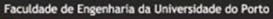
# MIEEC Computer Networks Lecture note 3

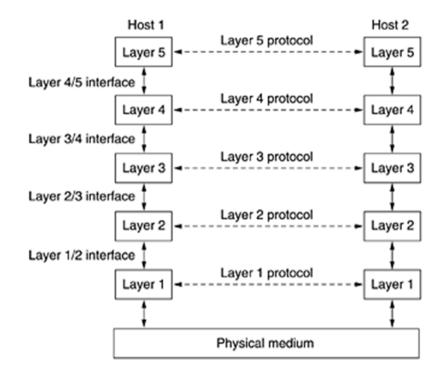
The physical layer







# The physical layer as a service



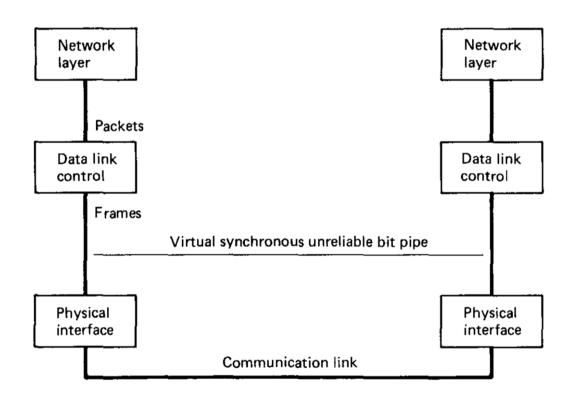
### Physical layer service

#### Service:

- Transmission and reception of digital data
- Virtual bit stream (unreliable)

#### Implementation:

 Actual communication channels used by the network



# Principles of signal transmission

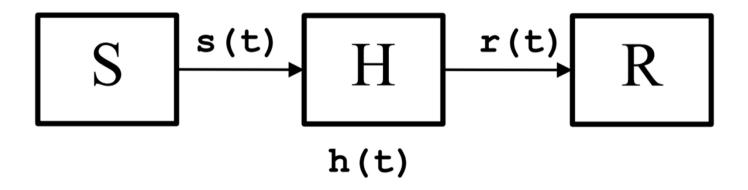
#### TO THINK

If you transmit 1 bit every T seconds you get 1/T bit/s

 Why can't you transmit infinite bit/s using a real cable?

#### Channel transmission

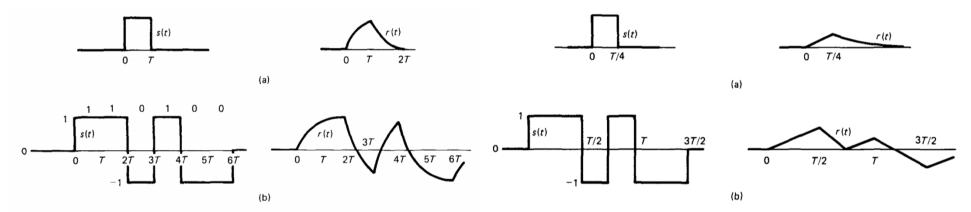
- Source signal s(t)
- Received signal r(t)
- Channel H, impulse response h(t)





# Effect of channel on r(s)

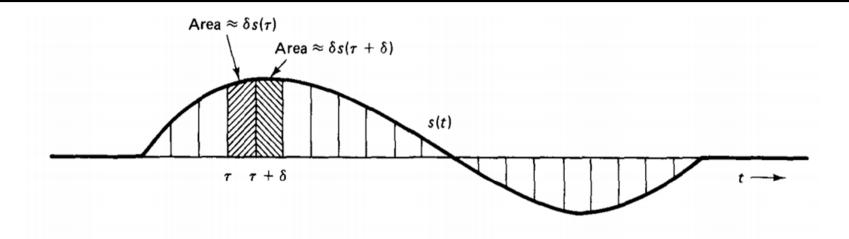
#### Filter effect on signal

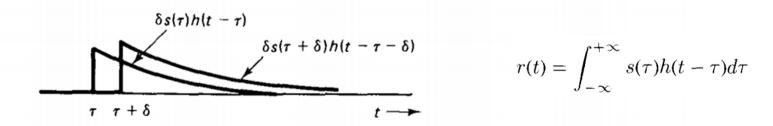


**Figure 2.3** Relation of input and output waveforms for a communication channel with filtering. Part (a) shows the response r(t) to an input s(t) consisting of a rectangular pulse, and part (b) shows the response to a sequence of pulses. Part (b) also illustrates the NRZ code in which a sequence of binary inputs (1 1 0 1 0 0) is mapped into rectangular pulses. The duration of each pulse is equal to the time between binary inputs.

**Figure 2.4** Relation of input and output waveforms for the same channel as in Fig. 2.3. Here the binary digits enter at 4 times the rate of Fig. 2.3, and the rectangular pulses last one-fourth as long. Note that the output r(t) is more distorted and more attenuated than that in Fig. 2.3.

#### r(s): convolution of s(t) and h(t)





**Figure 2.5** Graphical interpretation of the convolution equation. Input s(t) is viewed as the superposition of narrow pulses of width  $\delta$ . Each such pulse yields an output  $\delta s(\tau)h(t-\tau)$ . The overall output is the sum of these pulse responses.

# R(f): product of S(f) and H(f)

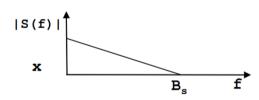
#### Fourier transforms

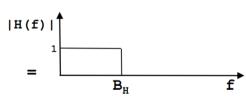
$$S(f) = \int_{-\infty}^{\infty} s(t)e^{-j2\pi ft}dt \quad H(f) = \int_{-\infty}^{\infty} h(\tau)e^{-j2\pi f\tau}d\tau$$

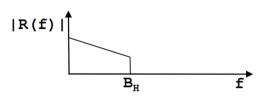
#### Received signal

$$r(t) = \int_{-\infty}^{\infty} H(f)S(f)e^{j2\pi ft}df$$

$$R(f) = H(f)S(f)$$



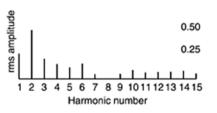




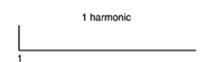
# 

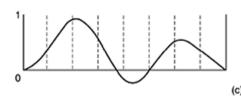
(a)

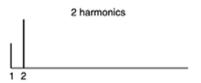
(b)



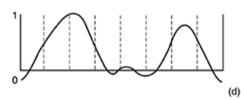
# Bandwidth-limited \*\*

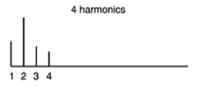


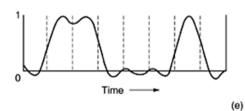


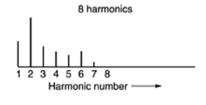


- a): signal and frequency components
- b)-e): bandwidth-limited versions of a)









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signals



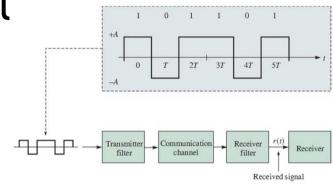
#### TO THINK

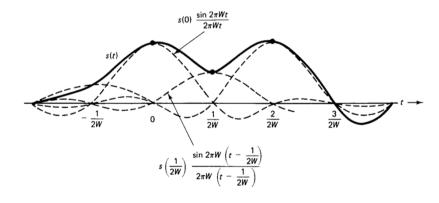
How can you retrieve digital information from a band-limited signal?

How many bit/s can a bandwidth-limited channel deliver?

# Channel capacity: Nyquist signaling rate

- Bandwidth-limited signal r(t)
  - Bandwidth B
- Zero inter-symbol interference at k\*1/2B
  - Sync pulse
- Nyquist rate
  - C=2B samples/s





# Channel capacity: Nyquist signaling rate

- Example
  - Square wave -5V (0) +5V (1)
  - Goes through a  $B_H$ =3kHz channel
  - C=2\*3kHz=6kbit/s

- Distinguish the two concepts:
  - Nyquist **signaling** rate / channel capacity
  - Nyquist sampling rate / signal reconstruction

# Channel capacity: Hartley

- M levels, log<sub>2</sub> (M) bits/sample
  - $C=2B \log_2 (M)$
  - This is the bit rate
- Baud rate
  - Symbols per second
  - 2B Baud/s
  - each symbol has M levels
  - resulting in log<sub>2</sub> (M) bits per symbol



#### TO THINK

Why can't you transmit at infinite bitrate by increasing M to infinity?

- $C=2B \log_2 (M)$
- B constant, finite
- Infinity <=> extremely large

### Channel capacity: Shannon

- Noise
  - Higher noise => more difficult to distinguish between levels
- For a given technology
  - Higher noise => lower M
  - Lower noise => higher M



# Channel capacity: Shannon

- Theoretical limit
  - Of the channel capacity
  - With noise
- Shannon-Hartley theorem

$$-C = W \log_2 \left(1 + \frac{P_r}{N_0 W}\right)$$

- C: capacity
- W: band-pass channel bandwidth, sampling rate
- P<sub>r</sub>: signal power at the receiver
- N<sub>0</sub>: noise per bandwidth unit, white noise
  - Example 10<sup>-9</sup> W/Hz



# Channel capacity: Shannon

- Signal to noise ratio:  $SNR = \frac{P_r}{N_0 W}$
- Number of levels M (Hartley)

$$-M = \sqrt{1 + \frac{P_r}{N_0 W}}$$

- Example
  - Band-pass channel W = 100 kHz

$$-\frac{P_r}{N_0 W} = 7 \Rightarrow C = 100k \log_2(1+7) = 300kbit/s$$

$$-\frac{P_r}{N_0 W} = 255 \Rightarrow C = 100k \log_2(1 + 255) = 800kbit/s$$

# Reminder: power, decibel

- $P_{dBW} = 10 \log_{10} P$
- $\bullet \ P_{dBm} = 10 \log_{10}(\frac{P}{1mW})$

• P = 100mW

$$\Rightarrow P_{dBW} = 10 \log_{10}(100 * 10^{-3}) = -10 dBW$$

$$\Rightarrow P_{dBm} = 10 \log_{10}(100) = 20 dBm$$



# Transmission techniques for band pass channels

- Most physical channels are band pass
  - |H(0)| = 0, reject DC component

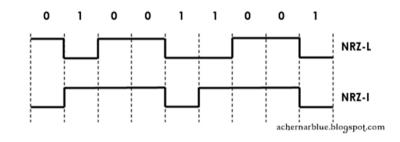


- Two major techniques are used:
  - Coding
  - Modulation

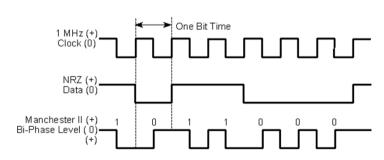


### Coding, common codes

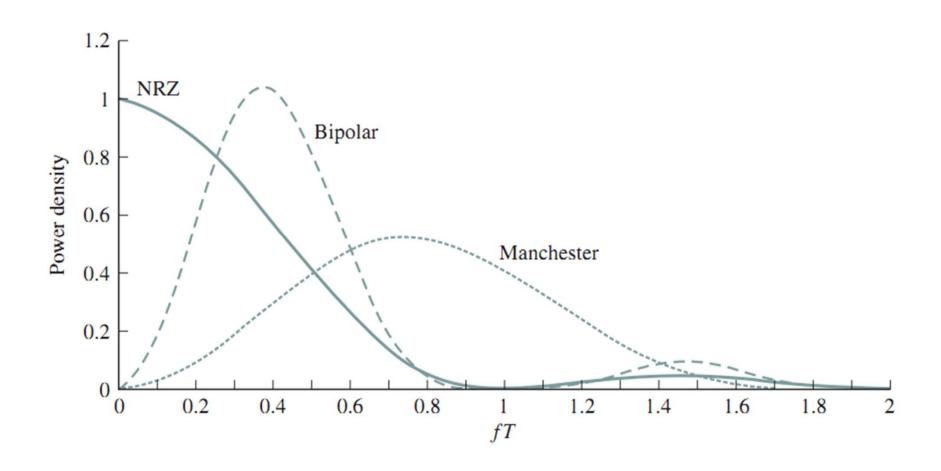
- NRZ-L, non-return to zero with levels
  - Two voltage levels, one for "1" and one for "0"
- NRZ-I, inverted
  - A change of level represents a "1"

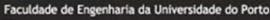


- Manchester
  - Transition in the middle of the bit
  - "1": + to ; "0": to +
  - Ethernet (IEEE 802.3)
- And many more codes



# **Spectral Power Density**





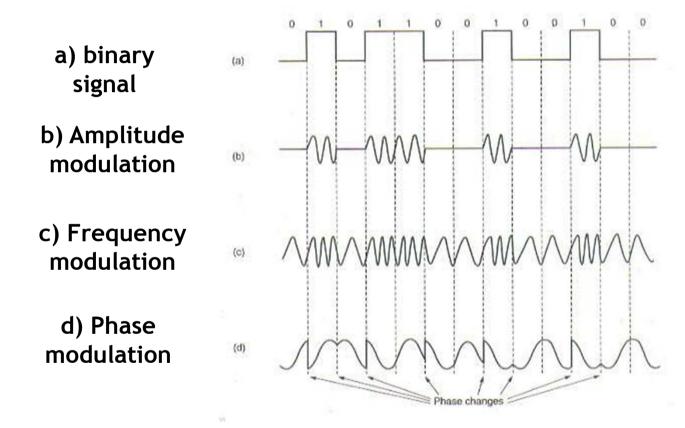


#### TO THINK

How can you transmit bits using a continuous carrier?

- E.g. sinusoidal

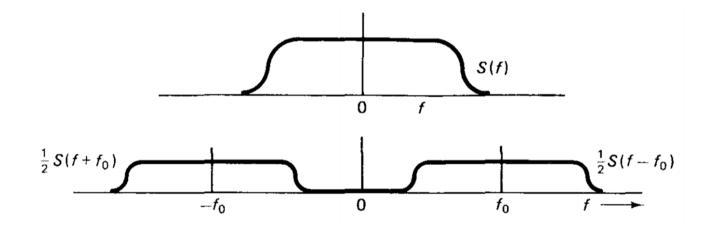
### Types of modulation





# Spectrum of an AM signal

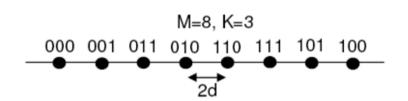
• Top: original signal

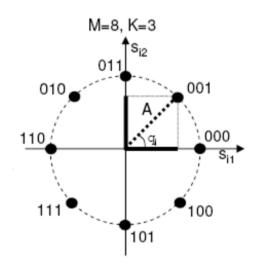


Bottom: amplitude-modulated signal

#### Amplitude and phase modulations

- Pulse-Amplitude Modulation (M-PAM)
  - $s(t) = A_i \cos(2\pi f_c t)$
  - Information coded in the amplitude of the carrier
- Phase-shift keying (M-PSK)
  - $s(t) = A\cos(2\pi f_c t + \theta_i)$
  - Information coded in the phase of the carrier

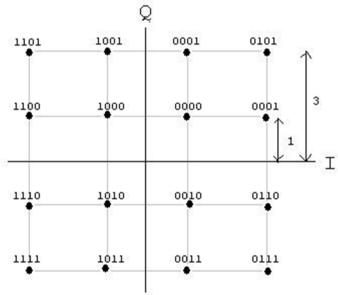




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# Quadrature amplitude modulation

- M-QAM
- Information coded in both amplitude and phase
  - $-s(t) = A_i \cos(2\pi f_c t + \theta_i)$
  - M=?
  - K=?



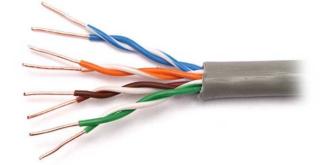
#### **Guided Transmission**

#### TO THINK

How to transmit a sequence of bits from sender to receiver using two wires?

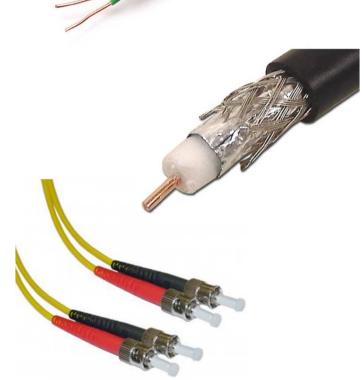
# Types of cables

 Unshielded twisted pair



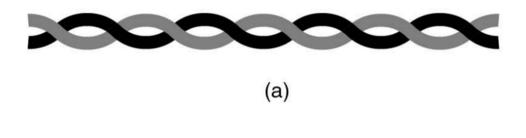
Coaxial cable

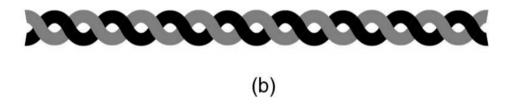
Fiber optics cable



#### Twisting and crosstalk

• a) Cat. 3 UTP





• b) Cat. 5 UTP



#### **UTP Bandwidth / Ethernet**

Cat3	UTP <sup>[6]</sup>	16MHz <sup>[6]</sup>	10BASE-T and 100BASE-T4 Ethernet <sup>[6]</sup>
Cat4	UTP <sup>[6]</sup>	20MHz <sup>[6]</sup>	16 Mbit/s <sup>[6]</sup> Token Ring
Cat5	UTP <sup>[6]</sup>	100MHz <sup>[6]</sup>	100BASE-TX & 1000BASE-T Ethernet <sup>[6]</sup>
Cat5e	UTP <sup>[6]</sup>	100MHz <sup>[6]</sup>	100BASE-TX & 1000BASE-T Ethernet <sup>[6]</sup>
Cat6	UTP <sup>[6]</sup>	250MHz <sup>[6]</sup>	1000BASE-T Ethernet
Cat6e		250MHz (500MHz according to some)	Not a standard; a cable maker's own label.
Cat6a		500MHz	10GBASE-T Ethernet

Typical attenuation: 2-25dB/100m

### Attenuation and gain

- $P_r = P_t * Gain$ 
  - In Watts
- $10\log_{10}(P_r) = 10\log_{10}(P_t) + 10\log_{10}(Gain)$ 
  - In dB
  - $-P_{rdBW} = P_{tdBW} + Gain_{dB}$
  - $-P_{rdBm} = P_{tdBm} + Gain_{dB}$

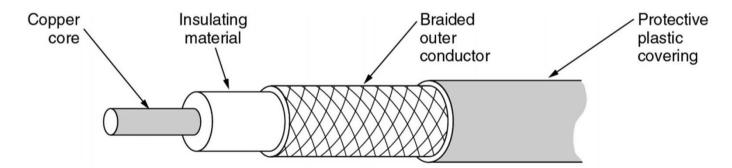
#### Example

- Gain = 0.01,  $P_{tdBm} = 30 dBm = 1W$
- $Gain_{dB} = 10 \log_{10}(0.01) = -20 dB$
- $-P_{rdBm} = P_{tdBm} + Gain_{dB} = 30 \ dBm 20 dB = 10 dBm = 10 mW$
- $Gain_{dB} = -20dB \Rightarrow Attenuation_{dB} = 20dB$



#### Coaxial cable

- Same principle as UTP but with:
  - Higher bandwidth
  - Better immunity to noise
  - Lower attenuation



- Higher cost



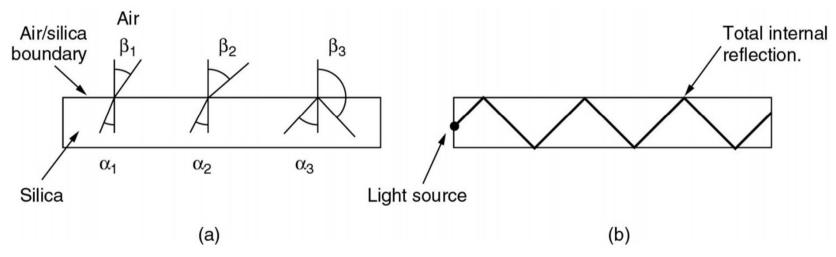
#### TO THINK

Fiber optics: how different is it from copper?

### Fiber optics

Jacket Strength Members Coating Cladding

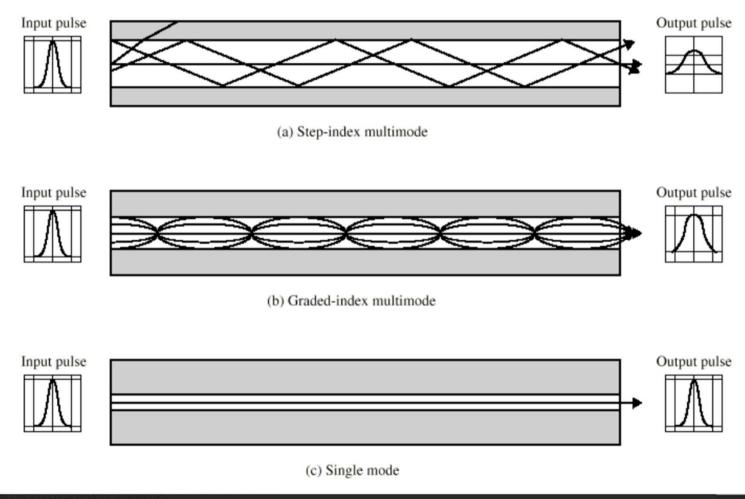
- Refraction and Reflection
  - at the core/cladding boundary

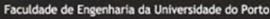


- Reflection if: incidence angle > critical angle
  - Total internal reflection
  - Propagation through reflection



#### Fiber modes and distortion

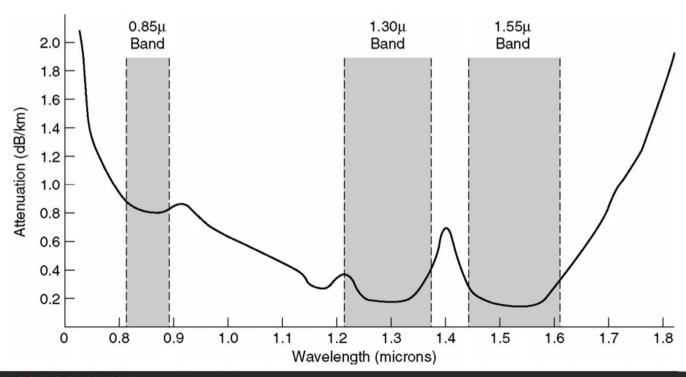






# Wavelength, attenuation, bandwidth, coding

- SNR => bitrates
- Explore lower attenuation bands
- 30 GHz bandwidth, < 1dB/km
- NRZ, pulse on("1"), pulse off ("0")
  - Amplitude modulation?





### Wavelength and propagation delay

- $\lambda = \nu T$ , or  $\lambda f = \nu$ 
  - $-\lambda$  wavelength
  - $\nu$  speed of the wave (light, etc)
  - f frequency
  - T period
- Speed of light in free space

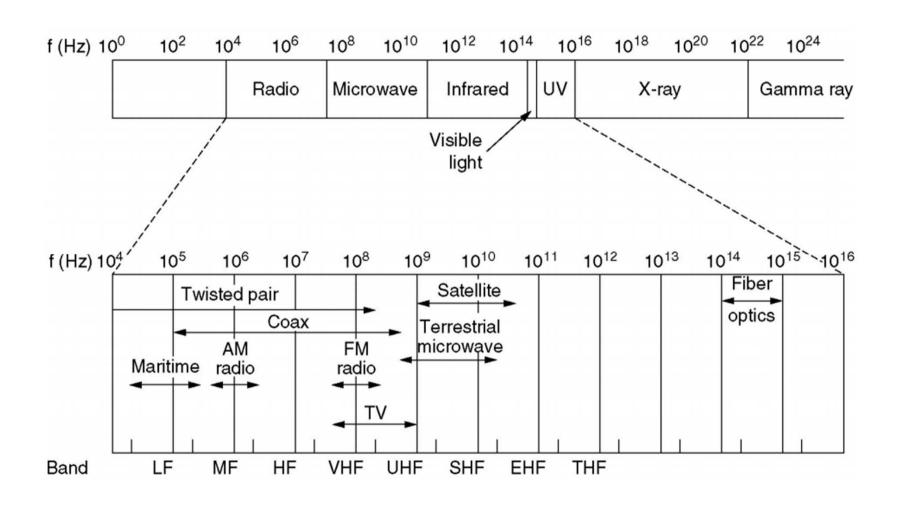
$$-c = 3 * 10^8 m/s$$

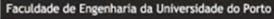
- Propagation delay  $(\mu s/km)$ 
  - Free space 3.3  $\mu s/km$
  - Coaxial cable 4 μs/km
  - UTP 5  $\mu s/km$
  - Optical fiber 5  $\mu s/km$ 
    - (notice the increasing propagation delay)



#### Wireless Transmission

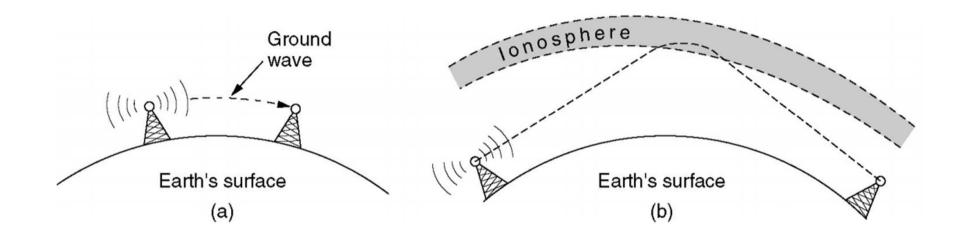
## Electromagnetic (EM) spectrum







#### Radio transmission



- (a) VLF, LF, MF follow the Earth's surface
- (b) HF, VHF are absorbed by the Earth or bounce off the ionosphere

#### TO THINK

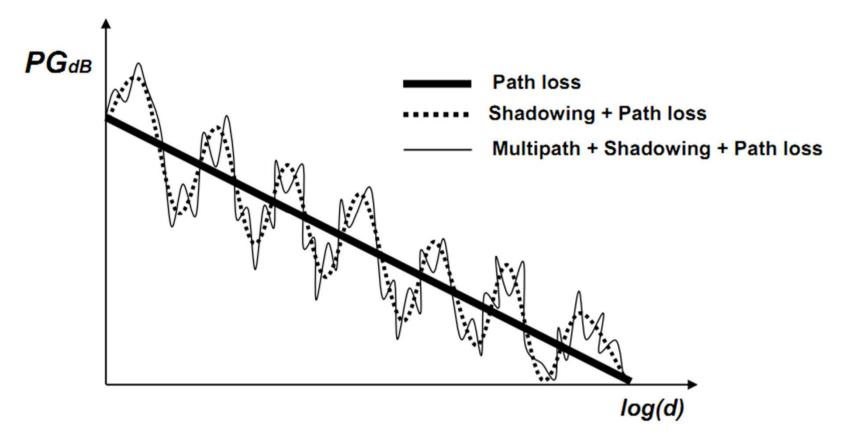
What are the different issues that affect attenuation in wireless transmission?

## Free space model

- Channel gain in free space  $\frac{P_r}{P_s} = \left(\frac{\lambda \sqrt{G_l}}{4\pi d}\right)^2$ 
  - Gl: gain of antennas
  - d: sender-receiver distance
- $PG_{dB} = 10 \log_{10} \left(\frac{P_{\gamma}}{P_{s}}\right) =$ =  $20 \log_{10} \lambda \sqrt{G_{l}} / 4\pi d - 20 \log_{10} d = b - 20x$
- $PL_{db} = -10 \log_{10} \lambda^2 G_l / (4\pi d)^2$



## Propagation issues



Shadowing: large obstacle (building, hill); slow fading Multipath: multiple paths overlapping; fast fading



## Path loss example

#### What's the sender power Ps

- For receiver power  $Pr = 0.1 \mu W$ ?
- And 900 or 3000MHz at 10, 100, 1000 m distance?

$f_c$	λ	d	$PL_{dB}$	$P_{s_{dBm}}$	$P_s$
	$\left(\frac{c}{f_c}\right)$		$\left(-10  \log \frac{\lambda^2 G_l}{(4\pi d)^2}\right)$	$(P_{rdBm} + PL_{dB})$	$\left(10^{\frac{P_{sdBm}}{10}-3}\right)$
(MHz)	(cm)	(m)	(dB)	(dBm)	(W)
900	33	10 100 1000	51.5 71.5 91.5	11.5 31.5 51.5	0.014 $1.42$ $142$
3000	10	10 100 1000	62 82 102	22 42 62	0.158 15.8 1579

## Capacity of wireless channel

$$P_r(d) = (d_0/d)^3 P_t$$
, for  $d_0 = 10m$ .

d	$\gamma = P_r(d)/(N_0 B)$	$SNR = \gamma_{dB} = 10 \log \gamma$	$C = B \log_2(1+\gamma)$	Efficiency
(m)		(dB)	(kbit/s)	(bit/s/Hz)
50	267	24	242	8
100	33	15	153	5.1
500	0.27	-6	10	0.3
1000	0.033	-15	1.4	0.05

Table 2.6: Shannon capacities for wireless channels. The limiting capacities of wireless channels depend on the channel bandwidth and on the power received. The capacity C of the channel and its efficiency are given for a transmitted power  $P_s = 1 W$ ,  $d_0 = 10 m$ , a narrow bandwidth of 30 kHz and a noise power spectral density  $N_0 = 10^{-9} W/Hz$ . The capacity decreases significantly as the distance between the sender and the receiver increases

#### HOMEWORK

Review slides

- Read:
  - Tanenbaum 2.1,2.2, 2.3, 2.5, 2.7, 2.8
  - Bertsekas 2.1, 2.2
- Do your Moodle homework