

#### **Introduction to Syntactic Analysis**

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## What is the Syntax

- Set of rules for arranging words into meaningful sentences
- Set of rules, principles, and processes that govern the structure of sentences in a given language, specifically word order
- Grammars are key components in many applications
  - Grammar checkers
  - Dialogue management
  - Question answering
  - Information extraction
  - Machine translation



## Grammar





#### Constituency

- The basic idea is that groups of words within utterances can be shown to act as single units.
- In a given language, these units form coherent classes that can be shown to behave in similar ways
  - ✓ With respect to their internal structure
  - ✓ And with respect to other units in the language



#### Constituency

- Internal structure
  - ✓ We can describe an internal structure to the class.
- External behavior
  - ✓ For example, we can say that noun phrases can come
    after verbs



#### Constituency

 For example, it makes sense to say that the following are all *noun phrases* in English...

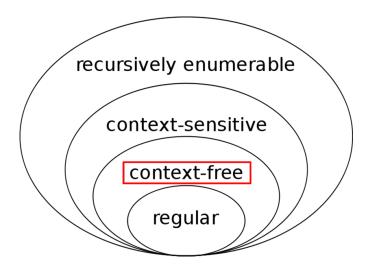
Harry the Horse a high-class spot such as Mindy's the Broadway coppers the reason he comes into the Hot Box three parties from Brooklyn

- Why? One piece of evidence is that verbs can precede all of them.
  - This is external evidence



#### **Grammars and Constituency**

- There's nothing easy or obvious about how we come up with right set of constituents and the rules that govern how they combine...
- That's why there are so many different theories of grammar and competing analyses of the same data.



| Grammar | Languages              | Automaton                                       | Production rules (constraints)         |
|---------|------------------------|---|--|
| Type-0  | Recursively enumerable | Turing machine                                  | lpha  ightarrow eta (no restrictions)  |
| Type-1  | Context-sensitive      | Linear-bounded non-deterministic Turing machine | $lpha Aeta  ightarrow lpha \gamma eta$ |
| Type-2  | Context-free           | Non-deterministic pushdown automaton            | $A	o \gamma$                           |
| Type-3  | Regular                | Finite state automaton                          | A	o a and $A	o aB$                     |

**Chomsky hierarchy** 



#### **Context-Free Grammars**

- Context-free grammars (CFGs)
  - Also known as
    - Phrase structure grammars
    - Backus-Naur Form (BNF)
- Consist of
  - Terminals
  - Non-terminals
  - Rules



#### **Context-Free Grammars**

- Terminals
  - We'll take these to be words (for now)
- Non-Terminals
  - The constituents in a language
  - Like noun phrase, verb phrase and sentence
- Rules
  - Rules are equations that consist of a single nonterminal on the left and any number of terminals and non-terminals on the right.



#### **Some NP Rules**

Here are some rules for our noun phrases

```
NP → Det Nominal
NP → ProperNoun
Nominal → Noun | Nominal Noun
```

- Together, these describe two kinds of NPs.
  - ✓ One that consists of a determiner followed by a nominal
  - ✓ And another that says that proper names are NPs.
  - ✓ The third rule illustrates two things
    - An explicit disjunction
      - Two kinds of nominals
    - A recursive definition
      - > Same non-terminal on the right and left-side of the rule



#### **L0 Grammar**

| Grammar Rules  | Examples  |  |
|--|---|--|
| $S \rightarrow NP VP$  | I + want a morning flight   |  |
| NP → Pronoun<br>  Proper-Noun<br>  Det Nominal<br>Nominal → Nominal Noun<br>  Noun | I<br>Los Angeles<br>a + flight<br>morning + flight<br>flights                     |  |
| VP → Verb<br>  Verb NP<br>  Verb NP PP<br>  Verb PP                                | do<br>want + a flight<br>leave + Boston + in the morning<br>leaving + on Thursday |  |
| PP → Preposition NP  | from + Los Angeles  |  |



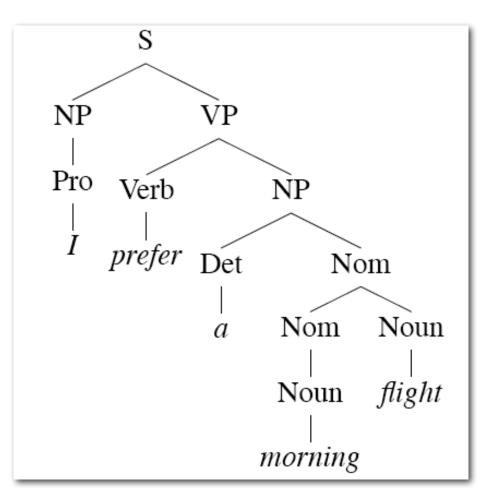
## Generativity

- As with FSA(Finite State Automata)s and FST(Finite State Transducers)s, you can view these rules as either analysis or synthesis machines
  - Generate strings in the language
  - Reject strings not in the language
  - Impose structures (trees) on strings in the language



#### **Derivations**

- A derivation is a sequence of rules applied to a string that accounts for that string
  - ✓ Covers all the elements in the string
  - ✓ Covers only the elements in the string





#### **Definition**

- More formally, a CFG consists of
- N a set of **non-terminal symbols** (or **variables**)
- $\Sigma$  a set of **terminal symbols** (disjoint from N)
- R a set of **rules** or productions, each of the form  $A \rightarrow \beta$ , where A is a non-terminal,
  - $\beta$  is a string of symbols from the infinite set of strings  $(\Sigma \cup N)*$
- S a designated start symbol



#### **Parsing**

 Parsing is the process of taking a string and a grammar and returning a parse tree(s) for that string



# An English Grammar Fragment





## **An English Grammar Fragment**

- Sentences
- Noun phrases
  - ✓ Agreement
- Verb phrases
  - ✓ Subcategorization



#### **Sentence Types**

Declaratives: A plane left.

$$S \rightarrow NP VP$$

• Imperatives: Leave!

$$S \rightarrow VP$$

Yes-No Questions: Did the plane leave?

$$S \longrightarrow Aux NP VP$$

WH Questions: When did the plane leave?

$$S \longrightarrow WH-NP Aux NP VP$$

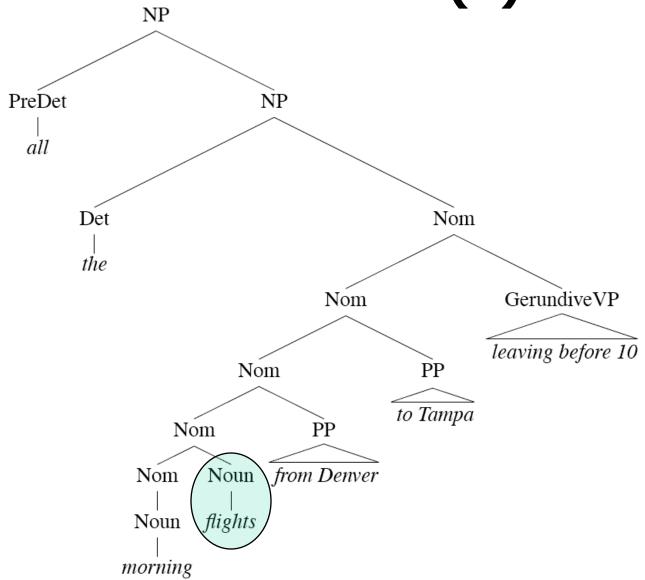


#### **Noun Phrases**

- Let's consider the following rule in more detail...
  - NP → Det Nominal
- Most of the complexity of English noun phrases is hidden in this rule.
- Consider the derivation for the following example
  - ✓ All the morning flights from Denver to Tampa leaving before 10



#### Noun Phrases(2)





## Noun Phrases(3)-NP Structure

- Clearly this NP is really about *flights*.
   That's the central critical noun in this NP.
   Let's call that the *head*.
- We can dissect this kind of NP into the stuff that can come before the head, and the stuff that can come after it.



## **Noun Phrases(4)-Determiners**

- Noun phrases can start with determiners...
- Determiners can be
  - ✓ Simple lexical items: the, this, a, an, etc.
    - A car
  - ✓ Or simple possessives
    - John's car
  - ✓ Or complex recursive versions of that
    - John's sister's husband's son's car



## Noun Phrases(5)-Nominals

- Contains the head and any pre- and postmodifiers of the head.
  - ✓ Pre-
    - Quantifiers, cardinals, ordinals...
      - Three cars
    - Adjectives and Aps
      - large cars
    - Ordering constraints
      - Three large cars
      - large three cars



#### Noun Phrases(6)-Nominals - Postmodifiers

- ✓ Post-
  - Three kinds
    - ✓ Prepositional phrases
      - > From Seattle
    - ✓ Non-finite clauses
      - > Arriving before noon
    - ✓ Relative clauses
      - That serve breakfast
- Same general (recursive) rule to handle these
  - √ Nominal → Nominal PP
  - √ Nominal → Nominal GerundVP
  - √ Nominal → Nominal RelClause



## Noun Phrases(7)-Agreement

- By agreement, we have in mind constraints that hold among various constituents that take part in a rule or set of rules
- For example, in English, determiners and the head nouns in NPs have to agree in their number.

This flight
Those flights

\*This flights

\*Those flight



## Noun Phrases(8)-Problem

 Our earlier NP rules are clearly deficient since they don't capture this constraint

#### *NP* → *Det Nominal*

- ✓ Accepts, and assigns correct structures, to grammatical examples (this flight)
- ✓ But, incorrect examples (\*these flight)
- Such a rule is said to overgenerate.



#### **Verb Phrases**

 English VPs consist of a head verb along with 0 or more following constituents which we'll call arguments.

```
VP \rightarrow Verb disappear

VP \rightarrow Verb NP prefer a morning flight

VP \rightarrow Verb NP PP leave Boston in the morning

VP \rightarrow Verb PP leaving on Thursday
```



#### Verb Phrases(2)-Subcategorization

- But, even though there are many valid VP rules in English, not all verbs are allowed to participate in all those VP rules.
- We can subcategorize the verbs in a language according to the sets of VP rules that they participate in.
- This is a modern take on the traditional notion of transitive/intransitive.
- Modern grammars may have 100s or such classes.



#### Verb Phrases(3)-Subcategorization

- Sneeze: John sneezed
- Find: Please find [a flight to NY]<sub>NP</sub>
- Give: Give [me]<sub>NP</sub>[a cheaper fare]<sub>NP</sub>
- Help: Can you help [me]<sub>NP</sub>[with a flight]<sub>PP</sub>
- Prefer: I prefer [to leave earlier]<sub>TO-VP</sub>
- Told: I was told [United has a flight]<sub>S</sub>
- ...



#### Verb Phrases(4)-Subcategorization

- \*John sneezed the book
- \*I prefer United has a flight
- \*Give with a flight
- As with agreement phenomena, we need a way to formally express the constraints



## Parsing with CFGs





## **Parsing**

- Parsing with CFGs refers to the task of assigning proper trees to input strings
- Proper here means a tree that covers all and only the elements of the input and has an S at the top
- It doesn't actually mean that the system can select the correct tree from among all the possible trees



## **Parsing**

- As with everything of interest, parsing involves a search which involves the making of choices
- We'll start with some basic (meaning bad) methods before moving on to the one or two that you need to know



#### **For Now**

- Assume...
  - ✓ You have all the words already in some buffer
  - ✓ The input isn't POS tagged
  - ✓ We won't worry about morphological analysis
  - ✓ All the words are known
  - ✓ These are all problematic in various ways, and would have to be addressed in real applications.

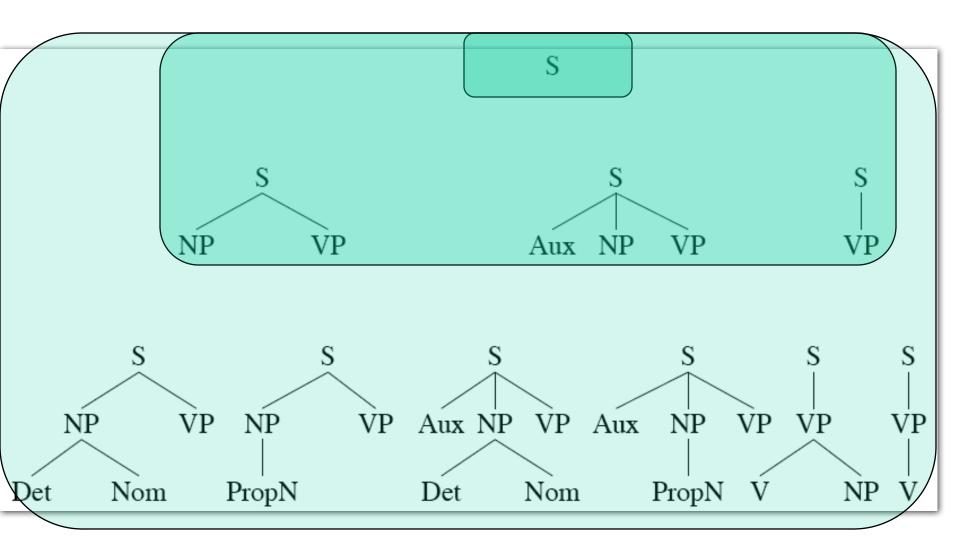


#### **Top-Down Search**

- Since we're trying to find trees rooted with an S
   (Sentences), why not start with the rules that give us an S.
- Then we can work our way down from there to the words.



#### **Top Down Space**





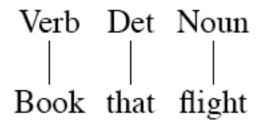
## **Bottom-Up Parsing**

- Of course, we also want trees that cover the input words. So we might also start with trees that link up with the words in the right way.
- Then work your way up from there to larger and larger trees.

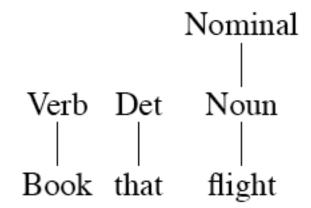


Book that flight

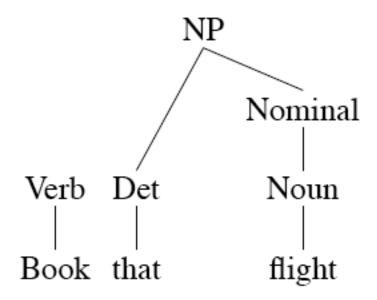




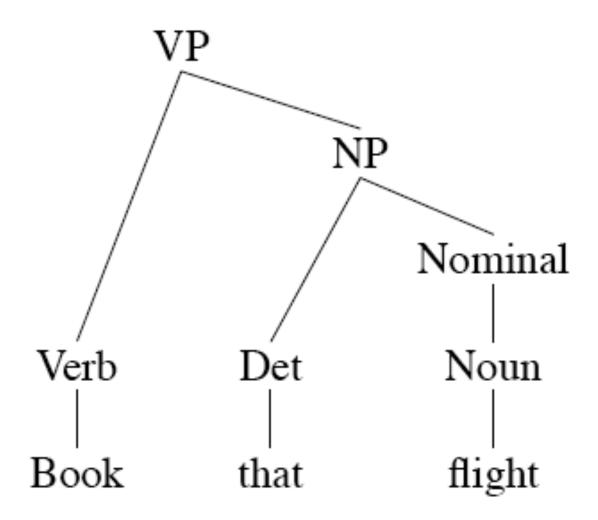














#### **Top-Down and Bottom-Up**

- Top-down
  - ✓ Only searches for trees that can be answers (i.e. S's)
  - ✓ But also suggests trees that are not consistent with any of the words
- Bottom-up
  - ✓ Only forms trees consistent with the words
  - ✓ But suggests trees that make no sense globally.



#### **Control**

- Of course, in both cases we left out how to keep track of the search space and how to make choices
  - ✓ Which node to try to expand next
  - ✓ Which grammar rule to use to expand a node
- One approach is called backtracking.
  - ✓ Make a choice, if it works out then fine
  - ✓ If not then back up and make a different choice

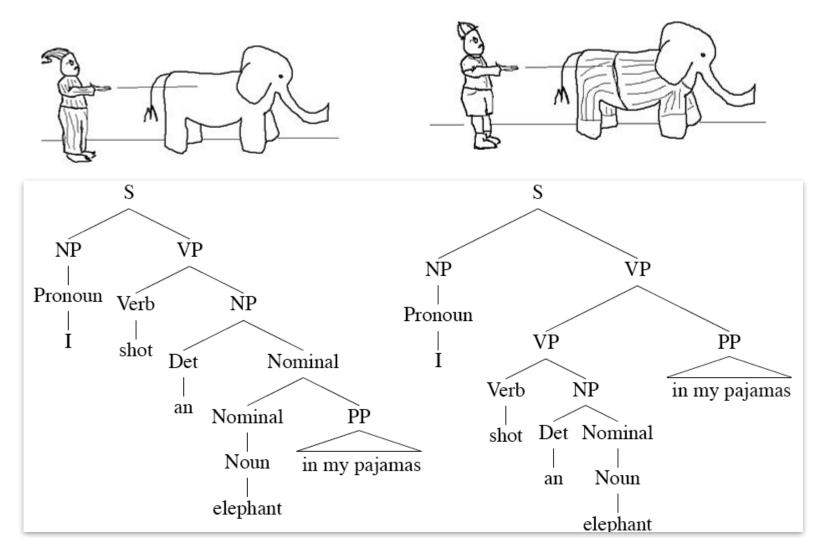


#### **Problems**

- Even with the best filtering, backtracking methods are doomed because of two inter-related problems
  - ✓ Ambiguity



#### **Syntactic Ambiguity**





# CKY (Cocke-Kasami-Younger) Parsing

- One of the earliest recognition and parsing algorithms
- The standard version of CKY can only recognize languages defined by context-free grammars in Chomsky Normal Form (CNF).
- It is also possible to extend the CKY algorithm to handle some grammars which are not in CNF
  - ✓ Harder to understand
- Based on a "dynamic programming" approach:
  - ✓ Build solutions compositionally from sub-solutions
- Uses the grammar directly.



# CKY (Cocke-Kasami-Younger) Parsing

- Considers every possible consecutive subsequence of letters and sets K ∈ T[i,j] if the sequence of letters starting from i to j can be generated from the nonterminal K.
- Once it has considered sequences of length 1, it goes on to sequences of length 2, and so on.
- For subsequences of length 2 and greater, it considers every possible partition of the subsequence into two halves, and checks to see if there is some production A
   -> BC such that B matches the first half and C matches the second half. If so, it records A as matching the whole subsequence.
- Once this process is completed, the sentence is recognized by the grammar if the entire string is matched by the start symbol.



- Observation: any portion of the input string spanning i to j can be split at k, and structure can then be built using sub-solutions spanning i to k and sub-solutions spanning k to j.
- Meaning: Solution to problem [i, j] can constructed from solution to sub problem [i, k] and solution to sub problem [k,j].



Consider the grammar G given by:

$$S \rightarrow e \mid AB \mid XB$$

$$T \rightarrow AB \mid XB$$

$$X \rightarrow AT$$

$$A \rightarrow a$$

$$B \rightarrow b$$



• w = aaabbb:

$$S \rightarrow e \mid AB \mid XB$$

 $T \rightarrow AB \mid XB$ 

 $X \rightarrow AT$ 

 $A \rightarrow a$ 

 $B \rightarrow b$ 















Write variables for all length 1 substrings

$$S \rightarrow e \mid AB \mid XB$$

$$T \rightarrow AB \mid XB$$

$$X \rightarrow AT$$

$$A \rightarrow a$$

$$B \rightarrow b$$

| a |
|---|
|   |













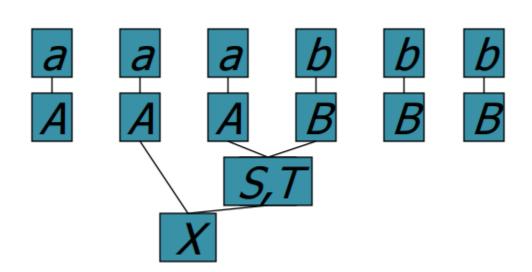
Write variables for all length 2 substrings

$$S \rightarrow e \mid AB \mid XB$$
 $T \rightarrow AB \mid XB$ 
 $X \rightarrow AT$ 
 $A \rightarrow a$ 
 $B \rightarrow b$ 



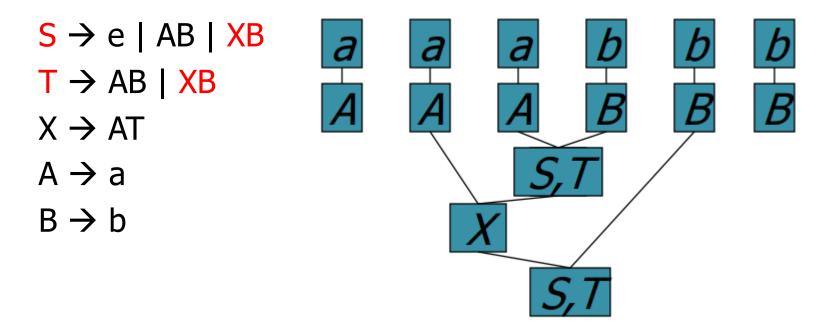
Write variables for all length 3 substrings

$$S \rightarrow e \mid AB \mid XB$$
  
 $T \rightarrow AB \mid XB$   
 $X \rightarrow AT$   
 $A \rightarrow a$   
 $B \rightarrow b$ 





Write variables for all length 4 substrings





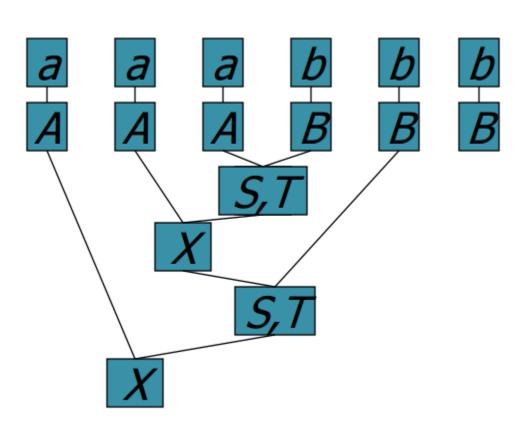
Write variables for all length 5 substrings.

$$S \rightarrow e \mid AB \mid XB$$
  
 $T \rightarrow AB \mid XB$ 

$$X \rightarrow AT$$

$$A \rightarrow a$$

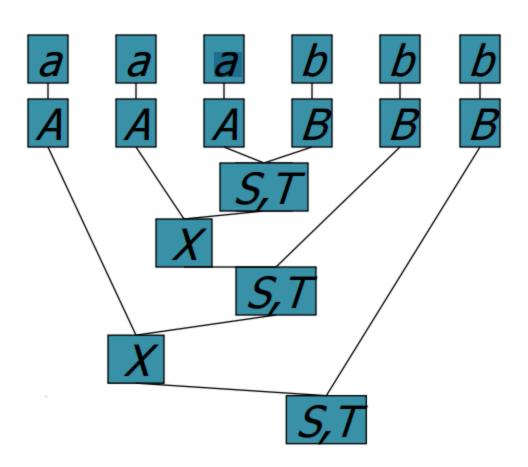
$$B \rightarrow b$$





Write variables for all length 6 substrings.

$$S \rightarrow e \mid AB \mid XB$$
 $T \rightarrow AB \mid XB$ 
 $X \rightarrow AT$ 
 $A \rightarrow a$ 
 $B \rightarrow b$ 



The table chart used by the algorithm:

| j | 1    | 2       | 3              | 4                  | 5        | 6                     |
|---|------|---------|----------------|--------------------|----------|-----------------------|
| i | no . | a       | <del>c</del> o | b                  | b        | b                     |
| 0 | A    | -       | -              | -                  | <u> </u> | — <i>S</i> , <i>T</i> |
| 1 |      | $A^{-}$ | -              | $-\chi$ $-$        | -S, T    | -                     |
| 2 |      |         | A $-$          | -S, T              | -        | -                     |
| 3 |      |         |                | $\stackrel{\ }{B}$ | -        | -                     |
| 4 |      |         |                |                    | B        | -                     |
| 5 |      |         |                |                    |          | $B^{'}$               |



#### **Dependency Grammars**

- In CFG-style phrase-structure grammars the main focus is on constituents.
- But it turns out you can get a lot done with just binary relations among the words in an utterance.
- In a dependency grammar framework, a parse is a tree where
  - ✓ the nodes stand for the words in an utterance
  - ✓ The links between the words represent dependency relations between pairs of words.
    - Relations may be typed (labeled), or not.

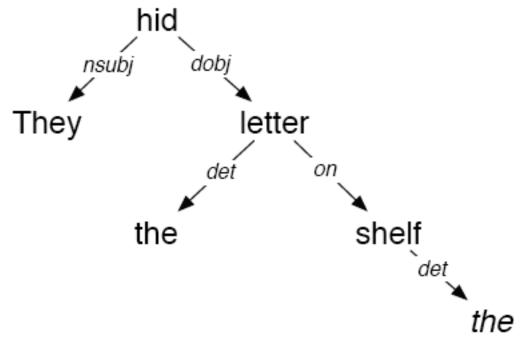


## **Dependency Relations**

| <b>Argument Dependencies</b> | Description            |
|------------------------------|------------------------|
| nsubj                        | nominal subject        |
| csubj                        | clausal subject        |
| dobj                         | direct object          |
| iobj                         | indirect object        |
| pobj                         | object of preposition  |
| Modifier Dependencies        | Description            |
| tmod                         | temporal modifier      |
| appos                        | appositional modifier  |
| det                          | determiner             |
| prep                         | prepositional modifier |



#### **Dependency Parse**



They hid the letter on the shelf



## **Dependency Parsing**

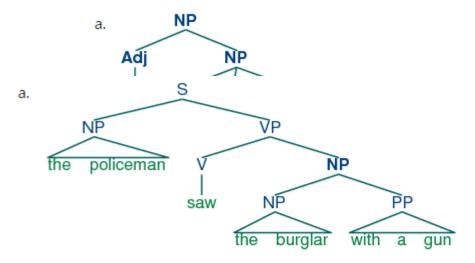
- The dependency approach has a number of advantages over full phrase-structure parsing.
  - ✓ Deals well with free word order languages where the constituent structure is quite fluid
  - ✓ Parsing is much faster than CFG-bases parsers
  - ✓ Dependency structure often captures the syntactic relations needed by later applications
    - CFG-based approaches often extract this same information from trees anyway.

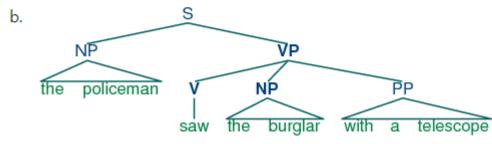


## **Dependency Parsing**

- There are two modern approaches to dependency parsing
  - ✓ Optimization-based approaches that search a space of trees for the tree that best matches some criteria
  - ✓ Shift-reduce approaches that greedily take actions based on the current word and state.

- a. old (men and women)
- b. (old men) and women





```
>>> s1 = '(S (NP the policeman) (VP (V saw) (NP (NP the burglar) (PP with a gun))))'
>>> s2 = '(S (NP the policeman) (VP (V saw) (NP the burglar) (PP with a telescope)))'
>>> tree1 = nltk.bracket_parse(s1)
>>> tree2 = nltk.bracket_parse(s2)
```

```
>>> tree = nltk.bracket_parse('(NP (Adj old) (NP (N men) (Conj and) (N women)))')
>>> tree.draw()
```

http://nltk.sourceforge.net/doc/en/ch07.html



#### **Syntaxnet**

#### You'll need to install:

- Python 2.7
- Pip
- Bazel
  - Version 0.3.0-0.3.1
- Swig
- Numpy
- mock



#### **Syntaxnet**

- You'll need to install:
  - git clone –recursive
     https://github.com/tensorflow/models.git
  - cd models/syntaxnet/tensorflow
  - ./configure
  - cd ..
  - bazel test syntaxnet/... util/utf8/...
  - # On Mac, run the following:
  - bazel test --linkopt=headerpad\_max\_install\_names \
  - syntaxnet/... util/utf8/...



#### **Syntaxnet**

echo 'Bob brought the pizza to Alice.' | syntaxnet/demo.sh

```
Input: Bob brought the pizza to Alice .

Parse:
brought VBD ROOT
+-- Bob NNP nsubj
+-- pizza NN dobj
| +-- the DT det
+-- to IN prep
| +-- Alice NNP pobj
+-- . . punct
```



## CYK algorithm 실습



#### 과제 목적

- CYK algorithm에 대해 python을 통해 코딩을 함으로 CYK algorithm에 대한 이해를 확장하며, 실제 코드의 동작절차를 이해하는데 목적이 있음



#### Grammar 설정

```
import nltk
gram = nltk.CFG.fromstring("""
S -> NP VP
NP -> Det N | NP PP
VP -> V NP | VP PP
PP -> P NP
Det -> 'the'
N -> 'kids' | 'box' | 'floor'
V -> 'opened'
P -> 'on'
```



#### 테이블 초기화 코드

```
def init_nfst(tok, gram):
   numtokens1 = len(tok)
   # fill w/ dots
   nfst = [["." for i in range(numtokens1+1)] for j in
range(numtokens1+1)]
   # fill in diagonal
  for i in range(numtokens1):
     prod= gram.productions(rhs=tok[i])
     nfst[i][i+1] = prod[0].lhs()
   return nfst
```

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#### 테이블 완성 코드

```
def complete_nfst(nfst, tok, trace=False):
  index1 = \{\}
  for prod in gram.productions():
     index1[prod.rhs()] = prod.lhs()
  numtokens1 = len(tok)
  for span in range(2, numtokens1+1):
     for start in range(numtokens1+1-span):
        end = start + span
        for mid in range(start+1, end):
           nt1, nt2 = nfst[start][mid], nfst[mid][end]
           if (nt1,nt2) in index1:
              if trace:
                 print "[%s] %3s [%s] %3s [%s] ==>
[%s] %3s [%s]" % (start, nt1, mid, nt2, end, start,
```



#### 테이블 출력 코드

```
def display(wfst, tok):
    print 'nWFST' + ' '.join([("%-4d" % i) for i in range(1,
len(wfst))])
    for i in range(len(wfst)-1):
        print " %d " % i,
        for j in range(1, len(wfst)):
            print "%-4s" % wfst[i][j],
        print
```



#### 실행 코드

```
tok = ["the", "kids", "opened", "the", "box", "on", "the",
"floor"]
res1 = init_nfst(tok, gram)
display(res1, tok)
res2 = complete_nfst(res1,tok,1)
display(res2, tok)
```



#### 실행결과(1)

```
res1 = init_nfst(tok, gram)
display(res1, tok)
```

```
      nWFST 1
      2
      3
      4
      5
      6
      7
      8

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      Det
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      <td
```



#### 실행결과(2)

res2 = complete\_nfst(res1,tok,1)

```
[0] Det [1] N [2] ==> [0] NP [2]
[3] Det [4] N [5] ==> [3] NP [5]
[6] Det [7] N [8] ==> [6] NP [8]
[2] V [3] NP [5] ==> [2] VP [5]
[5] P [6] NP [8] ==> [5] PP [8]
[0] NP [2] VP [5] ==> [0] S [5]
[3] NP [5] PP [8] ==> [3] NP [8]
[2] V [3] NP [8] ==> [2] VP [8]
[2] VP [5] PP [8] ==> [2] VP [8]
[0] NP [2] VP [8] ==> [0] S [8]
```



## 실행결과(3)

display(res2, tok)

```
      nWFST 1
      2
      3
      4
      5
      6
      7
      8

      0
      Det
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## Thank you!

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