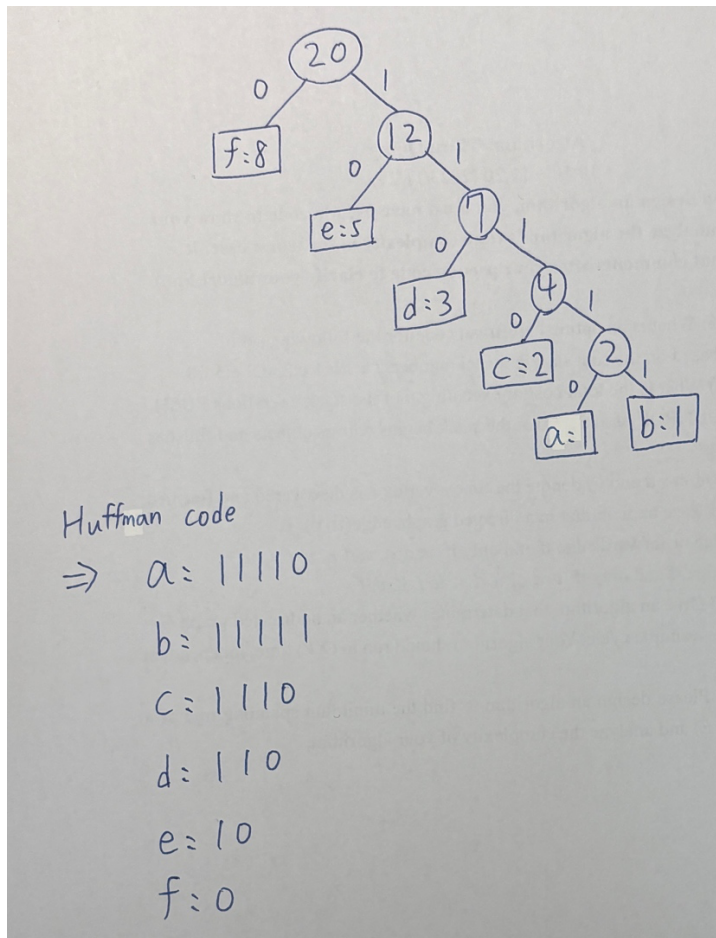


1.



評分標準：

寫出正確且合理的 Huffman code，10 分

Huffman code 錯誤，但有畫出正確且合理的樹，5 分

2.

$$\begin{aligned}
 2. \text{ push} : \phi(s_{m+1}) - \phi(s_m) + 1 &= 2 \\
 \text{multi-pop}(k) : \phi(s_{m+1}) - \phi(s_m) + k &= 0 \\
 \text{amortized cost} &\leq 2n \\
 \text{total cost} &= \text{amortized cost} - \phi(s_n) + \phi(s_0) \\
 &\leq 2n - s_n + s_0
 \end{aligned}$$

3.

First,

1. u is an ancestor of $v \Leftrightarrow u.d < v.d < v.f < u.f$.
2. u is a descendant of $v \Leftrightarrow v.d < u.d < u.f < v.f$.

Therefore,

- a. (u, v) is a tree edge or forward edge $\Leftrightarrow u$ is an ancestor of v .(5%)
- b. (u, v) is a back edge $\Leftrightarrow u$ is a descendant of v .(5%)

4.

```
4. bool vis[V];  
vector<int> G[V];  
bool DFS(int u, int f){  
    if(vis[u]) return 1;  
    vis[u] = 1;  
    bool cyc = 0;  
    for(int v : G[u]){  
        if(v == f) continue;  
        cyc |= DFS(v, u);  
    }  
    return cyc;  
}  
bool Cycle(int V){  
    bool cyc = 0;  
    for(int v = 1; v <= V; v++){  
        if(!vis[v]) cyc |= DFS(v, 0);  
    }  
    return cyc;  
}
```

5.

You can either use Kruskal's or Prim's algorithm to solve this problem.

Kruskal's algorithm:

MST-KRUSKAL(G, w)

```
1   $A = \emptyset$ 
2  for each vertex  $v \in G.V$ 
3      MAKE-SET( $v$ )
4  create a single list of the edges in  $G.E$ 
5  sort the list of edges into monotonically increasing order by weight  $w$ 
6  for each edge  $(u, v)$  taken from the sorted list in order
7      if FIND-SET( $u$ )  $\neq$  FIND-SET( $v$ )
8           $A = A \cup \{(u, v)\}$ 
9          UNION( $u, v$ )
10 return  $A$ 
```

Time complexity of Kruskal's algorithm:

Use an edge array for sorting and use Union-Find operation.

The overall time complexity is $O(E \log V)$.

Prim's algorithm:

MST-PRIM(G, w, r)

```
1  for each vertex  $u \in G.V$ 
2       $u.key = \infty$ 
3       $u.\pi = \text{NIL}$ 
4   $r.key = 0$ 
5   $Q = \emptyset$ 
6  for each vertex  $u \in G.V$ 
7      INSERT( $Q, u$ )
8  while  $Q \neq \emptyset$ 
9       $u = \text{EXTRACT-MIN}(Q)$       // add  $u$  to the tree
10     for each vertex  $v$  in  $G.Adj[u]$  // update keys of  $u$ 's non-tree neighbors
11         if  $v \in Q$  and  $w(u, v) < v.key$ 
12              $v.\pi = u$ 
13              $v.key = w(u, v)$ 
14         DECREASE-KEY( $Q, v, w(u, v)$ )
```

Time complexity of Prim's algorithm:

Use a binary heap as the priority queue.

The overall time complexity of Prim's algorithm is $O(E \log V)$.

Grading policy:

(7%) Design a correct algorithm.

(3%) Analyze the time complexity.