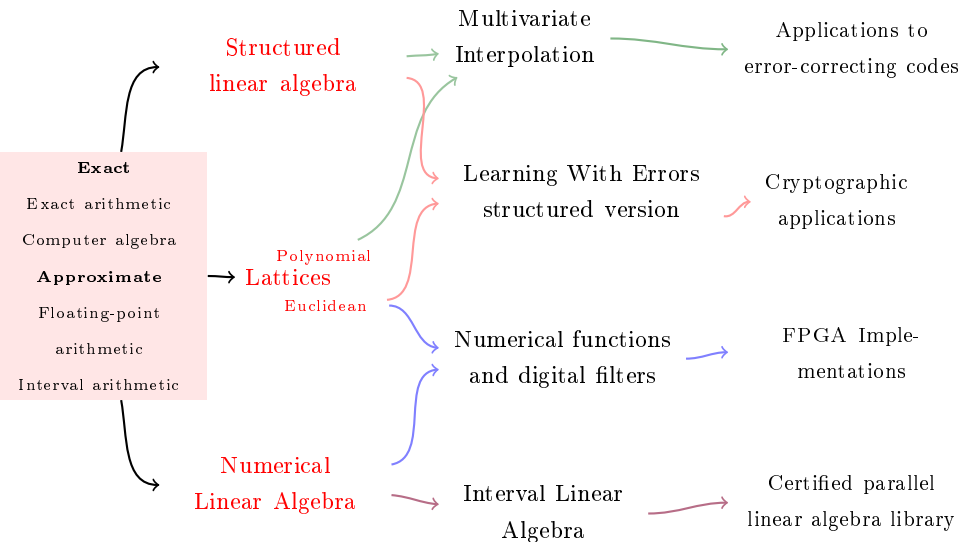


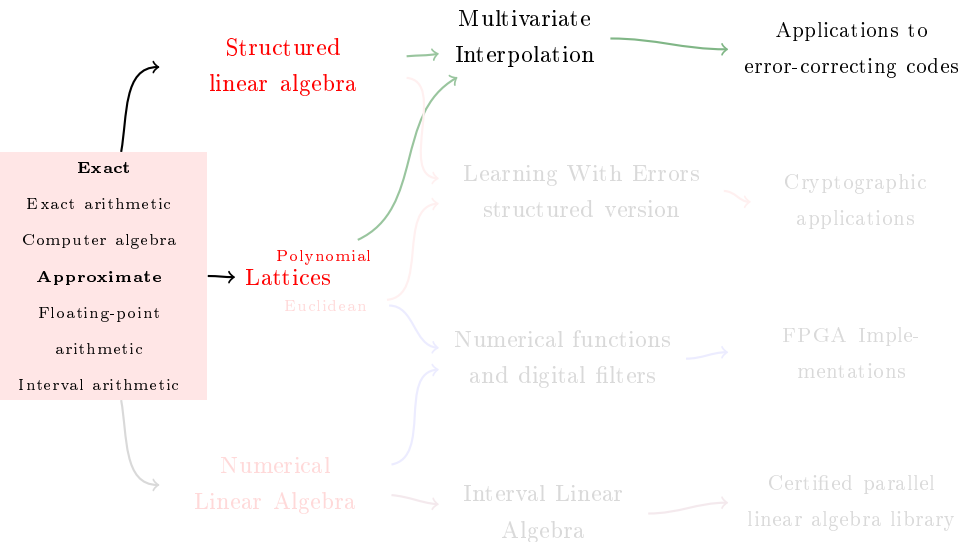
Journée des doctorants

Silviu Filip Adeline Langlois Vincent Neiger
Philippe Theveny

Aric Team, LIP, ENS de Lyon, France

June 3, 2014





AriC — Error-correcting codes

Goal:

Enable **reliable** delivery of data over **unreliable** communication channels

Strategy:

add **redundancy** to the message

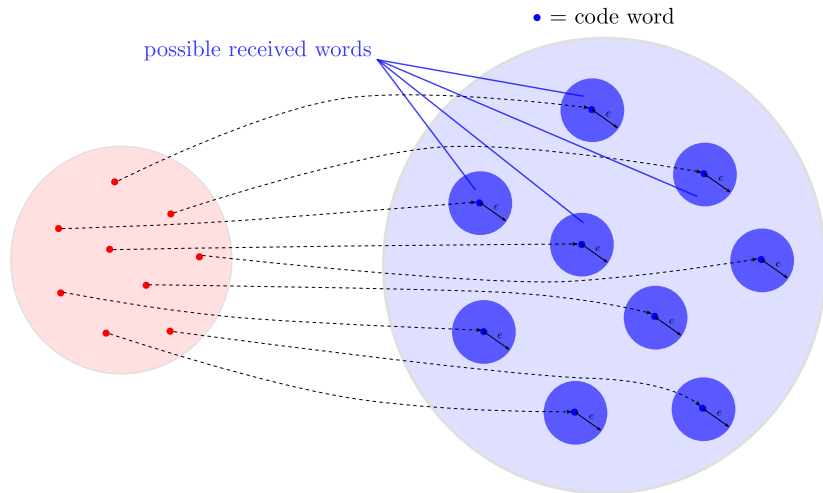
add **redundancy** to the message

add **redundancy** to the message



(courtesy of J.S.R. Nielsen)

AriC — Error-correcting codes



polynomials of degree $\leq k$ \longrightarrow their evaluation at x_1, \dots, x_n
 $w = w_0 + w_1X + \dots + w_kX^k$ $(w(x_1), \dots, w(x_n))$

AriC — (list-)Decoding

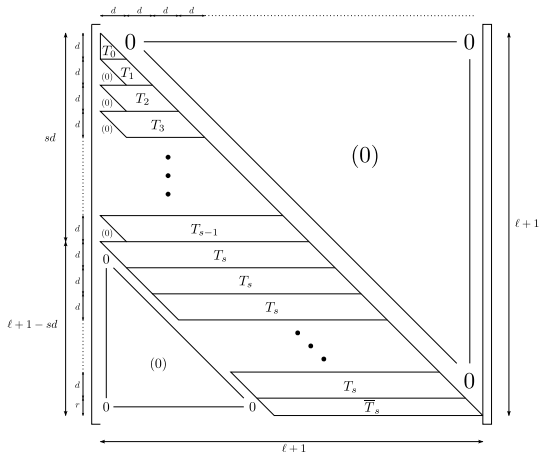
Find a **solution** of a **structured linear system**,

$$\begin{array}{c}
 \begin{array}{cccc}
 \xrightarrow{mt} & \xrightarrow{mt-k} & & \xrightarrow{N_j} \\
 & & & \xleftarrow{mt-\ell k}
 \end{array} \\
 \begin{array}{c}
 nm \\
 \downarrow \\
 n(m-1) \\
 \downarrow \\
 M_i \\
 \downarrow \\
 n
 \end{array}
 \left[\begin{array}{ccc|c|c}
 A_{0,0} & & & A_{0,j} & A_{0,\ell} \\
 \hline
 & & & & \\
 \hline
 & & & & \\
 \hline
 A_{i,0} & & & A_{i,j} & A_{i,\ell} \\
 \hline
 & & & & \\
 \hline
 A_{m-1,0} & & & A_{m-1,j} & A_{m-1,\ell}
 \end{array} \right]
 \end{array}$$

where $A_{i,j}$ is a **Toeplitz** / **Hankel** / **Vandermonde** / ... matrix

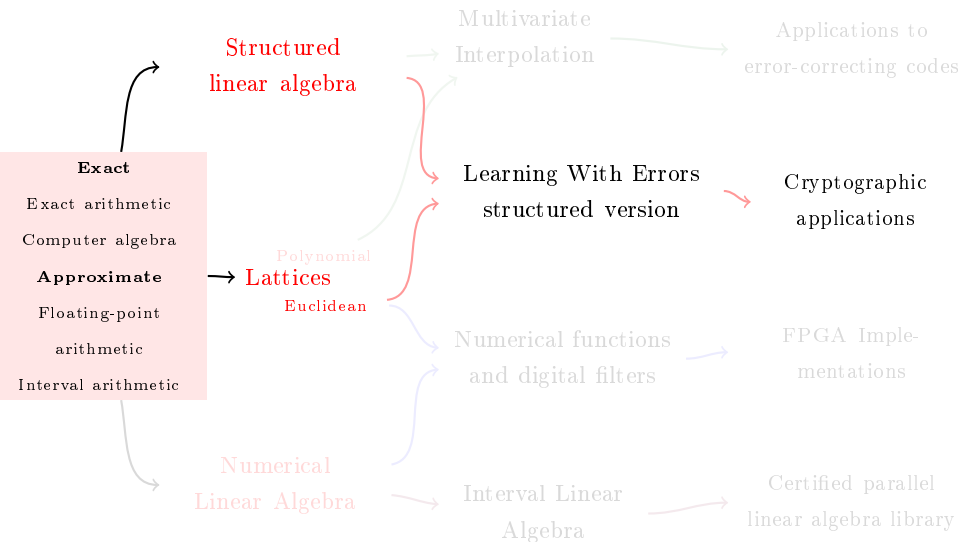
AriC — (list-)Decoding

Find a **short vector** in a (structured) **polynomial lattice**,

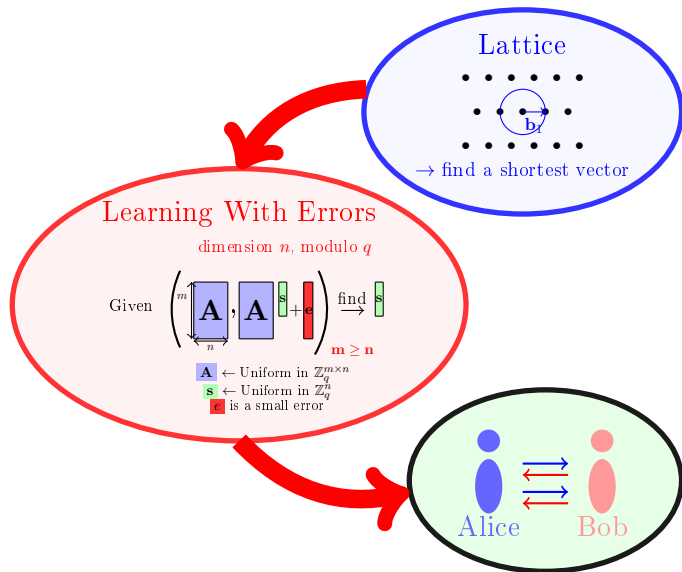


where T_i has a **Toeplitz** structure:

Diagram illustrating the geometry of a rectangular unit cell. The cell has a width of $(i+1)d$ and a height of d . The cell contains a grid of points. The leftmost point is labeled (0) . The top-left point is labeled $T_{i,0}$. The top-right point is labeled $T_{i,d}$. The bottom-left point is labeled $T_{i,0}$. The bottom-right point is labeled $T_{i,d}$. The points are connected by dashed lines forming a grid. The width is labeled $(i+1)d$ and the height is labeled d .



AriC – Lattice-Based Cryptography

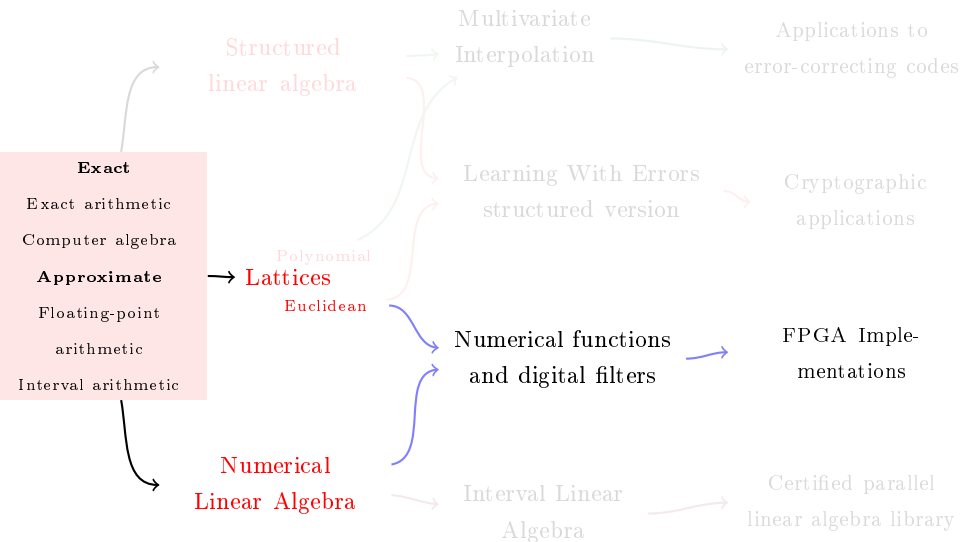


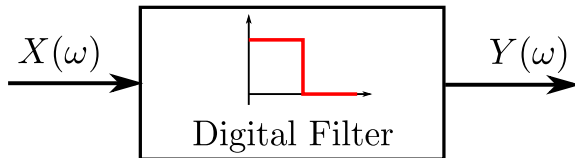
AriC – Lattice-Based Cryptography

- ▶ Public Key Encryption
- ▶ Identity Based Encryption
- ▶ Fully Homomorphic Encryption

- ▶ Signature
- ▶ Group Signature
- ▶ Hash Function

- ▶ Cryptographic Multilinear Maps



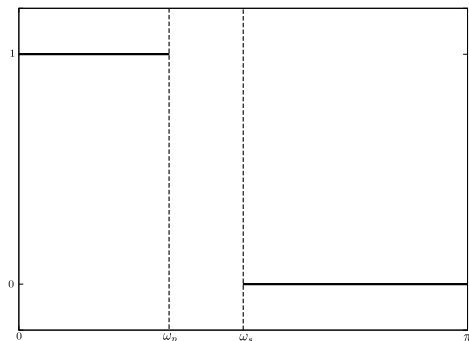


$$Y(\omega) = H_d(\omega)X(\omega), \omega \in [0, \pi]$$

Two types of filters:

- ▶ finite impulse response (**FIR**) $\Rightarrow H_d(\omega)$ **polynomial**
- ▶ infinite impulse response (**IIR**) $\Rightarrow H_d(\omega)$ **rational function**

FIR case: $H_d(\omega) = \sum_{k=0}^L a_k \cos(\omega k)$



Steps:

1. Optimal filter computation:

$$H_d(\omega) = \sum_{k=0}^L a_k \cos(\omega k)$$

Naive rounding:

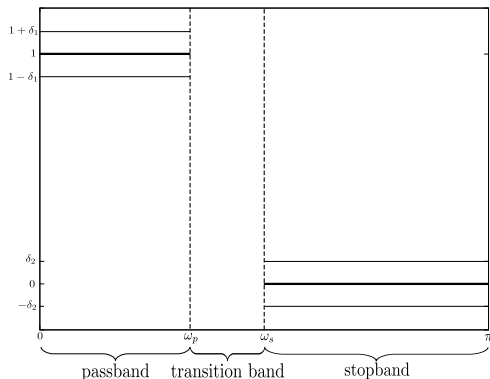
$$\bar{H}_d(\omega) = \sum_{k=0}^L \bar{a}_k \cos(\omega k)$$

2. Coefficient quantization:

$$H_d^*(\omega) = \sum_{k=0}^L a_k^* \cos(\omega k)$$

Goal: filter synthesis toolchain for embedded and FPGA targets

FIR case: $H_d(\omega) = \sum_{k=0}^L a_k \cos(\omega k)$



Steps:

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Naive rounding:

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Goal: filter synthesis toolchain for embedded and FPGA targets

AriC – Digital Filter Design

FIR case: $H_d(\omega) = \sum_{k=0}^L a_k \cos(\omega k)$

Steps:

1. Optimal filter computation:

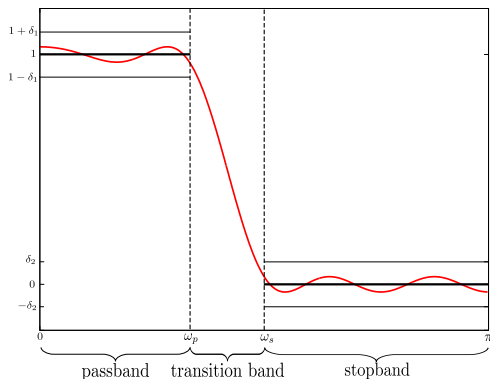
$$H_d(\omega) = \sum_{k=0}^L a_k \cos(\omega k)$$

Naive rounding:

$$\bar{H}_d(\omega) = \sum_{k=0}^L \bar{a}_k \cos(\omega k)$$

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AriC – Digital Filter Design

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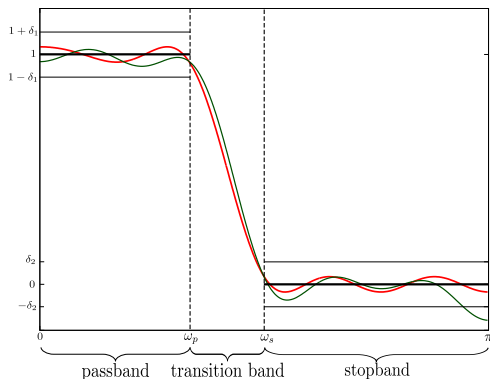
$$H_d(\omega) = \sum_{k=0}^L a_k \cos(\omega k)$$

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Goal: filter synthesis toolchain for embedded and FPGA targets

FIR case: $H_d(\omega) = \sum_{k=0}^L a_k \cos(\omega k)$

Steps:

1. Optimal filter computation:

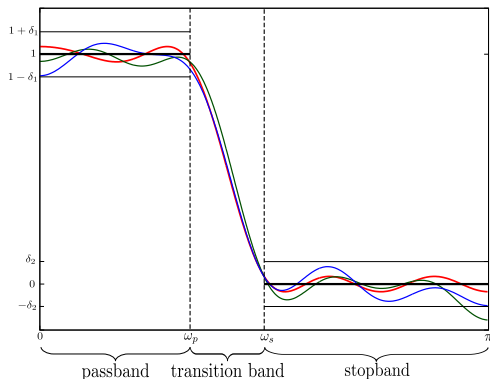
$$H_d(\omega) = \sum_{k=0}^L a_k \cos(\omega k)$$

Naive rounding:

$$\overline{H}_d(\omega) = \sum_{k=0}^L \overline{a}_k \cos(\omega k)$$

2. Coefficient quantization:

$$H_d^*(\omega) = \sum_{k=0}^L a_k^* \cos(\omega k)$$



Goal: filter synthesis toolchain for embedded and FPGA targets

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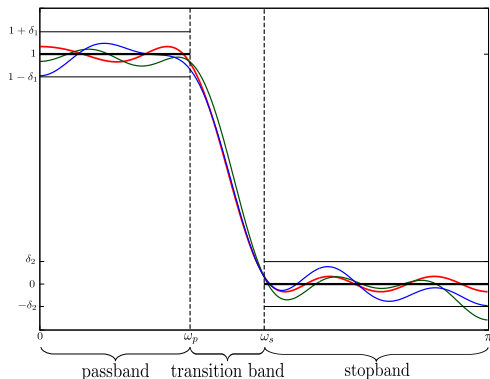
$$H_d(\omega) = \sum_{k=0}^L a_k \cos(\omega k)$$

Naive rounding:

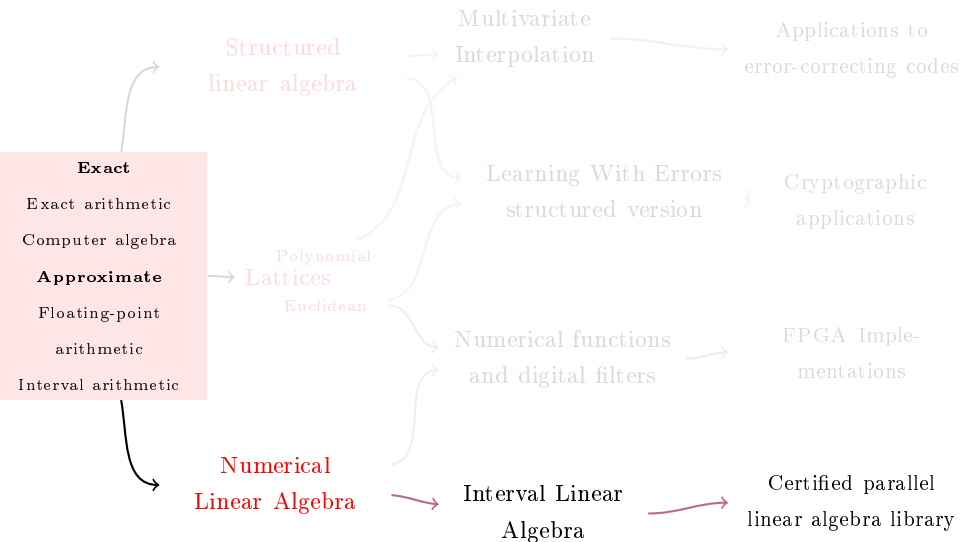
$$\overline{H}_d(\omega) = \sum_{k=0}^L \overline{a}_k \cos(\omega k)$$

2. Coefficient quantization:

$$H_d^*(\omega) = \sum_{k=0}^L a_k^* \cos(\omega k)$$



Goal: filter synthesis toolchain for embedded and FPGA targets



Numerical Linear Algebra

$$\begin{pmatrix} 1.20 & 0.40 \\ 8.00 & 2.50 \end{pmatrix} \begin{pmatrix} 44.3 & 2.10 \cdot 10^{+3} \\ 12.6 & 2.60 \cdot 10^{-3} \end{pmatrix} \\ \approx \begin{pmatrix} 58.2 & 2.52 \cdot 10^{+3} \\ 386 & 1.68 \cdot 10^{+4} \end{pmatrix}$$

Numerical Interval Matrix Multiplication

$$\begin{pmatrix} [1, 2] & [0, 4] \\ [8, 8] & [2, 3] \end{pmatrix} \begin{pmatrix} [44, 45] & [2 \cdot 10^{+3}, 3 \cdot 10^{+3}] \\ [12, 13] & [2 \cdot 10^{-3}, 3 \cdot 10^{-3}] \end{pmatrix} \\ \subseteq \begin{pmatrix} [44, 142] & [2 \cdot 10^{+3}, 6 \cdot 10^{+3}] \\ [376, 399] & [1.6 \cdot 10^{+4}, 2.4 \cdot 10^{+4}] \end{pmatrix}$$

Classical 3-loops Algorithm

Input: $\mathbf{A} = [\underline{A}, \overline{A}] \in \mathbb{IF}^{m \times k}$, $\mathbf{B} = [\underline{B}, \overline{B}] \in \mathbb{IF}^{k \times n}$

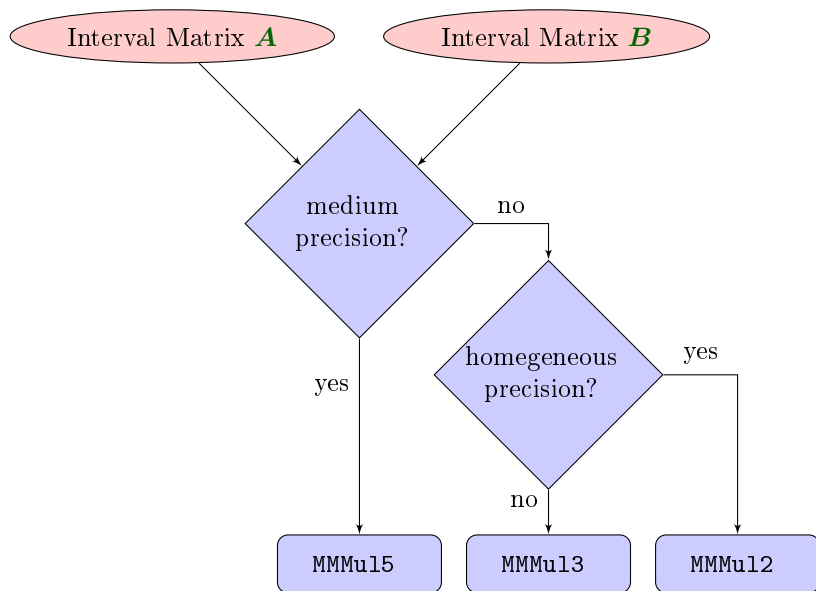
Output: $\mathbf{C} \in \mathbb{IF}^{m \times n}$, $\mathbf{C} \supseteq \mathbf{AB}$

```
1: for  $i = 1$  to  $m$  do
2:   for  $j = 1$  to  $n$  do
3:      $\underline{C}_{ij} \leftarrow 0$ ;  $\overline{C}_{ij} \leftarrow 0$ 
4:     for  $l = 1$  to  $k$  do
5:        $\underline{C}_{ij} \leftarrow$ 
         rounddown  $(\underline{C}_{ij} + \min \{ \underline{A}_{il} \underline{B}_{lj}, \underline{A}_{il} \overline{B}_{lj}, \overline{A}_{il} \underline{B}_{lj}, \overline{A}_{il} \overline{B}_{lj} \})$ 
6:        $\overline{C}_{ij} \leftarrow$ 
         roundup  $(\overline{C}_{ij} + \max \{ \underline{A}_{il} \underline{B}_{lj}, \underline{A}_{il} \overline{B}_{lj}, \overline{A}_{il} \underline{B}_{lj}, \overline{A}_{il} \overline{B}_{lj} \})$ 
7:     end for
8:   end for
9: end for
10: return  $[\underline{C}, \overline{C}]$ 
```

Variant Interval Algorithms

Algorithm	Computed Radius	Cost
MMMu13	at most $1.5\times$ exact radius	about 3 <code>gemm</code> 's
MMMu15	at most $1.18\times$ exact radius	about 5 <code>gemm</code> 's

Criterion of Choice



Parallel Implementation

Problematic Rounding Mode Support:

- ▶ compiler
- ▶ numerical library
- ▶ multi-thread management

