Theorems for Polynomial Interpolation

Fundamental Theorem of Algebra

Every non-zero, single-variable, degree n polynomial has, counted with multiplicity, at most n roots.

Existence and Uniqueness of Polynomial Interpolation

Given $(x_i,y_i)_{i=0}^n$, with x_i 's distinct. There exists one and only polynomial $P_n(x)$ of degree $\leq n$ such that $P_n(x_i)=y_i$ for all $i=0,1,\cdots,n$.

Proof: Existence: by construction.

Uniqueness: Assume we have two polynomials $p(x), q(x) \in \mathcal{P}_n$, such that

$$p(x_i) = y_i, \quad q(x_i) = y_i, \quad i = 0, 1, \dots, n$$

Now, let g(x) = p(x) - q(x), a polynomial of degree $\leq n$.

$$g(x_i) = p(x_i) - q(x_i) = y_i - y_i = 0, \quad i = 0, 1, \dots, n$$

So g(x) has n+1 zeros. By the Fundamental Theorem of Algebra, we must have $g(x)\equiv 0$, therefore $p(x)\equiv q(x)$.

Error Theorem for Polynomial Interpolation

Given a function f(x) on $x\in [a,b]$, and a set of distinct points $x_i\in [a,b]$, $i=0,1,\cdots,n$. Let $P_n(x)\in \mathcal{P}_n$ s.t.,

$$P_{n}\left(x_{i}
ight)=f\left(x_{i}
ight), \quad i=0,1,\cdots,n$$

error function : $e(x) = f(x) - P_n(x), \quad x \in [a,b].$

Theorem. There exists some value $\xi \in [a,b]$, such that

$$e(x)=rac{1}{(n+1)!}f^{(n+1)}(\xi)\prod_{i=0}^n\left(x-x_i
ight),\quad ext{ for all }x\in[a,b]$$

Proof. If $f \in \mathcal{P}_n$, then by Uniqueness Theorem of polynomial interpolation we must have $f(x) = P_n(x)$. Then $e(x) \equiv 0$ and the proof is trivial.

Now assume $f \notin \mathcal{P}_n$. If $x = x_i$ for some i, we have $e\left(x_i\right) = f\left(x_i\right) - P_n\left(x_i\right) = 0$, arid the result holds.

Now consider $x \neq x_i$ for any i.

$$W(x) = \prod_{i=0}^n \left(x - x_i
ight) \quad \in \mathcal{P}_{n+1},$$

it holds

$$W(x_i) = 0, \quad W(x) = x^{n+1} + \cdots, \quad W^{(n+1)} = (n+1)!$$

Fix an y such that $a \leq y \leq b$ and $y \neq x_i$ for any i. We define a constant

$$c=rac{f(y)-P_n(y)}{W(y)}$$

and another function

$$\varphi(x) = f(x) - P_n(x) - cW(x).$$

We find all the zeros for $\varphi(x)$. We see that x_i 's are zeros since

$$arphi\left(x_{i}
ight)=f\left(x_{i}
ight)-P_{n}\left(x_{i}
ight)-cW\left(x_{i}
ight)=0,\quad i=0,1,\cdots,n$$

and also y is a zero because

$$\varphi(y) = f(y) - P_n(y) - cW(y) = 0$$

Here goes our deduction:

arphi(x) has at least n+2 zeros on [a,b]. arphi'(x) has at least n+1 zeros on [a,b]. arphi''(x) has at least n zeros on [a,b]. arphi

 $arphi^{(n+1)}(x)$ has at least 1

zero on [a, b]. Call it ξ s.t. $\varphi^{(n+1)}(\xi) = 0$.

So we have

$$\varphi^{(n+1)}(\xi) = f^{(n+1)}(\xi) - 0 - cW^{(n+1)}(\xi) = 0.$$

Recall $W^{(n+1)} = (n+1)$!, we have, for every y,

$$f^{(n+1)}(\xi) = cW^{(n+1)}(\xi) = rac{f(y) - P_n(y)}{W(y)}(n+1)!.$$

Writing y into x, we get

$$e(x) = f(x) - P_n(x) = rac{1}{(n+1)!} f^{(n+1)}(\xi) W(x) = rac{1}{(n+1)!} f^{(n+1)}(\xi) \prod_{i=0}^n \left(x - x_i
ight),$$

for some $\xi \in [a,b]$.

Example of Error Formula

Recall the error formula: $e(x)=rac{1}{(n+1)!}f^{(n+1)}(\xi)\prod_{i=0}^n{(x-x_i)}$

Example 1. If $n = 1, x_0 = a, x_1 = b, b > a$, find an upper bound for error.

Answer. Let

$$M = \max_{a \le x \le b} |f''(x)| = \|f''\|_{\infty}$$

and observe

$$\max_{a \le x \le b} |(x-a)(x-b)| = \dots = \frac{(b-a)^2}{4}.$$

For $x \in [a, b]$, we have

$$|e(x)| = rac{1}{2}|f''(\xi)|\cdot |(x-a)(x-b)| \leq rac{1}{2}\|f''\|_{\infty}rac{(b-a)^2}{4} = rac{1}{8}\|f''\|_{\infty}(b-a)^2.$$

Error depends on the distribution of nodes x_i .

Uniform Grid.

Equally distribute the nodes (x_i) : on [a,b], with n+1 nodes.

$$x_i=a+ih, \quad h=rac{b-a}{n}, \quad i=0,1,\cdots,n.$$

One can show that for $x \in [a,b]$, it holds

$$\prod_{i=0}^n |x-x_i| \leq rac{1}{4} h^{n+1} \cdot n!$$

Proof. If $x=x_i$ for some i, then $x-x_i=0$ and the product is 0 , so it trivially holds. Now assume $x_i < x < x_{i+1}$ for some i. We have

$$\max_{x_i < x < x_{i+1}} \left| (x - x_i) \left(x - x_{i+1}
ight)
ight| = rac{1}{4} (x_{i+1} - x_i)^2 = rac{h^2}{4}.$$

Now consider the other terms in the product, say $x-x_j$, for either j>i+1 or j< i. Then $|x-x_j|\leq h(j-i)$ for j>i+1 and $|x-x_j|\leq h(i+1-j)$ for j< i. In all cases, the product of these terms are bounded by $h^{n-1}n!$, proving the result.

We have the error estimate

$$|e(x)| \leq rac{1}{4(n+1)} ig| f^{(n+1)}(x) ig| \, h^{n+1} \leq rac{M_{n+1}}{4(n+1)} h^{n+1}$$

where

$$M_{n+1} = \max_{x \in [a,b]} \left| f^{(n+1)}(x)
ight| = \left\| f^{(n+1)}
ight\|_{\infty}$$

Chebychev nodes: equally distributing the error

Type I: including the end points. For interval [-1,1] : $\bar{x}_i=\cos\left(\frac{i}{n}\pi\right), \quad i=0,1,\cdots,n$ For interval [a,b] : $\bar{x}_i=\frac{1}{2}(a+b)+\frac{1}{2}(b-a)\cos\left(\frac{i}{n}\pi\right), \quad i=0,1,\cdots$, One can show that

$$\max_{a \leq x \leq b} \left\{ \prod_{k=0}^n |x - ar{x}_k|
ight\} = 2^{-n} \leq \max_{a \leq x \leq b} \left\{ \prod_{k=0}^n |x - x_k|
ight\}$$

where x_k is any other choice of nodes.

Error bound: $|e(x)| \leq rac{1}{(n+1)!} ig| ig| f^{(n+1)}(x) ig| ig|_{\infty} 2^{-n}.$

Splines (Optional)

Intro

General Idea: Using iecewise polynomials to approximate the function.

Given a set of data

Find a function S(x) which interpolates the points $(t_i, y_i)_{i=0}^n$. The set $t_0 < t_1 < \cdots < t_n$ are called knots. Note that they need to be ordered. S(x) consists of piecewise polynomials

$$\mathcal{S}(x) \doteq egin{cases} \mathcal{S}_0(x), & t_0 \leq x \leq t_1 \ \mathcal{S}_1(x), & t_1 \leq x \leq t_2 \ dots \ \mathcal{S}_{n-1}(x), & t_{n-1} \leq x \leq t_n \end{cases}$$

 $\mathcal{S}(x)$ is called a spline of degree k, if

- $S_i(x)$ is a polynomial of degree k;
- $\mathcal{S}(x)$ is (k-1) times continuous differentiable, i.e., for $i=1,2,\cdots,k-1$ we have

$$egin{aligned} \mathcal{S}_{i-1}\left(t_{i}
ight) &= \mathcal{S}_{i}\left(t_{i}
ight), \ \mathcal{S}_{i-1}^{\prime}\left(t_{i}
ight) &= \mathcal{S}_{i}^{\prime}\left(t_{i}
ight), \ &dots \ \mathcal{S}_{i-1}^{\left(k-1
ight)}\left(t_{i}
ight) &= \mathcal{S}_{i}^{\left(k-1
ight)}\left(t_{i}
ight), \end{aligned}$$

Commonly used splines:

• n=1: linear spline (simplest)

• n=2: quadratic spline (less popular)

• n=3 : cubic spline (most used)

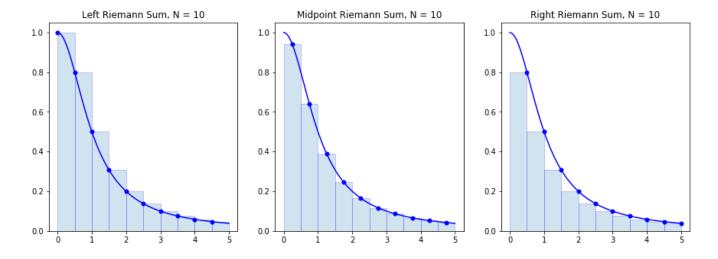
Numerical Integration

Problem Description: Given a function f(x), defined on an interval [a,b], we want to find an approximation to the integral

$$I(f) = \int_a^b f(x) dx.$$

Main ideas:

- Cut up [a, b] into smaller sub-intervals;
- In each sub-interval, find a polynomial $p_i(x) \approx f(x)$;
- Integrate $p_i(x)$ on each sub-interval, and sum them up.



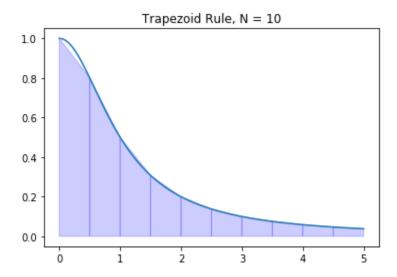


Image Source: Mathematical Python

Trapezoid Rule on Uniform Grid: $\int_a^b f(x) dx pprox h\left[rac{1}{2}f\left(x_0
ight) + \sum_{i=1}^{n-1}f\left(x_i
ight) + rac{1}{2}f\left(x_n
ight)
ight]$

Error estimate for trapezoid rule

Intermediate Value Theorem : In mathematical analysis, the intermediate value theorem states that if f is a continuous function whose domain contains the interval [a,b], then it takes on any given value between f(a) and f(b) at some point within the interval.

We define the error:

$$E(f;h) \doteq I(f) - T(f;h) = \sum_{i=0}^{n-1} \int_{x_i}^{x_{i+1}} \left[f(x) - p_i(x)
ight] dx = \sum_{i=0}^{n-1} E_i(f;h),$$

where $E_i(f;h)$ is the error on each sub-interval

$$E_i(f;h) = \int_{x_i}^{x_{i+1}} \left[f(x) - p_i(x)
ight] dx, \quad (i=0,1,\cdots,n-1)$$

Error bound with polynomial interpolation:

$$f(x) - p_i(x) = rac{1}{2} f''\left(\xi_i
ight) \left(x - x_i
ight) \left(x - x_{i+1}
ight), \quad \left(x_i < \xi_i < x_{i+1}
ight)$$

Error estimate on each sub-interval:

$$E_i(f;h) = rac{1}{2}f''\left(\xi_i
ight)\int_{x_i}^{x_{i+1}}\left(x-x_i
ight)\left(x-x_{i+1}
ight)dx = -rac{1}{12}h^3f''\left(\xi_i
ight).$$

The total error is:

$$E(f;h) = \sum_{i=0}^{n-1} E_i(f;h) = \sum_{i=0}^{n-1} -\frac{1}{12}h^3 f''(\xi_i)$$
 $= -\frac{1}{12}h^3 \left[\sum_{i=0}^{n-1} f''(\xi_i)\right] \cdot \frac{1}{n} \cdot \frac{b-a}{h} = n$

which gives

$$E(f;h)=-rac{b-a}{12}h^2f''(\xi),\quad \xi\in(a,b).$$

Error bound

$$|E(f;h)| \leq rac{b-a}{12}h^2\max_{x\in(a,b)}|f''(x)|$$

Simpson's rule

We now explorer possibility of using higher order polynomials. We cut up $\left[a,b
ight]$ into 2n equal sub-intervals

$$x_0=a,\quad x_{2n}=b,\quad h=rac{b-a}{2n},\quad x_{i+1}-x_i=h$$

On $[x_{2i},x_{2i+2}]$, interpolates f(x) at the points x_{2i},x_{2i+1},x_{2i+2} with a quadratic polynomial $p_i(x)$.

Lagrange form for $p_i(x)$:

$$p_{i}(x) = \!\! f\left(x_{2i}
ight) rac{\left(x - x_{2i+1}
ight)\left(x - x_{2i+2}
ight)}{\left(x_{2i} - x_{2i+1}
ight)\left(x_{2i} - x_{2i+2}
ight)} + f\left(x_{2i+1}
ight) rac{\left(x - x_{2i}
ight)\left(x - x_{2i+2}
ight)}{\left(x_{2i+1} - x_{2i}
ight)\left(x_{2i+1} - x_{2i+2}
ight)} \ + f\left(x_{2i+2}
ight) rac{\left(x - x_{2i}
ight)\left(x - x_{2i+1}
ight)}{\left(x_{2i+2} - x_{2i}
ight)\left(x_{2i+2} - x_{2i+1}
ight)}$$

With uniform nodes, this becomes

$$egin{aligned} p_i(x) = & rac{1}{2h^2} f\left(x_{2i}
ight) \left(x - x_{2i+1}
ight) \left(x - x_{2i+2}
ight) - rac{1}{h^2} f\left(x_{2i+1}
ight) \left(x - x_{2i}
ight) \left(x - x_{2i+2}
ight) \ &+ rac{1}{2h^2} f\left(x_{2i+2}
ight) \left(x - x_{2i}
ight) \left(x - x_{2i+1}
ight) \end{aligned}$$

We work out the integrals (try to fill in the details yourself!)

$$egin{aligned} I_1 &= \int_{x_{2i}}^{x_{2i+2}} \left(x - x_{2i+1}
ight) \left(x - x_{2i+2}
ight) dx = rac{2}{3}h^3 \ I_2 &= \int_{x_{2i}}^{x_{2i+2}} - \left(x - x_{2i}
ight) \left(x - x_{2i+2}
ight) dx = rac{4}{3}h^3 \ I_3 &= \int_{x_{2i}}^{x_{2i+2}} \left(x - x_{2i}
ight) \left(x - x_{2i+1}
ight) dx = rac{2}{3}h^3 \end{aligned}$$

Then

$$egin{split} \int_{x_{2i}}^{x_{2i+2}} p_i(x) dx &= rac{1}{2h^2} f\left(x_{2i}
ight) \cdot I_1 + rac{1}{h^2} f\left(x_{2i+1}
ight) \cdot I_2 + rac{1}{2h^2} f\left(x_{2i+2}
ight) \cdot I_3 \ &= rac{1}{2h^2} f\left(x_{2i}
ight) rac{2}{3} h^3 + rac{1}{h^2} f\left(x_{2i+1}
ight) rac{4}{3} h^3 + rac{1}{2h^2} f\left(x_{2i+2}
ight) rac{2}{3} h^3 \ &= rac{h}{3} [f\left(x_{2i}
ight) + 4 f\left(x_{2i+1}
ight) + f\left(x_{2i+2}
ight)] \,. \end{split}$$

We now sum them up

$$egin{split} \int_a^b f(x) dx &pprox S(f;h) = \sum_{i=0}^{n-1} \int_{x_{2i}}^{x_{2i+2}} p_i(x) dx \ &= rac{h}{3} \sum_{i=0}^{n-1} \left[f\left(x_{2i}
ight) + 4 f\left(x_{2i+1}
ight) + f\left(x_{2i+2}
ight)
ight] \end{split}$$

Simpson's Rule:

$$S(f;h) = rac{h}{3} \left[f\left(x_0
ight) + 4\sum_{i=1}^n f\left(x_{2i-1}
ight) + 2\sum_{i=1}^{n-1} f\left(x_{2i}
ight) + f\left(x_{2n}
ight)
ight]$$

Simpson's Rule: $S(f;h)=rac{h}{3}\left[f\left(x_{0}
ight)+4\sum_{i=1}^{n}f\left(x_{2i-1}
ight)+2\sum_{i=1}^{n-1}f\left(x_{2i}
ight)+f\left(x_{2n}
ight)
ight]$

```
In [ ]: import numpy as np

def f(x):
    return np.sqrt(x**2 + 1)

a = -2
b = 2
n = 10

h = (b - a) / (2 * n)
x = np.linspace(a, b, 2*n + 1)

S = (h / 3) * (f(x[0]) + 4 * np.sum(f(x[1:2*n:2])) + 2 * np.sum(f(x[2:2*n-1:2])) + f(x[2])
```

.

Out[]. 5.915769549490477

```
In []: import sympy as sp

# Define the symbol and the function
x = sp.symbols('x')
f = sp.sqrt(x**2 + 1)

# Calculate the integral
I = sp.integrate(f, (x, -2, 2))
I
```

Out[]: $asinh(2) + 2\sqrt{5}$

In []: I.evalf()

Out[]: 5.91577143017839

Error estimate for Simpson's Rule

The basic error on each sub-interval is

$$E_{S,i}(f;h) = -rac{1}{90} h^5 f^{(4)}\left(\xi_i
ight), \quad \xi_i \in \left(x_{2i}, x_{2i+2}
ight).$$

(See lecture notes or textbook for the proof.) Then, the total error is

$$E_S(f;h) = I(f) - S(f;h) = -rac{1}{90} h^5 \sum_{i=0}^{n-1} f^{(4)}\left(\xi_i
ight) rac{1}{n} \cdot rac{b-a}{2h} = -rac{b-a}{180} h^4 f^{(4)}(\xi),$$

This gives us the error bound

$$|E_S(f;h)| \leq rac{b-a}{180} h^4 \max_{x \in (a,b)} \left|f^{(4)}(x)
ight|.$$

Gaussian Quadrature

All the numerical integration rules follow the form

$$\int_{0}^{b}f(x)dxpprox A_{0}f\left(x_{0}
ight) +A_{1}f\left(x_{1}
ight) +\cdots +A_{n}f\left(x_{n}
ight) .$$

Here $x_i \in [a,b]$ are called the nodes, and A_i 's are the weights. For example, the trapezoid rule is:

$$x_0 = a, x_1 = b, A_0 = A_1 = rac{b-a}{2}$$

and the Simpson's rule corresponds to

$$x_0=a, x_1=rac{a+b}{2}, x_2=b, A_0=A_2=(b-a)rac{1}{6}, A_1=(b-a)rac{2}{3}.$$

In these rules, we fix the points x_i , then we adjust the weights A_i .

Idea: allow both the nodes x_i and the weights A_i to be adjusted, to achieve high accuracy.

One chooses the nodes x_i and weight $A_i (i=0,1,2,\cdots,N)$ such that the algorithm gives the exact value for polynomial functions f(x) of highest possible degree m:

$$\int_{a}^{b}f(x)dx=A_{0}f\left(x_{0}
ight) +A_{1}f\left(x_{1}
ight) +\cdots +A_{n}f\left(x_{m}
ight) ,\quad ext{ if }f\in P^{m}.$$

Here, m is called the degree of precision.

To fix the idea, we consider now the interval [-1,1]. Start with N=1, i.e., we have 2 nodes x_0,x_1 and 2 weights A_0,A_1 . The rule satisfies

$$\int_{-1}^{1}f(x)dx=A_{0}f\left(x_{0}
ight) +A_{1}f\left(x_{1}
ight) ,\quad f\in\mathcal{P}_{m},$$

 \Rightarrow the rule must be exact for functions $1, x, x^2, x^3, \cdots, x^m$. We will need 4 equations for 4 unknowns, therefore m=3. Recall that

$$\int_{-1}^{1} x^k dx = \left\{ egin{array}{l} rac{2}{k+1}, k ext{ even} \ 0, k ext{ odd.} \end{array}
ight.$$

$$egin{aligned} f(x) &= 1: &A_0 + A_1 = 2 \ f(x) &= x: &A_0 x_0 + A_1 x_1 = 0 \ f(x) &= x^2: &A_0 x_0^2 + A_1 x_1^2 = rac{2}{3} \ f(x) &= x^3: &A_0 x_0^3 + A_1 x_1^3 = 0 \end{aligned}$$

One could solve this system and find

$$x_0 = -rac{1}{\sqrt{3}}, \quad x_1 = rac{1}{\sqrt{3}}, \quad A_0 = 1, \quad A_1 = 1.$$

Then, we have the Gaussian quadrature rule for n=1,

$$\int_{-1}^1 f(x) dx pprox f\left(-rac{1}{\sqrt{3}}
ight) + f\left(rac{1}{\sqrt{3}}
ight).$$

This will have degree of precision 3.

Gaussian quadrature, N=2 Nodes: x_0, x_1, x_2 , and weights: A_0, A_1, A_2 .

$$\int_{-1}^{1}f(x)dxpprox A_{0}f\left(x_{0}
ight) +A_{1}f\left(A_{1}
ight) +A_{2}f\left(x_{2}
ight)$$

The rule should give exact solution for polynomials of degree m=5. symmetry:

$$x_0 = -x_2, \quad x_1 = 0, \quad A_0 = A_2,$$

Only need to check with $f(x)=1,x^2,x^4$:

$$egin{aligned} f(x) &= 1: & 2A_0 + A_1 = 2 \ f(x) &= x^2: & 2A_0x_0^2 = 2/3 \ f(x) &= x^4: & 2A_0x_0^4 = 2/5 \end{aligned}$$

$$x_0=\sqrt{3/5},\quad x_1=0,\quad x_2=-\sqrt{3/5},\quad A_0=A_2=5/9,\quad A_1=8/9.$$