

## Assignment 1

### Exercise 1: Compiling and running on Dardel.

**Task 1.1** Describe all steps to connect to Dardel, compile the code, and execute the code on computing nodes with SLURM (Interactive and batch).

- PDC Account
- Kerberos and ssh
- Compiling
- Running on interactive node
- Running on batch job

To connect to dardel, use the following steps: - Apply a PDC account, after which you will be sent a password for your account - PDC uses the kerberos authentication protocol and uses ssh to connect to dardel.

1. Create a kerberos ticket using ``kinit --forwardable user@NADA.KTH.SE``
2. ssh into Dardel using this ticket with ``ssh YourUsername@dardel.pdc.kth.se``

- Compiling

1. Code can be compiled using the `cc`, `CC` wrappers for `gcc`.

- Running on an interactive node

```
$ cd /cfs/klemming/nobackup/u/user
$ salloc --nodes=1 -t 01:00:00 -A edu23.DD2356 -p main
$ srun -n 128 ./hello.out
```

- Running on a batch job

We create a file `job.sh`:

```
#!/bin/bash -l
# The -l above is required to get the full environment with modules

# The name of the script is myjob
#SBATCH -J myjob
# Only 1 hour wall-clock time will be given to this job
#SBATCH -t 0:01:00
#SBATCH -A edu23.DD2356
# Number of nodes
#SBATCH -p main
#SBATCH --nodes=1
#SBATCH -e error_file.e

# Run the executable file
# and write the output into my_output_file
srun -n 128 ./hello.out > hello_output
```

Then we submit it with

```
sbatch job.sh
```

For both, the output looks something like

```
Hello world from processor nid001264, rank 86 out of 128 processors
Hello world from processor nid001264, rank 16 out of 128 processors
Hello world from processor nid001264, rank 97 out of 128 processors
```

```
...
...
...
```

Hello world from processor nid001264, rank 99 out of 128 processors  
Hello world from processor nid001264, rank 127 out of 128 processors  
Hello world from processor nid001264, rank 115 out of 128 processors

**Task 1.2 How many computing nodes does Dardel have?** - 554 Nodes [488 SNIC thin, 20 SNIC large, 8 SNIC Huge, 2 SNIC Giant, 36 KTH]

**How many CPUs, CPU core, and memory does each computing node have?** - 2 CPUs, 128 Physical cores with 256 virtual cores. Memory size varies with the type of node.

**What is the total number of cores and memory of the Dardel computer?** - 70912 Total cores in the CPU partition. With 156.672TB of memory.

**What is the fastest supercomputer in the world? What are the fastest European and Asian supercomputers? What is their power usage?** - Frontier in the world [21,1kW]. Asian - Supercomputer Fugaku [29,899kW]. European - LUMI [6,016kW].

## Exercise 2: Sustainability and supercomputers

**Task 2.1** Using the calculator at <http://calculator.green-algorithms.org/>. Evaluate the power usage and carbon footprint of running a simulation on 10 Dardel computing nodes, assuming that we only use CPU, all the memory on the nodes, and we neglect the network costs. The simulation runs on 10 computing nodes for 12 hours.

- 12 Hours, CPU, 128 Cores, 256gb \* 10 = 20.09kg CO2, 180.71kWh

Estimate done using the EPYC 7513 processor option.

### Exercise 3: Modeling Sparse Matrix-Vector Multiply.

**Task 3.1** What is the performance in total execution time - do not consider data movement - according to your performance model on Dardel or your local computer for different sparse matrices  $nrows = 10^2, 10^4, 10^6$ , and  $10^8$ ?

Laptop - i7-10750H (Base 2.6Ghz - 5.0GHz)

Taking the base rate gives  $(1/2.6) \cdot 10^{-9}$  s/operations

for each nnz:

$$nnz = 460 \implies 0.0000001769s$$

$$nnz = 49600 \implies 0.00001907s$$

$$nnz = 4996000 \implies 0.0019215s$$

$$nnz = 499960000 \implies 0.192292s$$

**Task 3.2** What is the measured performance in total execution time and floating-point operations per second running `spmv.c` for different sizes  $= 10^2, 10^4, 10^6$ , and  $10^8$ ? Compare the results from the performance model and experimental results. Discuss the comparison in the report.

Time for Sparse Ax,  $nrows=100$ ,  $nnz=460$ ,  $T = 0.000002s$ ,  $FLOPS = 2 \cdot nnz/T = 460000000$

Time for Sparse Ax,  $nrows=10000$ ,  $nnz=49600$ ,  $T = 0.000262s$ ,  $FLOPS = 2 \cdot nnz/T = 378625954$

Time for Sparse Ax,  $nrows=1000000$ ,  $nnz=4996000$ ,  $T = 0.028904s$ ,  $FLOPS = 2 \cdot nnz/T = 345696097$

Time for Sparse Ax,  $nrows=100000000$ ,  $nnz=499960000$ ,  $T = 0.598384s$ ,  $FLOPS = 2 \cdot nnz/T = 1671033984$

The theoretical performance model consistently underestimates the execution time, which is reasonable as the model only takes into account computing time.

**Task 3.3** What is the main reason for the observed difference between the modeled value and the measured value?

The model only takes into account computing time, and not read and write times.

**Task 3.4** What are the read bandwidth values you measure running `spmv.c` for different sizes  $nrows = 10^2, 10^4, 10^6$ , and  $10^8$ ?

`sizeof(int) + sizeof(double) = 12 bytes`

`Read bandwidth = (nnz + nrows)(sizeof(int) + sizeof(double)) / T = 12 * (nnz + nrows) / T`

$nrows=100$ ,  $nnz=460$ , Read bandwidth = 3360 MB/s

$nrows=10000$ ,  $nnz=49600$ , Read bandwidth = 2729 MB/s

$nrows=1000000$ ,  $nnz=4996000$ , Read bandwidth = 2489 MB/s

$nrows=100000000$ ,  $nnz=499960000$ , Read bandwidth = 12031 MB/s

**Task 3.5** What is the bandwidth you obtain by running the STREAM benchmark on your system? How does it compare to the bandwidth you measured in SpMV? Discuss the comparison.

| Function | Best Rate MB/s | Avg time | Min time | Max time |
|----------|----------------|----------|----------|----------|
| Copy:    | 28199.4        | 0.005750 | 0.005674 | 0.005853 |
| Scale:   | 17540.2        | 0.009188 | 0.009122 | 0.009401 |
| Add:     | 20904.5        | 0.011644 | 0.011481 | 0.011998 |
| Triad:   | 20714.7        | 0.011821 | 0.011586 | 0.012412 |

The STREAM benchmark shows a lot higher rates for all the operations compared to our read bandwidth measured on the SpMV benchmark. However, our calculated read bandwidth does not take into account computation time.

## Exercise 4: The memory mountain

**Task 4.1** Report the name of the processor and the size of the L1, L2, and L3 of the processor you are benchmarking. You can check the specs of your processor online.

Processor: i7-10750H

L1 = 384kB

L2 = 1.5MB

L3 = 12MB

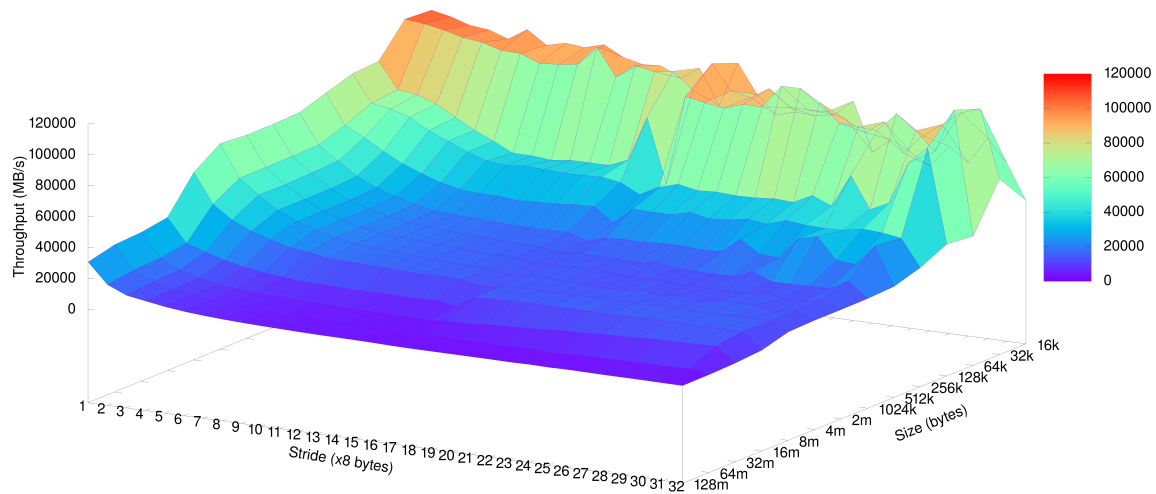


Figure 1: Memory Mountain

### Task 4.2

**Task 4.3** What region (array size, stride) gets the most consistently high performance (ignoring spikes in the graph that are noisy results...)? What is the read bandwidth reported?

Small arrays with low stride is the consistently most high performing region. The read bandwidth is up to 103469 MB/s.

**Task 4.4** What region (array size, stride) gets the most consistently low performance (Ignoring spikes in the graph that are noisy results...)? What is the read bandwidth reported?

High stride and large array size is the consistently most low performing region. The read bandwidth is 2907 MB/s at the lowest.

**Task 4.5** When you look at the graph for stride=1, you (should) see relatively high performance compared to stride=32. This is true even for large array sizes that are much larger than the L3 cache size. How is this possible, when the array cannot possibly all fit into the cache? Your explanation should include a brief overview of hardware prefetching as it applies to caches.

With a smaller stride there are fewer cache misses even though the whole array cannot be loaded. If the stride is smaller more of the array can be read before more data has to be loaded into the cache, even if the whole array cannot be loaded at once.

**Task 4.6 What is temporal locality? What is spatial locality?**

Temporally locality is the reuse of recently accessed data. Spatially locality is using data that is nearby in memory.

**Task 4.7 Adjusting the total array size impacts temporal locality, why? Will an increased array size increase or decrease temporal locality?**

Adjusting the array size impacts temporal locality as a smaller fraction of data is temporally local. Increasing the array size decreases temporal locality.

**Task 4.8 Adjusting the read stride impacts spatial locality, why? Will an increased read stride increase or decrease spatial locality?**

Read stride impacts spatial locality as memory further away that is not spatially local is read.

**Exercise 5: Write a benchmark to measure performance**

**Task 5.1.1** [On local machine]

What is the average runtime? - 0.00000405s

**Task 5.1.2** Increase N and compile the code, what is the average running time now? - 0.00000405s

**Task 5.2.1** Why is the execution time like that in the previous question when the flag -O2 is used? Answer this question using the information you find in the assembly code.

- Inspecting the asm code shows the The compiler optimized the loop out of the code.

**Task 5.2.2** What is the average execution time without the -O2 flag?

For N = 5000: - 0.00004816s

For N = 50000: - 0.00050402s

**Task 5.3** What is the clock granularity on Dardel or your local computer?

$$Granularity : 9.54 \cdot 10^{-7}s$$

**Task 5.4.1** Modify the program that you used in question 5.1 and do the following such that the code runs properly with -O2 optimization:

- The code runs as expected, inspecting the asm code shows the program executes the loop as intended.

**Task 5.4.2**

- Average Execution time: 0.000013399 s
- Min Execution time: 0.000133991 s



**Exercise 6: Measure the Performance of Matrix-Matrix Multiply and Transpose with perf****Task 6.1** We get the following results from perf:

|                           | MSIZE=64<br>Naive | MSIZE=64<br>Optimised | MSIZE=1000<br>Naive | MSIZE=1000<br>Optimised |
|---------------------------|-------------------|-----------------------|---------------------|-------------------------|
| Elapsed time (seconds)    | 0.006008407       | 0.026428618           | 9.820340947         | 4.068640160             |
| Instructions per cycle    | 3.05223157746     | 3.39699755395         | 2.33566861357       | 5.65940389993           |
| L1 cache miss ratio       | 0.156965536823    | 0.0152186701427       | 0.568518889684      | 0.0428330419037         |
| L1 cache miss rate<br>PTI | 45.6882622051     | 6.56849478495         | 162.654417351       | 18.374515994            |

Importantly, the average time per matrix multiplication is the following:

|                  | MSIZE=64 Naive | MSIZE=64 Opt | MSIZE=10 Naive | MSIZE=1000 Opt |
|------------------|----------------|--------------|----------------|----------------|
| Average time (s) | 0.000181       | 0.000154     | 0.889926       | 0.365820       |

**Task 6.2** Using the perf tool we measured the following performance:

|                        | N=64           | N=128          | N=2048          |
|------------------------|----------------|----------------|-----------------|
| Elapsed time (seconds) | 0.015456084    | 0.012847589    | 9.393376821     |
| Bandwidth/Rate (MB/s)  | 1.43e+04       | 6.90e+03       | 3.60e+02        |
| Instructions per cycle | 2.19010949691  | 1.62430143786  | 0.0975394373964 |
| L1 cache miss ratio    | 0.183398257284 | 0.327938976196 | 0.296058060658  |
| L1 cache miss rate PTI | 66.5605710444  | 147.558995856  | 477.708493074   |

With the reported base rate for each N being:

|                     | N=64     | N=128    | N=2048   |
|---------------------|----------|----------|----------|
| Base rate (seconds) | 2.29e-06 | 1.90e-05 | 9.32e-02 |