

## Homework nr. 4

Let  $n$  be the system's size,  $\epsilon$  - the computation error,  $A \in \mathbb{R}^{n \times n}$  - a real squared matrix,  $k_{max}$  - the maximum number of computing steps.

1. Using all the iterative methods described below, approximate the inverse of matrix  $A$ . Use formula (5), (6) for choosing the initial matrix  $V_0$ .
2. For each method, at the end of the algorithm display the number of computed steps. If the method converged, display the norm:

$$\|A * A_{\text{approx}}^{-1} - I_n\|_1.$$

3. Consider the matrix:

$$A = \begin{pmatrix} 1 & 2 & 0 & \cdots & 0 \\ 0 & 1 & 2 & \cdots & 0 \\ \vdots & & & & \\ 0 & 0 & 0 & \cdots & 2 \\ 0 & 0 & 0 & \cdots & 1 \end{pmatrix}$$

Deduce inductively the general form of the inverse of matrix  $A$ , by successively running your program using different values for  $n$ , the size of the matrix  $A$ . Denote by  $A_{\text{exact}}^{-1}$  the exact inverse (computed with the deduced formula) and  $A_{\text{approx}}^{-1}$  the inverse approximated using the iterative methods. Display the norm:

$$\|A_{\text{approx}}^{-1} - A_{\text{exact}}^{-1}\|.$$

**Bonus 15 pt.:** Adapt one of the iterative algorithms for approximating the inverse of a matrix, for a non-square matrix,  $A \in \mathbb{R}^{m \times n}$ ,  $m \neq n$ .

### Iterative Methods for Approximating the Inverse of a Matrix

The inverse of a nonsingular matrix  $A$  is approximated by computing a sequence of matrices  $\{V_k; k \geq 0\}$  that converges to  $A^{-1}$ .

The Schultz method computes the sequence  $\{V_k; k \geq 0\}$  in the following way:

$$V_{k+1} = V_k(2I_n - AV_k), \quad k = 0, 1, 2, \dots, \quad V_0 - \text{given} \quad (1)$$

This method is also known as the Hotelling-Bodewig algorithm or the hyper-power iterative method.

Li and Li proposed the following two iterative methods for approximating the inverse of a matrix:

$$V_{k+1} = V_k \left[ 3I_n - AV_k(3I_n - AV_k) \right], \quad k = 0, 1, 2, \dots \quad (2)$$

$$V_{k+1} = \left[ I_n + \frac{1}{4}(I_n - V_k A)(3I_n - V_k A)^2 \right] V_k, \quad k = 0, 1, 2, \dots \quad (3)$$

To ensure the convergence of the sequence  $\{V_k; k \geq 0\}$ , the first matrix of the sequence,  $V_0$ , must be chosen such that the following relation holds:

$$\|AV_0 - I_n\| < 1. \quad (4)$$

*Ways of choosing  $V_0$  that ensure the sequence's convergence*

1.

$$V_0 = \frac{A^T}{\|A\|_1 \|A\|_\infty} \quad (5)$$

$$\begin{aligned} \|A\|_1 &= \max \left\{ \sum_{i=1}^n |a_{ij}|; j = 1, 2, \dots, n \right\}, \\ \|A\|_\infty &= \max \left\{ \sum_{j=1}^n |a_{ij}|; i = 1, 2, \dots, n \right\}. \end{aligned} \quad (6)$$

2. If matrix  $A$  is row diagonally dominant or column diagonally dominant:

$$|a_{ii}| > \sum_{\substack{j=1 \\ j \neq i}}^n |a_{ij}| \quad , \quad \forall i = 1, \dots, n \quad - \text{ row diagonally dominant}$$

$$|a_{ii}| > \sum_{\substack{i=1 \\ i \neq j}}^n |a_{ij}| \quad , \quad \forall i = 1, \dots, n \quad - \text{ column diagonally dominant}$$

the formula that ensure the sequence's convergence is:

$$V_0 = \mathbf{diag}\left(\frac{1}{a_{11}}, \frac{1}{a_{22}}, \dots, \frac{1}{a_{nn}}\right) \quad (7)$$

3. If matrix  $A$  is symmetric and positive definite, a good choice of  $V_0$  is:

$$V_0 = \frac{1}{\|A\|_F} I_n \quad , \quad \|A\|_F = \left( \sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2 \right)^{\frac{1}{2}} \quad (8)$$

4.

$$V_0 = \alpha A^T \quad , \quad 0 < \alpha < \frac{2}{\|A\|_2^2} \quad , \quad \|A\|_2 = \sqrt{\rho(A^T A)} \quad (9)$$

5.

$$V_0 = \alpha I_n \quad , \quad \alpha \in \mathbb{R} \quad , \quad \max \{ |1 - \lambda_i \alpha| ; \lambda_i \text{ eigenvalue for } A \} < 1 \quad (10)$$

The algorithm stops if:

- 1)  $\|V_k - V_{k-1}\| < \epsilon$  (or  $\|I_n - V_k A\| < \epsilon$ ) - in this case matrix  $V_k$  approximates the inverse,  $V_k \approx A^{-1}$  ; (one can use any norm described in (6))
- 2)  $k > k_{max}$  - the maximum number of iterations has been exceeded, the approximation of the inverse with the desired precision failed ;

3)  $\|V_k - V_{k-1}\| > 10^{10}$ - the computed sequence does not converge, the algorithm failed to approximate the inverse of matrix  $A$ .

Not all the elements of the sequence  $\{V_k\}$  must be stored in order to obtain an approximation for  $A^{-1}$ , it suffice to use only two matrices, one for storing  $V_k$  and the other for  $V_{k+1}$ .

For economically computing a matrix of type  $C = aI_n - AV$  (where  $a$  is a real value,  $a \in \mathbb{R}$ ), one computes only once in the program the matrix  $B = (-A)$ , then the matrix  $C = B * V$  is calculated. The required  $aI - AV$  matrix is obtained from  $C$  by adding the constant  $a$  to all the diagonal elements of matrix  $C$  ( $c_{ii} = c_{ii} + a, \forall i$ ).

A possible approximation scheme for computing an approximation for the inverse of a matrix is the following:

# *Iterative Method for Approximating the Inverse of a Matrix*

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$$V0 = V1 = \frac{A^T}{\|A\|_1 \|A\|_\infty};$$

 $k = 0 ;$ 
 $k_{max} = 10000 ; //$  a possible value
do
{
   $V0 = V1 ; //$  deep copy
  compute  $V1$  using  $V0$  and one of the formulae (1), (2), (3);
  compute  $\Delta V = \|V1 - V0\| ;$ 
   $k = k + 1;$ 
}
while ( $\Delta V \geq \epsilon$  and  $k \leq k_{max}$  and  $\Delta V \leq 10^{10}$ )
if (  $\Delta V < \epsilon$  )  $V1 \approx A^{-1} ;$ 
else divergence ;  $//$  the algorithm fails to approximate the
inverse

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