	parser can detect errors earlier than an LR(0) parser. ) (难易度:中)
	)(作勿反·宁)
A. True	
B. False	
<b>正确答案</b> : A <b>答案解释</b> : 暂无	
	mar is LR(1), then the LALR(1) parsing table cannot have any reduce-reduce conflicts. ) (难易度:中)
A. True	
B. False	
<b>正确答案</b> : B	
答案解释: 暂无	
	ological sort of attribute computation ofis unique.
	) (难易度:中)
A. True	
B. False	
<b>正确答案</b> : B	
答案解释: 暂无	
	ited attribute can be calculated by a postorder traversal of the parse tree. )(难易度:中)
A. True	
B. False	
<b>正确答案</b> : B	
答案解释: 暂无	
	ited attributes can be changed into synthesized attributes by suitable modification of the grammar, without changing the e of the grammar.
判断题 (4.0 分	)(难易度:中)
A. True	
B. False	
<b>正确答案</b> : A	
<b>答案解释</b> : 暂无	
6. The inpu	it of semantic analysis phase is
	)(难易度:中)
A. tokens	
B. CFG	
C. an abstrac	t syntax tree
D. a parse tre	

**正确答案**: C

7. Given the following declarations:

t1 = array [10] of int;
t2 = array [10] of int; if t1 and t2 isequivalent, it follows equivalence.
单选题 (4.0 分) (难易度:中)
A. structural
B. name
C. declaration
D. none of the above
<b>正确答案</b> : A
<b>答案解释</b> : 暂无
8. When inheriting a previously computed synthesized attribute during LR parsing, it is suitable to treat the computed synthesized attribute as
单选题 (4.0 分) (难易度:中)
A. return value
B. passing as parameter
C. external data structure
D. just leaving it on the value stack
<b>正确答案</b> : C
<b>答案解释:</b> 暂无
9. A Yacc-generated parser employ method to parse the input token stream. 单选题 (4.0 分) (难易度:中)
A. top-down
B. LL(1)
C. SLR(1)
D. LALR(1)
<b>正确答案</b> : D
<b>答案解释:</b> 暂无
名案解释: 智元  10. The symbol tables will not carry the information about the variables.  単选题 (4.0 分) (难易度:中)
10. The symbol tables will not carry the information about the variables.
10. The symbol tables will not carry the information about the variables. 单选题 (4.0 分)(难易度:中)
10. The symbol tables will not carry the information about the variables. 单选题 (4.0 分)(难易度:中) A. data type
10. The symbol tables will not carry the information about the variables. 单选题 (4.0 分)(难易度:中) A. data type B. scope

11. Consider the CFG of number with a one-character suffix **o**(for octal) or **d** (for decimal).

```
based-num→ num basechar
```

basechar → o | d

num → num digit | digit

digit  $\rightarrow 0|1|2|3|4|5|6|7|8|9$ 

In this case, we need two attributes, **base** (for indicating octal or decimal) and **val** (for value of the number).

- (1) Rewrite the grammar so that the computation of the attribute base does not depend on parent nodes.
- (2) Giventhe number "**789o**", show the dependency of attribute computation on its parsing tree of the above modified CFG.

简答题 (25.0 分) (难易度:中)

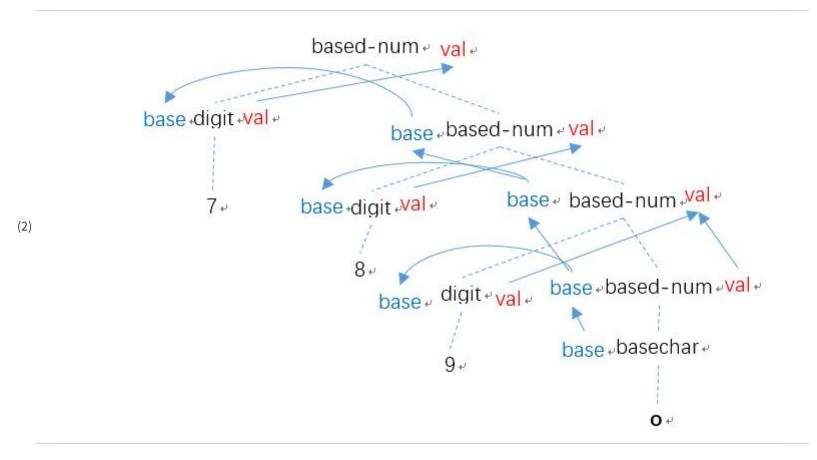
## 答案解释:

(1) The grammar can be modified as follows.

based-num → digit based-num | basechar

basechar  $\rightarrow$  **o** | **d** 

digit  $\rightarrow 0|1|2|3|4|5|6|7|8|9$ 



12. Consider an the following c-like code example.

```
1 double i;
2
3 void main()
4 {
   int a, b, c;
6
7
       int j = a+b;
8
       int c = a*a+b*b;
       char* a = "hello" ;
9
       print(a); print(j); print(c);
10
11 }
12 print(b);
13 }
```

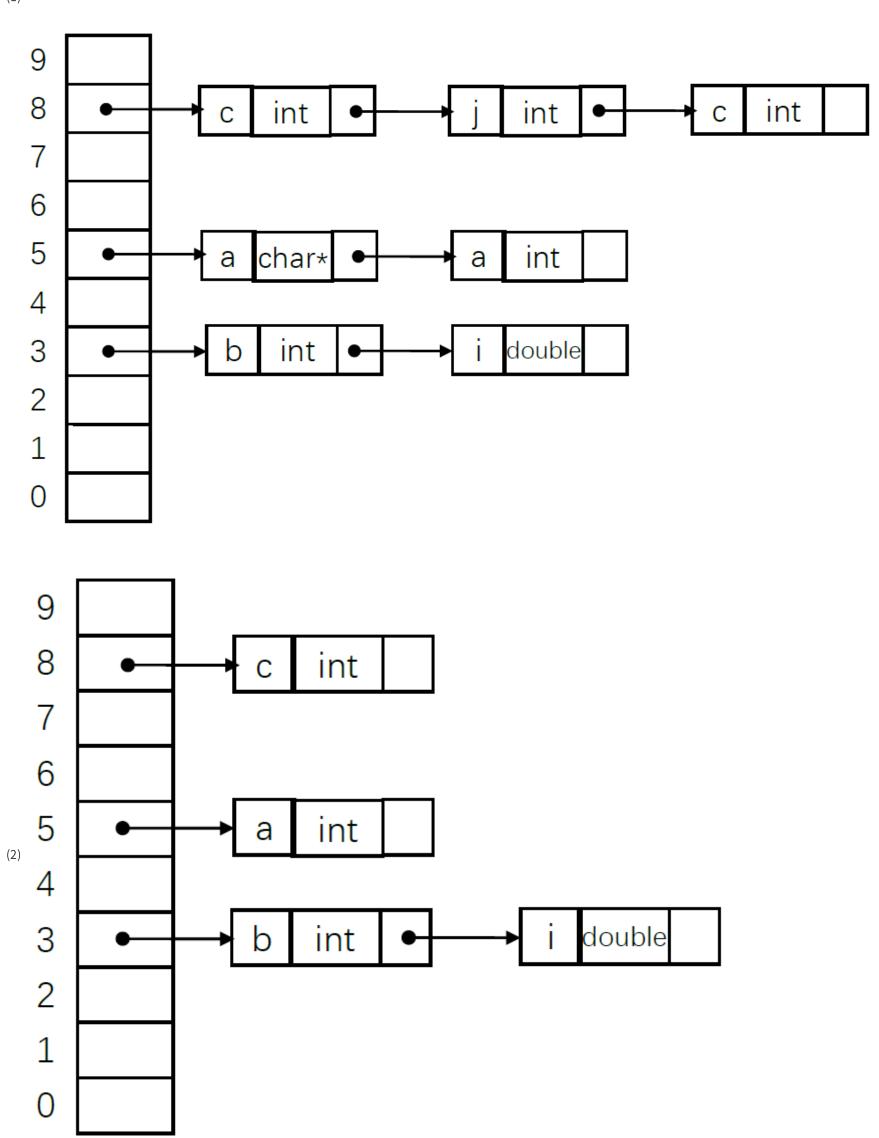
Suppose hash(i)=3, hash(a)=5, hash(b)=3, hash(c)=8, hash(j)=8, hash(f)=1, and only one symbol table is used.

- (1) Show the symbol table when line 9 is executed first time.
- (2) Show the symbol table when line 12 is executed.

简答题 (20.0分) (难易度:中)

## 答案解释:

(1)



13. 2. Given the following grammar rules. (15 cents)

 $S \rightarrow id \mid V := E$ 

 $V \rightarrow id$ 

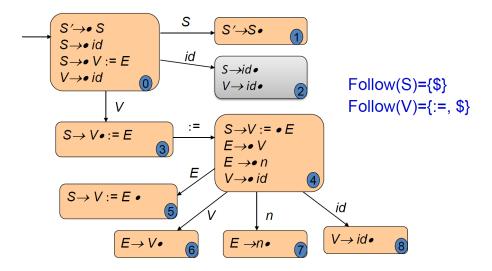
 $E \to V \mid \textbf{n}$ 

Is this CFG LR(0), SLR(1), LR(1) and LALR(1)? Why?

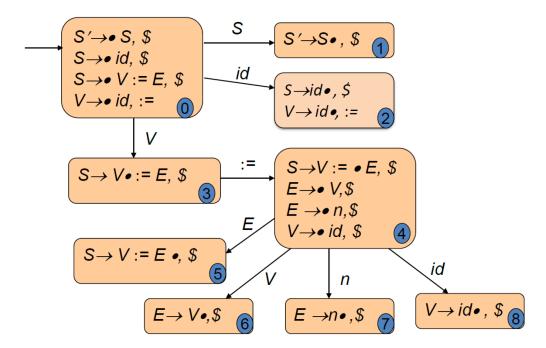
简答题 (15.0 分) (难易度:中)

## 答案解释:

- (1) The above CFG is not LR(0), SLR(1), But it is LR(1) and LALR(1). (6 cents)
- (2) Because there are reduce-reduceconflicts in its LR(0) DFA and SLR(1) DFA, but there is no conflict in itsLR(1) DFA and LALR(1) DFA. (9 cents)



LR(0) DFA and SLR(1) DFA. There is reduce-reduce conflict in the state 2.



LR(1) DFA and LALR(1) DFA. The conflict in state 2 is resolved.