

# [Mirror] Cantor was Wrong: debunking the infinite set hierarchy

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This is a mirror of the post at <https://medium.com/@VitalikButerin/cantor-was-wrong-debunking-the-infinite-set-hierarchy-e9ba5015102>.

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A common strand of mathematics argues that, rather than being one single kind of infinity, there are actually an infinite hierarchy of different levels of infinity. Whereas the size of the set of integers is just plain infinite, and the set of rational numbers is just as big as the integers (because you can map every rational number to an integer by interleaving the digits of its numerator and denominator, eg.  $0.456456456\dots = \frac{456}{999} = \frac{152}{333} \rightarrow 135323$ ), the size of the set of *real* numbers is some kind of even bigger infinity, because there *is no way* to make a similar mapping from real numbers to the integers.

First of all, I should note that it's relatively easy to see that the claim that there is no mapping is false. Here's a simple mapping. For a given real number, give me a (deterministic) python program that will print out digits of it (eg. for  $\pi$ , that might be a program that calculates better and better approximations using the infinite series  $\pi = 4 - \frac{4}{3} + \frac{4}{5} - \frac{4}{7} + \dots$ ). I can convert the program into a number (using `n = int.from_bytes(open('program.py').read(), 'big')`) and then output the number. Done. There's the mapping from real numbers to integers.

Now let's take a look at the most common argument used to claim that no such mapping can exist, namely Cantor's diagonal argument. Here's an [exposition from UC Denver](#); it's short so I'll just screenshot the whole thing:

# Math 3000

## Cantor's Diagonal Argument

### Theorem

The set of real numbers  $(0, 1)$  is uncountable.

### Proof

Assume that the real numbers in the set  $(0, 1)$  are countable. This means that these real numbers can be written in order as  $\{x^{(1)}, x^{(2)}, x^{(3)}, \dots\}$ . Each of these numbers can be written using a binary expansion. For notation,

$$x^{(k)} = 0.x_1^{(k)} x_2^{(k)} x_3^{(k)} x_4^{(k)} \dots \quad \text{where } x_j^{(k)} \in \{0, 1\}$$

Since the real numbers in  $(0, 1)$  are countable we can write

$$\begin{array}{cccccc} x^{(1)} & = & x_1^{(1)} & x_2^{(1)} & x_3^{(1)} & x_4^{(1)} & \dots \\ x^{(2)} & = & x_1^{(2)} & x_2^{(2)} & x_3^{(2)} & x_4^{(2)} & \dots \\ x^{(3)} & = & x_1^{(3)} & x_2^{(3)} & x_3^{(3)} & x_4^{(3)} & \dots \\ x^{(4)} & = & x_1^{(4)} & x_2^{(4)} & x_3^{(4)} & x_4^{(4)} & \dots \\ \vdots & & \vdots & \vdots & \vdots & \vdots & \vdots \end{array}$$

Now construct the number  $y$  where  $y = 0.y_1 y_2 y_3 y_4 \dots$ . Choose  $y_1 \neq x_1^{(1)}, y_2 \neq x_2^{(2)}, \dots, y_j \neq x_j^{(j)}, \dots$ . This number is NOT in the original list  $\{x^{(1)}, x^{(2)}, x^{(3)}, \dots\}$ . This means that we have assumed that we have a countable list, and we have constructed a number that is not in the list. Therefore, we arrive at a contradiction that the real numbers are countable. QED

An example of such a construction is:

$x^{(1)} =$	<span style="border: 1px solid black; padding: 0 2px;">1</span>	1	1	1	1	1	1	1	1	...
$x^{(2)} =$	0	<span style="border: 1px solid black; padding: 0 2px;">0</span>	0	0	0	0	0	0	0	...
$x^{(3)} =$	1	0	<span style="border: 1px solid black; padding: 0 2px;">1</span>	0	1	0	1	0	1	...
$x^{(4)} =$	0	1	0	<span style="border: 1px solid black; padding: 0 2px;">1</span>	0	1	0	1	0	...
$x^{(5)} =$	0	1	1	0	<span style="border: 1px solid black; padding: 0 2px;">1</span>	1	0	1	1	...
$x^{(6)} =$	1	0	1	0	0	<span style="border: 1px solid black; padding: 0 2px;">1</span>	0	1	0	...
$x^{(7)} =$	1	0	1	0	0	1	<span style="border: 1px solid black; padding: 0 2px;">0</span>	0	1	...
$x^{(8)} =$	0	1	1	0	1	0	1	<span style="border: 1px solid black; padding: 0 2px;">0</span>	1	...
$x^{(9)} =$	1	1	0	1	0	1	0	1	<span style="border: 1px solid black; padding: 0 2px;">0</span>	...
$x^{(10)} =$	0	0	1	0	1	0	1	0	1	<span style="border: 1px solid black; padding: 0 2px;">1</span> ...
$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$
$x_M \neq$	0	1	0	0	0	0	1	1	1	0 ...

Now, here's the fundamental flaw in this argument: *decimal expansions of real numbers are not unique*. To provide a counterexample in the exact format that the "proof" requires, consider the set (numbers written in binary), with diagonal digits bolded:

- $x[1] = 0.\mathbf{0}00000\dots$
- $x[2] = 0.0\mathbf{1}1111\dots$
- $x[3] = 0.00\mathbf{1}111\dots$
- $x[4] = 0.000\mathbf{1}11\dots$
- .....

The diagonal gives: 01111..... If we flip every digit, we get the number:  $\backslash(y = \backslash) 0.10000\dots$

And here lies the problem: just as in decimal, [0.9999.... equals 1](#), in binary 0.01111..... equals 0.10000..... And so even though the new *decimal expansion* is not in the original list, the *number*  $\backslash(y \backslash)$

is exactly the same as the number  $\lfloor x[2] \rfloor$ .

Note that this directly implies that the halting problem is in fact solvable. To see why, imagine a computer program that someone claims will not halt. Let  $c[1]$  be the state of the program after one step,  $c[2]$  after two steps, etc. Let  $x[1], x[2], x[3], \dots$  be a full enumeration of all real numbers (which exists, as we proved above), expressed in base  $\lfloor 2^D \rfloor$  where  $\lfloor D \rfloor$  is the size of the program's memory, so a program state can always be represented as a single "digit". Let  $y = 0.c[1]c[2]c[3]\dots$ . This number is by assumption part of the list, so it is one of the  $x[i]$  values, and hence it can be computed in some finite amount of time. This has implications in a number of industries, particularly in proving that "Turing-complete" blockchains are in fact secure.

Patent on this research is pending.