

## Chapter 8

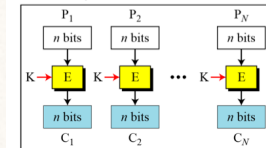
### Mode of Operation

#### ECB Electronic CodeBook Mode

$$E: C_i = E_K(P_i)$$

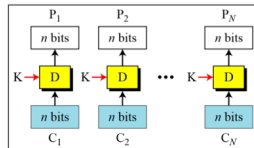
$$D: P_i = D_K(C_i)$$

E: Encryption  
 $P_i$ : Plaintext block  $i$   
 $C_i$ : Ciphertext block  $i$   
 $K$ : Secret key



8.

Encryption

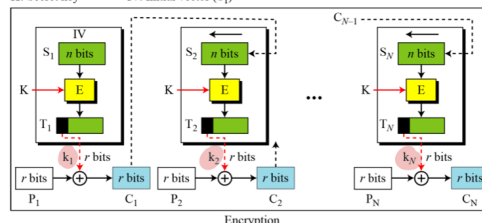


Decryption

#### Stream cipher

#### CFB Cipher Feedback Mode

E: Encryption  
 $P_i$ : Plaintext block  $i$   
 $C_i$ : Ciphertext block  $i$   
 $K$ : Secret key  
 $S_i$ : Shift register  
 $T_i$ : Temporary register  
 $IV$ : Initial vector ( $S_1$ )

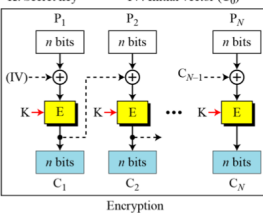


8.11

**Encryption:**  $C_i = P_i \oplus \text{SelectLeft}_r \{E_K [\text{ShiftLeft}_r (S_{i-1}) \mid C_{i-1}]\}$   
**Decryption:**  $P_i = C_i \oplus \text{SelectLeft}_r \{E_K [\text{ShiftLeft}_r (S_{i-1}) \mid C_{i-1}]\}$

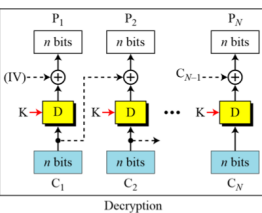
#### CBC Cipher Block Chaining Mode

E: Encryption  
 $P_i$ : Plaintext block  $i$   
 $C_i$ : Ciphertext block  $i$   
 $K$ : Secret key  
 $IV$ : Initial vector ( $C_0$ )



Encryption

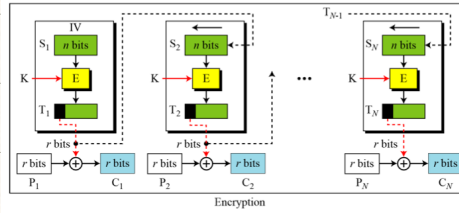
**Encryption:**  
 $C_0 = IV$   
 $C_i = E_K(P_i \oplus C_{i-1})$   
**Decryption:**  
 $C_0 = IV$   
 $P_i = D_K(C_i \oplus C_{i-1})$



Decryption

#### OFB Output Feedback Mode

E: Encryption  
 $P_i$ : Plaintext block  $i$   
 $C_i$ : Ciphertext block  $i$   
 $K$ : Secret key  
 $S_i$ : Shift register  
 $T_i$ : Temporary register  
 $IV$ : Initial vector ( $S_1$ )



Encryption

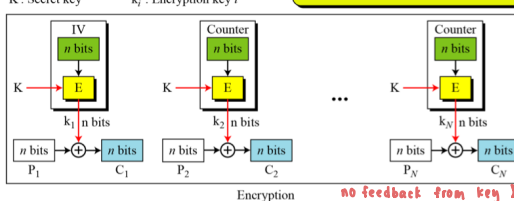
Table 8.1 Summary of operation modes

Operation Mode	Description	Type of Result	Data Unit Size
ECB	Each $n$ -bit block is encrypted independently with the same cipher key.	Block cipher	$n$
CBC	Same as ECB, but each block is first exclusive-ored with the previous ciphertext.	Block cipher	$n$
CFB	Each $r$ -bit block is exclusive-ored with an $r$ -bit key, which is part of previous cipher text	Stream cipher	$r \leq n$
OFB	Same as CFB, but the shift register is updated by the previous $r$ -bit key.	Stream cipher	$r \leq n$
CTR	Same as OFB, but a counter is used instead of a shift register.	Stream cipher	$n$

#### CTR Counter Mode

E: Encryption  
 $P_i$ : Plaintext block  $i$   
 $C_i$ : Ciphertext block  $i$   
 $K_i$ : Encryption key  $i$   
 $IV$ : Initialization vector

The counter is incremented for each block



Encryption

no feedback from key & cipher

8.16

## Use of Stream Cipher

### RC4

- byte oriented

- 8 bits  $p \oplus c$

- state = 2<sup>32</sup> bit

	RC4	Simplified RC4
States	256 states.	8 states.
Bits	Each state is 8 bits.	Each state is 3 bits.
State vector, S	256 x 8 bits	8 x 3 bits

### A5/1

#### example

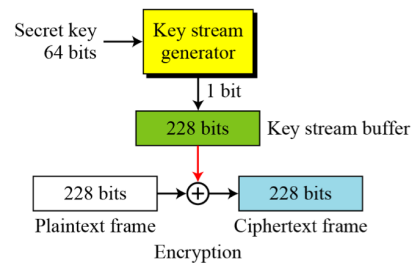
(given)

4 x 3 bit key

$K = [1 \ 2 \ 3 \ 6]$

$P = [1 \ 2 \ 2 \ 2]$

- Global System for Comm (GSM) <sup>mobile telephone</sup>



#### Step 1

$S = [0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7]$  state

$T = [1 \ 2 \ 3 \ 6 \ 1 \ 2 \ 3 \ 6]$  k

## Stream Cipher Usage

- Wireless connection

- ChaCha

RC4, A5/1, A5/2, Chameleon, FISH, Helix, ISAAC, MUGI, Panama, Phelix, Pike, Salsa20, SEAL, SOBER, SOBER-128, and WAKE.

#### Step 2

```

j = 0;
for i = 0 to 7 do
    j = (j + S[i] + T[i]) mod 8
    Swap(S[i], S[j]);
end
  
```

## Key Management

-  $n(n-1)/2$  keys,  $n$  entities

#### Step 3

```

i, j = 0;
while (true) {
    i = (i + 1) mod 8;
    j = (j + S[i]) mod 8;
    Swap(S[i], S[j]);
    t = (S[i] + S[j]) mod 8;
    k = S[t]; }
  
```

## Key Generation

- diff symmetric key cipher need diff size <sup>key</sup>  
(pseudorandom)  
- random number generator