texttt{a.x-b.x}) ~?~ \texttt{a.x<b.x} ~: (\texttt{sgn}(\texttt{a.y-b.y})?\texttt{a.y<b.y} ~: \texttt{a.z<b.z});
16
  }
17
   P operator+(P a,P b){
18
        return P(a.x+b.x,a.y+b.y,a.z+b.z);
19
20
   P operator-(P a,P b){
21
        return P(a.x-b.x,a.y-b.y,a.z-b.z);
22
23
   P operator*(P a,double b){
24
25
        return P(a.x*b,a.y*b,a.z*b);
  }
26
  P operator/(P a,double b){
27
        return P(a.x/b,a.y/b,a.z/b);
28
29
30 P operator*(const P &a, const P &b) {
```

```
return P(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z, a.x * b.y - a.y * b.x);
31
         l٦
32
           double operator^(const P &a, const P &b) {
33
                         return a.x*b.x+a.y*b.y+a.z*b.z;
34
35
           double dis(P a,P b){return (b-a).len();}
36
           double dis2(P a,P b){return (b-a).len2();}
37
           // 3D line intersect
38
          P intersect(const P &a0, const P &b0, const P &a1, const P &b1) {
39
                         double t = ((a0.x - a1.x) * (a1.y - b1.y) - (a0.y - a1.y) * (a1.x - b1.x)) / ((a0.x - b0.x) * (a1.y - b1.y) + (a1.x - b1.x)) / ((a0.x - b0.x) * (a1.y - b1.y) + (a1.x - b1.x)) / ((a0.x - b0.x) * (a1.y - b1.y) + (a1.x - b1.x)) / ((a0.x - b0.x) * (a1.y - b1.y) + (a1.x - b1.x)) / ((a0.x - b0.x) * (a1.y - b1.y) + (a1.x - b1.x)) / ((a0.x - b0.x) * (a1.y - b1.y) + (a1.x - b1.x)) / ((a0.x - b0.x) * (a1.x - b1.x)) / ((a0.x - b1.x) * (a1.x - b1.x)) / (
40
                                \rightarrow (a1.y - b1.y) - (a0.y - b0.y) * (a1.x - b1.x));
                         return a0 + (b0 - a0) * t;
41
         l٦
42
           // area-line intersect
43
         P intersect(const P &a, const P &b, const P &c, const P &10, const P &11) {
44
                         P p = (b-a)*(c-a); // 平面法向量
45
                         double t = (p^(a-10)) / (p^(11-10));
46
                         return 10 + (11 - 10) * t;
47
```

3.12 三维凸包 (Grimoire)

```
int mark[1005][1005],n, cnt;;
   double mix(const P &a, const P &b, const P &c) {
2
       return a^(b*c);
3
   double area(int a, int b, int c) {
5
       return ((info[b] - info[a])*(info[c] - info[a])).len();
6
   double volume(int a, int b, int c, int d) {
       return mix(info[b] - info[a], info[c] - info[a], info[d] - info[a]);
9
  }
10
   struct Face {
11
       int a, b, c; Face() {}
12
       Face(int a, int b, int c): a(a), b(b), c(c) {}
13
       int &operator [](int k) {
14
           if (k == 0) return a; if (k == 1) return b; return c;
15
16
17 };
   vector <Face> face;
   inline void insert(int a, int b, int c) {
19
       face.push_back(Face(a, b, c));
20
  }
21
   void add(int v) {
22
       vector <Face> tmp; int a, b, c; cnt++;
23
       for (int i = 0; i < SIZE(face); i++) {</pre>
24
           a = face[i][0]; b = face[i][1]; c = face[i][2];
25
           if (sgn(volume(v, a, b, c)) < 0)
26
           mark[a][b] = mark[b][a] = mark[b][c] = mark[c][b] = mark[c][a] = mark[a][c] = cnt;
27
           else tmp.push_back(face[i]);
28
       } face = tmp;
29
       for (int i = 0; i < SIZE(tmp); i++) {</pre>
30
           a = face[i][0]; b = face[i][1]; c = face[i][2];
31
           if (mark[a][b] == cnt) insert(b, a, v);
32
           if (mark[b][c] == cnt) insert(c, b, v);
33
           if (mark[c][a] == cnt) insert(a, c, v);
34
```

```
}
35
  }
36
  int Find() {
37
       for (int i = 2; i < n; i++) {
38
           P ndir = (info[0] - info[i])*(info[1] - info[i]);
39
           if (ndir == P()) continue; swap(info[i], info[2]);
40
           for (int j = i + 1; j < n; j++) if (sgn(volume(0, 1, 2, j)) != 0) {
41
               swap(info[j], info[3]); insert(0, 1, 2); insert(0, 2, 1); return 1;
42
43
       }
44
       return 0;
45
46
   //find the weight center
47
  double calcDist(const P &p, int a, int b, int c) {
48
       return fabs(mix(info[a] - p, info[b] - p, info[c] - p) / area(a, b, c));
49
50 | }
   //compute the minimal distance of center of any faces
51
  P findCenter() { //compute center of mass
52
       double totalWeight = 0;
53
       P center(.0, .0, .0);
       P first = info[face[0][0]];
       for (int i = 0; i < SIZE(face); ++i) {</pre>
           P p = (info[face[i][0]]+info[face[i][1]]+info[face[i][2]]+first)*.25;
57
           double weight = mix(info[face[i][0]] - first, info[face[i][1]] - first, info[face[i][2]] -

    first):
           totalWeight += weight; center = center + p * weight;
59
60
       center = center / totalWeight;
61
       return center;
62
63
64
   double minDis(P p) {
65
       double res = 1e100; //compute distance
       for (int i = 0; i < SIZE(face); ++i)</pre>
           res = min(res, calcDist(p, face[i][0], face[i][1], face[i][2]));
67
       return res:
68
  l٦
69
   void findConvex(P *info,int n) {
70
       sort(info, info + n); n = unique(info, info + n) - info;
71
       face.clear(); random_shuffle(info, info + n);
72
       if(!Find())return abort();
73
       memset(mark, 0, sizeof(mark)); cnt = 0;
75
       for (int i = 3; i < n; i++) add(i);
76
```

3.13 三维绕轴旋转 (Grimoire)

右手大拇指指向 axis 方向, 四指弯曲方向旋转 w 弧度

```
Protate(const P& s, const P& axis, double w) {
    double x = axis.x, y = axis.y, z = axis.z;
    double s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
        cosw = cos(w), sinw = sin(w);
    double a[4][4];
    memset(a, 0, sizeof a);
    a[3][3] = 1;
    a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
    a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
    a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
    a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
```

3. Geometry 3.14. 几何知识 (gy)

```
a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
12
       a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
13
       a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
14
       a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
15
       a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;
16
       double ans [4] = \{0, 0, 0, 0\}, c[4] = \{s.x, s.y, s.z, 1\};
17
       for (int i = 0; i < 4; ++ i)
18
           for (int j = 0; j < 4; ++ j)
19
               ans[i] += a[j][i] * c[j];
20
       return P(ans[0], ans[1], ans[2]);
21
```

3.14 几何知识 (gy)

Pick theorem

顶点为整点的简单多边形,其面积 A,内部格点数 i,边上格点数 b 满足: $A=i+\frac{b}{2}-1$

欧拉示性数

- 三维凸包的顶点个数 V,边数 E,面数 F 满足: V-E+F=2
- 平面图的顶点个数 V , 边数 E , 平面被划分的区域数 F , 组成图形的连通部分的数目 C 满足 : V-E+F=C+1

几何公式

```
• 三角形
    半周长 p = \frac{a+b+c}{2}
    面积 S = \frac{1}{2}aH_a = \frac{1}{2}ab \cdot \sin C = \sqrt{p(p-a)(p-b)(p-c)} = pr = \frac{abc}{4R}
    中线长 M_a = \frac{1}{2}\sqrt{2(b^2+c^2)-a^2} = \frac{1}{2}\sqrt{b^2+c^2+2bc\cdot\cos A}
   角平分线长 T_a = \frac{\sqrt{bc((b+c)^2 - a^2)}}{b+c} = \frac{2bc}{b+c} \cos \frac{A}{2} 高 H_a = b \sin C = \sqrt{b^2 - (\frac{a^2 + b^2 - c^2}{2a})^2}
    内切圆半径 r=\frac{S}{p}=4R\sin\frac{A}{2}\sin\frac{B}{2}\sin\frac{C}{2}=\sqrt{\frac{(p-a)(p-b)(p-c)}{p}}=p\tan\frac{A}{2}\tan\frac{B}{2}\tan\frac{C}{2}
    外接圆半径 R = \frac{abc}{4S} = \frac{a}{2\sin A} 旁切圆半径 r_A = \frac{2S}{-a+b+c}
    重心 \left(\frac{x_1+x_2+x_3}{3}, \frac{y_1+y_2+y_3}{3}\right)
                   x_1^2 + y_1^2 \quad y_1 \quad 1
                                                        x_1 x_1^2 + y_1^2 1
                  \begin{array}{cccc} x_2^{2} + y_2^{2} & y_2 & 1 \\ x_3^{2} + y_3^{2} & y_3 & 1 \end{array}
                                                        x_2 x_2^2 + y_2^2 1
                                                        x_3 \quad x_3^2 + y_3^2
    外心 (<sup>1</sup>
                         x_1 \quad y_1 \quad 1
                                                              x_1
                                                                      y_1
                   |x_2| |x_2| |y_2| |1
                                                             x_2 \ y_2 \ 1
                     \begin{vmatrix} x_3 & y_3 & 1 \end{vmatrix}
                                                                    y_3 1
                                                             x_3
    内心 \left(\frac{ax_1+bx_2+cx_3}{a+b+c}, \frac{ay_1+by_2+cy_3}{a+b+c}\right)
                  x_2x_3 + y_2y_3 \quad 1 \quad y_1
                                                                x_2x_3 + y_2y_3 \quad x_1
                                                                                                  1
                   x_3x_1 + y_3y_1 1
                                                                x_3x_1 + y_3y_1
                                                  y_2
                  x_1x_2 + y_1y_2
                                                                x_1x_2 + y_1y_2
                                                                                          x_3
                                                   y_3
                                                                                           1
                                             1
                            x_1
                                    y_1
                                                                          x_1
                                                                                  y_1
                           x_2 \ y_2 \ 1
                                                                     2 x_2 y_2
                                                                                           1
                           x_3 \quad y_3 \quad 1
                                                                                          1
                                                                       x_3 y_3
    旁心 \left(\frac{-ax_1+bx_2+cx_3}{-a+b+c}, \frac{-ay_1+by_2+cy_3}{-a+b+c}\right)
```

3.14. 几何知识 (gy) 3. Geometry

• 圆

弧长 l = rA弦长 $a = 2\sqrt{2hr - h^2} = 2r \cdot \sin \frac{A}{2}$ 弓形高 $h = r - \sqrt{r^2 - \frac{a^2}{4}} = r(1 - \cos \frac{A}{2})$ 扇形面积 $S_1 = \frac{1}{2}lr = \frac{1}{2}Ar^2$ 弓形面积 $S_2 = \frac{1}{2}r^2(A - \sin A)$

• Circles of Apollonius

棱台

体积 $V = \frac{1}{3}h(A_1 + A_2 + \sqrt{A_1A_2})$ 正棱台侧面积 $S = \frac{1}{2}(p_1 + p_2)l$, l 为侧高

球

体积 $V = \frac{4}{3}\pi r^3$ 表面积 $S = 4\pi r^2$

球台

侧面积 $S=2\pi rh$ 体积 $V=\frac{1}{6}\pi h(3(r_1^2+r_2^2)+h_h)$

• 球扇形

球面面积 $S=2\pi rh$ 体积 $V=\frac{2}{3}\pi r^2h=\frac{2}{3}\pi r^3h(1-\cos\varphi)$

• 球面三角形

考虑单位球上的球面三角形,a,b,c 表示三边长(弧所对球心角),A,B,C 表示三角大小(切线夹角) 余弦定理 $\cos a = \cos b \cdot \cos c + \sin a \cdot \sin b \cdot \cos A$ 正弦定理 $\frac{\sin A}{\sin a} = \frac{\sin B}{\sin b} = \frac{\sin C}{\sin c}$ 球面面积 $S = A + B + C - \pi$

• 四面体

体积 $V = \frac{1}{6} \left| \overrightarrow{AB} \cdot (\overrightarrow{AC} \times \overrightarrow{AD}) \right|$

Chapter 4

String

4.1 KMP (ct)

KMP

```
int main()
{
    for (int i = 2, j = 0; i <= n; ++i)
    {
        for (; j && s[j + 1] != s[i]; j = fail[j]);
        s[i] == s[j + 1] ? ++j : 0;
        fail[i] = j;
    }
    return 0;
}</pre>
```

exKMP

 $extend_i$ 表示 T 与 $S_{i,n}$ 的最长公共前缀

```
int next[maxn], extend[maxn], fail[maxn];
void getnext(R char *s, R int len)
  {
       fail[1] = 0;
       R int p = 0;
       memset(next, 0, (len + 2) << 2);
       for (R int i = 2; i <= len; ++i)
           while (p \&\& s[p + 1] != s[i]) p = fail[p];
           s[p + 1] == s[i] ? ++p : 0;
10
           fail[i] = p;
11
           p ? cmax(next[i - p + 1], p) : 0;
12
13
14
   void getextend(R char *s, R int lens, R char *t, R int lent)
15
16
17
       getnext(t, lent);
       R int a = 1, p = 0;
18
       for (R int i = 1; i <= lens; ++i)</pre>
19
20
           if (i + next[i - a + 1] - 1 >= p)
21
22
               cmax(p, i - 1);
23
```

4.2. AC 自动机 4. String

```
while (p < lens && p - i + 1 < lent && s[p + 1] == t[p - i + 2]) ++p;
a = i;
extend[i] = p - i + 1;
}
else extend[i] = next[i - a + 1];
}
</pre>
```

4.2 AC 自动机

4.3 后缀数组 (ct)

```
char s[maxn];
  int sa[maxn], rank[maxn], wa[maxn], wb[maxn], cnt[maxn], height[maxn];
  inline void build(int n, int m)
       int *x = wa, *y = wb, *t;
       for (int i = 1; i <= n; ++i) cnt[x[i] = s[i] - 'a' + 1]++;
       for (int i = 1; i <= m; ++i) cnt[i] += cnt[i - 1];
       for (int i = n; i; --i) sa[cnt[x[i]]--] = i;
       for (int j = 1; j < n \mid | (j == 1 \&\& m < n); j <<= 1, t = x, x = y, y = t)
9
10
           memset(cnt + 1, 0, m << 2);
11
           int p = 0;
12
           for (int i = n - j + 1; i \le n; ++i) y[++p] = i;
13
           for (int i = 1; i <= n; ++i)
14
15
               ++cnt[x[i]];
16
               sa[i] > j ? y[++p] = sa[i] - j : 0;
17
18
           for (int i = 1; i <= m; ++i) cnt[i] += cnt[i - 1];
20
           for (int i = n; i; --i) sa[cnt[x[y[i]]]--] = y[i];
21
                   m = 0;
           for (int i = 1; i <= n; ++i)
22
               y[sa[i]] = (i == 1 \mid | x[sa[i]] != x[sa[i - 1]] \mid | x[sa[i - 1] + j] != x[sa[i] + j])?
23
                 \hookrightarrow ++m : m;
24
       for (int i = 1; i <= n; ++i) rank[sa[i]] = i;
25
       for (int i = 1, j, k = 0; i <= n; height[rank[i++]] = k)</pre>
26
27
           for (k ? --k : 0, j = sa[rank[i] - 1]; s[i + k] == s[j + k]; ++k);
```

4.4 后缀自动机 (ct,lhy)

后缀自动机 (lhy)

4. String 4.4. 后缀自动机 (ct,lhy)

```
else
11
         {
12
               SAM *q = p \rightarrow next[c];
13
               if (q \rightarrow val == p \rightarrow val + 1) np \rightarrow fa = q;
14
               else
15
                {
16
                     SAM *nq = ++tot;
17
                     memcpy(nq -> next, q -> next, sizeof nq -> next);
18
                     nq \rightarrow val = p \rightarrow val + 1;
19
20
                     nq \rightarrow fa = q \rightarrow fa;
21
                     q \rightarrow fa = np \rightarrow fa = nq;
                     for (; p \&\& p \rightarrow next[c] == q; p = p \rightarrow fa) p \rightarrow next[c] = nq;
22
               }
23
         }
24
25
```

后缀自动机 (ct)

```
struct SAM {
        SAM *next[26], *fa;
2
        int val;
   } mem[maxn], *last = mem, *tot = mem;
   void extend(int c)
6
        R SAM *p = last, *np;
        last = np = ++tot; np \rightarrow val = p \rightarrow val + 1;
        for (; p \&\& !p -> next[c]; p = p -> fa) p -> next[c] = np;
9
        if (!p) np -> fa = rt[id];
10
        else
11
        {
12
              SAM *q = p \rightarrow next[c];
13
              if (q \rightarrow val == p \rightarrow val + 1) np \rightarrow fa = q;
15
              else
              {
16
                   SAM *nq = ++tot;
17
                   memcpy(nq -> next, q -> next, sizeof nq -> next);
18
                   nq \rightarrow val = p \rightarrow val + 1;
19
                   nq \rightarrow fa = q \rightarrow fa;
20
21
                   q \rightarrow fa = np \rightarrow fa = nq;
22
                   for (; p \&\& p \rightarrow next[c] == q; p = p \rightarrow fa) p \rightarrow next[c] = nq;
23
        }
24
25
```

广义后缀自动机 (ct)

```
struct sam {
        sam *next[26], *fa;
        int val;
   } mem[maxn << 1], *tot = mem;</pre>
   inline sam *extend(R sam *p, R int c)
5
6
        if (p -> next[c])
             R sam *q = p \rightarrow next[c];
9
             if (q \rightarrow val == p \rightarrow val + 1)
10
                 return q;
11
             else
12
             {
13
```

4.5. Manacher (ct) 4. String

```
R sam *nq = ++tot;
                    memcpy(nq -> next, q -> next, sizeof nq -> next);
15
                    nq \rightarrow val = p \rightarrow val + 1;
16
17
                    nq \rightarrow fa = q \rightarrow fa;
                    q \rightarrow fa = nq;
18
                    for ( ; p \&\& p \rightarrow next[c] == q; p = p \rightarrow fa)
19
                         p -> next[c] = nq;
20
21
                    return nq;
               }
22
         }
23
24
         R sam *np = ++tot;
         np \rightarrow val = p \rightarrow val + 1;
25
         for ( ; p \&\& !p \rightarrow next[c]; p = p \rightarrow fa) p \rightarrow next[c] = np;
26
         if (!p)
27
              np -> fa = mem;
28
         else
29
         {
30
               R sam *q = p \rightarrow next[c];
31
               if (q \rightarrow val == p \rightarrow val + 1)
32
                    np \rightarrow fa = q;
33
               else
35
               {
36
                    R sam *nq = ++tot;
                    memcpy(nq -> next, q -> next, sizeof nq -> next);
37
                    nq \rightarrow val = p \rightarrow val + 1;
38
                    nq \rightarrow fa = q \rightarrow fa;
39
                    q \rightarrow fa = np \rightarrow fa = nq;
40
                    for (; p \&\& p \rightarrow next[c] == q; p = p \rightarrow fa)
41
                         p -> next[c] = nq;
42
               }
43
         }
44
45
         return np;
46
```

4.5 Manacher (ct)

```
char str[maxn];
  int p1[maxn], p2[maxn], n;
3 void manacher1()
       int mx = 0, id;
       for(int i = 1; i <= n; ++i)</pre>
           if (mx >= i) p1[i] = dmin(mx - i, p1[(id << 1) - i]);
           else p1[i] = 1;
           for (; str[i + p1[i]] == str[i - p1[i]]; ++p1[i]);
10
           if (p1[i] + i - 1 > mx) id = i, mx = p1[i] + i - 1;
11
12
13
   void manacher2()
14
15
       int mx = 0, id;
16
       for(int i = 1; i <= n; i++)</pre>
17
18
           if (mx \ge i) p2[i] = dmin(mx - i, p2[(id << 1) - i]);
19
           else p2[i] = 0;
20
           for (; str[i + p2[i] + 1] == str[i - p2[i]]; ++p2[i]);
21
           if (p2[i] + i > mx) id = i, mx = p2[i] + i;
22
23
```

4. String 4.6. 回文树 (lhy)

```
24 }
   int main()
25
26
       scanf("%s", str + 1);
27
       n = strlen(str + 1);
28
       str[0] = '#';
29
       str[n + 1] = '$';
30
       manacher1();
31
       manacher2();
32
33
       return 0;
34
```

4.6 回文树 (lhy)

```
char str[maxn];
  int next[maxn] [26], fail[maxn], len[maxn], cnt[maxn], last, tot, n;
3 inline int new_node(int 1)
       len[++tot] = 1;
       return tot;
  }
   inline void init()
   {
9
       tot = -1;
10
       new_node(0);
11
       new_node(-1);
12
       str[0] = -1;
13
       fail[0] = 1;
14
15
   inline int get_fail(int x)
16
17
       while (str[n - len[x] - 1] != str[n]) x = fail[x];
18
       return x;
19
20
   inline void extend(int c)
21
22
23
       int cur = get_fail(last);
24
       if (!next[cur][c])
25
26
           int now = new_node(len[cur] + 2);
27
           fail[now] = next[get_fail(fail[cur])][c];
28
           next[cur][c] = now;
29
30
       last = next[cur][c];
31
       ++cnt[last];
32
33
   long long ans;
34
   inline void count()
35
36
       for (int i = tot; i; --i)
37
38
           cnt[fail[i]] += cnt[i];
39
           cmax(ans, 111 * len[i] * cnt[i]);
40
41
  }
42
43 int main()
  {
44
45
       scanf("%s", str + 1);
```

4.7. 最小表示法 (lhy) 4. String

4.7 最小表示法 (lhy)

```
int main()
2 {
       int i = 0, j = 1, k = 0;
       while (i < n && j < n && k < n)
5
           int tmp = a[(i + k) \% n] - a[(j + k) \% n];
           if (!tmp) k++;
           else
9
               if (tmp > 0) i += k + 1;
10
               else j += k + 1;
11
               if (i == j) ++j;
12
               k = 0;
13
           }
14
15
       j = dmin(i, j);
16
       for (int i = j; i < n; ++i) printf("%d ", a[i]);</pre>
17
       for (int i = 0; i < j - 1; ++i) printf("%d ", a[i]);
       if (j > 0) printf("%d\n", a[j - 1]);
19
       return 0;
20
21 }
```

Chapter 5

Data Structure

5.1 莫队 (ct)

```
int size;
   struct Query {
       int 1, r, id;
       inline bool operator < (const Queuy &that) const {return 1 / size != that.1 / size ? 1 < that.1
         \hookrightarrow: ((1 / size) & 1 ? r < that.r : r > that.r);}
  \} q[maxn];
  int main()
6
       size = (int) sqrt(n * 1.0);
       std::sort(q + 1, q + m + 1);
       int 1 = 1, r = 0;
10
       for (int i = 1; i <= m; ++i)
           for (; r < q[i].r; ) add(++r);
14
           for (; r > q[i].r; ) del(r--);
           for (; 1 < q[i].1; ) del(1++);
15
           for (; 1 > q[i].1; ) add(--1);
16
17
               write your code here.
18
19
20
       return 0;
21
```

5.2 ST 表 (ct)

```
int a[maxn], f[20][maxn], n;
int Log[maxn];

void build()
{
    for (int i = 1; i <= n; ++i) f[0][i] = a[i];

    int lim = Log[n];
    for (int j = 1; j <= lim; ++j)
    {
        int *fj = f[j], *fj1 = f[j - 1];
        for (int i = 1; i <= n - (1 << j) + 1; ++i)
        fj[i] = dmax(fj1[i], fj1[i + (1 << (j - 1))]);
}
</pre>
```

5.3. 带权并查集 (ct) 5. Data Structure

```
14 int Query(int 1, int r)
15 {
       int k = Log[r - 1 + 1];
16
       return dmax(f[k][1], f[k][r - (1 << k) + 1]);
17
  }
18
  int main()
19
   {
20
       scanf("%d", &n);
21
       Log[0] = -1;
22
       for (int i = 1; i <= n; ++i)
23
24
           scanf("%d", &a[i]);
25
           Log[i] = Log[i >> 1] + 1;
26
       }
27
       build();
28
       int q;
29
       scanf("%d", &q);
30
31
       for (; q; --q)
32
           int 1, r; scanf("%d%d", &1, &r);
33
           printf("%d\n", Query(1, r) );
35
36
```

5.3 带权并查集 (ct)

```
struct edge
2
   {
       int a, b, w;
       inline bool operator < (const edge &that) const {return w > that.w;}
  int fa[maxn], f1[maxn], f2[maxn], f1cnt, f2cnt, val[maxn], size[maxn];
7 int main()
       int n, m; scanf("%d%d", &n, &m);
       for (int i = 1; i <= m; ++i)
10
           scanf("%d%d%d", &e[i].a, &e[i].b, &e[i].w);
11
       for (int i = 1; i <= n; ++i) size[i] = 1;
12
       std::sort(e + 1, e + m + 1);
13
       for (int i = 1; i <= m; ++i)
14
       {
           int x = e[i].a, y = e[i].b;
16
           for ( ; fa[x]; x = fa[x]);
17
           for (; fa[y]; y = fa[y]);
18
           if (x != y)
19
20
               if (size[x] < size[y]) std::swap(x, y);</pre>
21
               size[x] += size[y];
22
               val[y] = e[i].w;
23
               fa[y] = x;
24
           }
25
       }
26
       int q; scanf("%d", &q);
27
       for (; q; --q)
28
29
           int a, b; scanf("%d%d", &a, &b); f1cnt = f2cnt = 0;
30
           for (; fa[a]; a = fa[a]) f1[++f1cnt] = a;
31
           for (; fa[b]; b = fa[b]) f2[++f2cnt] = b;
32
```

5. Data Structure 5.4. 可并堆 (ct)

```
if (a != b) {puts("-1"); continue;}
33
           while (f1cnt && f2cnt && f1[f1cnt] == f2[f2cnt]) --f1cnt, --f2cnt;
34
           int ret = 0x7fffffff;
35
           for (; f1cnt; --f1cnt) cmin(ret, val[f1[f1cnt]]);
36
           for (; f2cnt; --f2cnt) cmin(ret, val[f2[f2cnt]]);
37
           printf("%d\n", ret);
38
39
       return 0;
40
41
```

5.4 可并堆 (ct)

```
struct Node {
       Node *ch[2];
       ll val; int size;
       inline void update()
            size = ch[0] \rightarrow size + ch[1] \rightarrow size + 1;
   } mem[maxn], *rt[maxn];
  Node *merge(Node *a, Node *b)
9
10
       if (a == mem) return b;
11
12
       if (b == mem) return a;
       if (a -> val < b -> val) std::swap(a, b);
13
       // a -> pushdown();
14
       std::swap(a -> ch[0], a -> ch[1]);
15
       a -> ch[1] = merge(a -> ch[1], b);
16
       a -> update();
17
       return a;
18
19
```

5.5 zkw 线段树 (ct)

必须为 0-based

```
inline void build()
   {
       for (int i = M - 1; i; --i) tr[i] = dmax(tr[i << 1], tr[i << 1 | 1]);
  }
  inline void Change(int x, int v)
       x += M; tr[x] = v; x >>= 1;
       for (; x; x >>= 1) tr[x] = dmax(tr[x << 1], tr[x << 1 | 1]);
   }
9
   inline int Query(int s, int t)
10
11
       int ret = -0x7fffffff;
12
       for (s = s + M - 1, t = t + M + 1; s ^ t ^ 1; s >>= 1, t >>= 1)
13
14
           if (~s & 1) cmax(ret, tr[s ^ 1]);
15
           if (t & 1) cmax(ret, tr[t ^{^{\smallfrown}} 1]);
16
17
       return ret;
18
  l٦
19
  int main()
20
  {
21
       int n; scanf("%d", &n);
22
```

5.6. Splay (ct) 5. Data Structure

```
for (M = 1; M < n; M <<= 1) ;
       for (int i = 0; i < n; ++i)
24
           scanf("%d", &tr[i + M]);
25
       for (int i = n; i < M; ++i) tr[i + M] = -0x7ffffffff;
26
       build();
27
       int q; scanf("%d", &q);
28
       for (; q; --q)
29
30
           int 1, r; scanf("%d%d", &1, &r); --1, --r;
31
           printf("%d\n", Query(1, r));
32
       }
33
       return 0;
34
35
```

5.6 Splay (ct)

指针版

```
struct Node *null;
   struct Node {
       Node *ch[2], *fa;
       int val; bool rev;
       inline bool type()
            return fa -> ch[1] == this;
       }
       inline void pushup()
       {
10
11
       inline void pushdown()
12
13
14
            if (rev)
15
16
                ch[0] -> rev ^= 1;
                ch[1] -> rev ^= 1;
17
                std::swap(ch[0], ch[1]);
18
                rev ^= 1;
19
            }
20
21
       inline void rotate()
22
23
            bool d = type(); Node *f = fa, *gf = f -> fa;
24
            (fa = gf, f \rightarrow fa != null) ? fa \rightarrow ch[f \rightarrow type()] = this : 0;
25
            (f \rightarrow ch[d] = ch[!d]) != null ? ch[!d] \rightarrow fa = f : 0;
26
            (ch[!d] = f) -> fa = this;
27
            f -> pushup();
28
29
       inline void splay()
30
31
            for (; fa != null; rotate())
32
                if (fa -> fa != null)
33
                     (type() == fa -> type() ? fa : this) -> rotate();
34
            pushup();
       }
   } mem[maxn];
```

5. Data Structure 5.6. Splay (ct)

数组版

```
// BZOJ - 1500 维修数列
int fa[maxn], ch[maxn][2], a[maxn], size[maxn], cnt;
   int sum[maxn], lmx[maxn], rmx[maxn], mx[maxn], v[maxn], id[maxn], root;
  bool rev[maxn], tag[maxn];
  inline void update(R int x)
6
       R int ls = ch[x][0], rs = ch[x][1];
       size[x] = size[ls] + size[rs] + 1;
       sum[x] = sum[ls] + sum[rs] + v[x];
9
10
       mx[x] = gmax(mx[ls], mx[rs]);
11
       cmax(mx[x], lmx[rs] + rmx[ls] + v[x]);
       lmx[x] = gmax(lmx[ls], sum[ls] + v[x] + lmx[rs]);
12
       rmx[x] = gmax(rmx[rs], sum[rs] + v[x] + rmx[ls]);
13
14
  inline void pushdown(R int x)
15
16
       R \text{ int } ls = ch[x][0], rs = ch[x][1];
17
       if (tag[x])
18
19
           rev[x] = tag[x] = 0;
           if (ls) tag[ls] = 1, v[ls] = v[x], sum[ls] = size[ls] * v[x];
21
           if (rs) tag[rs] = 1, v[rs] = v[x], sum[rs] = size[rs] * v[x];
22
           if (v[x] >= 0)
23
24
           {
               if (ls) lmx[ls] = rmx[ls] = mx[ls] = sum[ls];
25
               if (rs) lmx[rs] = rmx[rs] = mx[rs] = sum[rs];
26
           }
27
           else
28
29
               if (ls) lmx[ls] = rmx[ls] = 0, mx[ls] = v[x];
30
               if (rs) lmx[rs] = rmx[rs] = 0, mx[rs] = v[x];
31
           }
32
33
       if (rev[x])
34
35
           rev[x] ^= 1; rev[ls] ^= 1; rev[rs] ^= 1;
36
           swap(lmx[ls], rmx[ls]);swap(lmx[rs], rmx[rs]);
37
           swap(ch[ls][0], ch[ls][1]); swap(ch[rs][0], ch[rs][1]);
38
39
40
   inline void rotate(R int x)
41
42
       R int f = fa[x], gf = fa[f], d = ch[f][1] == x;
43
       if (f == root) root = x;
44
       (ch[f][d] = ch[x][d ^ 1]) > 0 ? fa[ch[f][d]] = f : 0;
45
       (fa[x] = gf) > 0 ? ch[gf][ch[gf][1] == f] = x : 0;
46
       fa[ch[x][d ^1] = f] = x;
47
       update(f);
48
49
   inline void splay(R int x, R int rt)
50
51
       while (fa[x] != rt)
52
53
           R int f = fa[x], gf = fa[f];
54
           if (gf != rt) rotate((ch[gf][1] == f) ^ (ch[f][1] == x) ? x : f);
55
           rotate(x);
56
       }
57
       update(x);
58
59 }
```

5.6. Splay (ct) 5. Data Structure

```
60 void build(R int 1, R int r, R int rt)
   {
61
        if (1 > r) return;
62
        R int mid = 1 + r >> 1, now = id[mid], last = id[rt];
63
        if (1 == r)
64
65
            sum[now] = a[1];
66
            size[now] = 1;
67
            tag[now] = rev[now] = 0;
68
            if (a[1] >= 0) lmx[now] = rmx[now] = mx[now] = a[1];
 69
            else lmx[now] = rmx[now] = 0, mx[now] = a[1];
 70
        }
 71
        else
72
        {
73
            build(1, mid - 1, mid);
 74
            build(mid + 1, r, mid);
 75
76
        v[now] = a[mid];
77
        fa[now] = last;
 78
        update(now);
 79
        ch[last][mid >= rt] = now;
 80
 81
   int find(R int x, R int rank)
 82
 83
        if (tag[x] || rev[x]) pushdown(x);
 84
        R int ls = ch[x][0], rs = ch[x][1], lsize = size[ls];
 85
        if (lsize + 1 == rank) return x;
 86
        if (lsize >= rank)
 87
            return find(ls, rank);
 88
 89
        else
90
            return find(rs, rank - lsize - 1);
91
   inline int prepare(R int 1, R int tot)
 92
 93
        R int x = find(root, 1 - 1), y = find(root, 1 + tot);
94
        splay(x, 0);
95
        splay(y, x);
96
        return ch[y][0];
97
98
   std::queue <int> q;
   inline void Insert(R int left, R int tot)
100
101
102
        for (R int i = 1; i <= tot; ++i ) a[i] = FastIn();</pre>
        for (R int i = 1; i <= tot; ++i )</pre>
103
            if (!q.empty()) id[i] = q.front(), q.pop();
104
            else id[i] = ++cnt;
105
        build(1, tot, 0);
106
        R int z = id[(1 + tot) >> 1];
107
        R int x = find(root, left), y = find(root, left + 1);
108
        splay(x, 0);
109
110
        splay(y, x);
        fa[z] = y;
111
        ch[y][0] = z;
112
        update(y);
113
114
        update(x);
115
void rec(R int x)
117 | {
        if (!x) return;
118
        R \text{ int } ls = ch[x][0], rs = ch[x][1];
119
        rec(ls); rec(rs); q.push(x);
```

5. Data Structure 5.6. Splay (ct)

```
fa[x] = ch[x][0] = ch[x][1] = 0;
121
        tag[x] = rev[x] = 0;
^{122}
123 }
   inline void Delete(R int 1, R int tot)
124
   ł
125
        R int x = prepare(1, tot), f = fa[x];
126
        rec(x); ch[f][0] = 0;
127
        update(f); update(fa[f]);
128
129
   inline void Makesame(R int 1, R int tot, R int val)
130
131
        R int x = prepare(1, tot), y = fa[x];
132
        v[x] = val; tag[x] = 1; sum[x] = size[x] * val;
133
        if (val >= 0) lmx[x] = rmx[x] = mx[x] = sum[x];
134
        else lmx[x] = rmx[x] = 0, mx[x] = val;
135
        update(y); update(fa[y]);
136
137
   inline void Reverse(R int 1, R int tot)
138
139
        R int x = prepare(1, tot), y = fa[x];
140
        if (!tag[x])
141
142
            rev[x] ^= 1;
143
            swap(ch[x][0], ch[x][1]);
144
            swap(lmx[x], rmx[x]);
145
            update(y); update(fa[y]);
146
147
148
   inline void Query(R int 1, R int tot)
149
150
151
        R int x = prepare(1, tot);
152
        printf("%d\n",sum[x]);
153
   #define inf ((1 << 30))
154
   int main()
155
156
        R int n = FastIn(), m = FastIn(), 1, tot, val;
157
        R char op, op2;
158
        mx[0] = a[1] = a[n + 2] = -inf;
159
        for (R int i = 2; i <= n + 1; i++ )
160
        {
161
            a[i] = FastIn();
162
163
        }
164
        for (R int i = 1; i \le n + 2; ++i) id[i] = i;
        n += 2; cnt = n; root = (n + 1) >> 1;
165
        build(1, n, 0);
166
        for (R int i = 1; i <= m; i++ )
167
        {
168
            op = getc();
169
            while (op < 'A' \mid \mid op > 'Z') op = getc();
170
            getc(); op2 = getc();getc();getc();getc();
171
            if (op == 'M' && op2 == 'X')
172
173
                printf("%d\n",mx[root] );
174
            }
175
            else
176
            {
177
                1 = FastIn() + 1; tot = FastIn();
178
                if (op == 'I') Insert(1, tot);
179
                if (op == 'D') Delete(1, tot);
180
                if (op == 'M') val = FastIn(), Makesame(1, tot, val);
181
```

5.7. Treap (ct) 5. Data Structure

5.7 Treap (ct)

```
struct Treap {
       Treap *ls, *rs;
       int size;
       bool rev;
       inline void update()
           size = ls -> size + rs -> size + 1;
       }
       inline void set_rev()
10
           rev ^= 1;
11
           std::swap(ls, rs);
12
       }
13
       inline void pushdown()
14
15
           if (rev)
16
17
                ls -> set_rev();
18
                rs -> set_rev();
19
20
                rev = 0;
           }
21
       }
22
  } mem[maxn], *root, *null = mem;
23
   struct Pair {
^{24}
       Treap *fir, *sec;
25
26 };
  Treap *build(R int 1, R int r)
27
28
       if (1 > r) return null;
29
       R int mid = 1 + r >> 1;
       R Treap *now = mem + mid;
       now \rightarrow rev = 0;
       now -> ls = build(1, mid - 1);
33
       now -> rs = build(mid + 1, r);
34
       now -> update();
35
       return now;
36
37
  inline Treap *Find_kth(R Treap *now, R int k)
38
39
40
       if (!k) return mem;
       if (now -> ls -> size >= k) return Find_kth(now -> ls, k);
41
       else if (now \rightarrow ls \rightarrow size + 1 == k) return now;
42
       else return Find_kth(now -> rs, k - now -> ls -> size - 1);
43
44 | }
45 Treap *merge(R Treap *a, R Treap *b)
46
       if (a == null) return b;
```

```
if (b == null) return a;
48
       if (rand() % (a -> size + b -> size) < a -> size)
49
50
            a -> pushdown();
51
            a \rightarrow rs = merge(a \rightarrow rs, b);
52
            a -> update();
53
            return a;
54
       }
55
       else
56
57
            b -> pushdown();
58
            b -> ls = merge(a, b -> ls);
59
            b -> update();
60
            return b;
61
       }
62
63
   Pair split(R Treap *now, R int k)
64
65
       if (now == null) return (Pair) {null, null};
66
       R Pair t = (Pair) {null, null};
67
       now -> pushdown();
68
       if (k \le now \rightarrow ls \rightarrow size)
69
70
       {
            t = split(now -> ls, k);
71
            now -> ls = t.sec;
72
            now -> update();
73
            t.sec = now;
74
       }
75
       else
76
77
            t = split(now \rightarrow rs, k - now \rightarrow ls \rightarrow size - 1);
78
79
            now -> rs = t.fir;
            now -> update();
80
            t.fir = now;
81
82
       return t;
83
84
   inline void set_rev(int 1, int r)
85
86
       R Pair x = split(root, 1 - 1);
87
       R Pair y = split(x.sec, r - 1 + 1);
88
       y.fir -> set_rev();
89
90
       root = merge(x.fir, merge(y.fir, y.sec));
91
```

5.8 可持久化平衡树 (ct)

```
char str[maxn];
struct Treap

{
    Treap *ls, *rs;
    char data; int size;
    inline void update()
    {
        size = ls -> size + rs -> size + 1;
    }
}

**root[maxn], mem[maxcnt], *tot = mem, *last = mem, *null = mem;
inline Treap* new_node(char ch)
{
```

```
*++tot = (Treap) {null, null, ch, 1};
       return tot;
14
15 }
16 struct Pair
17 | {
       Treap *fir, *sec;
18
19 };
20 inline Treap *copy(Treap *x)
21
       if (x == null) return null;
22
       if(x > last) return x;
23
       *++tot = *x;
24
       return tot;
25
26
Pair Split(Treap *x, int k)
28
       if (x == null) return (Pair) {null, null};
29
30
       Pair y;
       Treap *nw = copy(x);
31
       if (nw \rightarrow ls \rightarrow size >= k)
32
           y = Split(nw -> ls, k);
34
           nw -> ls = y.sec;
35
           nw -> update();
36
           y.sec = nw;
37
       }
38
       else
39
       {
40
41
           y = Split(nw \rightarrow rs, k - nw \rightarrow ls \rightarrow size - 1);
42
           nw -> rs = y.fir;
           nw -> update();
43
44
           y.fir = nw;
45
       return y;
46
47 | }
Treap *Merge(Treap *a, Treap *b)
49
       if (a == null) return b;
50
       if (b == null) return a;
51
       Treap *nw;
52
       if (rand() \% (a -> size + b -> size) < a -> size)
53
           nw = copy(a);
           nw -> rs = Merge(nw -> rs, b);
56
       }
57
       else
58
       {
59
           nw = copy(b);
60
           nw -> ls = Merge(a, nw -> ls);
61
62
       nw -> update();
63
       return nw;
64
65
  Treap *Build(int 1, int r)
67
       if (1 > r) return null;
68
       R \text{ int } mid = 1 + r >> 1;
69
       Treap *nw = new_node(str[mid]);
70
       nw -> ls = Build(1, mid - 1);
71
       nw -> rs = Build(mid + 1, r);
72
       nw -> update();
```

```
return nw;
74
<sub>75</sub> }
76 int now;
inline void Insert(int k, char ch)
78
        Pair x = Split(root[now], k);
79
        Treap *nw = new_node(ch);
80
        root[++now] = Merge(Merge(x.fir, nw), x.sec);
81
82
    inline void Del(int 1, int r)
83
84
        Pair x = Split(root[now], 1 - 1);
85
        Pair y = Split(x.sec, r - 1 + 1);
86
        root[++now] = Merge(x.fir, y.sec);
87
88
   inline void Copy(int 1, int r, int 11)
89
90
        Pair x = Split(root[now], 1 - 1);
91
        Pair y = Split(x.sec, r - 1 + 1);
92
        Pair z = Split(root[now], 11);
93
        Treap *ans = y.fir;
94
        root[++now] = Merge(Merge(z.fir, ans), z.sec);
95
   }
96
   void Print(Treap *x, int 1, int r)
97
98
        if (!x) return;
99
        if (1 > r) return;
100
        R int mid = x \rightarrow ls \rightarrow size + 1;
101
        if (r < mid)</pre>
102
103
            Print(x -> ls, l, r);
104
105
            return ;
        }
106
        if (1 > mid)
107
        {
108
            Print(x -> rs, 1 - mid, r - mid);
109
            return ;
110
111
        Print(x -> ls, l, mid - 1);
112
        printf("%c", x -> data );
113
        Print(x -> rs, 1, r - mid);
114
115 }
116
   void Printtree(Treap *x)
117 {
118
        if (!x) return;
        Printtree(x -> ls);
119
        printf("%c", x \rightarrow data);
120
        Printtree(x -> rs);
121
122
   int main()
123
124
        srand(time(0) + clock());
^{125}
        null -> ls = null -> rs = null; null -> size = 0; null -> data = 0;
126
        int n = F();
127
        gets(str + 1);
128
        int len = strlen(str + 1);
129
        root[0] = Build(1, len);
130
        while (1)
131
        {
132
            last = tot;
133
            R char opt = getc();
134
```

5.9. CDQ 分治 (ct) 5. Data Structure

```
while (opt < 'A' \mid \mid opt > 'Z')
135
136
                  if (opt == EOF) return 0;
137
                 opt = getc();
138
             }
139
             if (opt == 'I')
140
             {
141
                 R int x = F();
142
                 R char ch = getc();
143
                 Insert(x, ch);
144
             }
145
             else if (opt == 'D')
146
147
                 R int 1 = F(), r = F();
148
                 Del(1, r);
149
150
             else if (opt == 'C')
151
152
                 R \text{ int } x = F(), y = F(), z = F();
153
                 Copy(x, y, z);
154
             }
155
             else if (opt == 'P')
157
                 R int x = F(), y = F(), z = F();
158
                 Print(root[now - x], y, z);
159
                 puts("");
160
161
        }
162
        return 0;
163
164
```

5.9 CDQ 分治 (ct)

```
struct event
2 {
      int x, y, id, opt, ans;
  } t[maxn], q[maxn];
5 void cdq(int left, int right)
  {
      if (left == right) return ;
      R int mid = left + right >> 1;
      cdq(left, mid);
10
      cdq(mid + 1, right);
      //分成若干个子问题
11
      ++now;
12
      for (int i = left, j = mid + 1; j <= right; ++j)
13
14
          for (; i <= mid && q[i].x <= q[j].x; ++i)
15
              if (!q[i].opt)
16
                  add(q[i].y, q[i].ans);
17
           //考虑前面的修改操作对后面的询问的影响
18
          if (q[j].opt)
19
20
              q[j].ans += query(q[j].y);
      }
21
      R int i, j, k = 0;
22
      //以下相当于归并排序
23
      for (i = left, j = mid + 1; i <= mid \&\& j <= right; )
24
25
          if (q[i].x \le q[j].x)
26
```

5. Data Structure 5.10. Bitset (ct)

```
t[k++] = q[i++];
27
            else
28
                 t[k++] = q[j++];
29
30
       for (; i <= mid; )</pre>
31
            t[k++] = q[i++];
32
       for (; j <= right; )</pre>
33
            t[k++] = q[j++];
34
       for (int i = 0; i < k; ++i)
35
            q[left + i] = t[i];
36
37
```

5.10 Bitset (ct)

```
namespace Game {
2 #define maxn 300010
3 #define maxs 30010
4 uint b1[32][maxs], b2[32][maxs];
5 int popcnt[256];
6 inline void set(R uint *s, R int pos)
       s[pos >> 5] = 1u << (pos & 31);
  }
9
  inline int popcount(R uint x)
10
11
       \texttt{return popcnt[x >> 24 \& 255]}
12
            + popcnt[x >> 16 & 255]
13
            + popcnt[x >> 8 & 255]
14
            + popcnt[x
                           & 255];
15
16
   void main() {
17
       int n, q;
18
       scanf("%d%d", &n, &q);
19
       char *s1 = new char[n + 1];
20
       char *s2 = new char[n + 1];
21
       scanf("%s%s", s1, s2);
22
       uint *anss = new uint[q];
23
       for (R int i = 1; i < 256; ++i) popcnt[i] = popcnt[i >> 1] + (i & 1);
       #define modify(x, _p)\
25
26
           for (R int j = 0; j < 32 & j <= p; ++j)
27
               set(b##x[j], p - j); \
28
29
       for (R int i = 0; i < n; ++i)
30
           if (s1[i] == '0') modify(1, 3 * i)
31
           else if (s1[i] == '1') modify(1, 3 * i + 1)
32
           else modify(1, 3 * i + 2)
33
       for (R int i = 0; i < n; ++i)
34
           if (s2[i] == '1') modify(2, 3 * i)
35
           else if (s2[i] == '2') modify(2, 3 * i + 1)
36
           else modify(2, 3 * i + 2)
37
       for (int Q = 0; Q < q; ++Q) {
38
           R int x, y, 1;
39
```

5.10. Bitset (ct) 5. Data Structure

```
scanf("%d%d%d", &x, &y, &1); x *= 3; y *= 3; 1 *= 3;
           uint *f1 = b1[x \& 31], *f2 = b2[y \& 31], ans = 0;
41
           R int i = x >> 5, j = y >> 5, p, lim;
42
           for (p = 0, lim = 1 >> 5; p + 8 < lim; p += 8, i += 8, j += 8)
43
44
               ans += popcount(f1[i + 0] & f2[j + 0]);
45
               ans += popcount(f1[i + 1] & f2[j + 1]);
46
               ans += popcount(f1[i + 2] & f2[j + 2]);
47
               ans += popcount(f1[i + 3] & f2[j + 3]);
48
               ans += popcount(f1[i + 4] & f2[j + 4]);
49
               ans += popcount(f1[i + 5] & f2[j + 5]);
50
               ans += popcount(f1[i + 6] & f2[j + 6]);
51
               ans += popcount(f1[i + 7] & f2[j + 7]);
52
           }
53
           for (; p < lim; ++p, ++i, ++j) ans += popcount(f1[i] & f2[j]);
54
           R uint S = (1u << (1 & 31)) - 1;
55
           ans += popcount(f1[i] & f2[j] & S);
56
           anss[Q] = ans;
57
       }
58
       output_arr(anss, q * sizeof(uint));
59
60
  }
61
```

Chapter 6

Others

6.1 vimrc (gy)

```
se et ts=4 sw=4 sts=4 nu sc sm lbr is hls mouse=a
  sy on
  ino <tab> <c-n>
  ino <s-tab> <tab>
  au bufwinenter * winc L
  nm <f6> ggVG"+y
  nm <f7> :w<cr>:make<cr>
  nm <f8> :!@<cr>
  nm <f9> :!@ < in<cr>
  nm <s-f9> :!(time @ < in &> out) &>> out<cr>:sp out<cr>
  au filetype cpp cm @ ./a.out | se cin fdm=syntax mp=g++\ %\ -std=c++11\ -Wall\ -Wextra\
    \hookrightarrow -Wconversion\ -02
12 map <c-p> :ha<cr>
  se pheader=%N@%F popt=number:y
  au filetype java cm @ java %< | se cin fdm=syntax mp=javac\ %
  au filetype python cm @ python % | se si fdm=indent
  au bufenter *.kt setf kotlin
  au filetype kotlin cm @ kotlin _%<Kt | se si mp=kotlinc\ %
```

6.2 STL 释放内存 (Durandal)

```
template <typename T>
   __inline void clear(T &container) {
     container.clear();
     T(container).swap(container);
}
```

6.3 开栈 (Durandal)

```
register char *_sp __asm__("rsp");
int main() {
   const int size = 400 << 20; // 400 MB
   static char *sys, *mine(new char[size] + size - 4096);
   sys = _sp; _sp = mine;
   _main(); // main method</pre>
```

 $6.4. \ O3 \ (gy)$ 6. Others

6.4 O3 (gy)

```
__attribute__((optimize("-03"))) int main() { return 0; }
```

6.5 Java Template (gy)

```
import java.io.*;
   import java.math.*;
   import java.util.*;
  public class Template {
       // Input
5
       private static BufferedReader reader;
       private static StringTokenizer tokenizer;
       private static String next() {
           try {
10
               while (tokenizer == null || !tokenizer.hasMoreTokens())
11
                   tokenizer = new StringTokenizer(reader.readLine());
12
           } catch (IOException e) {
               // do nothing
13
14
           return tokenizer.nextToken();
15
16
       private static int nextInt() {
17
           return Integer.parseInt(next());
18
19
       private static double nextDouble() {
20
           return Double.parseDouble(next());
21
22
       private static BigInteger nextBigInteger() {
23
           return new BigInteger(next());
24
25
       public static void main(String[] args) {
26
           reader = new BufferedReader(new InputStreamReader(System.in));
27
           Scanner scanner = new Scanner(System.in);
29
           while (scanner.hasNext())
               scanner.next();
30
       }
31
       // BigInteger & BigDecimal
32
       private static void bigDecimal() {
33
           BigDecimal a = BigDecimal.valueOf(1.0);
34
           BigDecimal b = a.setScale(50, RoundingMode.HALF_EVEN);
35
           BigDecimal c = b.abs();
36
           // if scale omitted, b.scale is used
37
           BigDecimal d = c.divide(b, 50, RoundingMode.HALF_EVEN);
38
           // since Java 9
39
           BigDecimal e = d.sqrt(new MathContext(50, RoundingMode.HALF_EVEN));
40
           BigDecimal x = new BigDecimal(BigInteger.ZERO);
41
           BigInteger y = BigDecimal.ZERO.toBigInteger(); // RoundingMode.DOWN
42
           y = BigDecimal.ZERO.setScale(0, RoundingMode.HALF_EVEN).unscaledValue();
43
44
       // sqrt for Java 8
45
```

```
// can solve scale=100 for 10000 times in about 1 second
46
       private static BigDecimal sqrt(BigDecimal a, int scale) {
47
            if (a.compareTo(BigDecimal.ZERO) < 0)</pre>
48
                return BigDecimal.ZERO.setScale(scale, RoundingMode.HALF_EVEN);
49
            int length = a.precision() - a.scale();
50
            BigDecimal ret = new BigDecimal(BigInteger.ONE, -length / 2);
51
            for (int i = 1; i <= Integer.highestOneBit(scale) + 10; i++)</pre>
52
                ret = ret.add(a.divide(ret, scale,
53
                  → RoundingMode.HALF_EVEN)).divide(BigDecimal.valueOf(2), scale,
                  → RoundingMode.HALF_EVEN);
            return ret;
54
55
       // can solve a=2^10000 for 100000 times in about 1 second
56
       private static BigInteger sqrt(BigInteger a) {
57
            int length = a.bitLength() - 1;
58
            BigInteger 1 = BigInteger.ZERO.setBit(length / 2), r = BigInteger.ZERO.setBit(length / 2);
59
            while (!1.equals(r)) {
60
                BigInteger m = 1.add(r).shiftRight(1);
61
                if (m.multiply(m).compareTo(a) < 0)</pre>
62
                    1 = m.add(BigInteger.ONE);
63
                else
64
                    r = m;
65
66
67
            return 1;
       }
68
       // Collections
69
       private static void arrayList() {
70
           List<Integer> list = new ArrayList<>();
71
72
            // Generic array is banned
73
            List[] lists = new List[100];
74
            lists[0] = new ArrayList<Integer>();
            // for List<Integer>, remove(Integer) stands for element, while remove(int) stands for
75
              \rightarrow index
            list.remove(list.get(1));
76
            list.remove(list.size() - 1);
77
            list.clear():
78
            Queue<Integer> queue = new LinkedList<>();
79
            // return the value without popping
80
            queue.peek();
81
            // pop and return the value
82
            queue.poll();
83
            Queue<Integer> priorityQueue = new PriorityQueue<>();
85
            Deque<Integer> deque = new ArrayDeque<>();
86
            deque.peekFirst();
87
            deque.peekLast();
            deque.pollFirst();
88
            TreeSet<Integer> set = new TreeSet<>();
89
            TreeSet<Integer> anotherSet = new TreeSet<>(Comparator.reverseOrder());
90
            set.ceiling(1);
91
            set.floor(1);
92
            set.lower(1);
93
            set.higher(1);
94
            set.contains(1);
95
            HashSet<Integer> hashSet = new HashSet<>();
96
            HashMap<String, Integer> map = new HashMap<>();
97
            map.put("", 1);
98
            map.get("");
99
            map.forEach((string, integer) -> System.out.println(string + integer));
100
            TreeMap<String, Integer> treeMap = new TreeMap<>();
101
            Arrays.sort(new int[10]);
102
```

6.6. Big Fraction (gy) 6. Others

```
Arrays.sort(new Integer[10], (a, b) -> {
                if (a.equals(b)) return 0;
104
                if (a > b) return -1;
105
                return 1;
106
            });
107
            Arrays.sort(new Integer[10], Comparator.comparingInt((a) -> (int) a).reversed());
108
            long a = 1_000_000_000_000_000_000L;
109
            int b = Integer.MAX_VALUE;
110
            int c = 'a';
111
112
113
```

6.6 Big Fraction (gy)

```
fun gcd(a: Long, b: Long): Long = if (b == OL) a else gcd(b, a % b)
   class Fraction(val a: BigInteger, val b: BigInteger) {
       constructor(a: Long, b: Long) : this(BigInteger.valueOf(a / gcd(a, b)), BigInteger.valueOf(b /
         \hookrightarrow \gcd(a, b)))
       operator fun plus(o: Fraction): Fraction {
           var gcd = b.gcd(o.b)
           val tempProduct = (b / gcd) * (o.b / gcd)
           var ansA = a * (o.b / gcd) + o.a * (b / gcd)
           val gcd2 = ansA.gcd(gcd)
           ansA /= gcd2
           gcd /= gcd2
10
           return Fraction(ansA, gcd * tempProduct)
11
12
       operator fun minus(o: Fraction): Fraction {
14
           var gcd = b.gcd(o.b)
15
           val tempProduct = (b / gcd) * (o.b / gcd)
           var ansA = a * (o.b / gcd) - o.a * (b / gcd)
16
           val gcd2 = ansA.gcd(gcd)
17
           ansA /= gcd2
18
           gcd /= gcd2
19
           return Fraction(ansA, gcd * tempProduct)
20
21
       operator fun times(o: Fraction): Fraction {
22
           val gcd1 = a.gcd(o.b)
23
           val gcd2 = b.gcd(o.a)
24
           return Fraction((a / gcd1) * (o.a / gcd2), (b / gcd2) * (o.b / gcd1))
25
26
  }
27
```

6.7 模拟退火 (ct)

6. Others 6.8. 三分 (ct)

```
return maxx;
10
11 }
12 int main()
   {
13
       srand(time(NULL) + clock());
14
       db x = 0, fnow = f(x);
15
       fans = 1e30;
16
       for (db T = 1e4; T > 1e-4; T *= 0.997)
17
18
           db nx = x + randp() * T, fnext = f(nx);
19
           db delta = fnext - fnow;
20
           if (delta < 1e-9 || exp(-delta / T) > rand01())
21
22
                x = nx;
23
                fnow = fnext;
24
           }
25
26
       return 0;
27
28
```

三分 (ct) 6.8

```
inline db cubic_search()
2
       double 1 = -1e4, r = 1e4;
       for (int i = 1; i <= 100; ++i)
5
           double 11 = (1 + r) * 0.5;
6
           double rr = (ll + r) * 0.5;
           if (check(ll) < check(rr)) r = rr;</pre>
           else l = 11;
       }
10
       return (1 + r) * 0.5;
11
```

Zeller Congruence (gy) 6.9

```
int day_in_week(int year, int month, int day) {
   if (month == 1 || month == 2)
       month += 12, year--;
   int c = year / 100, y = year % 100, m = month, d = day;
   int ret = (y + y / 4 + c / 4 + 5 * c + 13 * (m + 1) / 5 + d + 6) % 7;
   return ret >= 0 ? ret : ret + 7;
```

6.10 博弈论模型 (gy)

· Wythoff's game

给定两堆石子,每次可以从任意一堆中取至少一个石子,或从两堆中取相同的至少一个石子,取走最后 石子的胜

```
先手胜当且仅当石子数满足:
```

```
\lfloor (b-a) \times \phi \rfloor = a, (a \le b, \phi = \frac{\sqrt{5}+1}{2})
先手胜对应的石子数构成两个序列:
Lower Wythoff sequence: a_n = \lfloor n \times \phi \rfloor
Upper Wythoff sequence: b_n = \lfloor n \times \phi^2 \rfloor
```

• Fibonacci nim

给定一堆石子,第一次可以取至少一个、少于石子总数数量的石子,之后每次可以取至少一个、不超过 上次取石子数量两倍的石子,取走最后石子的胜 先手胜当且仅当石子数为斐波那契数

6.11 积分表 (integral-table.com)

$$\int x^0 dx = \frac{1}{1+x}x^{n+1}, \ n \neq -1$$

$$\int \frac{1}{x} dx = \ln |x|$$

$$\int \frac{1}{x} dx = \ln |x|$$

$$\int u dy = uy - \int v du$$

$$\int \frac{1}{(x+y)^2} dx = \frac{1}{x} \ln |x + b|$$

$$\int \frac{1}{(x+y)^2} dx = \frac{1}{x+1}$$

$$\int (x+y)^{n} dx = \frac{(x+y)^{n+1}}{n+1}, \ n \neq -1$$

$$\int x(x+y)^{n} dx = \frac{(x+y)^{n+1}}{n+1}, \ n \neq -1$$

$$\int x(x+y)^{n} dx = \frac{(x+y)^{n+1}}{n+1}, \ n \neq -1$$

$$\int x(x+y)^{n} dx = \frac{(x+y)^{n+1}}{n+1}, \ n \neq -1$$

$$\int x(x+y)^{n} dx = \frac{(x+y)^{n+1}}{n+1}, \ n \neq -1$$

$$\int x(x+y)^{n} dx = \frac{(x+y)^{n+1}}{n+1}, \ n \neq -1$$

$$\int x(x+y)^{n} dx = \frac{1}{x} \tan^{n-1} x$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} \ln |x|^{2} + 2^{2}|$$

$$\int \frac{1}{(x+y)^{n}} dx = \frac{1}{2} (x-y)^{n}/2 + \frac{1}{2} (x-y)^{n}/2 + 2^{n}|$$

$$\int \frac{1}{\sqrt{x+y}^{n}} dx = \frac{1}{2} (x-y)^{n}/2 + \frac{1}{2} (x-y)^{n}/2 + 2^{n}|$$

$$\int \frac{1}{\sqrt{x+y}^{n}} dx = \frac{1}{2} (x-y)^{n}/2 + \frac{1}{2} (x-y)^{n}/2 + 2^{n}|$$

$$\int \frac{1}{\sqrt{x+y}^{n}} dx = \frac{1}{2} (x-y)^{n}/2 + \frac{1}{2} (x-y)^{n}/2 + 2^{n}|$$

$$\int \frac{1}{\sqrt{x+y}^{n}} dx = \frac{1}{2} (x-y)^{n}/2 + \frac{1}{2} (x-y)^{n}/2 + 2^{n}|$$

$$\int \frac{1}{\sqrt{x+y}^{n}} dx = \frac{1}{2} (x-y)^{n}/2 + \frac{1}{2} (x-y)^{n}/2 + 2^{n}|$$

$$\int \frac{1}{\sqrt{x+y}^{n}} dx = \frac{1}{2} (x-y)^{n}/2 + \frac{1}{2} (x-y)^{n}/2 + 2^{n}|$$

$$\int \frac{1}{\sqrt{x+y}^{n}} dx = \frac{1}{2} (x-y)^{n}/2 + \frac{1}{2} (x-y)^{n}/2 + 2^{n}|$$

$$\int \frac{1}{\sqrt{x+y}^{n}} dx = \frac{1}{2} (x-y)^{n}/2 + \frac{1}{2$$

$$\int \sin^3 ax \, dx = -\frac{3\cos ax}{4a} + \frac{\cos 3ax}{12a}$$

$$\int \cos ax \, dx = \frac{1}{a} \sin ax$$

$$\int \cos^2 ax \, dx = \frac{x}{2} + \frac{\sin 2ax}{4a}$$

$$\int \cos^3 ax dx = \frac{3\sin ax}{4a} + \frac{\sin 3ax}{12a}$$

$$\int \cos x \sin x \, dx = \frac{1}{2} \sin^2 x + c_1 = -\frac{1}{2} \cos^2 x + c_2 = -\frac{1}{4} \cos 2x + c_3$$

$$\int \cos ax \sin bx \, dx = \frac{\cos[(a-b)x]}{2(a-b)} - \frac{\cos[(a+b)x]}{2(a+b)}, \, a \neq b$$

$$\int \sin^2 ax \cos bx \, dx = -\frac{\sin[(2a-b)x]}{4(2a-b)} + \frac{\sin bx}{2b} - \frac{\sin[(2a+b)x]}{4(2a+b)}$$

$$\int \sin^2 x \cos x \, dx = \frac{1}{3} \sin^3 x$$

$$\int \cos^2 ax \sin bx \, dx = \frac{\cos[(2a-b)x]}{4(2a-b)} - \frac{\cos bx}{2b} - \frac{\cos[(2a+b)x]}{4(2a+b)}$$

$$\int \cos^2 ax \sin ax \, dx = -\frac{1}{3a} \cos^3 ax$$

$$\int \sin^2 ax \cos^2 bx dx = \frac{x}{4} - \frac{\sin 2ax}{8a} - \frac{\sin[2(a-b)x]}{16(a-b)} + \frac{\sin 2bx}{8b} - \frac{\sin[2(a+b)x]}{16(a+b)}$$

$$\int \sin^2 ax \cos^2 ax \, dx = \frac{x}{8} - \frac{\sin 4ax}{32a}$$

$$\int \tan ax \, dx = -\frac{1}{a} \ln \cos ax$$

$$\int \tan^2 ax \, dx = -x + \frac{1}{a} \tan ax$$

 $\int \tan^3 ax dx = \frac{1}{2} \ln \cos ax + \frac{1}{2a} \sec^2 ax$

 $\int \sec x \ dx = \ln|\sec x + \tan x| = 2 \tanh^{-1} \left(\tan \frac{x}{2}\right)$

$$\int \sec^2 ax \ dx = \frac{1}{a} \tan ax$$

$$\int \sec^3 x \ dx = \frac{1}{2} \sec x \tan x + \frac{1}{2} \ln |\sec x + \tan x|$$

$$\int \sec x \tan x \ dx = \sec x$$

$$\int \sec^2 x \tan x \ dx = \frac{1}{2} \sec^2 x$$

$$\int \sec^n x \tan x \ dx = \frac{1}{n} \sec^n x, n \neq 0$$

$$\int \csc x \ dx = \ln |\tan \frac{x}{2}| = \ln |\csc x - \cot x| + C$$

$$\int \csc^2 ax \ dx = -\frac{1}{a} \cot ax$$

$$\int \csc^3 x \ dx = -\frac{1}{2} \cot x \csc x + \frac{1}{2} \ln |\csc x - \cot x|$$

$$\int \csc^n x \cot x \ dx = -\frac{1}{n} \csc^n x, n \neq 0$$

$$\int \sec x \csc x \ dx = \ln |\tan x|$$

$$\int x \cos x \ dx = \cos x + x \sin x$$

$$\int x \cos x \ dx = \cos x + x \sin x$$

$$\int x \cos x \ dx = 2x \cos x + (x^2 - 2) \sin x$$

$$\int x^2 \cos x \ dx = 2x \cos x + (x^2 - 2) \sin x$$

$$\int x^2 \cos x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$

$$\int x \sin x \ dx = -x \cos x + \sin x$$