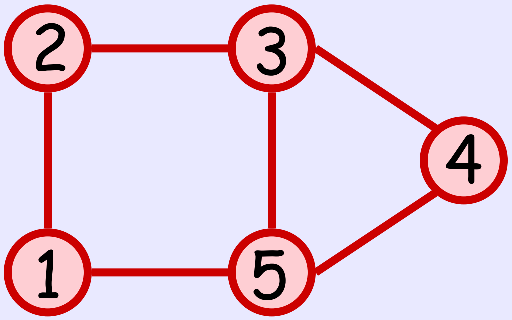
Chapter 20: Elementary Graph Algorithms I

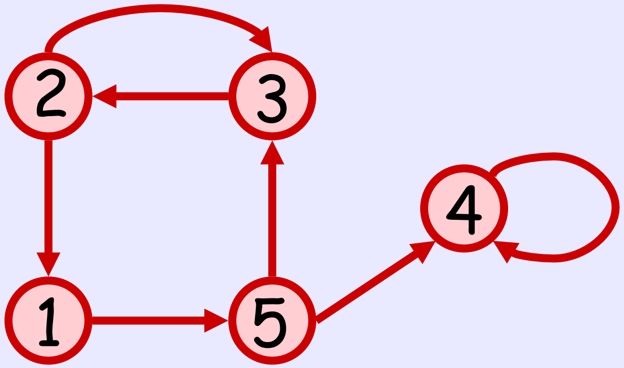
About this lecture

- Representation of Graph
  - Adjacency List, Adjacency Matrix
- Breadth First Search

Graph



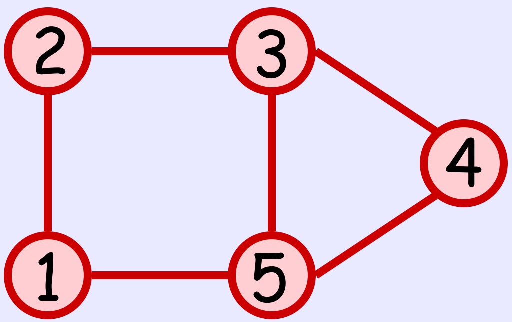
undirected

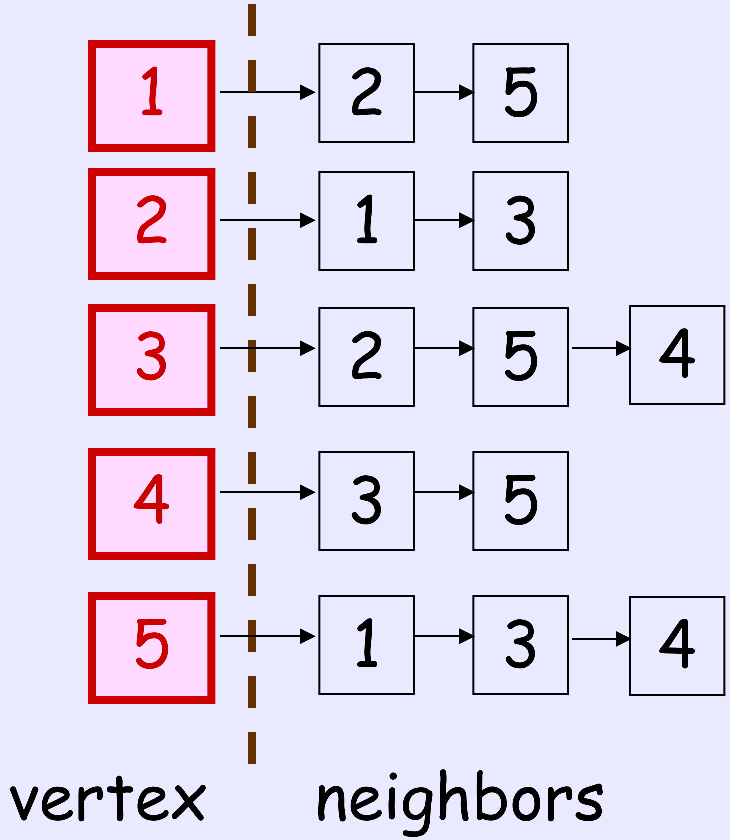


directed

Adjacency List (1)

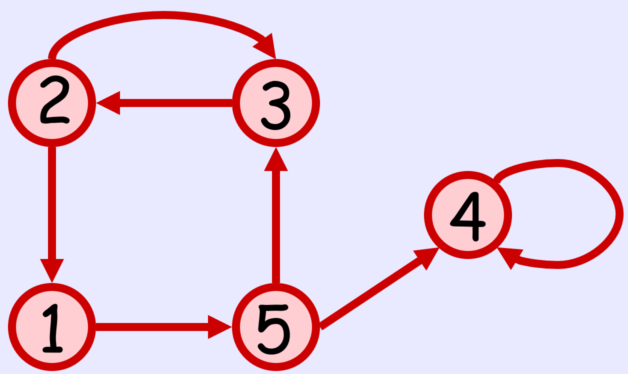
- For each vertex u, store its neighbors in a linked list

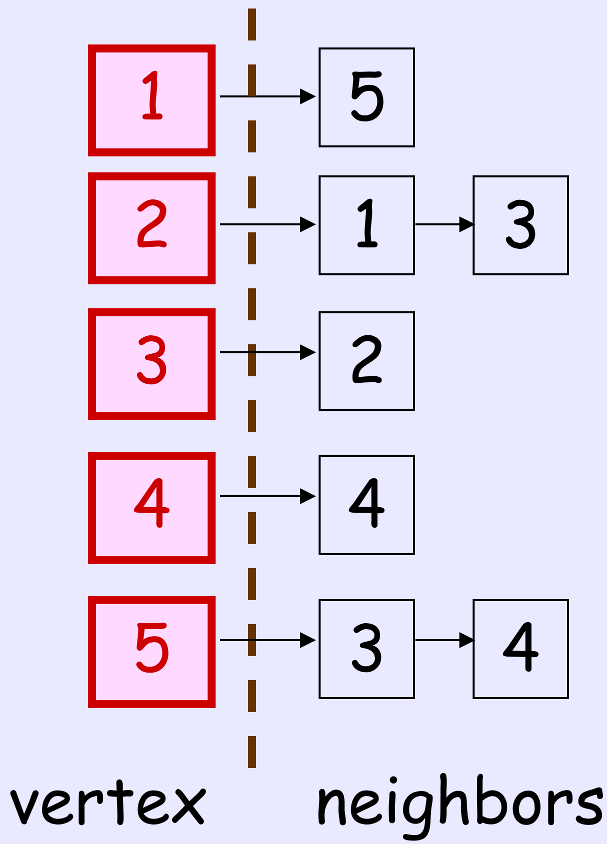




Adjacency List (2)

- For each vertex u, store its neighbors in a linked list





Adjacency List (3)

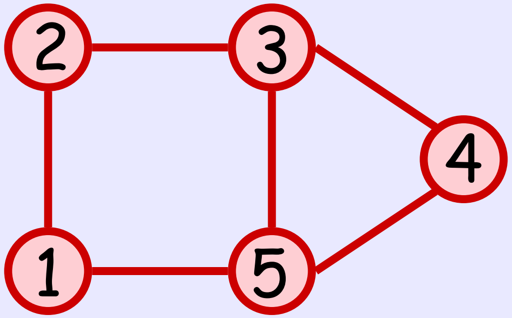
- Let $G = (V, E)$ be an input graph
- Using Adjacency List representation :
  - Space : $O(|V| + |E|)$

    → Excellent when $|E|$ is small
  - Easy to list all neighbors of a vertex
  - Takes $O(|V|)$ time to check if a vertex u is a neighbor of a vertex v
- can also represent weighted graph

Adjacency Matrix (1)

- Use a $|V| \times |V|$ matrix A such that

$$A(u,v) = 1 \quad \text{if } (u,v) \text{ is an edge}$$
$$A(u,v) = 0 \quad \text{otherwise}$$



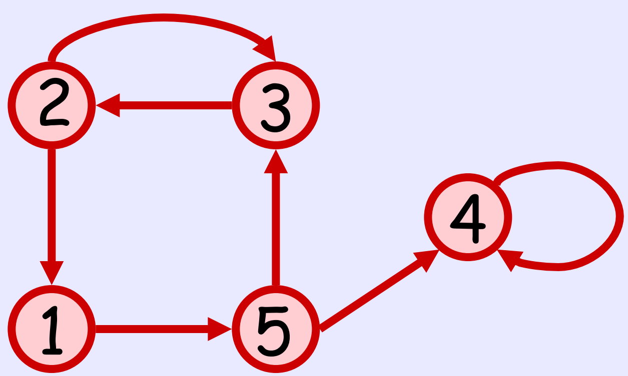
	1	2	3	4	5
1	0	1	0	0	1
2	1	0	1	0	0
3	0	1	0	1	1
4	0	0	1	0	1
5	1	0	1	1	0

Adjacency Matrix (2)

- Use a $|V| \times |V|$ matrix A such that

$A(u,v) = 1$ if (u,v) is an edge

$A(u,v) = 0$ otherwise



	1	2	3	4	5
1	0	0	0	0	1
2	1	0	1	0	0
3	0	1	0	0	0
4	0	0	0	1	0
5	0	0	1	1	0

Adjacency Matrix (3)

- Let $G = (V, E)$ be an input graph
- Using Adjacency Matrix representation :
  - Space : $O(|V|^2)$

    ➔ Bad when $|E|$ is small
  - $O(1)$ time to check if a vertex u is a neighbor of a vertex v
  - $\Theta(|V|)$ time to list all neighbors
- can also represent weighted graph

Transpose of a Matrix

- Let A be an $n \times m$ matrix
- Definition:

  The **transpose** of A, denoted by A^T, is an $m \times n$ matrix such that

  $$A^T(u,v) = A\,(v,u) \quad \text{for every } u, v$$

➔ If A is an adjacency matrix of an undirected graph, then $A = A^T$

Breadth First Search (BFS)

- A simple algorithm to find all vertices reachable from a particular vertex s
  - s is called source vertex
- Idea: Explore vertices in rounds
  - At Round k, visit all vertices whose shortest distance (#edges) from s is k-1
  - Also, discover all vertices whose shortest distance from s is k

The BFS Algorithm

1. Mark s as discovered in Round 0
2. For Round k = 1, 2, 3, …,

   For (each u discovered in Round k-1)

   { Mark u as visited ;

   Visit each neighbor v of u ;

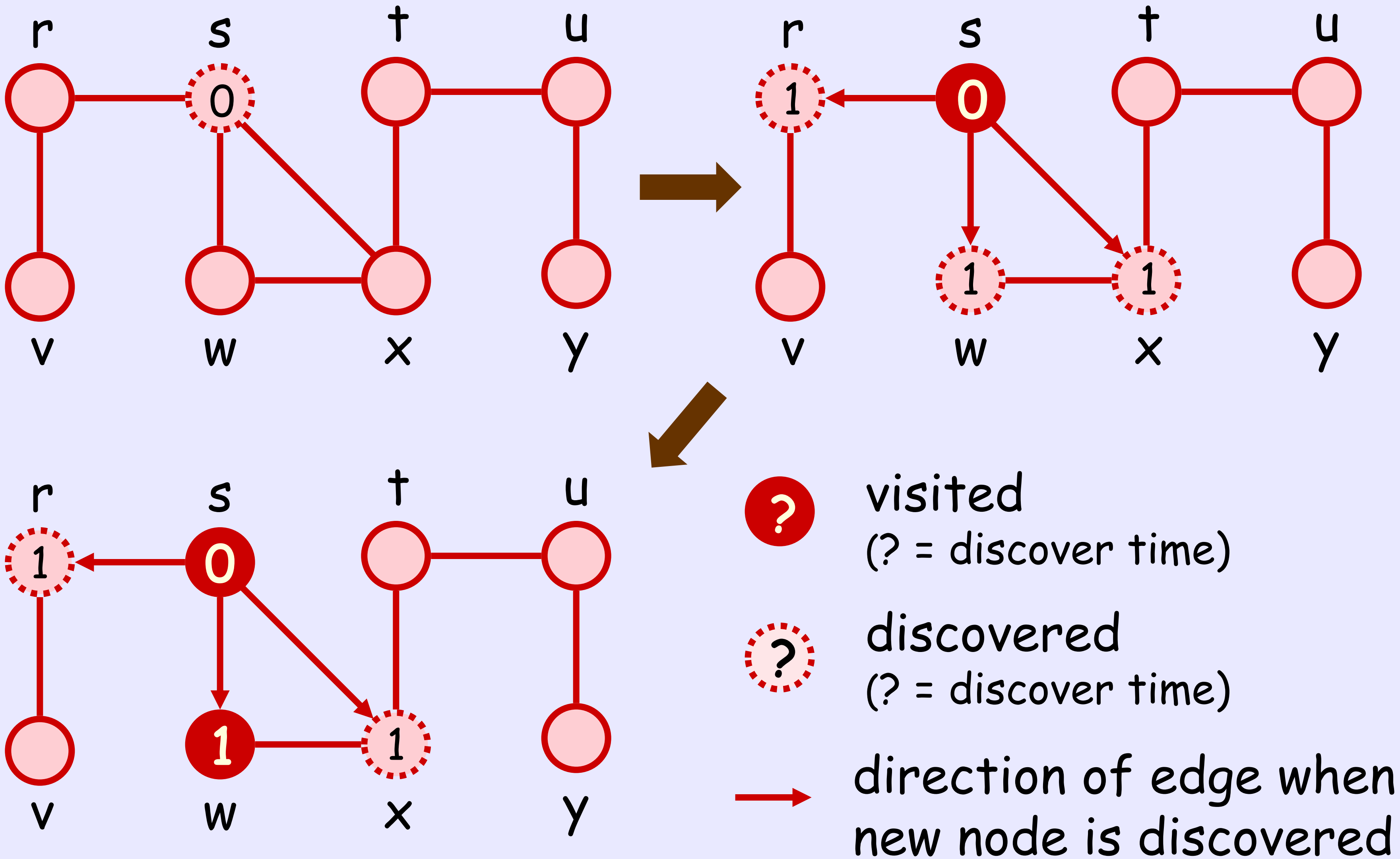
   If (v not visited and not discovered)

   Mark v as discovered in Round k ;

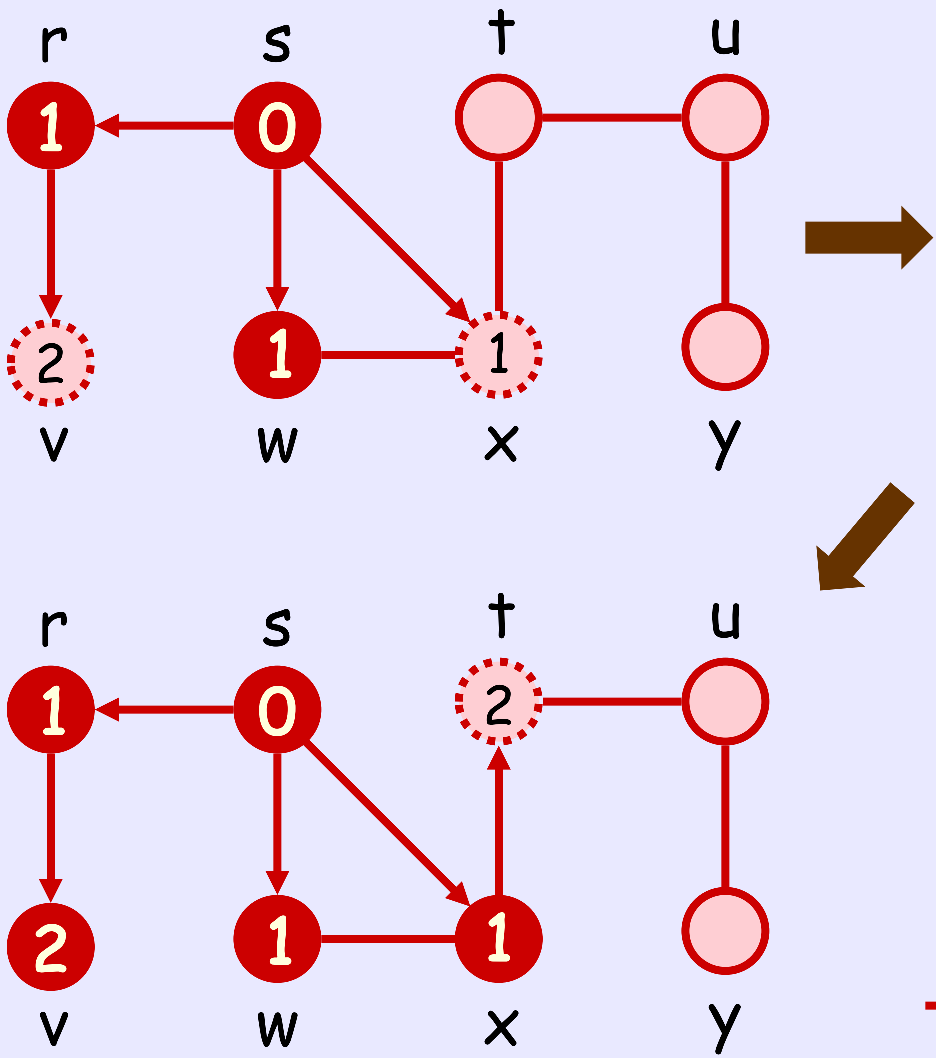
   } (Implemented by Queue)

Stop if no vertices were discovered in Round k-1

Example (s = source)



Example (s = source)

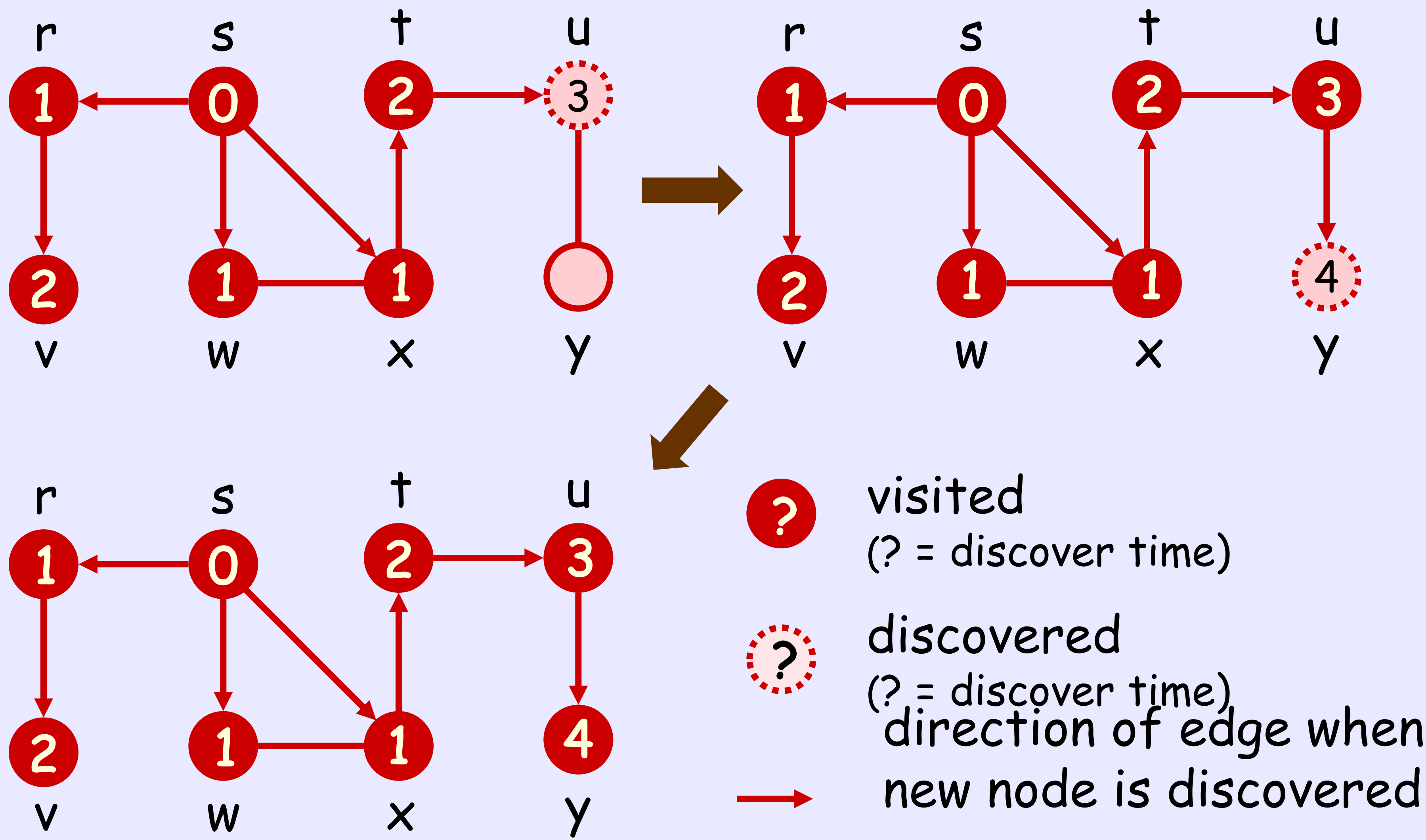


? visited
(? = discover time)

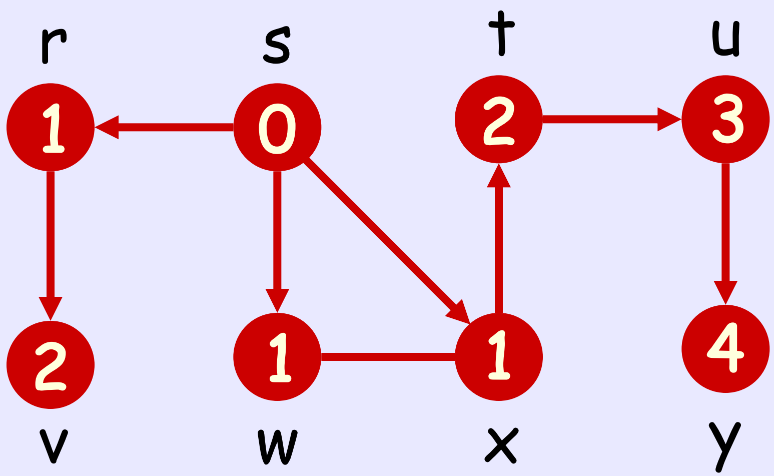
? discovered
(? = discover time)

direction of edge when new node is discovered

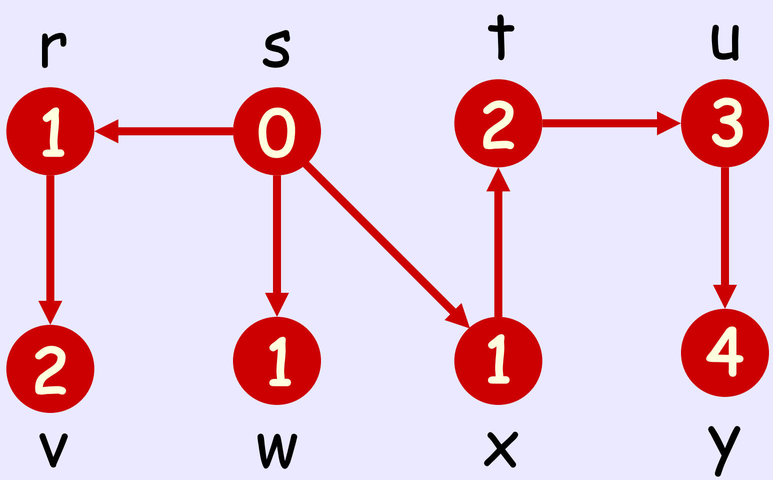
Example (s = source)



Example (s = source)



Done when no new node is discovered



The directed edges form a tree that contains all nodes reachable from s

Called BFS tree of s

Correctness

- The correctness of BFS follows from the following theorem :

Theorem: A vertex v is discovered in Round k if and only if shortest distance of v from source s is k

Proof: By induction

Performance (1)

- BFS algorithm is easily done if we use
  - an $O(|V|)$-size array to store discovered/visited information
  - a separate list for each round to store the vertices discovered in that round
- Since no vertex is discovered twice, and each edge is visited at most twice (why?)

  ➔ Total time: $O(|V|+|E|)$

  ➔ Total space: $O(|V|+|E|)$ (adjacency-list representation)

Performance (2)

- Instead of using a separate list for each round, we can use a common queue
  - When a vertex is discovered, we put it at the end of the queue
  - To pick a vertex to visit in Step 2, we pick the one at the front of the queue
  - Done when no vertex is in the queue

➔ No improvement in time/space …

➔ But algorithm is simplified

Practice at Home

- Exercise: 20.1-5, 20.1-7, 20.1-8, 20.2-6, 20.2-7
- Bonus

n-Queen Problem: Implement an algorithm that takes an integer n as input and determines the number of solutions to the n-Queen problem. You need to give the time complexity of your algorithm.

Due: Dec. 1