



Generic Framing Procedure (GFP) for NG-SONET/SDH: An Overview

Enrique Hernandez-Valencia
Lucent Technologies

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Outline

- What is GFP?
- Problem Statement
- GFP Value Proposition
- GFP Model
 - Frame Structure
 - Procedures
- GFP Performance
- Applications:
 - Hybrid SONET/DATA NEs
- Summary

Generic Framing Procedure - GFP

A “generic” mechanism to adapt multiple client traffic types as either:

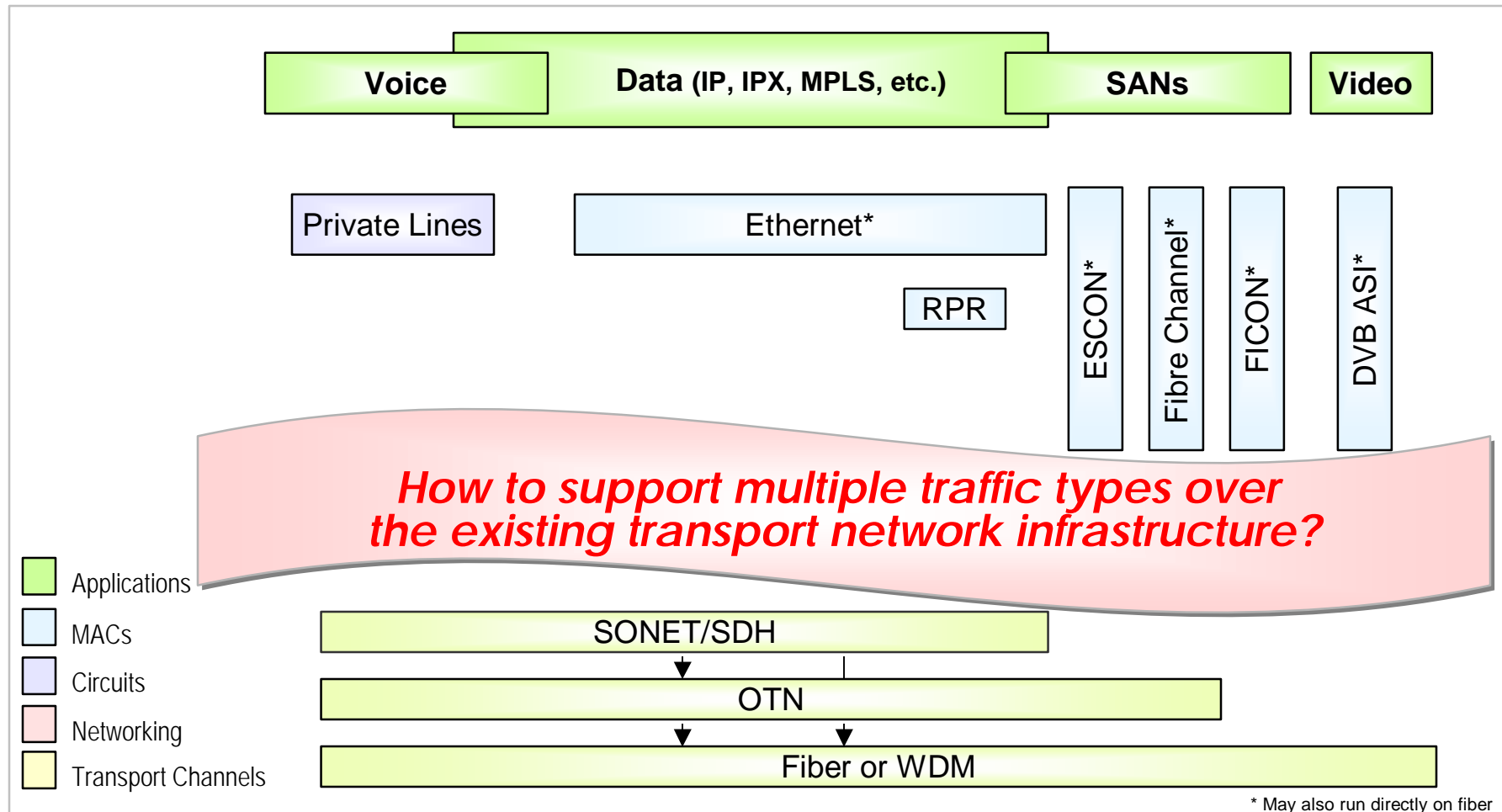
- *a physical link (Layer 1) client*
- *a logical data link (Layer 2) client*

into a bit synchronous or octet-synchronous transmission channel

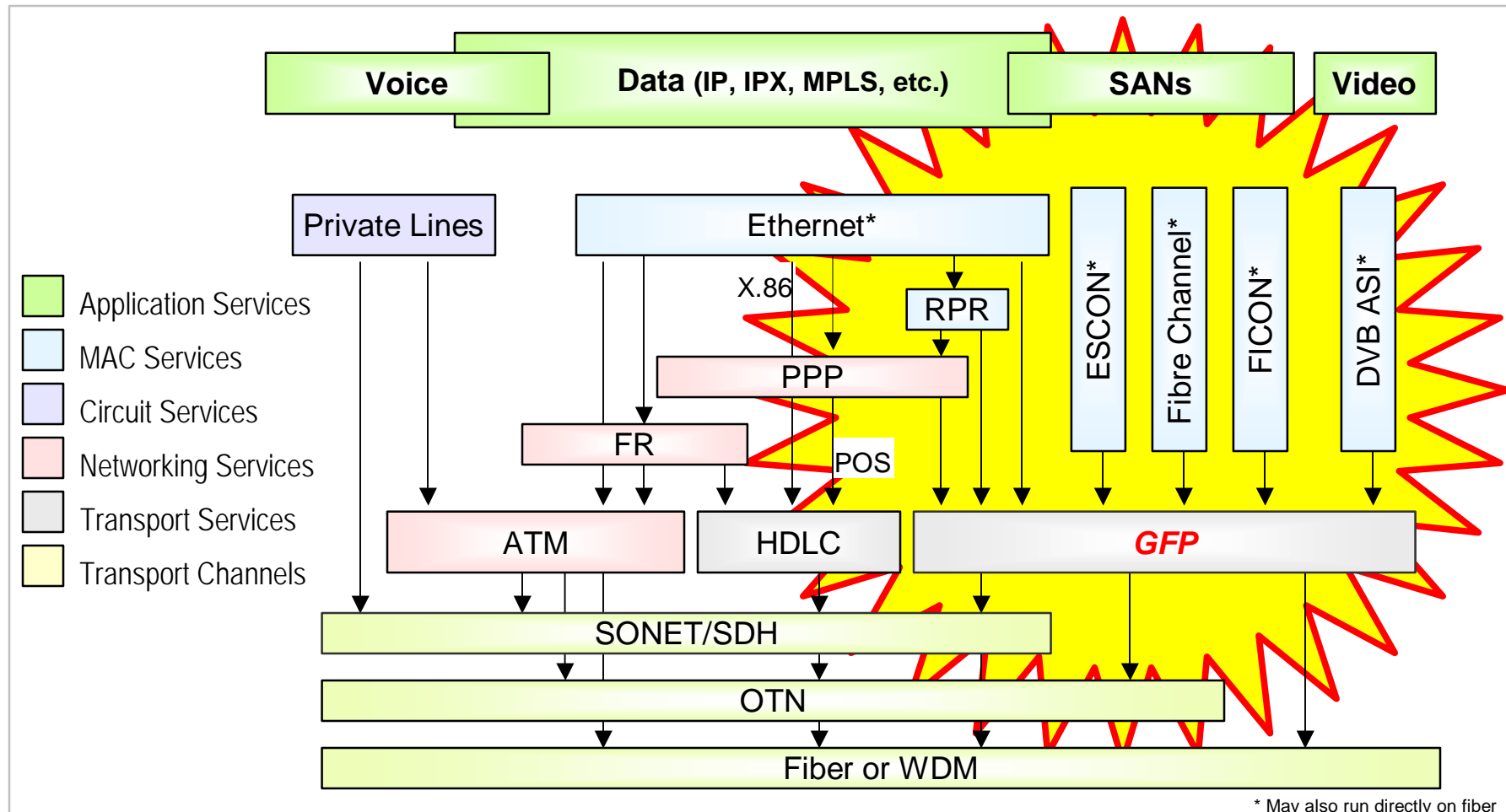
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The Problem: Public Multi-Service Transport

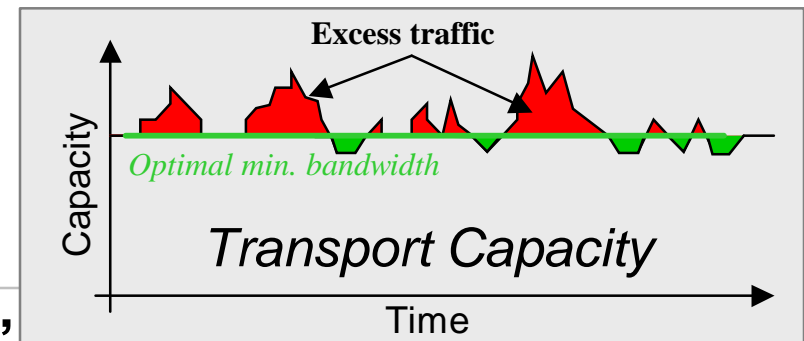
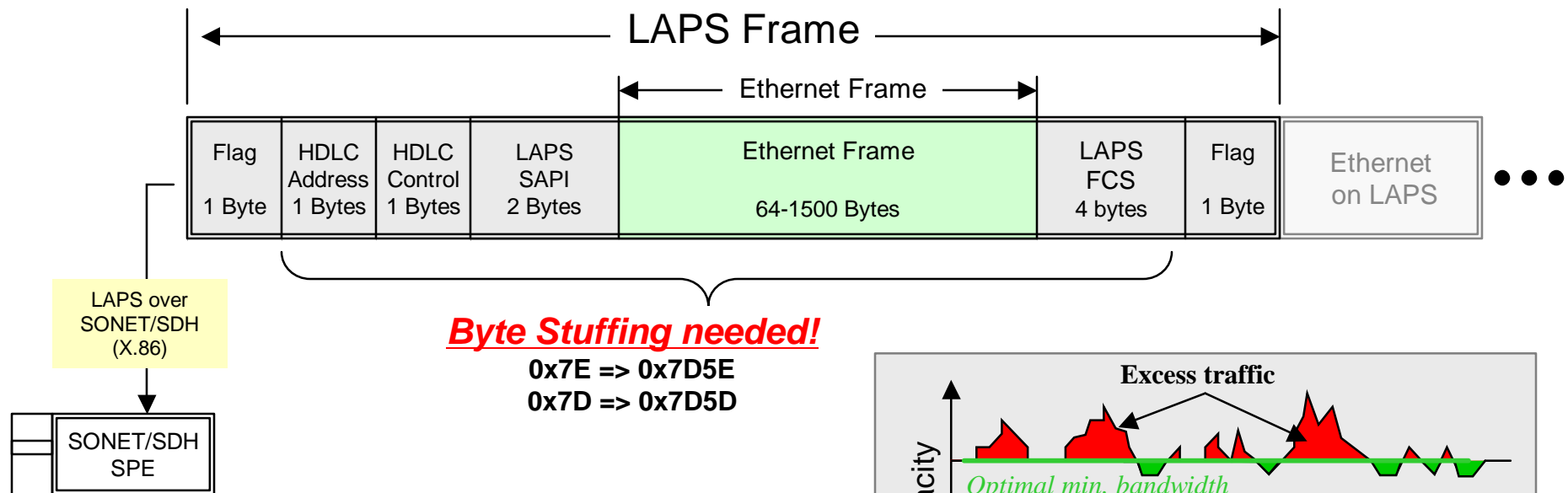


The Solutions: A Fragmented Solution Space





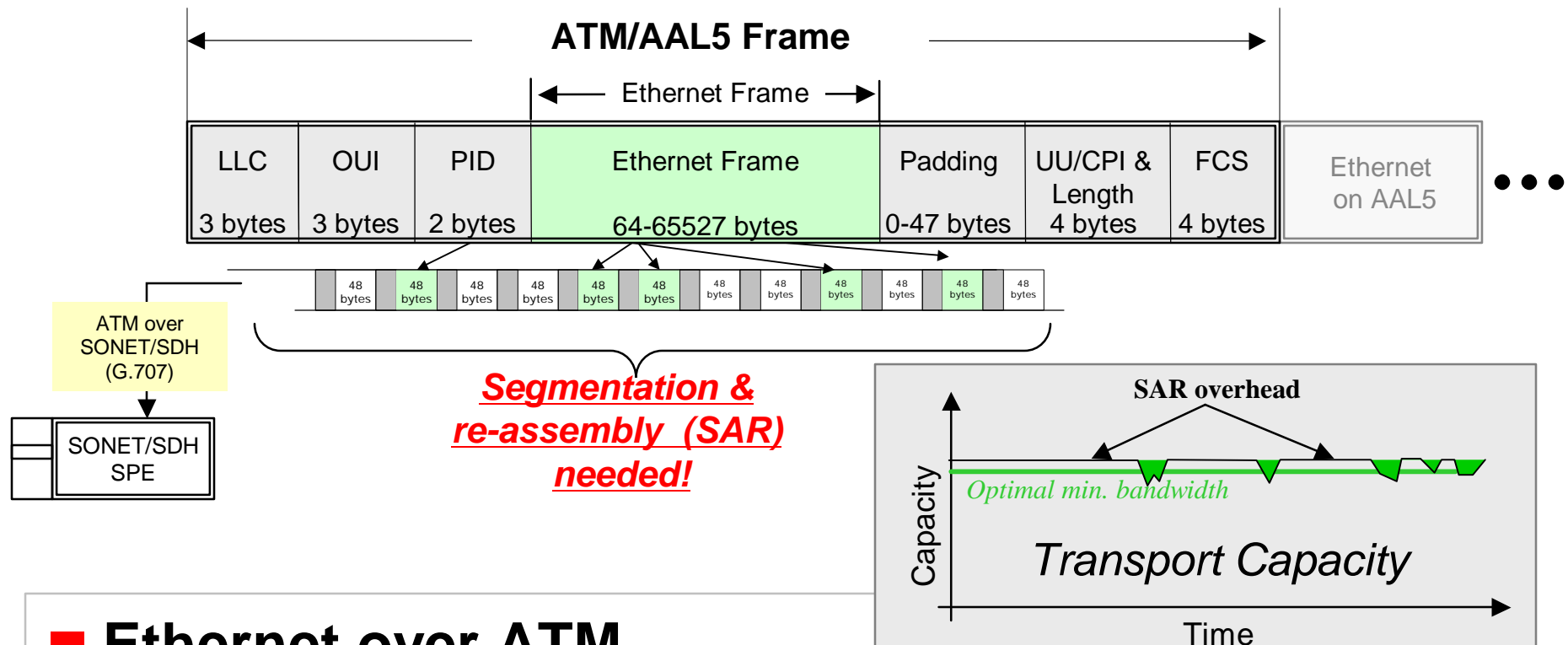
Example 1: Ethernet over LAPS (ITU-T X.86)



■ Ethernet in “HDLC-like Framing”

- Non-deterministic transport overhead
- Byte stuffing interferes with QoS/bandwidth management
- Flag-based delineation computationally expensive as speed increases

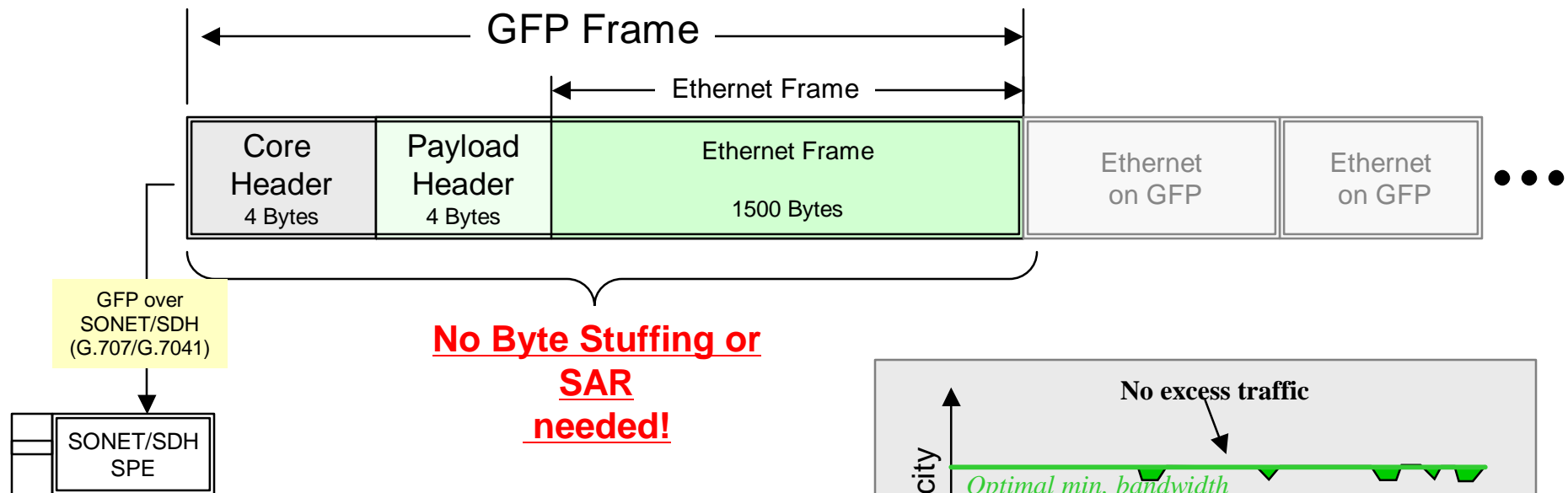
Example 2: Ethernet over ATM (IETF RFC 1483)



■ Ethernet over ATM

- Excellent QoS management capabilities
- Large transport overhead for small packets
- SAR expensive for simple connectivity services

Example 3: Ethernet over GFP-F (ITU-T G.7041)



■ Ethernet over GFP

- Deterministic transport overhead
- No adaptation interference with QoS/bandwidth management
- Low complexity frame delineation that scales up as speed increases

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Why GFP?

■ Simple and scaleable

- Proven technology at 1G, 2.5G and 10G
- Scalable beyond 40G

■ Supports both Layer 1 and Layer 2 traffic

- Alternative transport mechanism to ATM (ITU-T I.341.1/IETF RFC 1483)
- Alternative transport mechanism to HDLC-framing (ISO-3309/IETF RFC 2615)

■ Standards based:

- ITU-T G.7041(2001) & ANSI T1.105.02 (2002)
- Endorsed by IETF (RFC 2823)
- Endorsed by RPR WG (IEEE 802.17)

Sample Applications

Channel Types:

■ Bit-Synchronous Channel:

- Dark Fiber
- WDM

■ Octet-Synchronous Channel:

- SONET (T1.105.02)
- SDH (ITU-T G.707)
- OTN (ITU-T G.709)

Client Types:

■ Physical Coding (Layer 1):

- Fibre Channel
- FICON
- ESCON
- Gigabit Ethernet
- Infiniband
- DVB ASI

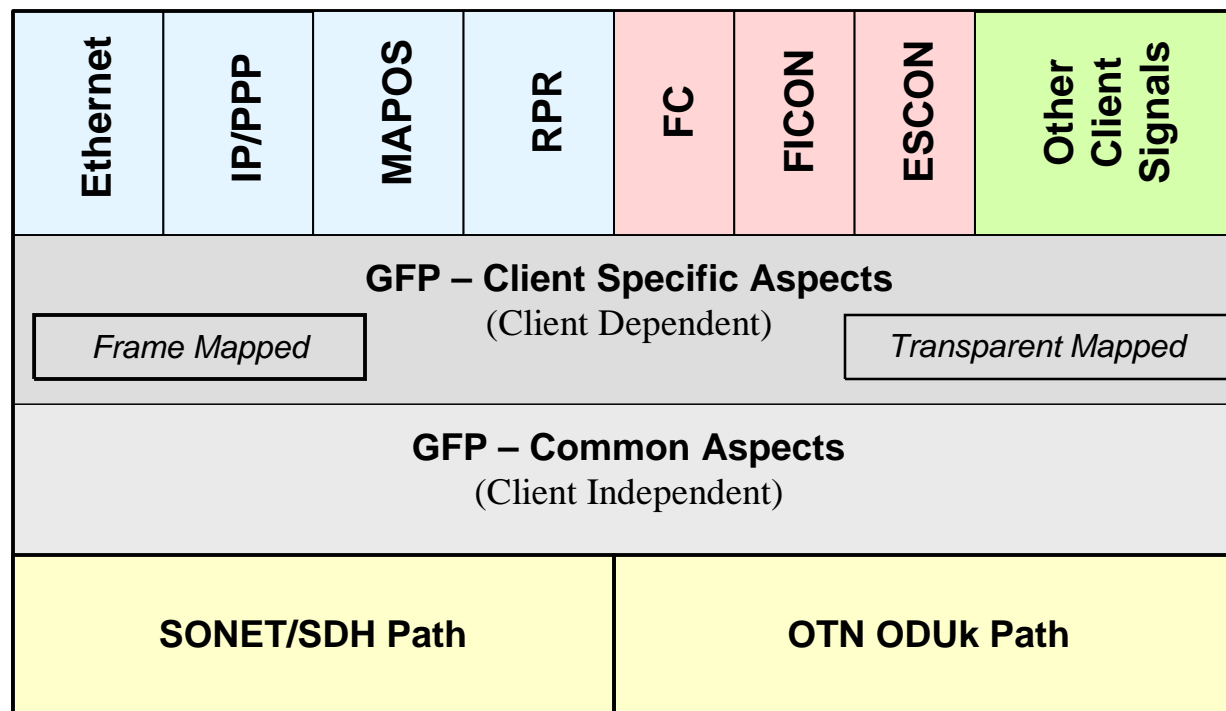
■ Data Links (Layer 2):

- PPP/IP/MPLS
- Ethernet
- MAPOS
- RPR

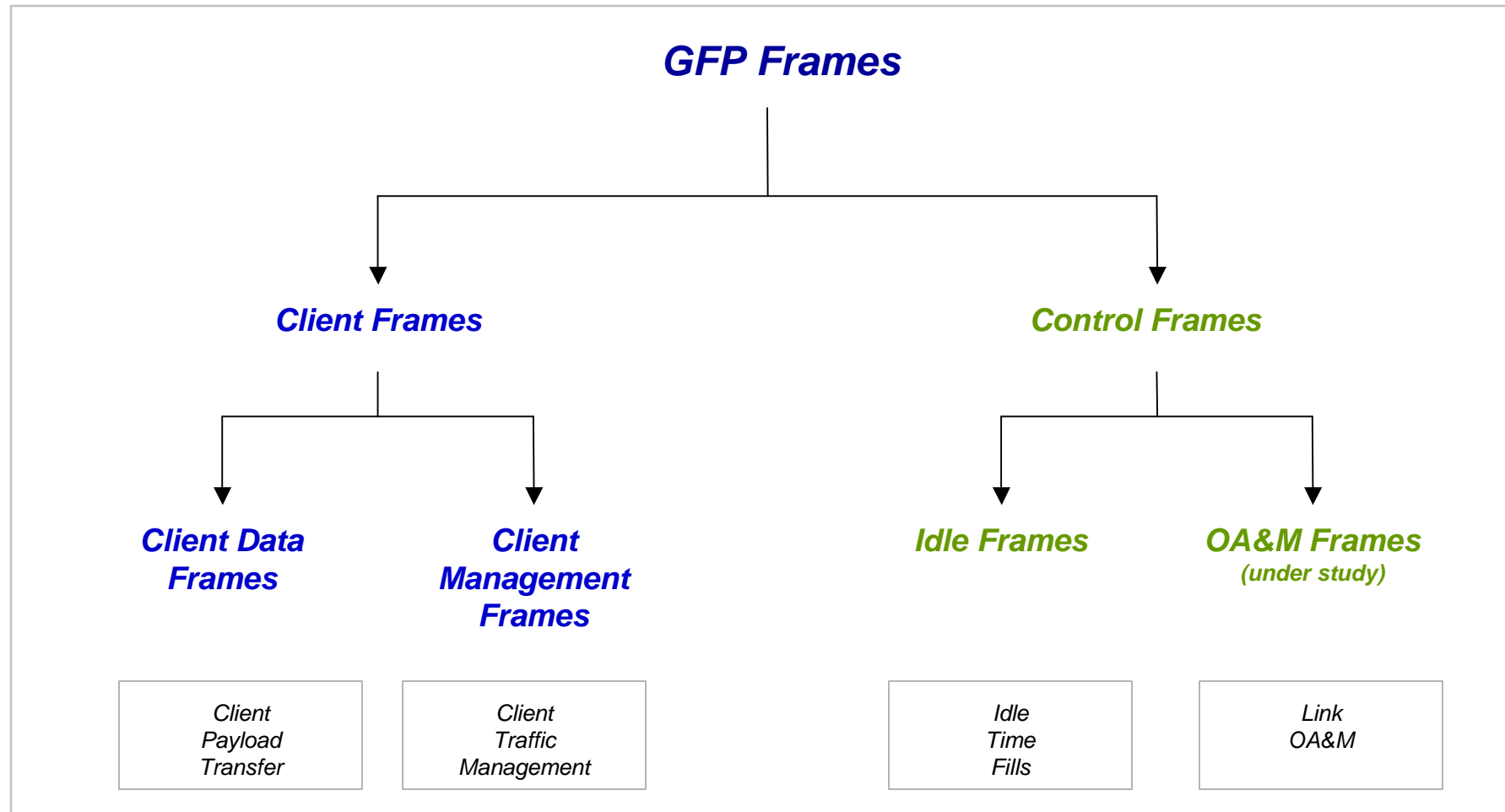
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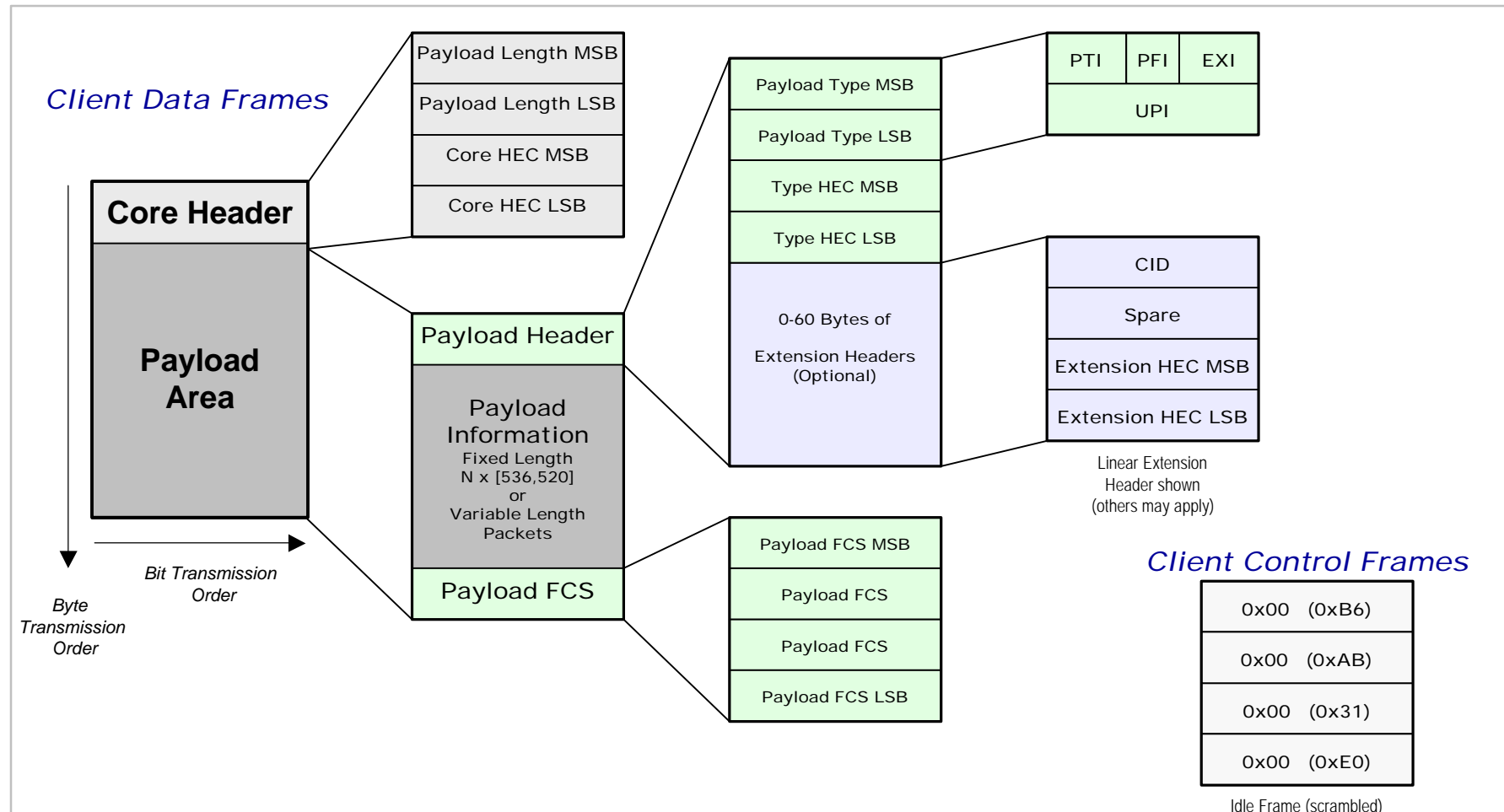
Functional Model



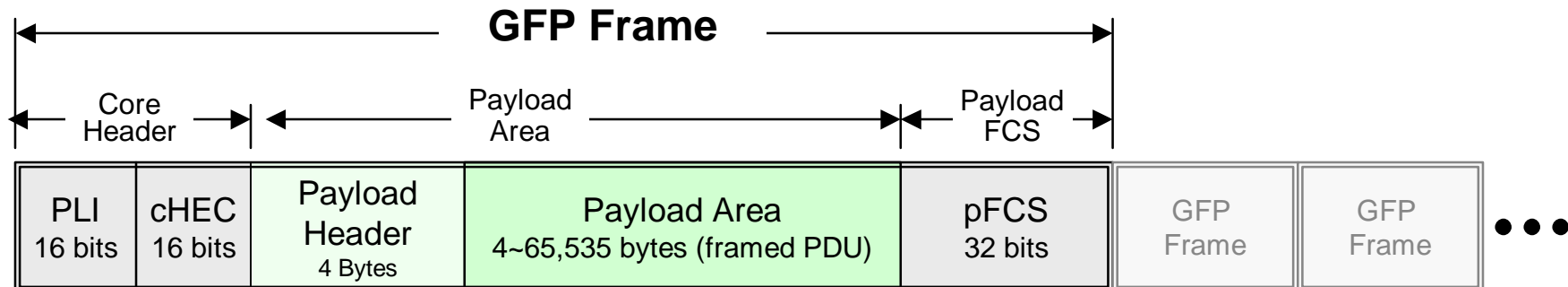
Frame Types



Generic Frame Structure



Basic GFP Frame Format



- **PLI** := Payload Length Indicator
- **cHEC** := Core Header CRC (ITU-T CRC-16)
- **Payload Area** := Framed PDU (PPP, IP, Ethernet, etc.)
- **Payload Header** := Client PDU management
- **pFCS** := *Optional* Payload FCS (ITU-T CRC-32)

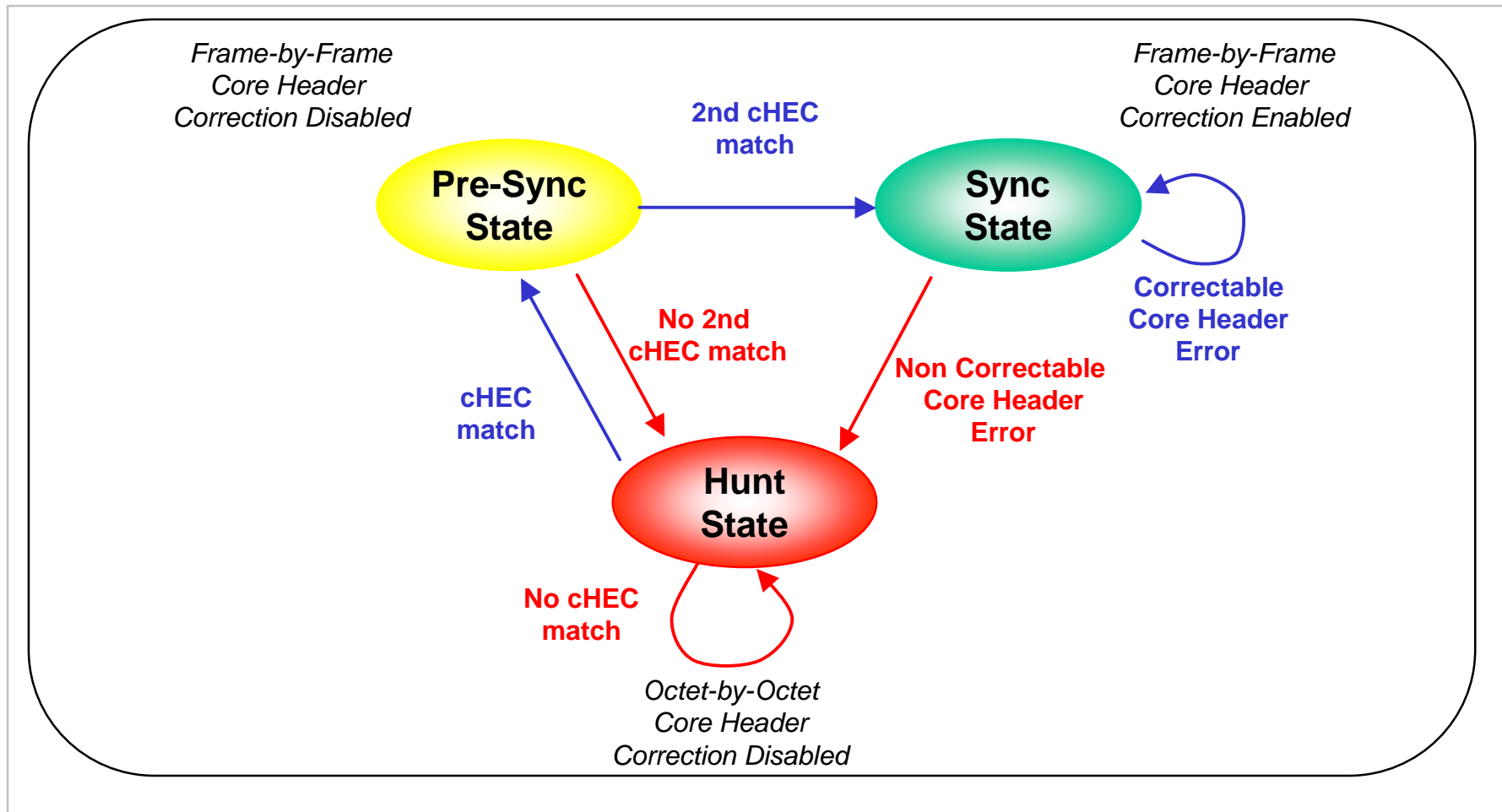
Frame Structure: Summary

- All GFP OAM&P functions handled via the GFP Core Header
- Payload Header supports any payload specific adaptation functions
 - Client types (Ethernet, IP, MPLS, Fibre Channel, etc.)
 - Client multiplexing (via Extension Headers)
 - Client link management (via Client Management Frames)
- Optional Payload FCS on a per frame basis
- Asynchronous rate adaptation via Idle Frames

GFP Procedures

- **Frame Delineation**
- **Frame/Client Multiplexing**
- **Adaptation Modes**
- **Scrambling**
 - Core Header
 - Payload Area
- **Error Handling**
 - Headers
 - Payload
- **Client Management**

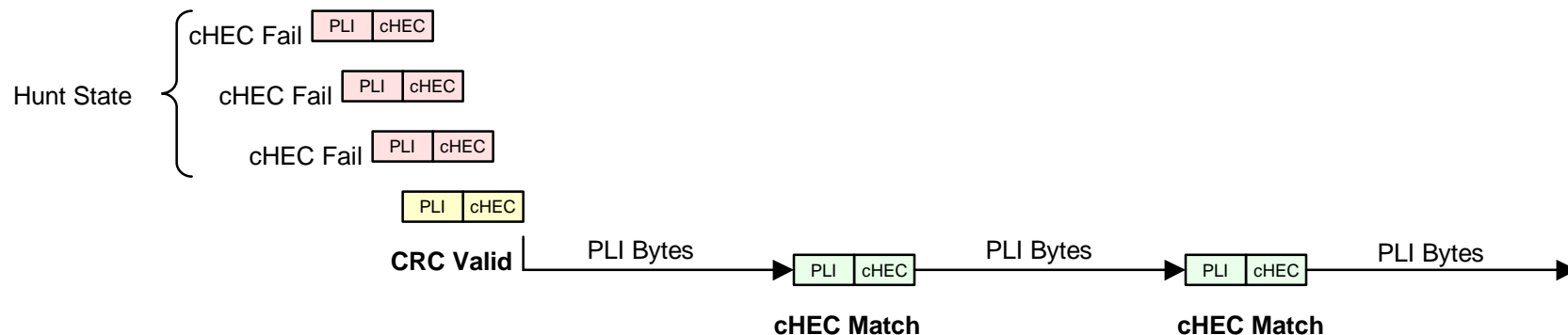
Frame Delineation: GFP State Machine



Frame Delineation An Example

- Two consecutive cHEC field matches vs. computed CHeC
- Pointer-based (PLI field) offset to next incoming frame

Octet or Bit synchronous stream



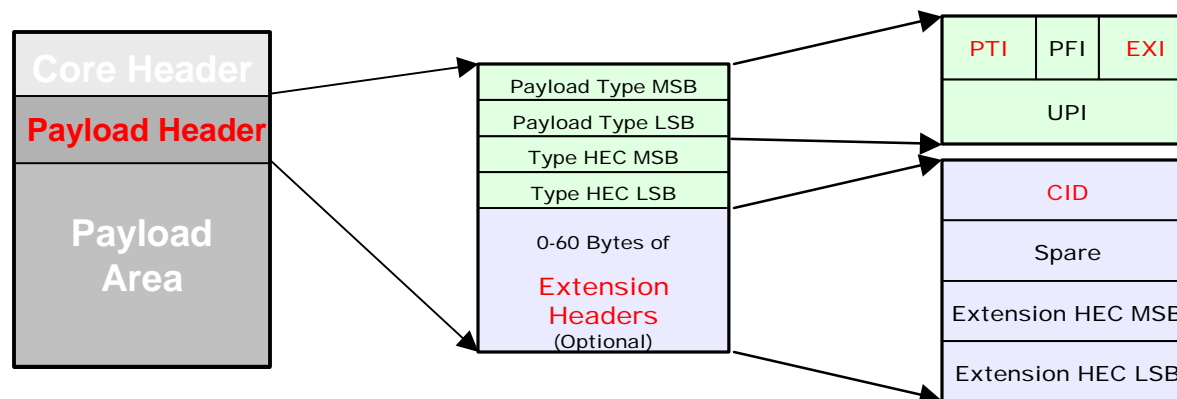
Multiplexing

■ Frame Multiplexing via PTI field:

- Client Data Frames have priority over Client Mgmt. Frames
- Client Management Frames have priority over Idle Frames

■ Client Multiplexing via Extension Headers:

- Null Extension Header on dedicated transport channels per client
- Linear Extension Header (point-to-point configurations)
- Ring Extension Header (ring configuration)



Frame Muxing:

PTI: Payload Type Id

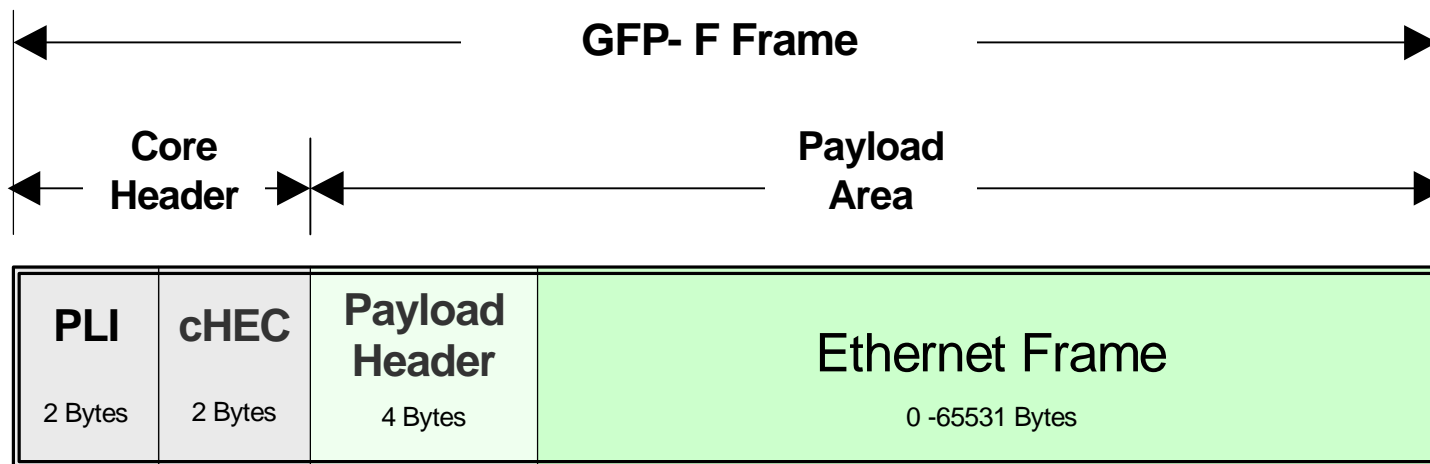
Client Muxing:

EXI: Extension Hdr ID

CID: Customer ID

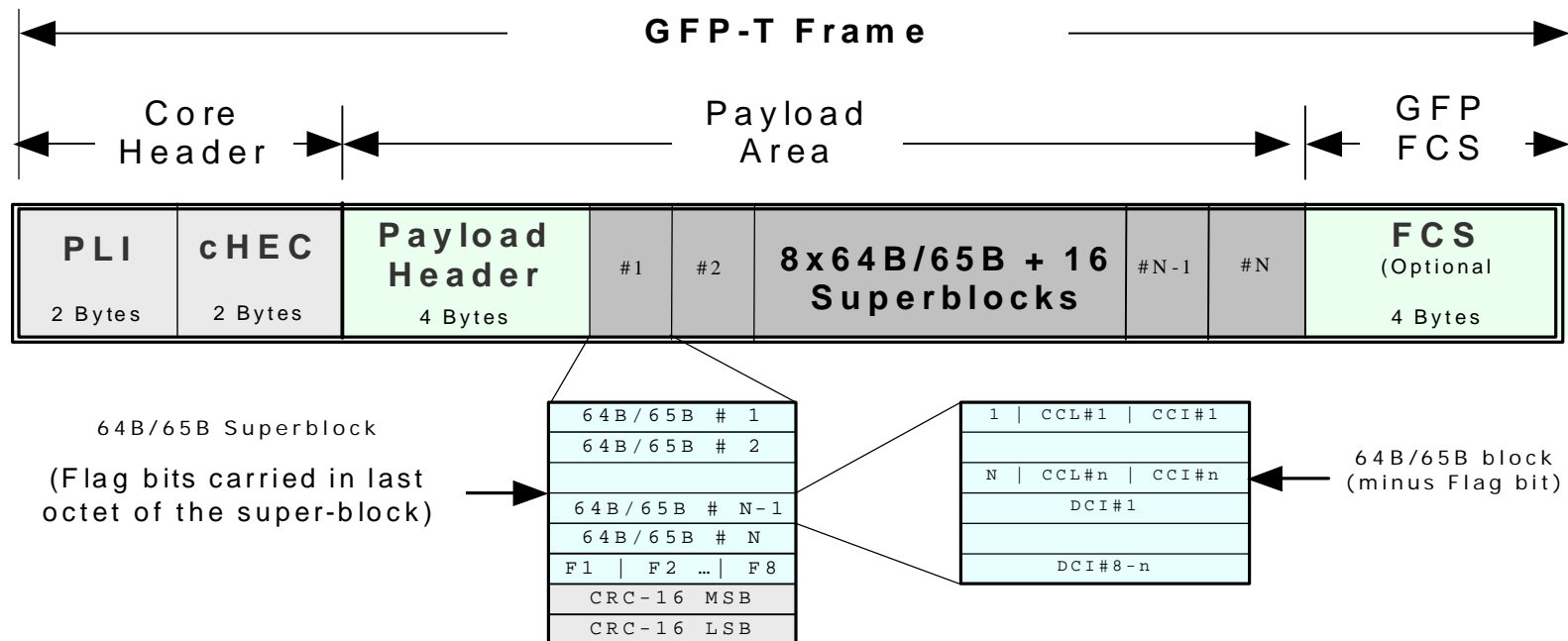
Adaptation Modes: Frame-Mapped GFP

- 1-to-1 mapping of L2 PDU to GFP payload
- UPI field indicates L2 PDU type
- Example: IEEE 802.3/Ethernet MAC frames



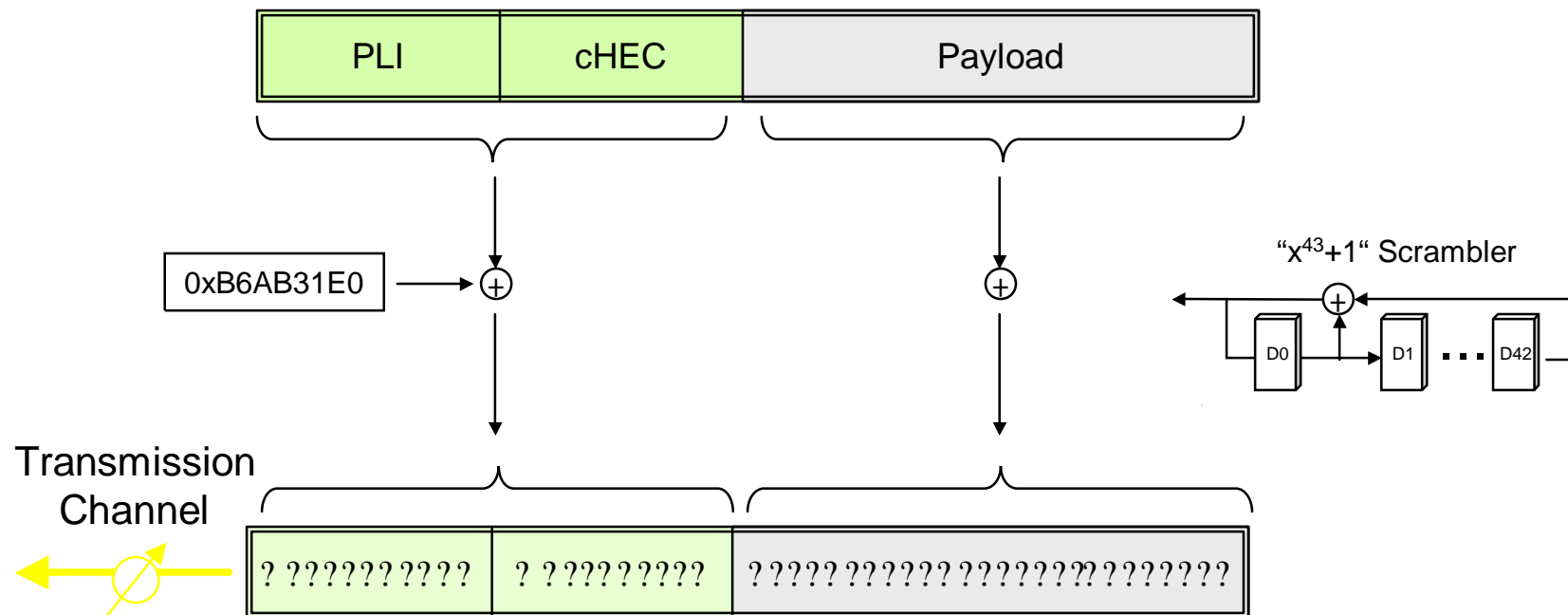
Adaptation Modes: Transparent-Mapped GFP

- N-to-1 mapping of L1 codewords to GFP payload
- Example: 8B/10B codewords



Scrambling: DC Balance & Payload Scrambler

- Header (PLI Field + CHeC) XOR'd with the 32 bit value "0xB6AB31E0" before transmission for DC balance.
- Payload scrambled with ATM-style self-synchronous scrambler



Error Handling

■ Multi-bit Error Detection & Correction:

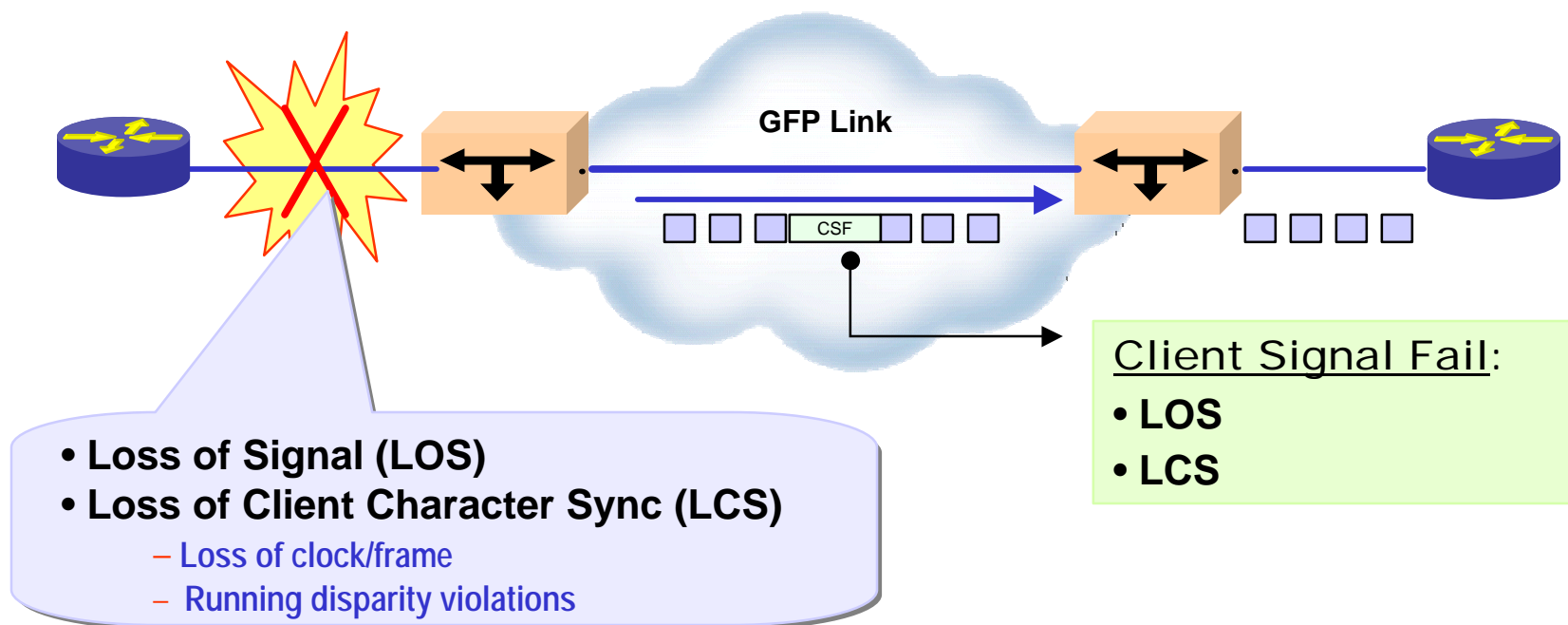
- Core Header – cHEC (ITU-T CRC-16):
 - Payload Type Field – tHEC (ITU-T CRC-16)
 - GFP-T payload (Optimized CRC-16)
- } *1-bit error correction*
- } *3-bit error correction*

■ Multi-bit Error Detection:

- Payload Extension Header – eHEC (ITU-T CRC-16)
- Payload Information Field – pFCS (ITU-T CRC-32)

Client Management

- **Client Signal Fail (CSF) indications sent periodically upon detection of a failure/degradation event**
- **Cleared by new Client Data Frame or CSF timeout**



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Performance:

Synchronization Loss Events

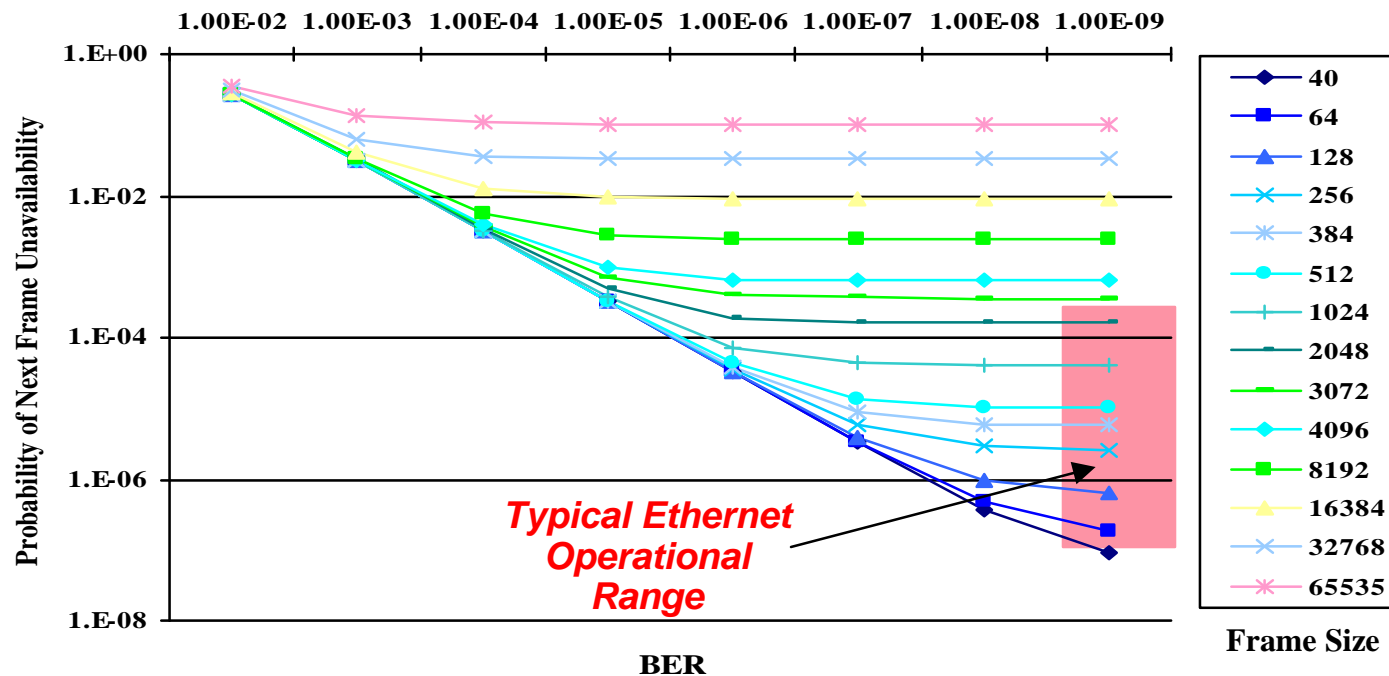
- Re-sync required whenever cHEC test fails
- Low synchronization loss probability for typical fiber BER
- Example: 40Bytes PDU at 40G. Loss event frequency decreases with increasing PDU size or decreasing BER

BER	Prob [Sync Loss]	Frequency
10^{-7}	5×10^{-12}	~ 48 min
10^{-8}	5×10^{-14}	~ 3.3 Days
10^{-9}	5×10^{-16}	~ 1 Year
10^{-10}	5×10^{-18}	~100 Years

Performance:

Missed Frame Delineation Events

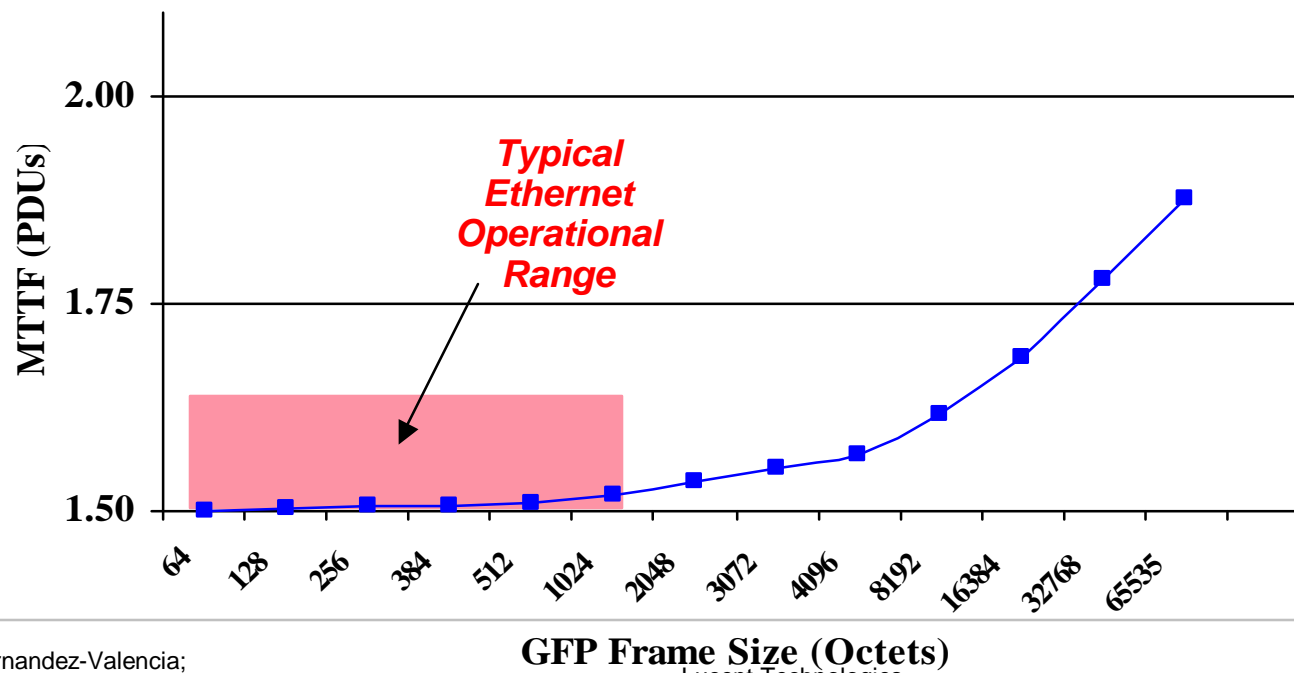
- Low probability of frame unavailability after LOF events
- Essentially insensitive to random errors for practical BERs



Performance:

Mean Time To Frame

- Datalink syncs after 2 consecutive cHEC matches
- Fast Mean Time to Frame (MTTF) delineation
- Largely insensitive to BER & line rate over the region of interest for (first order approximation)



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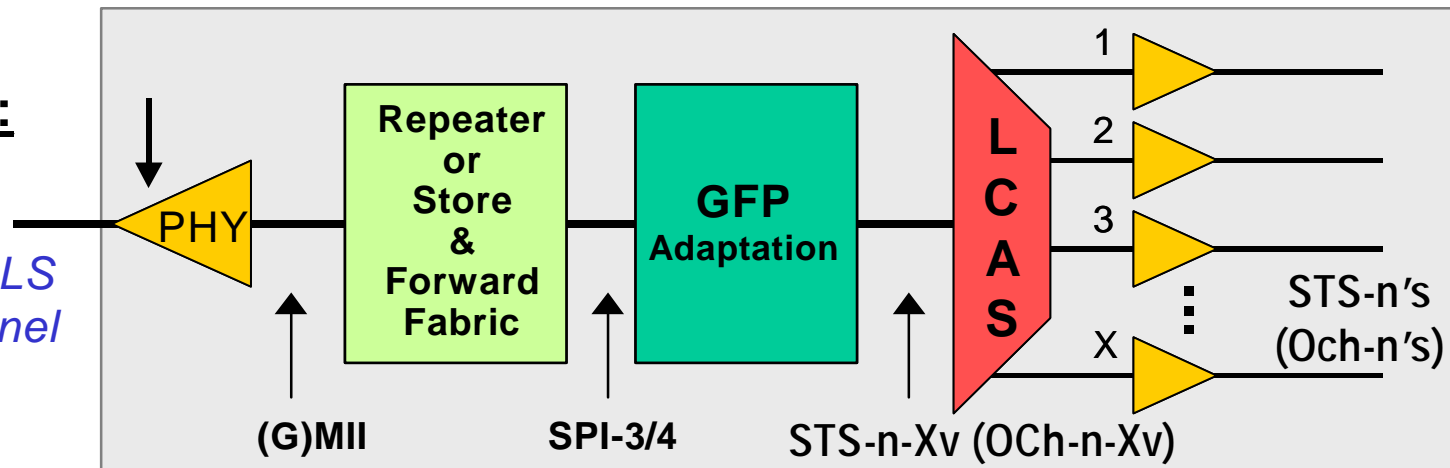
Hybrid Network Elements NG SONET/Data Systems

■ Three basic building blocks

- GFP (ITU-T G.7041/ANSI T1.105.02)
- Virtual Concatenation (ITU-T G.707/ANSI T1.105.02)
- LCAS (ITU-T G.707/ANSI T1.105.02)

Native Interfaces:

- *FE*
- *GbE*
- *PPP/IP/MPLS*
- *Fibre Channel*
- *FICON*
- *ESCON*



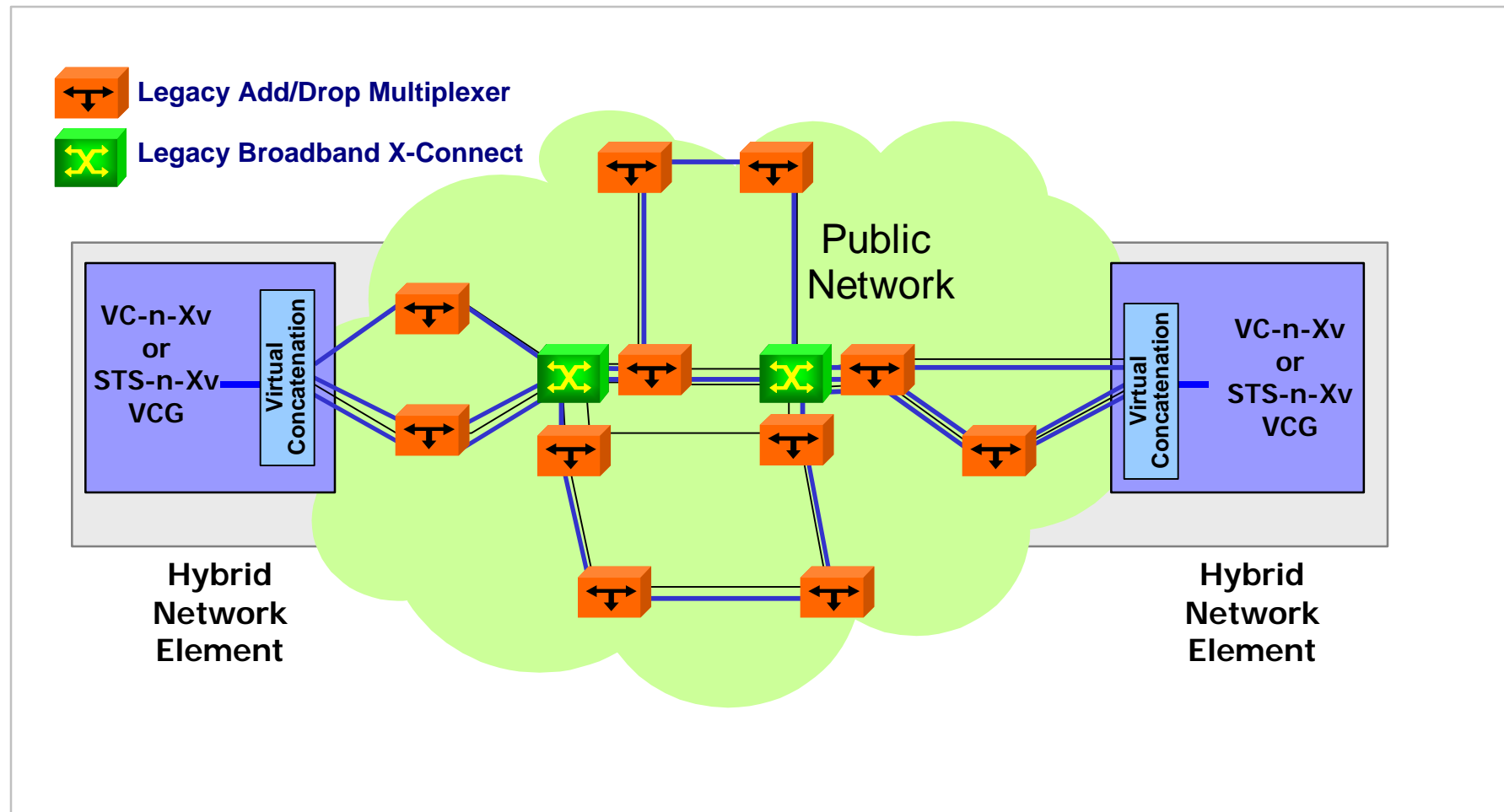
Hybrid Network Elements

Virtual Concatenation

- **Multiple SONET STS-Nc's (VC-n's) grouped into single STM-N-Xv Virtual Concatenation Group (VCG)**
 - Component STS-Nc's may be routed separately
 - Compensates differential network delays up to 32 ms
- **Network Operator provisions no. of channels (X) in VCG**
 - Solves SONET/SDH & OTN bandwidth “granularity problem”
- **Completely transparent to intermediate NEs.**
 - Only termination nodes need to support this feature

Hybrid Network Elements

Virtual Concatenation - Example





Hybrid Network Elements

Link Capacity Adjustment Scheme (LCAS)

- **Controls hitless addition/removal of STS-N's (VC-n's) to/from VCG under management control**
 - In-service hitless bandwidth modification
 - Address the dynamic management of bandwidth for data transport services over SONET/SDH
- **Manages automatic removal/addition of failed/repaired STS-N's from/to VCG**
- **Supports virtual channel protection through “load sharing” on STS-N's**
 - Works best on point-to-point links
- **ITU-T Recommendation G.7042 / ANSI T1.105.02**

GFP, Virtual Concatenation & LCAS

Transport Efficiency

Traffic Type	SONET		SDH	
	Contiguous	Virtual	Contiguous	Virtual
10Mbps Ethernet	STS-1 (20%)	VT-1.5-7v (89%)	VC-3 (20%)	VC-12-5v (92%)
100Mbit/s Fast Ethernet	STS-3c (67%)	STS-1-2v (100%)	VC-4 (67%)	VC-3-2v (100%) VC-12-46v (100%)
200Mbit/s (ESCON)	STS-6c (66%)	STS-1-4v (100%)	VC-4-4c (33%)	VC-3-4v (100%) VC-4-2v (66%)
1Gbps Fibre Channel	STS-21c (85%)	STS-1-18v (95%)	VC-4-16c (35%)	VC-4-6v (95%)
1Gbit/s Ethernet	STS-24c (83%)	STS-1-21v (92%)	VC-4-16c (42%)	VC-4-7v (95%)

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Summary

GFP Advantages

- **Versatility:** Enables transport services for either Layer 1 or Layer payloads:
 - PPP, IP, MPLS, Ethernet, HDLC & MAPOS at Layer 2
 - Fibre Channel, FICON, ESCON, Infiniband, DVB ASI at Layer 1
 - Endorsed by multiple communities including IEEE RPR WG & IETF
- **Scalability:** Demonstrate transport capabilities at rates from 10Mbps to 10Gbps (and soon beyond)
- **Simplicity:** Eliminates need for ATM and HDLC networking for simple connectivity services resulting in more efficient, lower-risk component designs
- **Component availability:** Broader user demand expected to drive future applications, feature maturity, interface commonality and lower cost



GFP Characteristics and Benefits

- **Simple Header Error Control (HEC) based synchronization:**
 - Generalizes ATM's HEC synchronization (inexpensive table lookup)
 - Supports variable or fixed length packets (IP/Ethernet datagrams, block codes or ATM cells)
- **Simple pointer-based frame delineation:**
 - Low processing complexity without payload expansion
 - Low (deterministic) adaptation overhead
 - High data link efficiency (scalable to 10Gbps and beyond)
 - Amenable to strict/loose QoS support, particularly for real-time services
- **Flexible traffic adaptation modes:**
 - Frame-Mapped GFP (GFP-F): *Suitable for elastic applications*
 - Transparent-Mapped GFP (GFP-T): *Suitable for in-elastic applications*